# Determinants Group 12

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#### Overview

- Why calculate determinants?
  - Invertibility of matrices
  - Cramer's Rule
  - Eigenvalues and Eigenvectors
  - Volume
  - Jacobian determinant
- Using LU decomposition to compute determinants
  - What is LU decomposition?
  - Is there always an LU decomposition?
  - PLU decomposition
  - How it helps us compute determinants
  - How do we find the PLU decomposition?
  - Runtime analysis
- 3 Other algorithms for computing determinants
  - Leibniz formula
  - Laplace expansion
  - Gaussian elimination



### Invertbililty of matrices

#### Theorem

An  $n \times n$  square matrix A is invertible if and only if

$$\det(A) \neq 0$$

#### Theorem

Given an equation  $A\mathbf{x} = \mathbf{b}$  The solutions to  $\mathbf{x}$  are given by

$$x_i = \frac{\det(A_i)}{\det(A)}$$

with A<sub>i</sub> being the matrix formed by replacing the ith column of A by **b**.

### Eigenvalues and Eigenvectors

#### Definition

The Eigenvalues of a matrix A are the roots of the characteristic polynomial as defined

$$\chi_A = \det(A - \lambda I) = \mathbf{0}$$

#### Volume

The absolute value of the determinant of real vectors is equal to the volume of the parallelepiped spanned by those vectors.  $f: \mathbb{R}^n \to \mathbb{R}^n$ : the linear map represented by the A. S: any measurable subset of  $\mathbb{R}^n$ .

$$\mathsf{volume}(f(S)) = \sqrt{\det(A^T A)} \times \mathsf{volume}(S)$$

The volume of any tetrahedron, given its vertices a, b, c, and d is

$$\frac{\det(a-b,b-c,c-d)}{6}$$

#### Jacobian determinant

For  $f: \mathbb{R}^n \to \mathbb{R}^n$ , the Jacobian matrix is the  $n \times n$  matrix whose entries are defined as

$$D(f) = \left(\frac{\delta f_i}{\delta x_j}\right)_{1 \le i, j \le n}$$

Its determinant is known as the Jacobian determinant.

If the determinant of a continuously differentiable function f at a point p is...

- Non-zero, f is invertible near a point p in  $\mathbb{R}^n$ .
- Positive, then f preserves orientation near p.
- Negative, then f reverses orientation near p.

#### Definition

An LU decomposition of an invertible matrix A is a factorization

$$A = LU$$

where L and U are lower and upper triangular matrices, respectively.

### Is there always an LU decomposition?

#### No.

An LU decomposition of A exists if and only if each of its *leading* principle minors (contiguous square submatrices in the top-left corner of A), are also invertible.

#### Example

$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

This matrix is invertible but has no LU decomposition.

What can we do?

### PLU decomposition

#### Partial pivoting.

We can pivot the matrix into the correct form by multiplication with an orthogonal, permutation matrix P (representing a permutation  $\sigma_P$ ) which gives us the PLU decomposition:

$$\sigma_P(A) = PA = LU$$

This technique works on any invertible matrix.

### How it helps us compute determinants

Now that we have PA = LU, it follows that

$$A = P^{-1}LU$$
$$= P^{T}LU$$

since  $P^{-1} = P^T$  by the definition of orthogonal matrices.

### How it helps us compute determinants (cont.)

Now that we have  $A = P^T L U$ , it follows that

$$det(A) = det(P^{T}LU)$$

$$= det(P^{T}) \cdot det(L) \cdot det(U)$$
 (Thm. 4.4.1)
$$= det(P) \cdot det(L) \cdot det(U)$$
 (Lem. 4.4.4)

#### Given that

- the determinant of a triangular matrix is the product of its diagonal elements
- the determinant of a permutation matrix (P) is the parity of the permutation it represents  $(\sigma_P)$

it follows that

$$\det(A) = \operatorname{sgn}(\sigma_P) \cdot \left(\prod_{i=1}^n I_{i,i}\right) \left(\prod_{i=1}^n u_{i,i}\right)$$

### How do we find the PLU decomposition?

ASDFGHJK

### Runtime analysis

How quick is it?

- The PLU decomposition can be computed in  $\mathcal{O}(n^3)$  time.
- The determinants of the triangular matrices computed in  $\mathcal{O}(n)$  time.
- The parity of the permutation matrix in  $\mathcal{O}(n^2)$  time.

Therefore the total runtime for computing the determinant using the method is

$$\mathcal{O}(n^3) + \mathcal{O}(n^2) + \mathcal{O}(n) = \mathcal{O}(n^3)$$

What are some other methods to compute determinants?

#### Leibniz formula

#### Definition

The Leibniz formula defines the determinant of  $A \in \mathbb{M}(n)$  as

$$\det(A) = \sum_{\sigma \in \mathfrak{S}_n} \left( \operatorname{sgn}(\sigma) \cdot \prod_{i=1}^n a_{i,\sigma(i)} \right)$$

where  $\mathfrak{S}_n$  is the set of permutations length n.

Computing the determinant using this method is slow with runtime  $\mathcal{O}((n+1)!)$ .

### Laplace expansion

The Laplace (1st row) expansion for computing determinants is usually the first method taught for computing determinants of  $3\times 3$  matrices and larger.

