## Numerical Methods Bootcamp

Lecture 3
Solving nonlinear rational expectations models

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- In yesterday's lecture I showed you how a more "advanced" method gave more accurate results using N = 5 grid points than the discretised method using N = 1,000 grid points.
- I also told you that you needed to know about functional approximations and nonlinear equations solvers in order to know how to do that yourself.
- So now that you know, how do you do it?



$$V(k) = \max_{k'} \{ u(f(k) + (1 - \delta) - k') + \beta V(k') \}, \quad \text{for } k \in \mathcal{K}.$$

- Notice that k' is not restricted to be in  $\mathcal{K}$ .
- One way to solve:
  - Guess for  $\hat{V}_0(k)$  for all  $k \in \mathcal{K}$ .
  - Use some approximation method to evaluate V<sub>0</sub>(k') for k ∉ K
  - Use some optimisation routine (e.g. fminunc in matlab) to solve the maximisation problem.
  - Rinse and repeat.

$$V(k) = \max_{k'} \{ u(f(k) + (1 - \delta) - k') + \beta V(k') \}, \quad \text{for } k \in \mathcal{K}.$$

- ▶ But for our problem, and for many other problems, we know that V(k) ought to be
  - 1. Concave
  - 2. Differentiable
- So first order conditions are necessary and sufficient.
- Optimality satisfies

$$u'(f(k) + (1 - \delta)k - k') = \beta V'(k'), \text{ for } k \in \mathcal{K}.$$



In addition, the envelope theorem tells us that

$$V'(k) = (1 + f'(k) - \delta)u'(c)$$

- So what I did was starting with a guess  $k' = g_0(k)$  for  $k \in \mathcal{K}$ .
- Found  $V_0'(k)$  using the envelope theorem

$$V_0'(k) = (1 + f'(k) - \delta)u'(f(k) + (1 - \delta) - g_0(k)), \text{ for } k \in \mathcal{K}.$$

- Used a linear approximation to create the function  $\hat{V}'_0(k)$   $\forall k$
- Found k' as the solution to

$$u'(f(k)+(1-\delta)k-k')=\beta\hat{V}'_0(k'), \quad \text{for } k\in\mathcal{K}.$$

- ▶ And updated to  $k' = g_1(k)$  for  $k \in \mathcal{K}$ .
- Then I repeated this until

$$||u'(f(k)+(1-\delta)k-g_n(k))-\beta \hat{V}'_n(g_n(k))||<\varepsilon, \quad \text{for } k\in\mathcal{K}.$$



- Notice that I never updated or computed the value function itself
  - But only its slope
- And slope information is all that I need to solve for the solution to the nonlinear equation
- And the solution to the nonlinear equation is all I need to update the slope of the value function!

- If I instead would have approximated the value function, I would have
  - 1. Approximated level information
  - Use the derivative of this level information (which could be wildly inaccurate) to find the policy function
  - 3. Used the policy function to update level information.
- That's an unnecessary and inaccurate procedure (for differentiable problems)

- Yet the contraction property still kicks in, and I will find k' = g(k)
- And if I would need the value function, I could simply iterate on

$$V_{n+1}(k) = u(f(k) + (1-\delta)k - g(k)) + \beta V_n(g(k)), \quad \text{for } k \in \mathcal{K}$$
 with  $V_0(k) = 0$ .

Until

$$||V_{n+1} - V_n|| < \varepsilon$$



Why does the contraction property kick in?

▶ We know that the sequence  $\{V_n\}$  defined as

$$V_{n+1}(k) = \max_{k'} \{ u(f(k) + (1 - \delta) - k') + \beta V_n(k') \}$$

converges to V.

• We also know that each associated policy function  $g_{n+1}(k)$  such that

$$V_{n+1}(k) = u(f(k) + (1-\delta) - \mathbf{g}(k)) + \beta V_n(\mathbf{g}(k)),$$

converges to g.

these should have subscript n+1



Why does the contraction property kick in?

