A puzzle

Alice and Bob are trying to send each other a message without Eve the Eavesdropper being able to read it.

There are some positive integers in each column. **Only Alice** can access the numbers in her column. **Only Bob** can access the numbers in his column. **Anyone** can access the numbers in the public column.

Alice
a
a is between
1 and n .

Public	
$g \qquad n$	
g is a small prime number	
n is a very big number.	

 $\begin{array}{c} b \\ b \\ \text{ is between} \\ 1 \text{ and } n. \end{array}$

If Eve can read anything Alice and Bob send to each other, how can Alice and Bob both know a number without Eve knowing that number as well?





Cryptography

Vincent Macri



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- 2 Primes
- 3 Factoring
- 4 Introduction to Cryptography
- 5 RSA





Quick review Modular Arithmetic

We define the **mod** operator as being the remainder when dividing two numbers. That is:

$$a \bmod b =$$
the remainder of $a \div b$

In some programming languages, modulo is written as % or rem. Use whichever notation you are most comfortable with.

Examples

$$4 \mod 2 = 0$$

$$7 \mod 3 = 1$$

$$5 \mod 2 = 1$$

$$9 \mod 5 = 4$$

The definition of modulo (mod for short) is a bit trickier with negative numbers. It also doesn't matter for today, as we're only looking at mod with positive numbers.

Calculating modulo Modular Arithmetic

While we could do long division to find the remainder when we want to calculate modulo, I prefer to use this formula:

$$a \mod b = a - b \left\lfloor \frac{a}{b} \right\rfloor$$

Where $\lfloor x \rfloor$ is the floor function, which rounds a number **down** to an integer.

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Where $\lfloor x \rfloor$ is the floor function, which rounds a number **down** to an integer.

There are other method, but I think this one is the hardest to mess up. Use whatever method you are most comfortable with.

How do we do division in modular arithmetic?



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Division is the inverse of multiplication.



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Recall from the ${\bf group\ theory}$ lesson that the identity element in multiplication is 1.



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So, for modulus n, b is the inverse of a when:

$$a \times b \mod n = 1 \mid 0 < a, b < n$$



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$$3 \times 7 \mod 20 = 1$$

Here, 7 is the inverse of 3, and 3 is the inverse of 7.





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2 does **not** have an inverse modulo 4.

$$2 \times b \mod 4 = 1$$

There is no integer value for b that satisfies this equation.



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2 does **not** have an inverse modulo 4.

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We will soon learn the algorithm to find the inverse to modular multiplication.



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What is a prime number? Primes

A **prime number** is a positive integer that is only divisible by 1 and itself.

Examples

$$\mathbb{P} = \{2, 3, 5, 7, 11, 13, 17, \dots\}$$

If an integer greater than 1 is not prime, it is called a **composite** number.

1 is special, and is called the **unit number**

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Proof 1 is not prime.

In the past, some mathematicians said that 1 is prime. All of them are dead now.

$$\therefore 1 \notin \mathbb{P}$$



The largest known prime number¹ is:

$$M_{77232917} = 2^{77232917} - 1$$

If you were to print this number out, it would be 6055 pages long! This prime was discovered by Jonathan Pace on December 26, 2017 after 6 days of continuous computer computations. The discovery was published on January 3, 2018.



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What's special and useful about Mersenne primes?



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This number is a **Mersenne prime**. These are prime numbers of the form $2^n - 1$, and we label these primes as M_n for short.

What's special and useful about Mersenne primes? Not much.



How many primes are there?

Primes

Is the number of primes finite?



How many primes are there? Primes

Is the number of primes finite?

No! There are infinite prime numbers!

This was proved thousands of years ago by Euclid.



Proof of infinite primes Primes

Assume the list of primes is finite, and there are only n prime numbers. We will call our list of prime numbers P.

$$P = \{p_1, p_2, \dots, p_{n-1}, p_n\}$$

Where p_k is the kth prime number.

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Where p_k is the kth prime number.

Now, let m be the product of all numbers in P plus 1.

$$m = (p_1 \times p_2 \times \dots \times p_{n-1} \times p_n) + 1 = \left(\sum_{i=1}^n p_i\right) + 1$$

m is either prime or not prime. Let's look at both cases.





First, let's consider the case that m is prime.



Proof of infinite primes: m is prime $_{\mathrm{Primes}}$

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Note that m is not in our original list, P.



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Note that m is not in our original list, P.

If m is prime, our original list is incomplete, and there are more prime numbers than we listed.





If m is not prime, then it must be divisible by a prime number. Notice that m cannot be divisible by any numbers in P, as they would not divide a number that is a multiple of themselves plus 1.

