Touist reference manual

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Contents

1. Introduction

Touist is a language that allows to express propositional logic [1, 2]. You are provided with two programs: a graphical interface, refered as `touist.jar` (as it is written in Java) and `touistc`, the command-line compiler and solver (written in Ocaml).

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The touist language aims at making the writing of propositional logic as approchable as possible. Propositions and connectors syntaxes are close to what would be written on paper. Propositions are words that can contain numbers, the character "_". Here are some example of formulas:

Propositional logic	touist language
$\neg p$	not p
$p \wedge q$	p and q
$p \vee q$	p or q
$p \oplus q$	p xor q
$p \rightarrow q$	p => q
$p \leftrightarrow q$	p <=> q

After typing a formula, you can ask `touist` to find a valuation (true or false) of each proposition so that the whole formula is true (such valuation, also called *interpretation*, is called *model*). When a model exists, your formula is *satisfiable*. For example, a model of $p \vee q$ is $\{p = true, q = false\}$. To check the models of this formula using `touist`, you can do

Graphical Java interface	Command-line interface[1]
1. Type `p and q`	1. Create a file `p and q`
2. Press "Solve"	2. Type `./touistc -sat --solve yourfile`
3. Press "Next" to see other models	3. The first model is displayed

1.1. Simple inference checking

Checking the truth of a given reasoning (or *inference*) may be handy. From a wikipedia example:

Premise 1:	If it's raining then it's cloudy
Premise 2:	It's raining
Conclusion:	It's cloudy

This inference can be written

$$\{raining \rightarrow cloudy, raining\} \models cloudy$$

The *infer* symbol ($\models$) does not belong to the `touist` language (we call it "meta-language"). This means that we have to transform this notation to an actual propositional formula.

Theorem 1.
Let H be a set of formulas (*hypotheses* or *premises*) and C a formula (*conclusion*). Then
$$H \models C \text{ if and only if } H \cup \{\neg C\} \text{ is unsatisfiable.}$$

From this theorem, we just have to check that the set of formulas

$$\{raining \rightarrow cloudy, raining, \neg cloudy\}$$

has no model. We can translate this set to `touist` language (comments begin with two semi-colon ";;"):

```
raining => cloudy        ;; Premise 1
raining                  ;; Premise 2
not cloudy               ;; Conclusion
```

Note. In touist, the premises are simply formulas separated by a new line. A new line is semantically equivalent to the **and** connector: the previous bit of **touist** code could be equivalently written

```
(raining => cloudy) and raining and not cloudy
```

2. Language reference

2.1. Structure of a touist file

```
<touist-file> ::= <affect> <touist-file>
                | <formula> <touist-file>
                | EOF
```

A touist file is a whitespace-separated list of affectations and formulas. Affectations are global and cannot be done in nested formulas. They can be anywhere in the file (at the beginning, at the end or interlaced with formulas). A whitespace is a space, tab or newline.

Comments begin with the ";;" sequence.

2.2. Affectations

Affectations have the following syntax:

```
<affect> ::= <var> "=" (<int>|<float>|<bool>|<prop>|<set>)
```

This kind of affectation is global and applies for the whole code, even if the affectation happens after the formulas when the variable appears. These affectations are evaluated before any formula.

The order of affectation is important when you want to use a variable in an affectation expression. For example:

```
$N = 10
$set = [1..$N]    ;; $N must be defined before $set
```

2.3. Variables

```
TERM = [_0-9]*[a-zA-Z][a-zA-Z_0-9]*
<var> ::=
  | "$" TERM                              (* simple-var *)
  | "$" TERM "(" <comma-list(<float>|<int>|<prop>)> ")" (* tuple-var *)
```

Simple variables ("simple-var")
    Of the form `$var`
    In a formula, it is expected to contain a proposition

In an expression, it can contain an integer, a float, a prop or a set

Tuple variable ("tuple-var")

Of the form `$var($i,a,4)`

The leading variable (e.g., `$var`) must contain a proposition

It will always be expanded to a proposition. For example, if `$var=p` and `$i=q`, then it will expand to `p(q,a,4)`

The nested indices (e.g., `$i`) can be either a integer, float, proposition or boolean

Here are some examples of variables:

Simple-var	Tuple-var
`$N`	`$place($number)`
`$time`	`$action($i,$j)`
`$SIZE`	
`$is_over`	

2.4. Propositions

```
TERM = [_0-9]*[a-zA-Z][a-zA-Z_0-9]*
<prop> ::=
    | <var>
    | TERM
    | TERM "(" <comma-list(<float> | <int> | <prop>)> ")"
```

A simple proposition is a simple word that can contain numbers and the underscore symbol ("_"). A tuple proposition (we can it as a *predicate*), of the form `prop(1,$i,abc)`, must have indices of type integer, float, boolean or set.