We also know that

$$u'(f(k) + (1 - \delta) - k') - \beta V'_n(k') = 0,$$

is necessary and sufficient for an optimum

Turns out that the envelope theorem also says that

$$V'_n(k) = (1 + f'(k) - \delta)u'(f(k) + (1 - \delta)k - g_n(k))$$



#### Why does the contraction property kick in?

▶ So the sequence {g<sub>n</sub>} defined as

$$u'(f(k) + (1 - \delta) - g_{n+1}(k)) - \beta(1 + f'(g_{n+1}(k)) - \delta)$$
  
  $\times u'(f(g_{n+1}(k)) + (1 - \delta)g_{n+1}(k) - g_n(g_{n+1}(k))) = 0,$ 

converges to g(k).

So we can also just iterate on the Euler equation



Why does the contraction property kick in?

- ▶ That is, suppose you have  $\hat{g}_n(k)$ .
- ▶ For each  $k \in \mathcal{K}$  find k' as

$$u'(f(k) + (1 - \delta) - k') - \beta(1 + f'(k') - \delta)$$
  
  $\times u'(f(k') + (1 - \delta)k' - \hat{g}_n(k')) = 0,$ 

and use the k' values to construct  $\hat{g}_{n+1}(k)$  using some functional approximation.

Why does the contraction property kick in?

- ▶ This is basically what I did.
- ▶ But I started with a guess  $\hat{V}'_n(k)$ .
- ▶ For each  $k \in \mathcal{K}$  I found k' as

$$u'(f(k) + (1 - \delta) - k') - \beta \hat{V}'_n(k') = 0,$$

and use the k' values to construct  $\hat{V}'_{n+1}(k)$  using

$$V'_{n+1}(k) = (1 + f'(k') - \delta)u'(f(k') + (1 - \delta)k' - \hat{g}_n(k')) = 0$$

Same same but different.

- One last thing
- How do I solve

$$u'(f(k)+(1-\delta)k-k')=\beta \hat{V}'_0(k'), \text{ for } k \in \mathcal{K}?$$

- Newton's method!
- Set the problem up as

$$-u'(f(k)+(1-\delta)k-k')+\beta\hat{V}'_n(k')=0, \quad \text{for } k\in\mathcal{K}$$

• We can call this  $F_n(k, k')$ .





And notice that

$$\frac{\partial F(k,k')}{\partial k'} = u''(f(k) + (1-\delta)k - k') + \beta \frac{\partial \hat{V}_n(k')}{\partial k'}$$

- We have an initial guess for F(k, k'), namely  $k' = g_n^0(k)$ .
- ▶ Using this initial guess  $k' = g_n^0(k)$ ,  $k \in \mathcal{K}$ , the last derivative can be approximated as

$$rac{\partial \hat{V}_n'(k')}{\partial k'} pprox rac{ ext{gradient}(\hat{V}_n'(g_n(k)))}{ ext{gradient}(g_n(k))}$$



So define

$$dE(k) = u''(f(k) + (1 - \delta)k - g_n^0(k)) +$$
 
$$\beta \frac{\operatorname{gradient}(\hat{V}_n'(g_n^0(k)))}{\operatorname{gradient}(g_n^0(k))}, \quad \text{for } k \in \mathcal{K}$$

And

$$E(k) = -u'(f(k) + (1 - \delta)k - k') + \beta \hat{V}'_n(k') = 0, \quad \text{for } k \in \mathcal{K}$$

And update

$$g_n^1(k) = g_n^0(k) - \frac{E(k)}{dE(k)}$$

Until (weak) convergence.



That is until

$$||g_n^i - g_n^{i-1}|| < 1e(-2)$$

or so (or perhaps just one iteration!)

- Doing this step at a tighter tolerance level is often unnecessary and takes up (a lot) of computer time.
- Notice that the real constraint is rather on

$$||g_n-g_{n-1}||<\varepsilon$$

▶ Which should be tight,  $\varepsilon = 1e(-6)$  is a good rule of thumb.