#### $\mathsf{Theorem}$

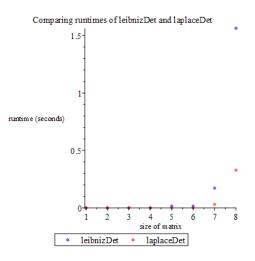
The formula for the (1st row) Laplace expansion of  $A \in \mathbb{M}(n)$  is given as:

$$\det(A) = \sum_{j=1}^n \mathsf{a}_{1,j} \; \mathsf{C}_{1,j}$$

where  $C_{i,j}$  is the (i,j) cofactor of A.

Its runtime complexity of  $\mathcal{O}(n!)$  is poor.

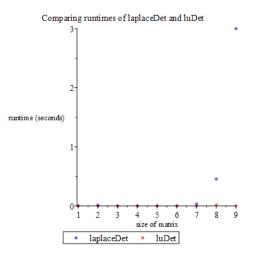
### Laplace expansion vs Leibniz formula



Runtimes are similar — both run in exponential time.



### Laplace expansion vs LU decomposition



The difference between the exponential and polynomial-time function is clear.

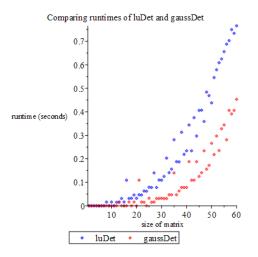


#### Gaussian elimination

- The determinant of a triangular matrix can be computed by taking the product of its diagonal entries (which is a quick  $\mathcal{O}(n)$  operation).
- Any invertible square matrix can be transformed into echelon form by performing Gaussian elimination, which takes  $\mathcal{O}(n^3)$  time.

So how does it compare to LU decomposition?

### Gaussian elimination vs LU decomposition



The difference in runtimes is small (a constant factor).



### Gaussian elimination (cont.)

Conventional Gaussian elimination requires division. This has two problems:

- Over  $\mathbb{M}(n; \mathbb{R})$  solutions maybe inexact, so precision is lost.
- Division is not a ring operation, so would not necessarily work on matrices over a ring.

This is can be addressed by using...

### Bareiss algorithm

- Addresses the issue of precision-loss by performing integer-preserving Gaussian elimination on integer matrices.
- The runtime complexity is  $\mathcal{O}(n^3)$  which is the same as conventional Gaussian Elimination, whilst preserving exactness.

### Bird's algorithm

Define  $\mu: \mathbb{M}(n) \to \mathbb{M}(n)$ :

$$\mu(X) = \begin{pmatrix} \mu_{2,2} - x_{2,2} & x_{1,2} & \cdots & x_{1,n-1} & x_{1,n} \\ 0 & \mu_{3,3} - x_{3,3} & \cdots & x_{2,n-1} & x_{2,n} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & \mu_{n,n} - x_{n,n} & x_{n-1,n} \\ 0 & 0 & \cdots & 0 & 0 \end{pmatrix}$$

 $F_{\Delta}^{n}(X) = \mu(F_{\Delta}^{n-1}(X)) \cdot A$ 

and 
$$F_A: \mathbb{M}(n) \to \mathbb{M}(n)$$
, with  $A \in \mathbb{M}(n)$ 

$$F_A(X) = \mu(X) \cdot A$$

$$F_A^2(X) = \mu(F_A(X)) \cdot A$$

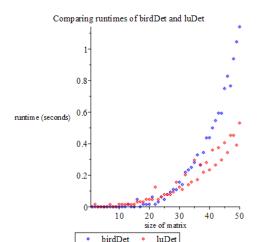
$$\vdots$$

### Bird's algorithm (cont.)

#### Bird's Theorem

$$F_A^{n-1}(A) = \begin{pmatrix} d & 0 & \cdots & 0 \\ 0 & 0 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & 0 \end{pmatrix} \text{ with } d = \begin{cases} \det(A) & \text{odd } n \\ -\det(A) & \text{even } n \end{cases}$$

- Enables the *division-free* computation of determinants in  $\mathcal{O}(n \cdot M(n))$  where M(n) is the runtime complexity of the matrix multiplication algorithm used.
- If the conventional  $\mathcal{O}(n^3)$  matrix multiplication algorithm is used, then Bird's algorithm will run in  $\mathcal{O}(n^4)$  time.
- But this can be reduced to  $\mathcal{O}(n^{3.8})$  by using the *Strassen* algorithm for matrix multiplication.

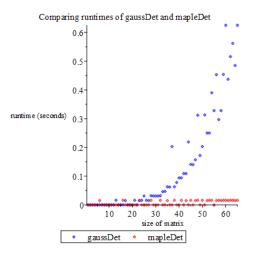


Bird's runtimes increase noticeably more rapidly than LU decomposition, but it's still polynomial. 4 D > 4 B > 4 B > 4 B > 9 Q P

### Summary of determinant algorithms

Algorithm	Runtime	Exact?
Leibniz formula	$\mathcal{O}((n+1)!)$	Yes
Laplace expansion	$\mathcal{O}(n!)$	Yes
LU decomposition	$\mathcal{O}(n^3)$	No
Gaussian elimination	$\mathcal{O}(n^3)$	No
Bareiss algorithm	$\mathcal{O}(n^3)$	Yes
Bird's algorithm	$\mathcal{O}(n^{3.8})$	Yes

### How fast is Maple's built-in determinant function?



Very. Maple's optimisation means a fair comparison cannot be made.



## Thanks!