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For example:

$$P = \{2, 3, 5, 7, 11, 13\}$$

$$m = 2 \times 3 \times 5 \times 7 \times 11 \times 13 + 1 = 30031$$

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For example:

$$P = \{2, 3, 5, 7, 11, 13\}$$
 $m = 2 \times 3 \times 5 \times 7 \times 11 \times 13 + 1 = 30\,031$
 $30\,031 \bmod 2 = 1$
 $30\,031 \bmod 3 = 1$
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Here, we can see that since $30\,031$ is a multiple plus 1 of every number in P, no numbers in P will divide it.





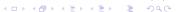
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Primality Primes

How do we check if a number is prime?



Primality Primes

How do we check if a number is prime?

How do we check if a number is prime quickly?



Primality Primes

How do we check if a number is prime?

How do we check if a number is prime quickly?

With a **very** fast computer. Algorithms exist (some of which run in polynomial time) but they are **very** slow.

Here, "quickly" means the computer will finish before we die.

This fits in the margins Primes

Theorem (Fermat's Little Theorem)

Let p be a prime number, and a an integer that does not divide p. Then:

$$a^{p-1} \bmod p = 1$$

This fits in the margins Primes

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Let p be a prime number, and a an integer that does not divide p. Then:

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Theorem (Euler-Fermat Generalization)

Fermat's Little Theorem can be generalized as:

$$a^{\phi(n)} \bmod n = 1$$

Where $\phi(n)$ is Euler's totient function, which gives us the number of integers less than or equal to n that are coprime to n.



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For an extra challenge, prove the **Euler-Fermat Generalization** using **Fermat's Little Theorem**.

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Divisibility Factoring

We will now introduce a new notation, which is more of a shortcut. If b divides a with no remainder, then we will write $b \mid a$. More formally:

$$b \mid a \equiv a \bmod b = 0$$

Or:

$$b \mid a \iff a = bc$$

Where a, b, and c and positive integers. If $b \mid a$, then b is a factor of a.



Theorem (The Unique Factorization Theorem)

Every positive integer has a unique representation as a product of prime numbers.

That is, for all numbers $n \in \mathbb{Z}^+$:

$$n = p_1^{a_1} \times p_2^{a_2} \times \dots \times p_k^{a_k}$$

Where p_i is prime, and a_i is a positive integer.

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Where p_i is prime, and a_i is a positive integer.

Example (180)

$$180 = 2^2 \times 3^2 \times 5$$
$$180 = 2 \times 2 \times 3 \times 3 \times 5$$



How do we factor a number?



How do we factor a number?

How do we factor a large number?





How do we factor a number?

How do we factor a large number?

Try this one:

RSA-2048

251959084756578934940271832400483985714292821262040320277771378360436620207075955562640185258807 844069182906412495150821892985591491761845028084891200728449926873928072877767359714183472702618 963750149718246911650776133798590957000973304597488084284017974291004624586918171951187461215151 726546322822168699875491824224336372590851418654620435767984233871847744479207399342365848238242 811981638150106748104516603773060562016196762561338441436038339044149526344321901146575444541784 240209246165157233507787077498171257724679629263863563732899121548314381678998850404453640235273 81951378636564391212010397122822120720357

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How do we factor a large number?

Try this one:

RSA-2048

261959084756578934940271832400483985714292821262040320277771378360436620207075955562640185258807 844069182906412495150821892985591491761845028084891200728449926873928072877767359714183472702618 963750149718246911650776133798590957000973304597488084284017974291006424586918171951187461215151 726546322822168669875491824224336372590851418654620435767984233871847744479207399342365848238242 811981638150106748104516603773060562016196762561338441436038339044149526344321901146575444541784 240209246165157233507787077498171257724679629263863563732899121548314381678998850404453640235273 81951378636564391212010397122822120720357

This number has two factors. Nobody knows what they are. There was a $\$200\,000$ prize to factor this number. People had over 15 years to factor it, but nobody was able to before the contest period ended.



What is the greatest common divisor? Factoring

The greatest common divisor (GCD) of two numbers is the largest number that divides both numbers.

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More formally:

$$x = \gcd(a, b) \mid a, b \in \mathbb{Z}$$

Where x is the largest number such that:

$$x \mid a \wedge x \mid b$$



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If two numbers are coprime, their gcd is 1.

How do we find the greatest common divisor? Factoring

How do we find the greatest common divisor?



How do we find the greatest common divisor? Factoring

How do we find the greatest common divisor?

We could list all the factors, and the biggest one would be the gcd. But, as we saw above, factoring is very hard. Is there a better way?



How do we find the greatest common divisor? Factoring

How do we find the greatest common divisor?