 In general, our functional equations boil down to something like

$$E_{z'}[f(x,x',x'',z,z')=0,\quad \text{all }x\in\mathcal{X} \text{ and }z\in\mathcal{Z}$$

And we would like to find a function x' = g(x, z) such that

$$E_{z'}[f(x,g(x,z),g(g(x,z),z'),z,z')=0$$

- There are several approaches to this
- Recall that what we will find is only an approximation  $\hat{g}(\cdot) \approx g(\cdot)$ .
- ► This approximation is, for a vector of X and Z, given by a vector of x'.
- So in principle we can just use a nonlinear equation solver to find x' for all  $x \in X$  and all  $z \in Z$  such that

$$E_{z'}[f(x,x',\hat{g}(x',z'),z,z')=0,\quad \text{all }x\in\mathcal{X} \text{ and }z\in\mathcal{Z}$$

▶ Where  $\hat{g}(\cdot)$  is consistent with x'.



- This is a "direct approach", and is normally very fast
- But it's only fast when it works.
- And my experience is that it doesn't work too often.
- Big systems of nonlinear equations are difficult to solve.
- Not very robust.
- Use only if speed is crucial.

#### Fixed point iteration

- ▶ Suppose we have a candidate policy function,  $g_n(k, z)$
- Sometimes we can rewrite the problem

$$E_{z'}[f(x,x',x'',z,z')]=0,\quad \text{all }x\in\mathcal{X} \text{ and }z\in\mathcal{Z}$$

as

$$E_{z'}[F(x,x',x'',z,z')]-x'=0$$
, all  $x\in\mathcal{X}$  and  $z\in\mathcal{Z}$ 

▶ Insert  $g_n(x, z)$  to get

$$E_{z'}[F(x,x',g_n(x,z),z,z')]-x'=0$$
, all  $x\in\mathcal{X}$  and  $z\in\mathcal{Z}$ 



#### Fixed point iteration

Problem again

$$E_{z'}[F(x,x',g_n(x',z'),z,z')]-x'=0,\quad ext{all }x\in\mathcal{X} ext{ and }z\in\mathcal{Z}$$

Fixed point iteration would then suggest that we find x' as

$$x' = E_{z'}[F(x, g_n(x, z), g_n(g_n(x, z), z'), z, z')],$$
  
all  $x \in X$  and  $z \in Z$ 

Until

$$||g_n(x,z)-E_{z'}[F(x,g_n(x,z),g_n(g_n(x,z),z'),z,z')]||<\varepsilon$$



#### Fixed point iteration

- Fixed point iteration is normally very fast
  - No derivatives to compute, just evaluate a function
- But it can have some serious convergence problems
  - A good initial guess is very helpful.
- It can be useful to first find

$$\hat{x}' = E_{z'}[F(x, g_n(x, z), g_n(g_n(x, z), z'), z, z')],$$
  
all  $x \in X$  and  $z \in Z$ 

And find x' as

$$x' = \rho \hat{x'} + (1 - \rho) g_n(x, z)$$
, all  $x \in \mathcal{X}$  and  $z \in \mathcal{Z}$ 

for a low value of  $\rho$ .



#### Time iteration

- ▶ Suppose we have a candidate policy function,  $g_n(k, z)$
- Our problem is

$$E_{z'}[f(x,x',x'',z,z')=0,\quad \text{all }x\in\mathcal{X} \text{ and }z\in\mathcal{Z}$$

▶ Insert  $g_n(k, z)$ 

$$E_{z'}[f(x,x',g_n(x',z'),z,z')=0,\quad \text{all }x\in\mathcal{X} \text{ and }z\in\mathcal{Z}$$

- ▶ Use a nonlinear equation solver to find x', and update to  $g_{n+1}(x, z)$ .
- Iterate until

$$||E_{z'}[f(x,g_n(x,z),g_n(g_n(x,z),z'),z,z')|| < \varepsilon$$



#### Time iteration

- This is actually what I did in the beginning of this lecture
- It's my favourite method as it is a contraction mapping!
  - So convergence is at least theoretically guaranteed
- But it's slower than fixed point iteration.

#### Getting the best of both worlds

Use time iteration until

$$||E_{z'}[f(x,g_n(x,z),g_n(g_n(x,z),z'),z,z')||<\varepsilon$$

for a large value of  $\varepsilon$ .

- Then switch to fixed point iteration.
- Possibly use the solution to a linearized version of your problem as an initial guess (if such a solution exist).

