

Formal Methods

Seminar 2

First Part of the Programming Assignment

Design in Ocaml the Abstract Syntax Tree for a small object-oriented language called CoreJava.

The syntax of CoreJava is described in the following slides.

Syntax of CoreJava

A CoreJava program P consists of a list of class declarations as follows:

$P ::= \text{def}^+$

$\text{def}^+ ::= \text{def} \mid \text{def1}; \text{def2}; \dots; \text{defn}$

The program contains at least one class declaration.

The last class declaration must contain a method called “**main**” which is the starting execution point for the CoreJava program.

Syntax of CoreJava

A class declaration def consists of the following:

def ::= class cn1 extends cn2 { field* # meth* }

where cn1 is the current class name and cn2 is the name of the parent class

The body of the class consists of a list of fields declarations followed by a list of methods declarations. The fields list is separated by “#” from the methods list. **The fields list is either empty or field+ ::= field1;field2;...;fieldn. The methods list is either empty or meth+ ::= meth1;meth2;...;methn.**

Syntax of CoreJava

A field declaration consists of the following:

$\text{field} ::= \text{typ fn}$

where typ is the type of the field and fn is the name of the field

A type typ can be either a primitive type, or a class name or the bottom (the type of null).

$\text{typ} ::= \text{prim} \mid \text{cn} \mid _ \mid _$

$\text{prim} ::= \text{int} \mid \text{float} \mid \text{bool} \mid \text{void}$

Syntax of CoreJava

A method declaration consists of the following:

$\text{meth} ::= \text{typ1 mn}((\text{typ pn})^*) \{e\}$

where typ1 is the type of the method result, mn is the method name, $(\text{typ pn})^*$ is the list of method parameters.

The list of method parameters can be either empty or $(\text{typ pn})^+ ::= \text{typ1 pn1, typ2 pn2, ..., typn pnn}$

The body of the method is given by the expression e .

Syntax of CoreJava

An expression is defined as follows:

$e ::= \text{null}$	(null value)
kint	(a constant value)
kfloat	
$\text{true} \mid \text{false}$	
$()$	(a value of type void)
v	(a variable name)
$v.f$	(a field f of an object denoted by v)

Syntax of CoreJava

An expression is defined as follows:

$e ::= \dots$

| $v = e$ (variable assignment)

| $v.f = e$ (object field assignment)

| $\{ (typ\ v)\ e1 \}$ (local variable declaration such that v has type typ and its scope is the expression $e1$)

| $e1; e2$ (a sequence such that $e2$ is executed only after $e1$ is executed)

| $\text{if } v \text{ then } \{e1\} \text{ else } \{e2\}$ (conditional where the condition is a variable)

Syntax of CoreJava

An expression is defined as follows:

$e ::= \dots$

| $e1 \text{ opint } e2$ (arithmetic expressions for ints
where $\text{opint} ::= +|-|*|/$)

| $e1 \text{ opfloat } e2$ (arithmetic expressions for floats
where $\text{opfloat} ::= +.|-.|*|/|.$)

| $e1 \ \&\& \ e2$ (logical expressions)

| $e1 \ || \ e2$

| $!e$ (negation)

Syntax of CoreJava

An expression is defined as follows:

$e ::= \dots$

| $\text{new cn}(v^*)$ (creation of an instance object of the class cn . The fields of the new object are initialized by the list of variables v^* . The fields are initialized in the order defined in the declaration of class cn)

| $v.\text{mn}(v^*)$ (method call where v is the method receiver and the list of variables v^* are the current arguments initializing the method parameters)

| $\text{while } v \{e\}$ (loop where the condition is a variable)

Syntax of CoreJava

An expression can also be

$e ::= \dots$

| $e1 \text{ opcmp } e2$ (relational expressions where
 $\text{opcm} ::= < | <= | = | != | > | >=$)

| $(cn) v$ (cast expression)

| $v \text{ instanceof } cn$

Discussion on the Ocaml implementation

The following slides show how you can represent the Abstract Syntax Tree of CoreJava in Ocaml.

A program can be represented as a dictionary of classes (association list or you may want to consult Chapter 13 from realworldocaml.org):

```
type progr= (string*classDecl) list
```

```
type classDecl =  
  (string*string*fldDeclList*mthDeclList)
```

```
type fldDeclList = fldDecl list
```

```
type fldDecl = typ * string
```

```
type mthDeclList = mthDecl list
```

```
type mthDecl =(typ*string*fPrmList*blkExp)
```

```
type fPrmList = fPrm list
```

```
type fPrm = (typ*string)
```

type typ = Tprim of tPrim

| Tclass of string | Tbot

type tPrim = Tint | Tfloat | Tbool | Tvoid

type blkExp = Bvar of typ*string*exp

| Bnvar of exp

type val = Vnull | Int of int | Float of float

| Bool of bool | Vvoid

type exp = Value of val

| Var of string | Vfld of string*string

|(see next slide)

type exp = ...

- | AsgnV of string*exp**
- | AsgnF of string*string*exp**
- | Blk of blkExp**
- | Seq of exp*exp**
- | If of string*blkExp*blkExp**
- | AddInt of exp*exp**
- | MulInt of exp*exp**
- | (please continue yourself the first assignment)**

AST Example

The following Core-Java expression:

if m then {z=z+1} else {z=2}

Is represented in Ocaml as the following AST:

```
Let ifexp1 =(IF ("m",  
                (Blk (Bnvar (AsgnV ("z", (AddInt (Var "z", (Value (Int  
1))))))))) ,  
                (Blk (Bnvar (AsgnV ("z", (Value (Int 2)) )))  
                )
```


**Please represent the following Core-Java
examples as ASTs in Ocaml:**

Core-Java Example

```
class A extends Object{  
int f1;  
#  
int m1(int a, int b) { (int c)  
                c=a+b;this.f1=this.f1+c;c};  
}
```

Core-Java Example

```
class B extends A{  
  A f2;  
  #  
  A m2(A x, A y) {(A z) { (int n)  
    n=x.m1(1,2)-y.m1(2,1);  
    {(bool m) m=(x.f1-y.f1)>n;  
    if m then {z=new A(m)} else {z=new A(n)}  
  }  
};this.f2=z;z  
}
```

Core-Java Example

```
Class Main extends Object{ #  
    Void main(){ (B o1) o1=new B(0,null);  
        { (A o2) o2=new A(2);  
            { (A o3) o3=new A(3);  
                o2 =o1.m2(o2,o3)  
            }  
        }  
    }  
}
```