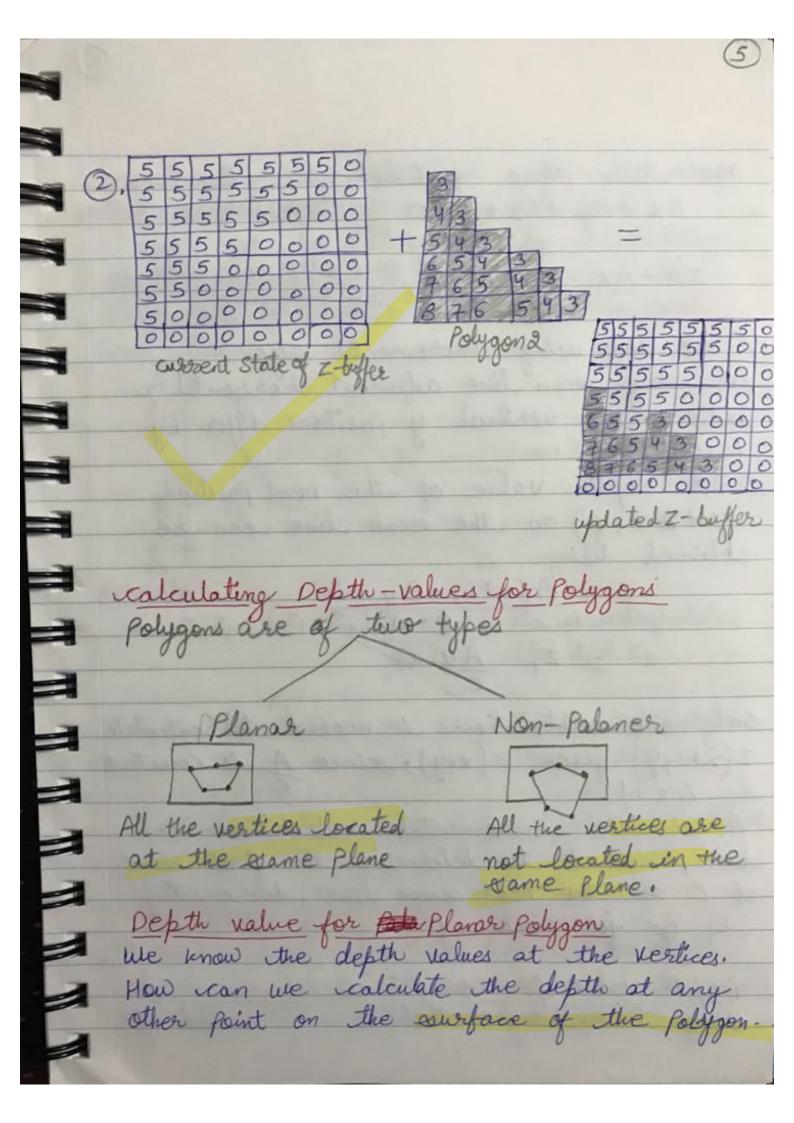
Reference Resource: J.D. Foley, A. Van Dan, Feiner, Hughes, Computer Graphic Principles & Practice, 2nd edition, Publication Addison 3 we do not want to draw surfaces that are hidden. In this chapter we will study three algorithms that compute which ourfoces are hidden, we can then bypass them & down only the ourfaces that are - August Visible. three Algo. Depth-Sort Algo. Z-buffer/ warnocks depth buffer Algo. (1) Z - buffer / Depth Buffer A gorithm
An example to understand the concept Frame Buffer: It stores the information about each & every fixel [--- | screen Array of Pixel x \*y Eg 10001011 1-) on monochrome display In case of colorer display it will have multiple value Z-buffer . Size X\*1 It will extere z-value of each fisel (x, y) we will maintain a z-buffer it will extere the depth of the fixel being displaced at each fix 11 11 11 11 Put in frame buffer. The Z-buffer will be (Polygon 2) with the Previous depth (stored value of Polygon 1). Pivel with greater depth OIt is an image -space approach. 2). Each Surface (Polygon) is processed one fixel fosition at a time. 3). The depth value for fixel are compared I the closest polygon determines the color to be displayed. (4). The Polygons can be exammed in any order (5), It requires that we have available not only a frame buffer t in which color values are actored, but also a z-buffer Z, with the same number of entries, in which a z- Value is stored for each fixel. Algorithm's Steps (1). Intelize frame buffer F to background color. (11). Intialize z-buffer z to zero representing the z-value at the back clipping Plane. 1 The largest value that can be extored in the z-buffer represents the zof the (ii). Scar convert each Polygon in arbitrary order 

