

Algorithms and Data Structures 1 CS 0445



Fall 2022
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(Slides are adapted from Dr. Ramirez's and Dr. Farnan's CS1501 slides.)

Announcements

- Upcoming Deadlines:
 - Homework 5: this Friday @ 11:59 pm
 - Lab 4: next Monday @ 11:59 pm
 - Programming Assignment 1: Friday Oct. 7th Monday Oct. 10th
 - Autograder is up on GradeScope
- If you think you lost points in a lab assignment because of the autograder or because of a simple mistake
 - please reach out to Grader TA over Piazza
- Live Remote Support Session for Assignment 1
 - Recording and slides on Canvas
- Student Support Hours of the teaching team are posted on the Syllabus page

Previous Lecture ...

- ADT Stack
 - Linked implementation
 - Implementation using ADT List
 - Application: Building a simple parser of Algebraic expressions

Today ...

- ADT Stack
 - Application: Building a simple parser of Algebraic expressions
 - Application: Runtime stack
- Recursion

Our Plan for Processing Algebraic Expressions

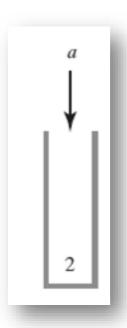
- 1. Check if input infix expression is balanced
- 2. Convert the expression from infix to postfix
- 3. Evaluate the postfix expression

Our Plan for Processing Algebraic Expressions

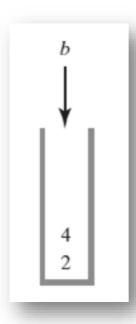
- 1. Check if input infix expression is balanced
- 2. Convert the expression from infix to postfix
- 3. Evaluate the postfix expression

- 1. Initialize an empty Stack
- 2. for each character in postfix expression
 - 1. if variable, push its value to Stack
 - 2. if operator
 - pop second operand
 - 2. pop first operand
 - 3. apply operator to two operands
 - 4. push result
- 3. Return the remaining value in Stack

The stack during the evaluation of the postfix expression



The stack during the evaluation of the postfix expression



The stack during the evaluation of the postfix expression



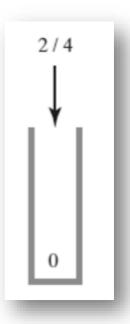
The stack during the evaluation of the postfix expression



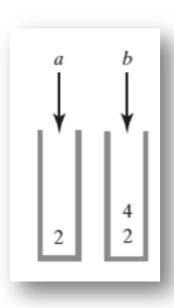
The stack during the evaluation of the postfix expression



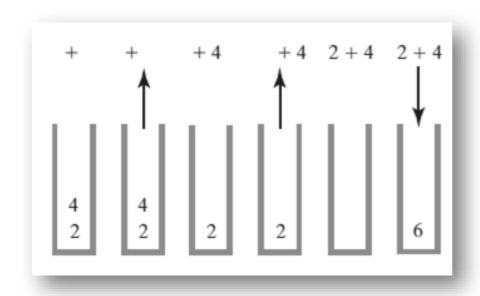
The stack during the evaluation of the postfix expression



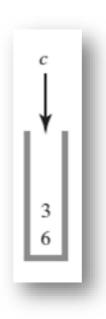
The stack during the evaluation of the postfix expression ab+c/ when a is 2, b is 4, and c is 3



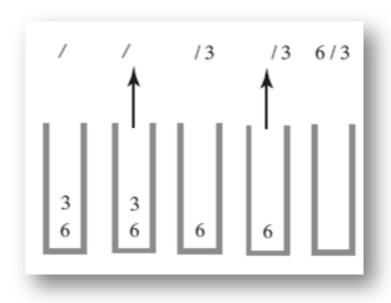
The stack during the evaluation of the postfix expression a b + c / when a is 2, b is 4, and c is 3



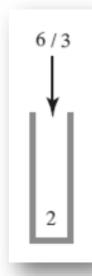
The stack during the evaluation of the postfix expression a b + c \prime when a is 2, b is 4, and c is 3



The stack during the evaluation of the postfix expression a b + c / when a is 2, b is 4, and c is 3



The stack during the evaluation of the postfix expression a b + c / when a is 2, b is 4, and c is 3