#### Improving on time iteration

- In the univariate case there exist an improvement technique for time iteration known as the method of endogenous gridpoints.
- It works like a charm
  - As fast as fixed point iteration
  - With sustained contraction property
- Only works for the univariate case, but still.



#### Method of Endogenous Gridpoints

- The idea is the following:
- It is impossible to get a closed form solution for x' from

$$E_{z'}[f(x, x', g_n(x', z'), z, z') = 0, \text{ all } x \in \mathcal{X} \text{ and } z \in \mathcal{Z}$$

- So for time iteration we used a nonlinear equation solver to find this.
- ► However, given a value of x', it is usually trivial to find a close for solution for x!

#### Method of Endogenous Gridpoints

- ▶ So create a grid for x' instead, X'.
- And solve for x from

$$E_{z'}[f(x,x',g_n(x',z'),z,z')=0,\quad \text{all }x'\in\mathcal{X}' \text{ and }z\in\mathcal{Z}$$

Then use your vector of x' together with your solutions x to update  $g_{n+1}(x,z)!$ 

#### Method of Endogenous Gridpoints

For instance, the stochastic growth model is given by

$$u'(m - k') = \beta E_{z'}[(1 + z'f(k') - \delta)u'(m' - g_n(m', z'))]$$
with  $m = zf(k) + (1 - \delta)k$ 

$$m' = z'f(k') + (1 - \delta)k'$$

Notice that m is as a legitimate state variable as k.



#### Method of Endogenous Gridpoints

$$u'(m-k') = \beta E_{z'}[(1+z'f(k')-\delta)u'(m'-g_n(m',z'))]$$

▶ So given a grid for k' and z, calculate m', and find m as

$$m = u'^{-1} \{ \beta E_{z'}[(1 + z'f(k') - \delta)u'(m' - g_n(m', z'))] \} + k'$$

▶ Then update  $g_{n+1}(m, z)$  until convergence.

## Occasionally binding constraints

- Many interesting economic problem involves various sorts of inequality constraints that "occasionally bind"
  - Borrowing constraints
  - Irreversibility constraints
  - Collateral constraints
  - Implementability constraints
- Linearisation techniques are (almost) useless to solve models with these types of constraints
- It is generally perceived as hard to solve such problems
- I disagree

# Occasionally binding constraints

#### Let's take two examples

- 1. A borrowing constraint
- Irreversible investment

# A borrowing constraint

Consider the following optimisation problem

$$egin{aligned} V(b_0,s_0) \max_{\{c_t(s^t),b_{t+1}(s^t)\}_{t=0}^\infty} \sum_{t=0}^\infty \sum_{s^t \in \mathcal{S}^{t+1}} eta^t u(c_t(s^t)) P(s^t,s_0) \ & ext{subject to} \quad c_t(s^t) + b_{t+1}(s^t) = s_t w + (1-s_t) \mu w + (1+r) b_t(s^t), \ & b_{t+1}(s^t) \geq \underline{b} \ & \forall t, orall s^t \in \mathcal{S}^{t+1} \quad b_0, s_0 ext{ are given} \end{aligned}$$

#### Bellman equation

$$v(b,s) = \max_{b' \ge \underline{b}} \{u(b(1+r) + w(s) - b') + \beta \sum_{s'=0}^{1} v(b',s')p(s',s)\}$$

# A borrowing constraint

$$V(b,s) = \max_{b' \ge \underline{b}} \{u(b(1+r) + w(s) - b') + \beta \sum_{s'=0}^{1} V(b',s')p(s',s)\}$$

First order condition

$$u'(b(1+r)+w(s)-b')-\mu(b,s)=\beta\sum_{s'=0}^{1}V_{b'}(b',s')p(s',s)$$

where  $\mu(b, s)$  is the Lagrange multiplier on the borrowing constraint.



