# Information Science Model Answers of Sorting and String

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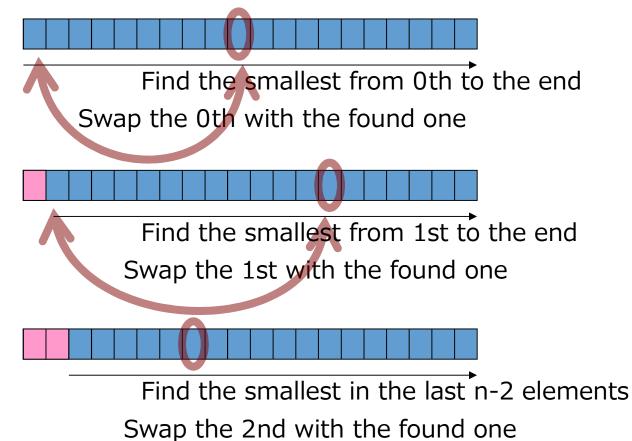


### Contents

- Sorting exercises
  - Behavior
  - Making a graph
- Dealing with String: Match and reverse

### Review of Sorting Exercises

- Simple Sort Algorithm
  - Repeat the following
    - Find the minimum index of the remaining elements
    - Swap it with the head and delete it



```
def simplesort(a)
  for i in 0..(a.length()-1)
    k = \min index(a,i)
    v = a[k]
    a[k] = a[i]
    a[i] = v
  end
  a
end
```

```
def simplesort(a)
  for i in 0..(a.length()-1)
    k = \min index(a,i)
    v = a[k]
                      Find the index of the
     a[k] = a[i]
                      minimum in the last
     a[i] = v
                       (n-i) elements in a
  end
  a
end
```

```
def simplesort(a)
  for i in 0..(a.length()-1)
    k = \min index(a,i)
    v = a[k]
    a[k] = a[i]
    a[i] = v
  end
                            Swap a[k] and a[i].
  a
                            Necessary to use a
                               variable v
end
```

# Exercise: Apply to a=[1,4,2,9,8,3,2,6,4]

#### It repeats 9 times(not 6)

Do nothing at i=0 (1 is at the beginning)

```
8 in this case
def simplesort(a)
  for i in 0..(a.length()-1)
 # display a at this point
    k = \min index(a,i)
    v = a[k]
    a[k] = a[i]
    a[i] = v
  end
  a
end
```

Swap min and head (don't change the others)

```
a=[1,4,2,9,8,3,2,6,4]
      1 a=[1,4,2,9,8,3,2,6,4]
      2 a=[1,2,4,9,8,3,2,6,4]
      3 a=[1,2,2,9,8,3,4,6,4]
      4 a=[1,2,2,3,8,9,4,6,4]
      5 a=[1,2,2,3,4,9,8,6,4]
      6 a=[1,2,2,3,4,4,8,6,9]
      7 a=[1,2,2,3,4,4,6,8,9]
      8 a=[1,2,2,3,4,4,6,8,9]
output a=[1,2,2,3,4,4,6,<mark>8</mark>,9]
```

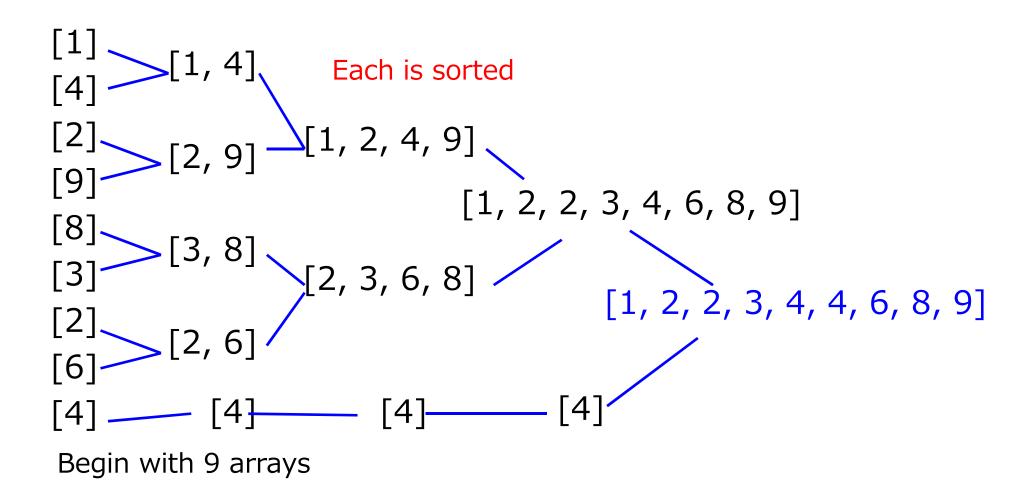
### min\_index(a, i)