Evaluating Postfix Expressions

Algorithm for evaluating postfix expressions.

```
Algorithm evaluatePostfix(postfix)
  // Evaluates a postfix expression.
  valueStack = a new empty stack
  while (postfix has characters left to parse)
      nextCharacter = next nonblank character of postfix
      switch (nextCharacter)
        case variable:
            valueStack.push(value of the variable nextCharacter)
            break
case + incase incase incase case case case incase
```

Evaluating Postfix Expressions

Algorithm for evaluating postfix expressions.

```
break

case '+': case '-': case '*': case '/': case '^':

operandTwo = valueStack.pop()

operandOne = valueStack.pop()

result = the result of the operation in nextCharacter and its operands

operandOne and operandTwo

valueStack.push(result)

break

default: break // Ignore unexpected characters

}
```

What is the running time?

- in terms of n, the length of the input prefix string
- Check balance
 - how many times does each character get pushed?
 - at most 1
 - how many times does each character get poped?
 - at most 1
 - What is the runtime of push and pop?
 - O(1)
 - O(n)
- Convert infix to postfix: O(n)
- Evaluate postfix: O(n)
- Total: O(3n) = O(n)
- Three passes!
- Can we do better?
- Yes! We can use two passes only
 - Expect to require more space
 - space-time tradeoff

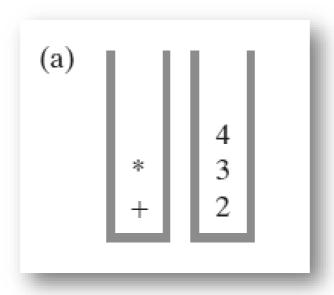
Evaluating Infix Expressions with 2 passes only

- We will use two stacks
 - Operator Stack
 - Operand stack
- Scan the expression once:
 - follow the steps of infix conversion to postfix,
 - except
 - instead of appending to postfix output, push to operand stack
 - when popping an operator, pop second then first operands, apply operator, push result to operand stack
- While operator stack not empty
 - pop an operator
 - pop second operand then first operand
 - apply the operator and push result to operand stack
- Result is the remaining value in the operand stack

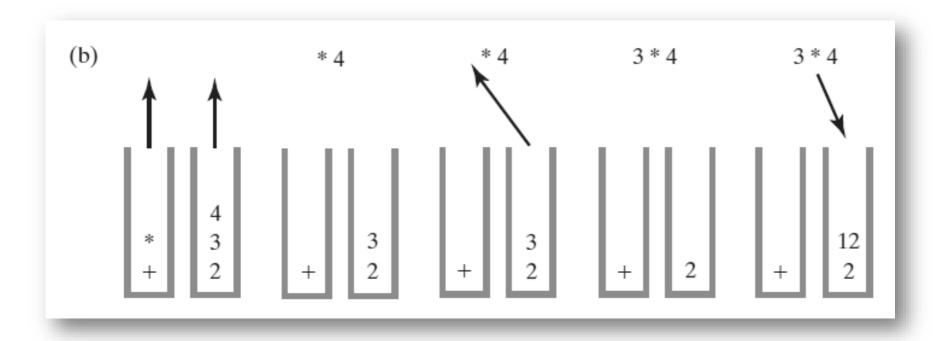
Evaluating Infix Expressions with 2 passes only

Two stacks during the evaluation of a + b * c when a is 2, b is 3, and c is 4:

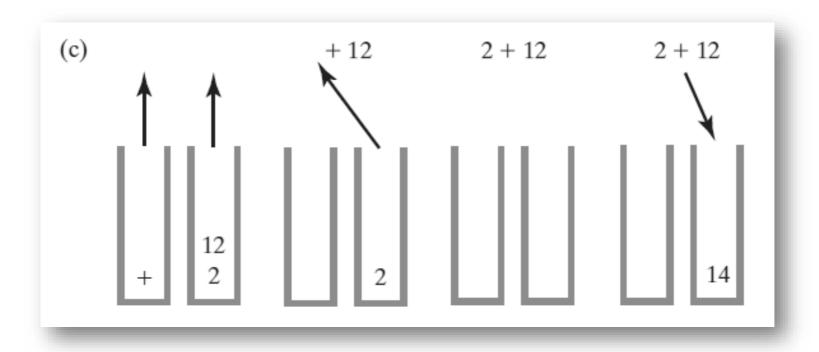
after reaching the end of the expression;