# A borrowing constraint

Using the envelope theorem

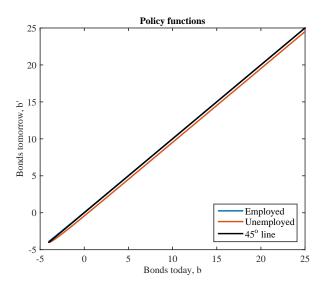
$$u'(b(1+r)+w(s)-b')-\mu(b,s) = \ eta(1+r)\sum_{s'=0}^1 u'(b'(1+r)+w(s')-g(b',s')) p(s',s)$$

▶ Suppose we have a guess  $g_n(b, s)$ . Then find  $\tilde{b}'$  as

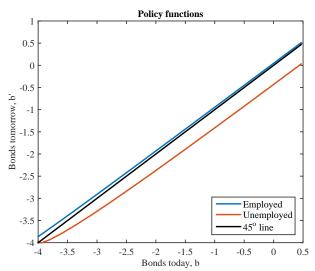
$$u'(b(1+r)+w(s)-\tilde{b}') = \ \beta(1+r)\sum_{s'=0}^{1}u'(\tilde{b}'(1+r)+w(s')-g_n(\tilde{b}',s'))p(s',s)$$

- Construct  $\tilde{g}_{n+1}(b,s)$  using  $\tilde{b}'$ .
- Update  $g_{n+1}(b, s) = \max{\{\tilde{g}_{n+1}(b, s), b\}}$

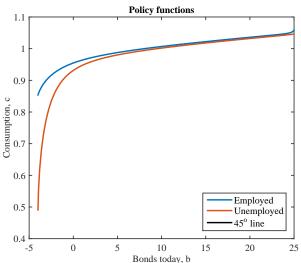




#### zoom in to see the non-linearities



#### Really affects consumption



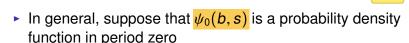
# Long run distribution

- When we analysed the problem solved by discretised value function iteration we could represent our policy function as a transition matrix
- From this we could calculate the transition matrix for the entire economy
- Find the long run distribution of capital
- And therefore calculate things like the average capital stock.

# Long run distribution

- In this problem, however, the policy function is of the type g(b, s).
- This cannot be represented as a transition matrix.
- So how can we find the long run distribution of bonds (and employment status) in this case?

# Long run distribution



Then

$$\psi_1(b',s') = \sum_{s \in \mathcal{S}} \sum_{\{b:b'=g(b,s)\}} \psi_0(b,s) p(s',s)$$

And in general

$$\psi_{t+1}(b',s') = \sum_{s \in \mathcal{S}} \sum_{\{b:b'=g(b,s)\}} \psi_t(b,s) p(s',s)$$



# Complications

 $\blacktriangleright$  A stationary cross-sectional distribution,  $\psi$ , is such that

$$\psi(b',s') = \sum_{s \in \mathcal{S}} \sum_{\{b:b'=q(b,s)\}} \psi(b,s) p(s',s)$$



This is very tricky to compute. The best practice I am aware of is to convert  $g(a, \theta)$  into a transition matrix!

#### Consider the following policy function

$$\begin{pmatrix}
1 \\
2 \\
3 \\
4 \\
5
\end{pmatrix}
\rightarrow
\begin{pmatrix}
1.8 \\
2.4 \\
3 \\
3.6 \\
4.2
\end{pmatrix}$$

#### Nearest neighbor interpolation

$$\begin{pmatrix} 1 \\ 2 \\ 3 \\ 4 \\ 5 \end{pmatrix} \rightarrow \begin{pmatrix} 2 \\ 2 \\ 3 \\ 4 \\ 4 \end{pmatrix}$$

#### Can be written as transition matrix

$$\begin{pmatrix}
0 & 1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 & 0 \\
0 & 0 & 1 & 0 & 0 \\
0 & 0 & 0 & 1 & 0 \\
0 & 0 & 0 & 1 & 0
\end{pmatrix}$$

But we can be a bit smarter. Policy function

$$\begin{pmatrix}
1 \\
2 \\
3 \\
4 \\
5
\end{pmatrix}
\rightarrow
\begin{pmatrix}
1.8 \\
2.4 \\
3 \\
3.6 \\
4.2
\end{pmatrix}$$



$$\begin{pmatrix} 0.2 & 0.8 & 0 & 0 & 0 \\ 0 & 0.6 & 0.4 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0.4 & 0.6 & 0 \\ 0 & 0 & 0 & 0.8 & 0.2 \end{pmatrix}$$