Find the index with minimum value among a[i]..a[a.length()-1]

```
def min index(a, i)
                                        m: index w/ minimum value
 m = i
                                           so far
 for p in i.. (a.length()-1)
                                        check all entries in
                                        p=i..the last
    if a[p] < a[m]
                                           If p-th one is smaller,
      m = p
                                           set m to be p.
    end
 end
 m
end
                             Find the smallest in the last n-i+1 elements
                              Swap the i-th with the found one
```

# Exercise: Apply merge Sort

$$= [1,4,2,9,8,3,2,6,4]$$



### Exercise: Apply merge Sort

> #arrays reduces by almost half in each iteration

$$\bullet \ 9(\cong n) \rightarrow 5(\cong n/2) \rightarrow 3(\cong n/4) \rightarrow \ 2(\cong n/8) \rightarrow 1(\cong n/16)$$

```
Min integer at least n/2
                          t: # iterations
                           = min integer w/ n/2<sup>t</sup> \leq 1
                               t ≅ log n
                   [1, 2, 2, 3, 4, 6, 8, 9]
                               [1, 2, \dot{2}, 3, 4, 4, 6, 8, 9]
```

Contents

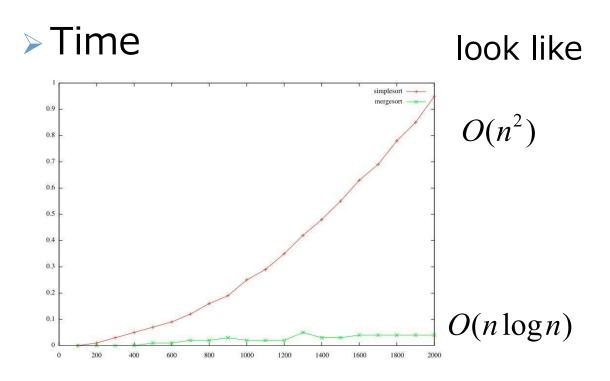
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### Remarks on Exercise 4

 Make a graph of the computational times for simple sort and merge sort

- > Remark
  - Function "bench" does not work at home
    - have to install new software "gnuplot"