We could list all the factors, and the biggest one would be the \gcd . But, as we saw above, factoring is very hard. Is there a better way? Of course. Otherwise I wouldn't ask.



Thousands of years ago, Euclid came up with an algorithm to find the $\gcd. \label{eq:control}$

Euclidean algorithm

To find gcd(a, b), do the following:

- **1** Let $r_0 = a$, $r_1 = b$, and i = 1.
- 2 If $r_i = 0$ then $gcd(a, b) = r_i$.
- Write $r_{i-1} = q_i r_i + r_{i+1}$ and increment i by 1. Here, $r_{i+1} = r_i \mod r_{i-1}$.
- 4 Go back to step 2.



Extended Euclidean algorithm Factoring

The extended Euclidean algorithm is essentially the Euclidean algorithm in reverse.

We use substitution while working backwards.

It allows us to find two integers, \boldsymbol{x} and \boldsymbol{y} , that satisfy:

$$\gcd(a,b) = ax + by$$

Extended Euclidean algorithm Factoring

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It allows us to find two integers, x and y, that satisfy:

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When gcd(a,b) = 1, $x = a^{-1} \mod b$, where a^{-1} is the inverse of a in modulus b.

Extended Euclidean algorithm example Factoring

Find the inverse of 3 in modulus 26.

$$26 = (8)3 + 2$$

$$3 = (1)2 + 1$$

$$2 = (2)1 + 0$$

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And so the inverse of 3 in $\mod 26$ is 9. We can verify this:

$$(3 \times 9) \bmod 26 = 1$$



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What are factors used for? Introduction to Cryptography

Factorization of numbers is very useful in cryptography. The reason for this is that factoring large numbers takes a **very** long time, but the maths for checking factorization are quick. We can use this to develop a way to encode messages so they can only be read by certain people. This is called cryptography.



What is cryptography Introduction to Cryptography

Simply put, cryptography is the study of ways to encrypt messages.

Encryption is when you transform a message so that it cannot easily be read by someone without a key. Encryption is like a lock, but instead of locking your house, it locks information.

The use of encryption goes back thousands of years.

One example of encryption was used by Julius Caesar to keep military messages protected from spies.

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Shift every letter in the message that many letters to the right, wrapping around when you reach z.



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For example, with n=3:

Plaintext message: 'Crypto is fun!' Encrypted message: 'Fubswr lv ixq!'

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Obviously, this isn't very secure. Why?





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RSA encryption RSA

The RSA encryption algorithm was originally classified before being rediscovered by three people with the initials R, S, and A in 1977. RSA encryption or a variant of it is used today in many online systems, especially for setting up more permanent encrypted sessions.

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Today, we will learn how to encrypt messages using the RSA system.

While RSA has not been broken by anyone, there are systems that are considered to be **more** secure because they provide something called "perfect forward secrecy". However, a lot of these systems are similar to, or even use RSA.





RSA overview RSA

What is RSA?



RSA overview RSA

What is RSA?

■ Public keys



What is RSA?

- Public keys
- Private keys



What is RSA?

- Public keys
- Private keys
- Padlocks



What is RSA?

- Public keys
- Private keys
- Padlocks
- Factoring is hard



What is RSA?

- Public keys
- Private keys
- Padlocks
- Factoring is hard

Alice and Bob.



The first step of RSA encryption is to generate a public-private keypair.

I Alice starts by picking two random prime numbers, p and q, that are similar in size. The bigger the better.

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Alice's public keypair is (n, e).

Alice's private key is d.



Encrypting RSA

Bob wants to send a message to Alice. He has turned his message into a number m, such that m < n (if $m \ge n$, then Bob will split the message up into multiple short messages).

If m and n are not coprime, then one could easily factorize n, thus breaking the encryption.

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Bob then downloads Alice's public key from somewhere he trusts and calculates the ${\bf ciphertext},\ c$:

$$c = m^e \bmod n$$

And he sends the encrypted message, \boldsymbol{c} to Alice.





Decrypting RSA

To decrypt the message, Alice calculates calculates r, which is equal to m, Bob's message:

$$r = c^d \bmod n$$



Decrypting RSA

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How do we know that r = m?



Euler's theorem RSA

Theorem (Euler's theorem)

When a and n are coprime positive integers:

$$a^{\phi(n)} \bmod n = 1$$



Proof that r=m RSA

Remember that $(m^e)^d = m^{ed}$. So:

$$r = m^{ed} \bmod n$$

We know that $ed=1 \bmod \phi(n)$ because we chose e and d to have that property.

This means that an integer, q, exists such that:

$$ed = q\phi(n) + 1$$