Two stacks during the evaluation of a + b * c when a is 2, b is 3, and c is 4: while performing the multiplication;



Two stacks during the evaluation of a + b * c when a is 2, b is 3, and c is 4: (c) while performing the addition



```
Algorithm evaluateInfix(infix)
    // Evaluates an infix expression.
    operatorStack = a new empty stack
    valueStack = a new empty stack
    while (infix has characters left to process)
       nextCharacter = next nonblank character of infix
       switch (nextCharacter)
          case variable:
            valueStack.push(value of the variable nextCharacter)
            break
         case 'A' :
            operatorStack.push(nextCharacter)
            break
         case '+' : case '-' : case '*' : case '/' :
```

```
VINGERAL CONTROL CONTR
                   case '+' : case '-' : case '*' : case '/' :
                                while (!operatorStack.isEmpty() and
                                                         precedence of nextCharacter <= precedence of operatorStack.peek())</pre>
                                             // Execute operator at top of operatorStack
                                             topOperator = operatorStack.pop()
                                             operandTwo = valueStack.pop()
                                             operandOne = valueStack.pop()
                                             result = the result of the operation in topOperator and its operands
                                                                                    operandOne and operandTwo
                                             valueStack.push(result)
                                operatorStack.push(nextCharacter)
                                break
                   case '(' :
                                operatorStack.push(nextCharacter)
                                break
.....Fase in initial histock is not empty if infix expression is walldown
```

```
case '('
    operatorStack.push(nextCharacter)
    break
  case ')': // Stack is not empty if infix expression is valid
    topOperator = operatorStack.pop()
    while (topOperator != '(')
       operandTwo = valueStack.pop()
       operandOne = valueStack.pop()
       result = the result of the operation in topOperator and its operands
               operandOne and operandTwo
       valueStack.push(result)
       topOperator = operatorStack.pop()
    break
```

The Runtime Stack (aka program stack)

Under the hood/behind the scenes alert!

The Runtime Stack (aka program stack)

- Under the hood/behind the scenes alert!
- A stack is created for each running program
 - called runtime stack
- The stack is used to hold data for each method call
 - in an activation record (aka activation frame or just frame)
- Activation record stores:
 - method parameters
 - local (method) variables
 - address of return point (i.e., next statement to execute after returning from call)
- When a method is called, its activation record is pushed to the runtime stack
- When a method returns, the top activation record is popped

Example

- The following code has three methods:
 - main
 - methodA
 - method
- main calls methodA
- methodA calls method
- Side note: methodA and method must be static
 - because they are called from main, which must static
- What are the local variables of each method?
- What are the parameters of each method?

```
public static
     void main(string[] arg)
        int x = 5;
        int y = methodA(x);
     } // end main
100
     public static
     int methodA(int a)
        int z = 2:
        methodB(z);
120
        return z:
     } // end methodA
150
     public static
     void methodB(int b)
     } // end methodB
```

Program

The Program Stack

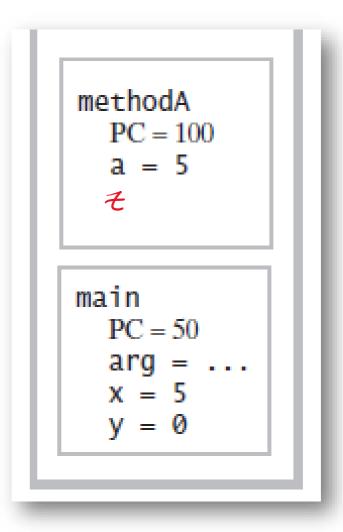
- The program stack when main begins execution
- PC is the Program Counter CPU register
 - it keeps track of the address of the next instruction to execute

```
public static
     void main(string[] arg)
        int x = 5;
        int y = methodA(x);
       // end main
     public static
100
     int methodA(int a)
        int z = 2;
        methodB(z);
120
        return z;
     } // end methodA
                                      main
     public static
150
                                        PC = 1
     void methodB(int b)
     } // end methodB
                                           (a)
                                        Program stack at th
           Program
```

The Program Stack

- The program stack when methodA begins execution
- Before methodA starts, the value of PC is stored in the activation record of main