Note. A tuple proposition that is in an expression and that contains at least one set in its indices will be expanded to a set of the cartesian product of the set indices. In the following table, the two right-columns show how the propositions are expanded whether they are in an expression or in a formula:

Proposition	is in a formula	is in an expression
`p([a])`	`p([a])`	`p(a)`
`p([a,b,c])`	`p([a,b,c])`	`[p(a),p(b),p(c)]`
`p([a,b],[1..2])`	`p([a,b],[1..2])`	`[p(a,1),p(b,1)`
		`p(a,2),p(b,2)]`

2.5. Numeric expression

The available operations on integers and floats are `+`, `-`, `*`, `/`, `mod` (modulo) and `abs()` (absolute value).

Here is the complete rules for numeric operators:

```
<num-operation(<T>)> ::=
    | <T> "+" <T>
    | <T> "-" <T>
```

```
        |     "-" <T>
        | <T> "*" <T>
        | <T> "/" <T>
<num-operation-others(<T>)> ::=
        | <T> "mod" <T>
        | "abs(" <T> ")"
```

Note. Integer and float expressions cannot be mixed. It is necessary to cast explicitly to the other type when the types are not matching. For example, the expression `1+2.0` is invalid and should be written `1+int(2.0)` (gives an integer) or `float(1)+2.0` (gives a float). Some operators are specific to integer or float types:

- `card([a,b])` returns an integer,
- `sqrt(3)` returns a float.

2.5.1. Integers

```
<int> ::=
        | "(" <int> ")"
        | <var>
        | INT
        | num-operation(<int>)
        | num-operation-others(<int>)
        | "if" <bool> "then" <int> "else" <int> "end"
        | "int(" (<int>|<float>) ")"
        | "card(" <set> ")"
```

2.5.2. Floats

```
<float> ::=
        | "(" <float> ")"
        | <var>
        | FLOAT
        | num-operation(<float>)
        | num-operation-others(<float>)
        | "if" <bool> "then" <float> "else" <float> "end"
        | "float(" (<int>|<float>)  ")"
        | "sqrt(" <float> ")"
```

2.6. Booleans

The constants are `true` and `false`. The boolean connectors are $>,<,\geq$ (`>=`),$\leq$ (`<=`),$=$ (`==`) and $\neq$ (`!=`). The operators that return a boolean are `subset(,)`, `empty()` and `in`.

```
<bool> ::= "(" <bool> ")"
    | <var>
    | "true"
    | "false"
    | (<int>|<float>|<prop>|<bool>) "in" <set>
    | "subset(" <set> "," <set> ")"
    | "empty(" <set> ")"
    | <equality(<int>|<float>|<prop>)>
    | <order(<int>|<float>)>
    | <connectors(<bool>)>
<equality(<T>)> ::=
    | <T> "!=" <T>
    | <T> "==" <T>
<order(<T>)> ::=
    | <T> ">" <T>
    | <T> "<" <T>
    | <T> "<=" <T>
    | <T> ">=" <T>
```

Note. Booleans cannot be mixed with formulas. In a formula, the evaluation (choosing true or false) is not done during the translation from touist to the "solver-friendly" language. Conversely, a boolean expression must be evaluable during the translation.

2.7. Sets

Sets can contain anything (propositions, integers, floats, booleans or even other sets) as long as all elements have the same type.

Along with the previously mentionned `card(x)`, `empty()`, `in` and `subset(,)`, common operations on sets are available: $P \cup Q$ (`inter($P,$Q)`), $P \cap Q$ (`inter($P,$Q)`), $P \setminus Q$ (`diff($P,$Q)`) and `powerset($P)`.

Here is a description of the non-obvious operators:

- `powerset($P)` computes the set of all possible subsets of the set P
- `subset($P,$Q)` returns `true` if $P \subset Q$
- `empty($P)` returns `true` if P is empty
- `card($P)` returns the cardinal of P

Here is the complete rule:

```
<set> ::= "(" <set> ")"
    | <var>
    | "[" <comma-list(<int>|<float>|<prop>|<bool>)> "]"
    | "[ <int> ".." <int> "]"
    | "union(" <set> "," <set> ")"
    | "inter(" <set> "," <set> ")"
    | "diff(" <set> "," <set> ")"
    | "powerset(" <set> ")"
```

2.8. Formulas

2.8.1. Connectors

A formula is a sequence of propositions (that can be variables) and connectors $\neg p$ (not), $\wedge$ (and), $\vee$ (or), $\oplus$ (xor), $\rightarrow$ (=>) or $\leftrightarrow$ (<=>).

```
<connectors(<T>)> ::=
    |       "not" <T>
    | <T> "and" <T>
    | <T> "or" <T>
    | <T> "xor" <T>
    | <T> "=>" <T>
    | <T> "<=>" <T>
```