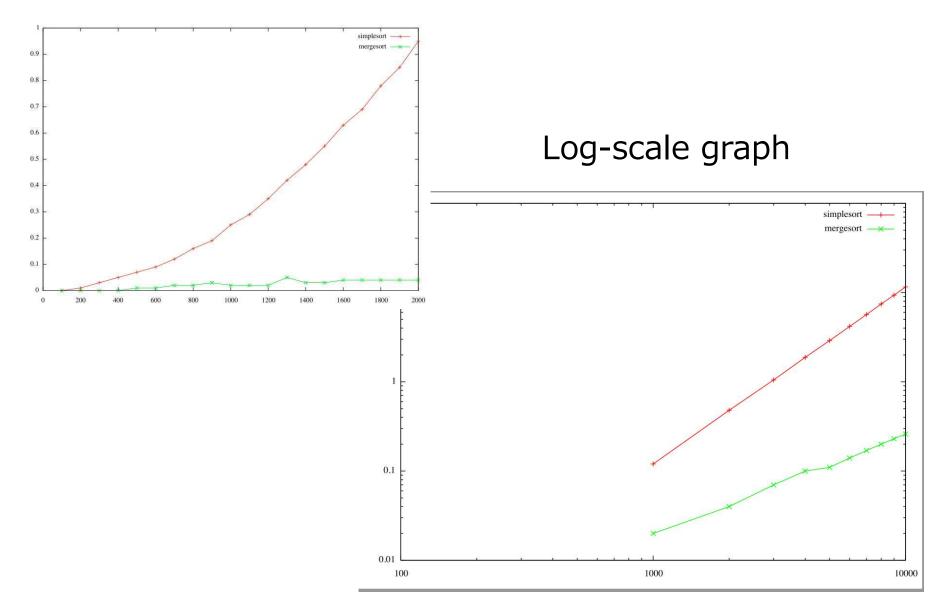
## How to Read Log Scale Graph



- How can we analyze more precisely?
  - → log-scale graph

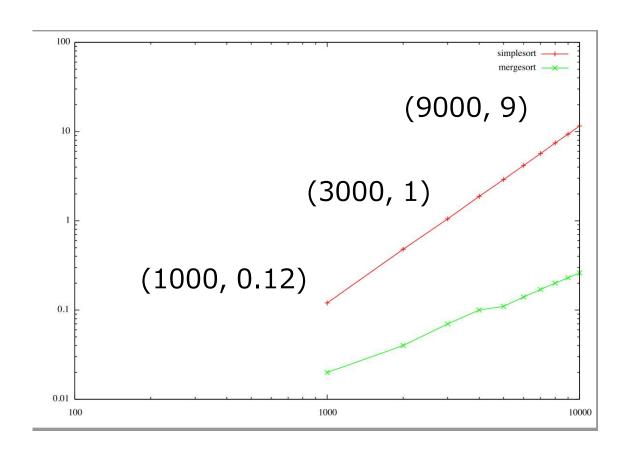
# How to Read Log Scale Graph

#### > Time



## Simple Sort Graph

Compute the slope using representatives



# Simple Sort Graph

➤ Slope

$$\frac{\log 1 - \log 0.12}{\log 3000 - \log 1000} = 1.92.. \qquad \frac{\log 1}{\log 1000} = 1.92..$$

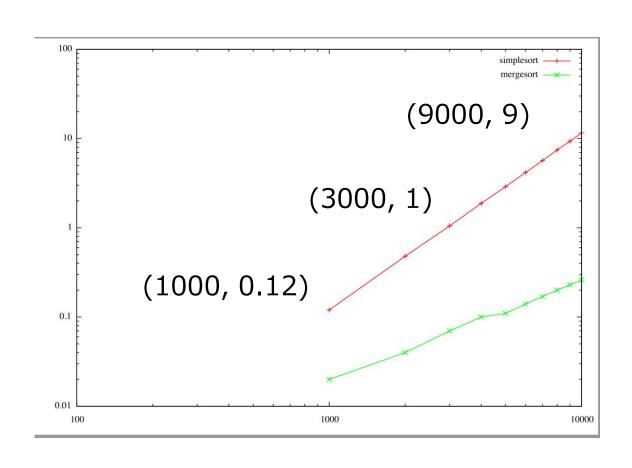
$$\frac{\log 9 - \log 1}{\log 9000 - \log 3000} = 2$$

> Implying that

 $\log(\text{time}) \propto 2 \log n$ 

time  $\propto n^2$ 

 $O(n^2)$ 



### Merge-Sort Graph

Unlikely to a straight line

$$\frac{\log 0.07 - \log 0.03}{\log 4000 - \log 3000} = 2.94.. \qquad \frac{\log 0.11 - \log 0.03}{\log 8000 - \log 3000} = 1.32..$$

