

Lambda Calculus

Now you can bring a computer to your tests!

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What is lambda calculus?

Introduction

- Created by Alonzo Church



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Introduction

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- A way of representing **pure** mathematical functions
- Can represent any computer program
- Equivalent to Turing machines



$x + 1$

Introduction

In math class, we would define a function that accepts an argument x and outputs $x + 1$ as so:

$$f(x)$$



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$$f(4) = 4 + 1 = 5$$



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You can think of λ as f , and $.$ as $=$.



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Currying

One Argument

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And we use it like this:

$$(\lambda x. \lambda y. x + y)(2 \ 3)$$



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Why one argument?

One Argument

- Very simple



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- Very powerful



Why one argument?

One Argument

- Very simple
- Very powerful
- Functions can only have one variable



But I'm lazy

One Argument

Me too!



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We have some shortcuts to help us write down lambda calculus expressions, but it's important to remember what they represent, without the shortcuts.



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Can be abbreviated as:

$$\lambda xyz. A$$



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$$\lambda x. \lambda y. \lambda z. A$$

Can be abbreviated as:

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Also, we assume that we evaluate a function with “multiple” arguments starting with the leftmost parameter.



This is stupid. It just makes everything harder.

One Argument

you're right!



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For those examples, you're right!



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One Argument

For those examples, you're right!
Let's get to the fun stuff now!



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Boolean logic

Booleans

Quote

“Any program can be written in lambda calculus.”



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— Me, 5 minutes ago



Boolean logic

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So, let's bring on the Booleans!



TRUE and FALSE

Booleans

Definition (TRUE)

$$\text{TRUE} = \lambda xy.x$$



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$$\text{TRUE} = \lambda xy.x$$

Definition (FALSE)

$$\text{FALSE} = \lambda xy.y$$



TRUE and FALSE

Booleans

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$$\text{TRUE} = \lambda xy.x$$

Definition (FALSE)

$$\text{FALSE} = \lambda xy.y$$

So TRUE returns the first value, and FALSE returns the second.



TRUE and FALSE

Booleans

Definition (TRUE)

$$\text{TRUE} = \lambda xy.x$$

Definition (FALSE)

$$\text{FALSE} = \lambda xy.y$$

So TRUE returns the first value, and FALSE returns the second.
We will use TRUE and FALSE as shorthand for these definitions.



NOT

Booleans

Definition (NOT)

$$\text{NOT} = \lambda b.b(\text{FALSE } \text{TRUE})$$



NOT

Booleans

Definition (NOT)

$$\text{NOT} = \lambda b.b(\text{FALSE } \text{TRUE})$$

NOT TRUE

$$(\lambda b.b(\text{FALSE } \text{TRUE})) \text{TRUE} =$$



NOT

Booleans

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$$\text{NOT} = \lambda b.b(\text{FALSE } \text{TRUE})$$

NOT TRUE

$$(\lambda b.b(\text{FALSE } \text{TRUE})) \text{TRUE} = \text{TRUE}(\text{FALSE } \text{TRUE})$$



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Definition (NOT)

$$\text{NOT} = \lambda b.b(\text{FALSE } \text{TRUE})$$

NOT TRUE

$$\begin{aligned} (\lambda b.b(\text{FALSE } \text{TRUE})) \text{TRUE} &= \text{TRUE}(\text{FALSE } \text{TRUE}) \\ &= \lambda xy.x(\text{FALSE } \text{TRUE}) \end{aligned}$$



NOT

Booleans

Definition (NOT)

$$\text{NOT} = \lambda b.b(\text{FALSE } \text{TRUE})$$

NOT TRUE

$$\begin{aligned}(\lambda b.b(\text{FALSE } \text{TRUE})) \text{ TRUE} &= \text{TRUE}(\text{FALSE } \text{TRUE}) \\&= \lambda xy.x(\text{FALSE } \text{TRUE}) \\&= \text{FALSE}\end{aligned}$$



AND

Booleans

Definition (AND)

$$\text{AND} = (\lambda pq.p)(q \text{ } p)$$



AND

Booleans

Definition (AND)

$$\text{AND} = (\lambda pq.p)(q\ p)$$

AND(TRUE FALSE)

$$((\lambda pq.p)(q\ p))\ (\text{TRUE}\ \text{FALSE})$$



Definition (AND)

$$\text{AND} = (\lambda pq.p)(q \text{ } p)$$

AND(TRUE FALSE)

$$\begin{aligned} & ((\lambda pq.p)(q \text{ } p)) (\text{TRUE FALSE}) \\ &= \text{TRUE}(\text{FALSE TRUE}) \end{aligned}$$



Definition (AND)

$$\text{AND} = (\lambda pq.p)(q \text{ } p)$$

AND(TRUE FALSE)

$$\begin{aligned} & ((\lambda pq.p)(q \text{ } p)) (\text{TRUE FALSE}) \\ &= \text{TRUE}(\text{FALSE TRUE}) \\ &= \text{FALSE} \end{aligned}$$



IFTHENELSE

Booleans

Definition (AND)

$$\text{IFTHENELSE} = (\lambda b t f. b)(t \ f)$$

Where b is a Boolean, t is the value to return if $b = \text{TRUE}$, and f is the value to return if $b = \text{FALSE}$.

$\text{IFTHENELSE}(\text{TRUE} \text{ "Math is great" "Math is kool"})$

$((\lambda b t f. b)(t \ f)) (\text{TRUE} \text{ "Math is kool" "Math is lit"})$