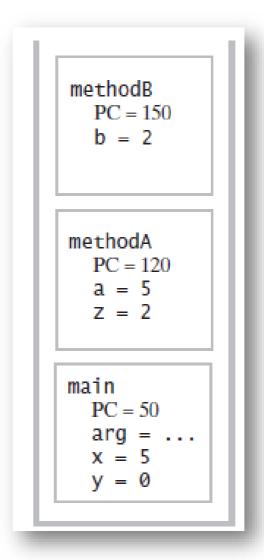
```
public static
     void main(string[] arg)
        int x = 5;
        int y = methodA(x);
     } // end main
     public static
100
     int methodA(int a)
        int z = 2;
120
        methodB(z);
        return z;
     } // end methodA
     public static
     void methodB(int b)
     } // end methodB
           Program
```



The Program Stack

- The program stack when methodB begins execution
- Before methodB starts, the value of PC is stored in the activation record of methodA

```
public static
     void main(string[] arg)
        int x = 5;
        int y = methodA(x);
     } // end main
     public static
100
     int methodA(int a)
        int z = 2;
120
        methodB(z);
        return z;
     } // end methodA
     public static
     void methodB(int b)
     } // end methodB
           Program
```

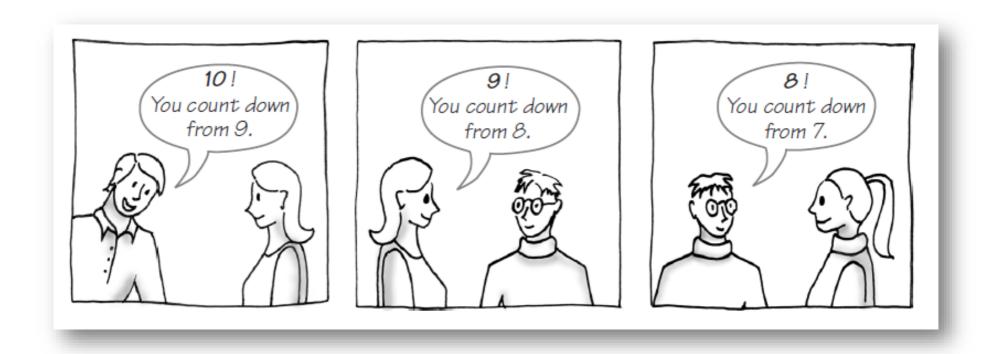


Recursion

What Is Recursion?

- Consider hiring a contractor to build
 - He hires a subcontractor for a portion of the job
 - That subcontractor hires a sub-subcontractor to do a smaller portion of job
- The last sub-sub- ... subcontractor finishes
 - Each one finishes and reports "done" up the line

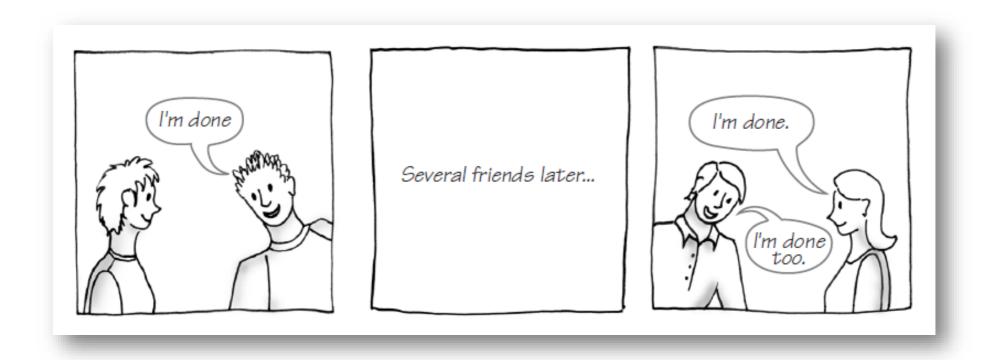
Counting down from 10



Counting down from 10



Counting down from 10



```
/** Counts down from a given positive integer.
@param integer An integer > 0. */
public static void countDown(int integer)
{
```

```
/** Counts down from a given positive integer.
    @param integer An integer > 0. */
public static void countDown(int integer)
{
    System.out.println(integer);
```

```
/** Counts down from a given positive integer.
    @param integer An integer > 0. */
public static void countDown(int integer)
{
    System.out.println(integer);
    if (integer > 1)
```

```
/** Counts down from a given positive integer.
    @param integer An integer > 0. */
public static void countDown(int integer)
   System.out.println(integer);
   if (integer > 1)
      countDown(integer - 1);
} // end countDown
```