2.8.2. Generalized connectors

Generalized connectors bigand, bigor, exact, atmost and atleast are also available for generalizing the formulas using sets. Here is the rule for these:

```
<generalized-connectors(<T>)> ::=
    | "bigand" <comma-list(<var>)> "in" <comma-list(<set>)>
                            ["when" <bool>] ":" <T> "end"
    | "bigor" <comma-list(<var>)> "in" <comma-list(<set>)>
                            ["when" <bool>] ":" <T> "end"
    | "exact(" <int> "," <set> ")"
    | "atmost(" <int> "," <set> ")"
    | "atleast(" <int> "," <set> ")"
```

Bigand and bigor   When multiple variables and sets are given, the bigand and bigor operators will produce the and/or sequence for each possible couple of value of each set (the set of couples is the cartesian product of the given sets). For example,

The formula	expands to...
$\bigwedge\limits_{\substack{i\in[1..2]\\j\in[a,b]}} p_{i,j}$	$p_{1,a} \wedge p_{1,b} \wedge p_{2,a} \wedge p_{2,b}$

`bigand $i,$j in [1..2],[a,b]:`
`p(i,j)`
`end`
	`p(1,a) and p(1,b)`
	`and p(2,a) and p(2,b)`

The when is optional and allows to apply a condition to each couple of valued variable.

On the following two examples, the math expression is given on the left and the matching touist code is given on the right:

$$\bigwedge_{\substack{i\in[1..n]\\ j\in[a,b,c]}} p_{i,j}$$

```
bigand $i,$j in [1..$n],[a,b,c]:
    p($i,$j)
end
```

$$\bigvee_{\substack{v\in[A,B,C]\\ x\in[1..9]\\ y\in[3..4]\\ x\neq y\\ x\neq A}} v_{x,y}$$

```
bigor $v,$x,$y
    in [A,B,C],[1..9],[3..4]
    when $v!=A and $x!=$y:
    $v($x)
end
```

Exact, atmost and atleast   The operator `exact(3,[a,b,c,d,e])` will produce the formula that ensures that, for any valuation, exactly 3 propositions can be `true` simultanously. The operator `atleast` ensures that at least N propositions are `true` simultanously, and `atmost` does the opposite.

Note. (TODO) explain the cases N=0 or empty set

2.8.3. Simple formulas

The constants $\top$ (`Top`) and $\bot$ (`Bot`) allows to express the "always true" and "always false". Here is the complete grammar:

```
<formula-simple> ::=
    | "Top"
    | "Bot"
    | <prop>
    | <var>
    | <formula(<formula-simple>)>

<formula(<T>)> ::=
    | "(" <T> ")"
    | "if" <bool> "then" <T> "else" <T> "end"
    | <connectors(<T>)>
    | <generalized-connectors(<T>)>
    | <let-affect(<T>)>
```

2.8.4. SMT formulas

Touist can also be given SMT formulas and output the SMT2-compliant file. (TODO)

2.8.5. Local variables

Sometimes, you want to use the same result in multiple places. You might not be able to use a global affectation (presented above) because you are in a nested formula. The `let` construct lets you create temporary variables inside formulas.

```
<let-affect<T>> ::=
    | "let" <var> "=" <int>|<float>|<bool>|<prop>|<set> ":" <formula<T>>
```

Example:

```
;; This piece of code has no actual meaning
$letters = [a,b,c,d,e]
bigand $letter,$number in $letters,[1..card($letters)]:
  has($letter,$number) =>
  let $without_letter = diff($letters,$letter): ;; keep temorary result
  bigand $l1 in $without_letter:
      p($letter)
  end
end
```