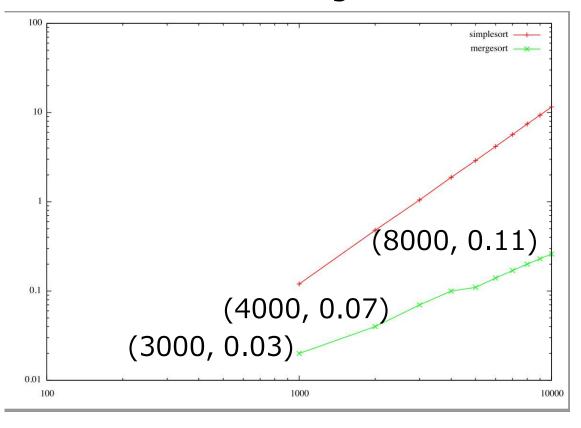
Slope is > 1, but < 2 when the size is large



worse than O(n) faster than  $O(n^2)$ 

Close to theory time  $\propto n \log n$ 

Might be better to plot "time"/n



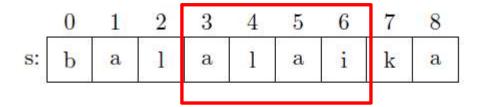
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> Dealing with String: Match and reverse

# Last Exercise1: Searching a Substring

- > Find some keyword in a string
  - Ex. Analysis of DNA sequences



Q. Is there "alai", denoted by p, in the sequence?

# (review) Searching Keywords

```
def match(s,p)
  i = 0
  w = p.length()
  while submatch(s,i,p,w) < w
    i = i + 1
  end
end
```

Function that counts how many characters are matched from the i-th char in s

Decide if # matched = length of p

Downloadable from the Ruby-program website

# (review) Procedure submatch(s,i,p,w)

```
def submatch (s,i,p,w)
   i = 0
   while j < w && s[(i+j)..(i+j)] == p[j..j]
      j = j + 1
                                     Compare
   end
                            the (i+j)th in s and the jth in p
end
                               a
                                     a
while they coincide
                                  p:
                                     a
we increment j by one
```

W

#### Last Exercise 1:

- What is the computational complexity of the function match?
  - Let n be the length of s, and m be the length of p,
  - Represent the complexity using n and m
- Hint: It suffices to consider the worst case:
  - Consider the following case & estimate #repetitions

```
□ s=[aa…aaab]

n-1

p=[aa…ab]

m-1
```

### Computational Time of Match

Repeatedly compare s[i..(i+w-1)] and p

s: a a a a a a a a a a a a a b

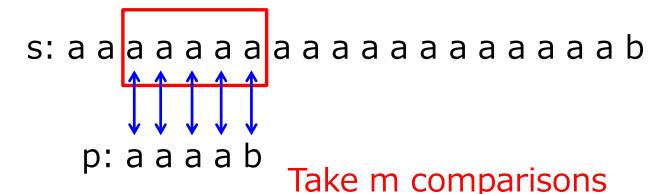
p: a a a a b

- ➤ When s=[aa···aaaab] and p=[aa···ab]
  - Need to consider s[i..(i+m-1)] for all i=0,···,n-m
    - $\square$  #(substrings of size m in s) = n m +1
  - Each time we call submatch

Computational time: (n-m+1)\*[time for submatch]

### Behavior of submatch (when fixing i)

Compare each element in s[i..(i+w-1)] and p



- ➤ When s=[aa···aaaab] and p=[aa···ab]
  - Need to consider s[i..(i+m-1)] for all i=0,···,n-m
    - $\square$  # substrings ([aaa···a] or [aaa···b] ) = n m +1
  - Each time we call submatch

```
total time: (n-m+1)*m (\leftarrowok as an answer)
= O(n*m) (leave only the most dominant term)
```

## Last Exercise 2: Reverse a String

- > Given a string, we want to reverse it
  - Ex. For s="abcdef", the output is "fedcba"

- Requirement
  - Use "while" (or "for")
  - Do not use s.reverse() nor s.split("").reverse().join()
    - Already implemented in Ruby