(iv). For each (x,y) pixel, calculate depth 'z' at that pixel. (v). compare the calculated value from step (iv) to the value stored in z-buffer. (Vi). If new Z(x,y) > = Zbuff (x,y), update Z-buffer to the new value of update frame (Vii). else do nothing. => The entire process is a search over each set of pairs ? Zi (xyy), Fi(xyy) for fixed x & y, to find the largest zi. The I the frame buffer record the information associated with the largest z encountered thus for each (xyy). Eg. A fixel's shade is shown by its color, z value is shown 5 5 5 5 00000 000 5555 0000 0000 000 0 0000000 0000 5550000 Z-buffer 5 5000 019 intralize to O. 0000 00 00000 0 9 00000 current state of 2-6-ffer



we have plane equation Ax + By + Cz + D = 0 Polygon surface equation Z = -Ax -By -D At each estep calculating z For any scan line adjacent horizontal x positions or vertical y positions differ by unit. The depth value of the next position  $(x+\Delta x, y)$  on the escan line can be obtained using z = -A(x+1) - By - D= Z, - A(D) only one subtraction is needed to calculate Z(x+1,y) given Z(x,y), since A is constant & DX=1. A similar incremental calculation can be performed to determine the first value of Z on the next scan line, decrementing by & for each sy. Depth Value for Non-Planar Polygons Using interpolating method (It is a method of

constructing new data points within the range of discrete set of known data points.) z(x,y) can be determined by interpolating the z co-ordinates of the Polygon's vertices along Pairs of edges and then across each Scan line. y1-ys Za is interpolated y1-y2 6/W Z1& Z2 Zb = Z1 - (Z1-Z3) y1-ys Zb is interpolated Zb - (Zb - Za) 26 - Xp Zp is interpolated Xb - Xa b/w Za & Zb Advantages of Z-buffer Algorithm

(D. The z-buffer also does not require that objects be polygons. It can be used to render any object if a shade and a z-value can

be determined for each point in its Projection. 2). It is easy to implement. 3). It can be saved along with the image and used later to merge in other Objects whoose z can be computed, (4). It Performs radio sort in it and y, requiring no comparisons & its z sort takes only one comparison for pivel for containing that fixel. 3). Total no. of Polygons in a ficture can be large, as it processes one object at a time Disadvantages of Z-buffer Algo 1). It requires a large amount of expace in for Z-buffer. D. It has an aliasing Problem, those fixels whered by the rendered edges may assigned slightly different Z values because of numerical inaccurocies in performing the z interpolation. 3). It is a time - consuming Process. Qwarnock's Algorithm

Area - auddivision algorithm follow the divide and conquer estrategy which estates, if it is easy to decide which folygons are visible in the area, they are diplayed Otherwise the area is subdivided into smaller areas to which the decision logic is applied rewrsively. Os ubdivides each area into four equal exqueres. D. It is a recursine subdivision process. intialarea first dividion second division If the visibility decision cannot be made then 탇 this area is further subdivided either until a visibility decision can be made or until the excreen area is a single fixel. This method is also known as quadtree method. 1. At each stage in recursion, each polygon has one of four Possible relationship with the area of interest (window) -: (1). Surrounding Polygons: completely contain the area of interest. (ii) Intersecting Polygons', Intersects the 

(ii). Contained Polygons: are completely inside the area. (iv). Disjoint Polygons: are completely outside the In four cases, a decision about an area can be made easily, no the area does not need to be further to be conquered: 1. All the Polygons are disjoint from the area. The background color can be diplayed in 2. There is only one intersecting or only one = 1 contained polygon. The area is first filled with the background color and then the fart of the Polygon contained in the area is \_ oscan converted. 3. There is a single surrounding folygon, the area is filled with the color of the surrounding 4. More than one polygon is intersecting, contained in or surrounding the area, but one 

is a sourrounding Polygon that is in front of all other Polygons. Determining whether a surrounding polygon is in front is done by computing the z-co-ordinate of the planes of all surrounding, folygon intersecting and contained polygons at the four Polygon whose four corner z co-ordinates are larger than those of any other polygons, then the entire area is filled with the color of surrounding Polygons. surrounding Polygon X: marks the intersection of surrounding polygon o: 4 4 4 interecting 4 In (a) the four intersections of the surrounding Polygon are all closer to the viemport than are any of the other intersections: the entire 1 area is filled by the surrounding polygon's color.

In (b) no decision can be made ? " on the left the plane of intersecting Polygon exems to be in front of the intersecting polygon exems surrounding Polygon.