2.9. Formal grammar

This section presents the grammar formatted in a BNF-like way. Some rules ("::=") are parametrized so that some parts of the grammar are "factorized" (the idea of parametrized rules come from the Menhir parser generator used for generating the touist parser).
Note. This grammar specification is not LL(1) and could not be implemented as such using Menhir; most of the type checking is made after the abstract syntaxic tree is produced. The only purpose of the present specification is to give a clear view of what is possible and not possible with this language.

```
INT    = [0-9]+
FLOAT  = [0-9]+\.[0-9]+
TERM   = [_0-9]*[a-zA-Z][a-zA-Z_0-9]*
```

```
<touist-file> ::= <affect> <touist-file>
                | <formula> <touist-file>
                | EOF

<var> ::= "$" TERM
    | "$" TERM "(" <comma-list(<float> | <int> | <prop>)> ")"

<prop> ::=
    | <var>
    | TERM
    | TERM "(" <comma-list(<float> | <int> | <prop>)> ")"

<affect> ::= <var> "=" (<int>|<float>|<bool>|<prop>|<set>)
```

```
<let-affect<T>> ::=
    | "let" <var> "=" <int>|<float>|<bool>|<prop>|<set> ":" <formula<T>>

<equality(<T>)> ::=
    | <T> "!=" <T>
    | <T> "==" <T>

<order(<T>)> ::=
    | <T> ">" <T>
    | <T> "<" <T>
    | <T> "<=" <T>
    | <T> ">=" <T>

<bool> ::= "(" <bool> ")"
    | <var>
    | "true"
    | "false"
    | (<int>|<float>|<prop>|<bool>) "in" <set>
    | "subset(" <set> "," <set> ")"
    | "empty(" <set> ")"
    | <equality(<int>|<float>|<prop>)>
    | <order(<int>|<float>)>
    | <connectors(<bool>)>

<num-operation(<T>)> ::=
    | <T> "+" <T>
    | <T> "-" <T>
    |     "-" <T>
    | <T> "*" <T>
    | <T> "/" <T>

<num-operation-others(<T>)> ::=
    | <T> "mod" <T>
    | "abs(" <T> ")"

<int> ::=
    | "(" <int> ")"
    | <var>
    | INT
    | num-operation(<int>)
    | num-operation-others(<int>)
    | "if" <bool> "then" <int> "else" <int> "end"
    | "int(" (<int>|<float>) ")"
    | "card(" <set> ")"
```

```
<float> ::=
    | "(" <float> ")"
    | <var>
    | FLOAT
    | num-operation(<float>)
    | num-operation-others(<float>)
    | "if" <bool> "then" <float> "else" <float> "end"
    | "float(" (<int>|<float>)  ")"
    | "sqrt(" <float> ")"

<set> ::= "(" <set> ")"
    | <var>
    | "[" <comma-list(<int>|<float>|<prop>|<bool>)> "]"
    | "[ <int> ".." <int> "]"
    | "union(" <set> "," <set> ")"
    | "inter(" <set> "," <set> ")"
    | "diff(" <set> "," <set> ")"
    | "powerset(" <set> ")"

<comma-list(<T>)> ::= <T> | <T> "," <comma-list(<T>)>

<generalized-connectors(<T>)> ::=
    | "bigand" <comma-list(<var>)> "in" <comma-list(<set>)>
                            ["when" <bool>] ":" <T> "end"
    | "bigor" <comma-list(<var>)> "in" <comma-list(<set>)>
                            ["when" <bool>] ":" <T> "end"
    | "exact(" <int> "," <set> ")"
    | "atmost(" <int> "," <set> ")"
    | "atleast(" <int> "," <set> ")"

<connectors(<T>)> ::=
    |     "not" <T>
    | <T> "and" <T>
    | <T> "or" <T>
    | <T> "xor" <T>
    | <T> "=>" <T>
    | <T> "<=>" <T>

<formula(<T>)> ::=
    | "(" <T> ")"
    | "if" <bool> "then" <T> "else" <T> "end"
    | <connectors(<T>)>
    | <generalized-connectors(<T>)>
    | <let-affect(<T>)>

<formula-simple> ::=
```

```
    | "Top"
    | "Bot"
    | <prop>
    | <var>
    | <formula(<formula-simple>)>

<formula-smt> ::=
    | <formula(<formula-smt>)>
    | <expr-smt>

<expr-smt> ::=
    | "Top"
    | "Bot"
    | <prop>
    | <var>
    | <int>
    | <float>
    | <order>(<expr-smt>)
    | <num-operations_standard(<expr-smt>)>
    | <equality(<expr-smt>)>
    | <in_parenthesis(<expr-smt>)>
```

References

[1] Khaled Skander Ben Slimane, Alexis Comte, Olivier Gasquet, Abdelwahab Heba, Olivier Lezaud, Frédéric Maris, and Maël Valais. "La Logique Facile Avec TouIST (formalisez et Résolvez Facilement Des Problèmes Du Monde Réel)." In *Actes Des 9es Journées d'Intelligence Artificielle Fondamentale (IAF 2015)*. 2015. http://pfia2015.inria.fr/actes/download.php?conf=IAF&file=Ben_Slimane_IAF_2015.pdf.

[2] Khaled Skander Ben Slimane, Alexis Comte, Olivier Gasquet, Abdelwahab Heba, Olivier Lezaud, Frederic Maris, and Mael Valais. "Twist Your Logic with TouIST." *CoRR* abs/1507.03663. 2015. http://arxiv.org/abs/1507.03663.