But in the income fluctuation problem (and many others) we normally have two (or many) policy functions

$$\begin{pmatrix} 1 \\ 2 \\ 3 \\ 4 \\ 5 \end{pmatrix} \xrightarrow{\text{if good state}} \begin{pmatrix} 2.2 \\ 2.8 \\ 3.4 \\ 4 \\ 4.6 \end{pmatrix}, \quad \text{and} \quad \begin{pmatrix} 1 \\ 2 \\ 3 \\ 4 \\ 5 \end{pmatrix} \xrightarrow{\text{if bad state}} \begin{pmatrix} 1.4 \\ 2 \\ 2.6 \\ 3.2 \\ 3.8 \end{pmatrix}$$

With some transition matrix for good and bad states

$$T = \begin{pmatrix} 0.8 & 0.2 \\ 0.3 & 0.7 \end{pmatrix}$$

Nearest neighbor interpolation

$$\begin{pmatrix} 1 \\ 2 \\ 3 \\ 4 \\ 5 \end{pmatrix} \xrightarrow{\text{if good state}} \begin{pmatrix} 2 \\ 3 \\ 3 \\ 4 \\ 5 \end{pmatrix}, \quad \text{and} \quad \begin{pmatrix} 1 \\ 2 \\ 3 \\ 4 \\ 5 \end{pmatrix} \xrightarrow{\text{if bad state}} \begin{pmatrix} 1 \\ 2 \\ 3 \\ 3 \\ 4 \end{pmatrix}$$

With some transition matrix for good and bad states

$$T = \begin{pmatrix} 0.8 & 0.2 \\ 0.3 & 0.7 \end{pmatrix}$$

Two transition matrices

$$M_g = \begin{pmatrix} 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{pmatrix}, \quad \text{and} \quad M_b = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \end{pmatrix}$$

With some transition matrix for good and bad states

$$T = \begin{pmatrix} 0.8 & 0.2 \\ 0.3 & 0.7 \end{pmatrix}$$

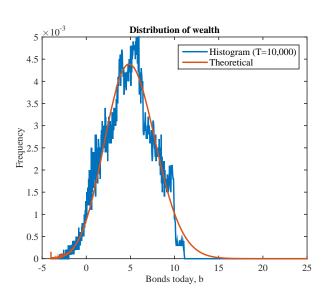
Full transition matrix is given by

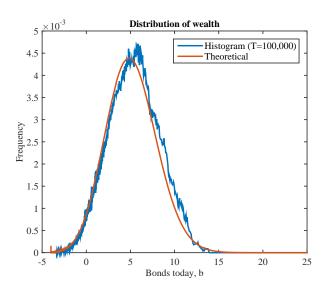
$$\begin{pmatrix} T(1,1) \cdot M_g & T(1,2) \cdot M_g \\ T(2,1) \cdot M_b & T(2,2) \cdot M_b \end{pmatrix}$$

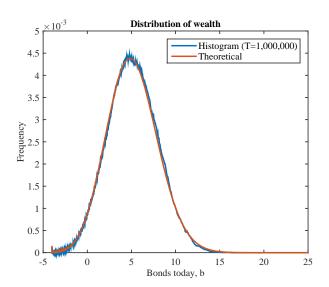
#### Full transition matrix is given by

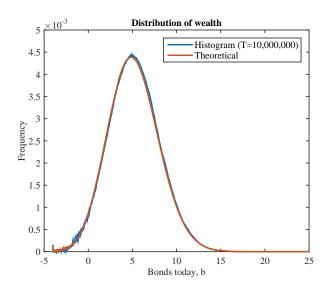
Doing the same thing with the smarter method gives