- Each call to the countdown method corresponds to one person
 - what did each person do?
 - say a number

```
/** Counts down from a given positive integer.
    @param integer An integer > 0. */
public static void countDown(int integer)
   System.out.println(integer);
   if (integer > 1)
      countDown(integer - 1);
} // end countDown
```

- Each call to the countdown method corresponds to one person
 - what did each person do?
 - say a number
 - check if not done

```
/** Counts down from a given positive integer.
    @param integer An integer > 0. */
public static void countDown(int integer)
   System.out.println(integer);
   if (integer > 1)
      countDown(integer - 1);
} // end countDown
```

- Each call to the countdown method corresponds to one person
 - what did each person do?
 - say a number
 - check if not done
 - ask a classmate to count down starting from the number before

```
/** Counts down from a given positive integer.
    @param integer An integer > 0. */
public static void countDown(int integer)
   System.out.println(integer);
   if (integer > 1)
      countDown(integer - 1);
} // end countDown
```

Definition

- Recursion is a problem-solving process
 - Breaks a problem into identical but smaller problems.
- A method that calls itself is a recursive method.
 - The invocation is a recursive call or recursive invocation.

Method must be given an input value

```
/** Counts down from a given positive integer.
    @param integer An integer > 0. */
public static void countDown(int integer)
{
    System.out.println(integer);
    if (integer > 1)
        countDown(integer - 1);
} // end countDown
```

- Method must be given an input value
- Method definition must contain logic that involves this input, leads to different cases

```
/** Counts down from a given positive integer.
    @param integer An integer > 0. */
public static void countDown(int integer)
{
    System.out.println(integer);
    if (integer > 1)
        countDown(integer - 1);
} // end countDown
```

- Method must be given an input value
- Method definition must contain logic that involves this input, leads to different cases
- One or more cases should provide solution that does not require recursion
 - otherwise, infinite recursion
 - if integer <= 1 → return

```
/** Counts down from a given positive integer.
    @param integer An integer > 0. */
public static void countDown(int integer)
{
    System.out.println(integer);
    if (integer > 1)
        countDown(integer - 1);
} // end countDown
```

- Method must be given an input value
- Method definition must contain logic that involves this input, leads to different cases
- One or more cases should provide solution that does not require recursion
- One or more cases must include a recursive invocation

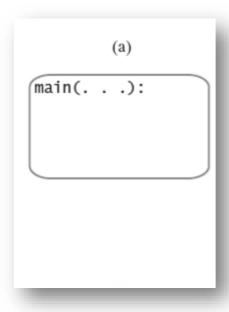
```
/** Counts down from a given positive integer.
    @param integer An integer > 0. */
public static void countDown(int integer)
{
    System.out.println(integer);
    if (integer > 1)
        countDown(integer - 1);
} // end countDown
```

Programming Tip

- Iterative method contains a loop
- Recursive method calls itself
- Some recursive methods contain a loop and call themselves
 - If the recursive method with loop uses while, make sure you did not mean to use an if statement

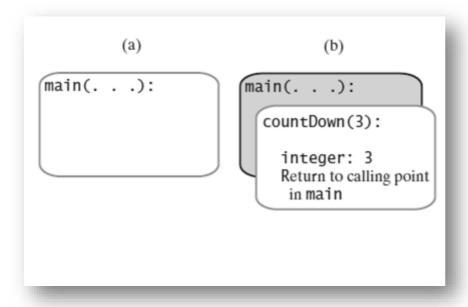
The stack of activation records during the execution of the call countDown (3)