### Last Exercise 2: Framework

> Fill in the question marks

```
def reverse(s)
  result = "" # empty string(length 0)
  i = ??
  while i >= 0 do
    ??
    i = i - 1
    end
  result # return the reversed string
end
```

### **Basic Behaviors**

#### > Idea

```
def reverse(s)
  result = ""
  i = s.length()-1
  while i >= 0 do
    ??
    i = i - 1
  end
  result
end
```

Read s one-by-one from the end

s: PEAK2013

# The first entry is the end index

# decrease by one

(final state)

result: 3 1 0 2 K E A P

#### > Idea

```
def reverse(s)
  result = ""
  i = s.length()-1
  while i >= 0 do
    ??
    i = i - 1
  end
  result
end
```

Read s one-by-one from the end

s: PEAK2013

result: 3 1 0···· (start from empty)

(final state)

result: 3 1 0 2 K E A P

#### > Idea

```
def reverse(s)
  result = ""
  i = s.length()-1
  while i >= 0 do
    ??
    i = i - 1
  end
  result
end
```

```
7(s.length-1)
  s: PEAK2013
result: 3 1
         Append s[i..i] at the end
         result = "3" + "1"
                    result
              in the previous iteration
```

#### > Idea

```
def reverse(s)
  result = ""
  i = s.length()-1
  while i >= 0 do
    ??
    i = i - 1
  end
  result
end
```

```
7(s.length-1)
  s: PEAK2013
result: 3 1
         Append s[i..i] at the end
         result = "31" + "0"
                    result
              in the previous iteration
```

### Exercise 2: Reverse a String

> Another solution: Using concatenation

```
def reverse(s)
                               s: PEAK2013
 result = ""
 i = s.length()-1
                                  "result"+"0"
 while i >= 0 do
                                   (="31")
  result = result + s[i..i]
                                  String so far obtained
  i = i - 1
 end
                        OK to be s[i]
 result
end
```

"result" contains the reversed string so far add s[i..i] at the end of "result"

#### Exercise 2: Recursive Ones

Based on recursion

```
def reverse_r(s)
    n = s.length()
    if n == 1
        s
    else
        s[n-1..n-1] + reverse_r(s[0..(n-2)])
    end
end
```

```
Ex. (Reverse of "PEAK2013") = 3 + (Reverse of "PEAK201")
```

```
def reverse_wrong(s)
  result = ""
```

s[i]=s[s.length-1-i]

s[s.length-1-i]=c

i = s.length()-1

while i >= 0 do

c=s[i]

i = i - 1

end

end

result

Prepare result

Because we return "result", need to change "result" not "s"

# Analysis of Program from some of you

7(s.length-1)

```
def reverse_wrong(s)
  result = ""
  i = s.length()-1
  while i >= 0 do
    c=s[i]
    s[i]=s[s.length-1-i]
                           S:
    s[s.length-1-i]=c
    i = i - 1
  end
  result
end
                           S:
```

```
PEAK2013
                    i=7
    3 E A K 2 0 1 P
                     i=6
    3 1 A K 2 0 E P
                     i=5
   3 1 0 K 2 A E P
                     i=4
    3 1 0 2 K A E P
                    Still going...
                     i=3
    3 1 0 K 2 A E P
                     i=2
    31AK20EP
                     i=1
   3 E A K 2 0 1 P
                     i=0
    PEAK2013
S:
                     output
```

### To make it correct: Run to half

7(s.length-1)

```
def reverse2(s)
  i = s.length() - 1
  k = s.length()/2
  while i > = k do
  # stop when i=k
    c=s[i]
                           S:
    s[i]=s[s.length-1-i]
                           S:
    s[s.length-1-i]=c
    i = i - 1
  end
  s # change output
end
```

```
PEAK2013
                 i=7
3 E A K 2 0 1 P
                  i=6
3 1 A K 2 0 E P
                 i=4
3 1 0 K 2 A E P
 3 1 0 2 K A E P
  Stop here (when i=3)
```