( 0	0.64	0.16	0	0	0	0.16	0.04	0	0 )
0	0.16	0.64	0	0	0	0.04	0.16	0	0
0	0	0.48	0.32	0	0	0	0.12	0.08	0
0	0	0	8.0	0	0	0	0	0.2	0
0	0	0	0.32	0.48	0	0	0	80.0	0.12
0.1	8 0.12	2 0	0	0	0.42	0.28	0	0	0
0	0.3	0	0	0	0	0.7	0	0	0
0	0.12	0.18	0	0	0	0.28	0.42	0	0
0	0	0.24	0.06	0	0	0	0.56	0.14	0
0	0	0.06	0.24	0	0	0	0.14	0.56	0 )









Consider the following optimisation problem

$$\begin{aligned} V(k_0, z_0) \max_{\{c_t(z^t), k_{t+1}(z^t)\}_{t=0}^{\infty}} \sum_{t=0}^{\infty} \sum_{z^t \in \mathcal{Z}^{t+1}} \beta^t u(c_t(z^t)) P(z^t, z_0) \\ \text{subject to} \quad c_t(z^t) + k_{t+1}(z^t) &= z_t f(k_t(z_{t-1})) + (1 - \delta) k_t(z^{t-1}), \\ k_{t+1}(z^t) &\geq (1 - \delta) k_t(z^{t-1}) \\ \forall t, \forall z^t \in \mathcal{Z}^{t+1} \quad k_0, z_0 \text{ are given} \end{aligned}$$

#### Bellman equation

$$v(k,z) = \max_{k' \geq (1-\delta)k} \{u(zf(k) + (1-\delta)k - k') + \beta \sum_{z' \in \mathcal{Z}} v(k',z')p(z',z)\}$$



$$v(k,z) = \max_{k' \ge (1-\delta)k} \{ u(zf(k) + (1-\delta)k - k') + \beta \sum_{z' \in \mathcal{Z}} v(k',z')p(z',z) \}$$

First order condition

$$U'(zf(k)+(1-\delta)k-k')-\mu(k,z)=\beta\sum_{z'\in\mathcal{Z}}V_{k'}(k',z')p(z',z)$$

where  $\mu(k, z)$  is the Lagrange multiplier on the irreversibility constraint.

The envelope theorem is slightly different here

$$V_k(k,z) = (1 + zf'(k) - \delta)u'(zf(k) + (1 - \delta)k - g(k,z)) - \mu(k,z)(1 - \delta)$$

So our first order condition is

$$u'(zf(k) + (1 - \delta)k - k') - \mu(k, z)$$

$$= \beta \sum_{z' \in \mathcal{Z}} [(1 + z'f'(k') - \delta)u'(z'f(k') + (1 - \delta)k' - g(k', z'))]$$

$$- \mu(k', z')(1 - \delta)[p(z', z)]$$





▶ Suppose we have guesses  $g_n(k, z)$  and  $\mu_n(k, z)$ . Then find  $\tilde{k}'$  as

$$u'(zf(k) + (1 - \delta)k - \tilde{k}')$$

$$= \beta \sum_{z' \in \mathcal{Z}} [(1 + z'f'(\tilde{k}') - \delta)u'(z'f(\tilde{k}') + (1 - \delta)\tilde{k}' - g_n(\tilde{k}', z'))$$

$$- \mu_n(\tilde{k}', z')(1 - \delta)]p(z', z)$$

- Construct  $\tilde{g}_{n+1}(k,z)$  using  $\tilde{k}'$ .
- ▶ Update  $g_{n+1}(k,z) = \max\{\tilde{g}_{n+1}(k,z), (1-\delta)k\}$
- Find  $\tilde{\mu}$  as

$$\begin{split} \widetilde{\mu} &= u'(zf(k) + (1-\delta)k - g_{n+1}(k,z)) \\ &- \beta \sum_{z' \in \mathcal{Z}} [(1+z'f'(g_{n+1}(k,z)) - \delta)u'(z'f(g_{n+1}(k,z)) \\ &+ (1-\delta)g_{n+1}(k,z) - g_n(g_{n+1}(k,z),z')) \\ &- \mu_n(g_{n+1}(k,z),z')(1-\delta)]p(z',z) \end{split}$$

- Construct  $\tilde{\mu}_{n+1}(k,z)$  using  $\tilde{\mu}$ .
- Update  $\mu_{n+1}(k, z) = \max\{\tilde{mu}_{n+1}(k, z), 0\}$