pushing → activation records pile up



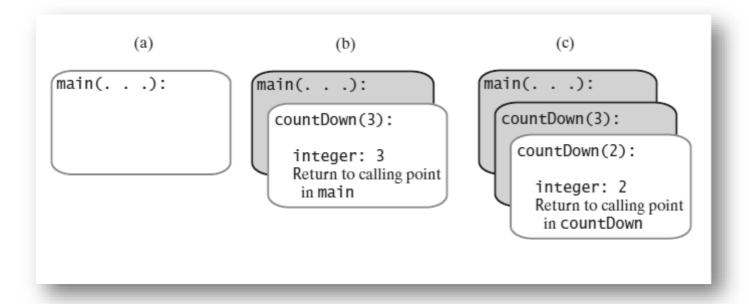
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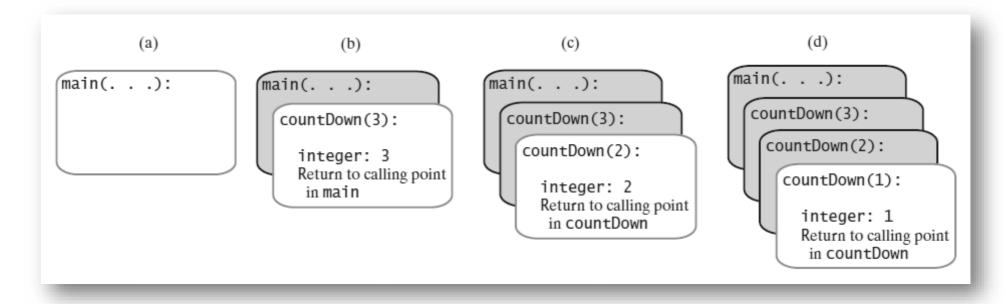
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The stack of activation records during the execution of the call countDown (3)

pushing -> activation records pile up



The stack of activation records during the execution of the call countDown (3)

popping → activation records tear down

```
(e)

main(. . .):

countDown(3):

countDown(2):

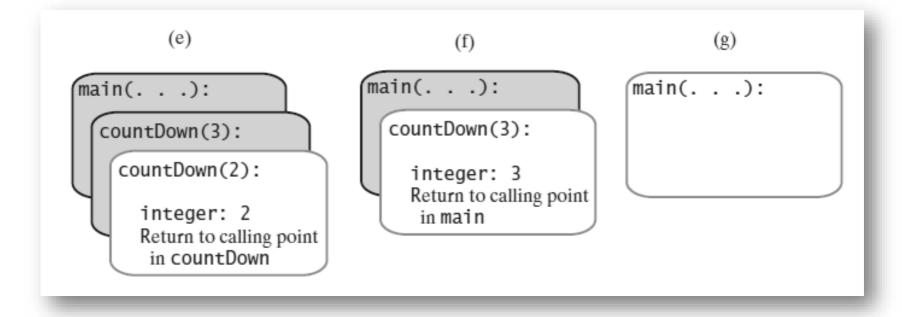
integer: 2
Return to calling point in countDown
```

The stack of activation records during the execution of the call countDown (3)

popping → activation records tear down

The stack of activation records during the execution of the call countDown (3)

popping → activation records tear down



Stack of Activation Records

- Each call to a method generates an activation record
- Recursive method uses more memory than an iterative method
 - Each recursive call generates an activation record
- If recursive call generates too many activation records, could cause stack overflow

$$\sum_{i=1}^{n} i$$

```
/** @param n An integer > 0.
    @return The sum 1 + 2 + ... + n. */
public static int sumOf(int n)
{
```

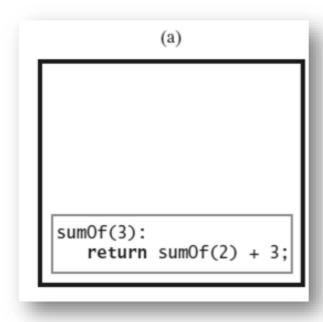
$$\sum_{i=1}^{n} i$$

$$sum = sumOf(n - 1) + n; // Recursive call$$

$$\sum_{i=1}^{n} i$$

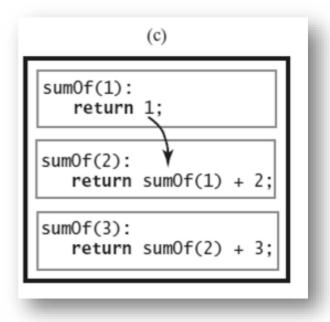
$$\sum_{i=1}^{n} i$$

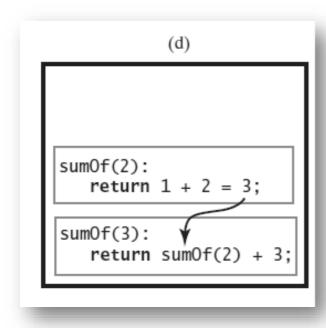
```
/** @param n An integer > 0.
    @return The sum 1 + 2 + ... + n. */
public static int sumOf(int n)
   int sum;
   if (n == 1)
      sum = 1;
                              // Base case
   else
      sum = sumOf(n - 1) + n; // Recursive call
   return sum;
} // end sumOf
```

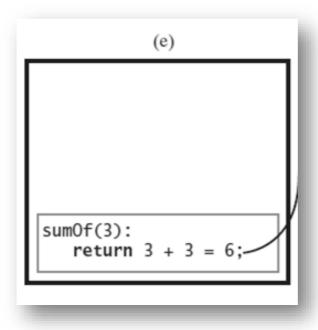


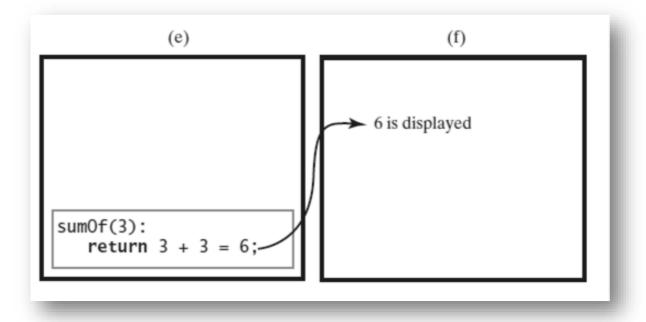
```
sumOf(2):
    return sumOf(1) + 2;

sumOf(3):
    return sumOf(2) + 3;
```









Recursively Processing an Array

A recursive method to display array.

Starting with array[first]

```
public static void displayArray(int array[], int first, int last)
{
    System.out.print(array[first] + " ");
```

Starting with array[first]

displayArray(array, first + 1, last);

Starting with array[first]

What is wrong with this method?

```
public static void displayArray(int array[], int first, int last)
{
    System.out.print(array[first] + " ");

    displayArray(array, first + 1, last);
} // end displayArray
```

Starting with array[first]

We need a base (non-recursive) case!

ask for help only when there is at least one array entry to display otherwise, return

```
public static void displayArray(int array[], int first, int last)
{
    System.out.print(array[first] + " ");
    if (first < last)
        displayArray(array, first + 1, last);
} // end displayArray</pre>
```

Alternatively, ...

```
public static void displayArray(int array[], int first, int last)
{
```

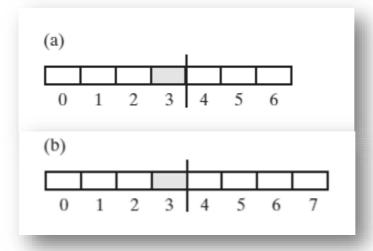
System.out.print (array[last] + " ");

Alternatively, ...

```
public static void displayArray(int array[], int first, int last)
{
    if (first <= last)
    {
        displayArray(array, first, last - 1);
        System.out.print (array[last] + " ");
    } // end if
} // end displayArray</pre>
```

How can we find the middle entry given first and last?

int mid = (first + last) / 2;



```
public static void displayArray(int array[], int first, int last)
{
```

```
int mid = (first + last) / 2;
displayArray(array, first, mid);
```

```
public static void displayArray(int array[], int first, int last)
{
```

```
int mid = (first + last) / 2;
displayArray(array, first, mid);
displayArray(array, mid + 1, last);
```

```
public static void displayArray(int array[], int first, int last)
{
   if (first == last)
      System.out.print(array[first] + " ");
   else
   {
      int mid = (first + last) / 2;
      displayArray(array, first, mid);
      displayArray(array, mid + 1, last);
   } // end if
} // end displayArray
```

```
public static void displayArray(int array[], int first, int last)
{
   if (first == last)
       System.out.print(array[first] + " ");
   else
   {
      int mid = (first + last) / 2;
      displayArray(array, first, mid);
      displayArray(array, mid + 1, last);
   } // end if
} // end displayArray
Consider
first + (last - first) / 2
   Why?
```

Displaying a Bag

Recursive method that is part of an implementation of an ADT is private

```
public void display()
{
    displayArray(0, numberOfEntries - 1);
} // end display

private void displayArray(int first, int last)
{
    System.out.println(bag[first]);
    if (first < last)
        displayArray(first + 1, last);
} // end displayArray</pre>
```