Warnock's algo, subdivides the area to simply the Pooblem.

There is no point in examining case I (Surrounding Polygon) & case 4 (Disjoint Polygon) only intersecting and contained Polygons need to be re-examined.

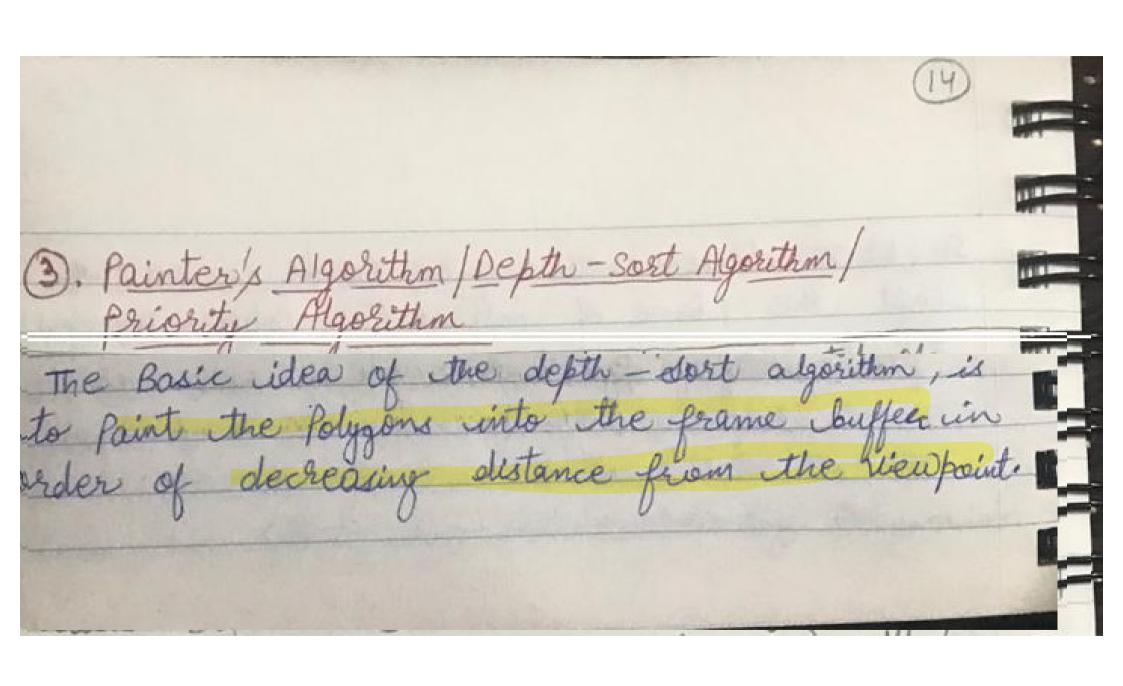
Since it follows the divide and confuer strategy, no parallel computer can be be used to speed up the process.

Disaduatage Test are complex & slow.

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A worst cose ocenario occurs when the case is not one of the four.

Please refer Pg 638 example.



(5) Three conceptual esteps are ferformed !-1. Sort all Polygons according to the smallest (farthest) Z coordinate of each. 2. Resolve any ambiguities this may cause when the Polygon's z extents overlap applitting Polygons 3. Scan convert each polygon in ascending order of semallest z-coordinate. To resolve the ambiguities in setep 2, the algorithm Performs 5 tests. Let the Polygon currently at the far end of the exorted list of Polygons be called P. Before this Polygon is ecan-converted into the frame buffer, it must be tested each Polygon a whose z-exent overlaps the z-extent of P, to Prove I cannot hide of and that Pican is be written before of upto 5 tests are performed in order of Ting complexitiy (The 5 test are in order of test fails enter in testa & seo on. As soon as one succedes, stop Performing the test let a way one of test Pass, Phas been shown not to obscure of the next polygon of overlapping P in Z is tested! If all such folygons pass, then P is scan converted & the next polygon on the list becomes the new P)

Five test (order matters) 1. Do the Polygons' x-extents not overlap?
2. Po the Polygons' y-extents notonerlap? 3. Is P entirely on the opposite side of Or's Plane from the viewpoint? stame side of P's Plane 4. Is of entirely on the as the viewpoint! 111 Viewpoint (Viewpoint 7117 test 3 is false test 4 fails) (: Pis not behind a) but test 4 is true 5. Do the Projections of the Polygons onto the (x) y) plane not overlap? (This ican be determined by comparing the edges.)

If all of the 5 test fails of then we have to do the se ording re-ordering of polygon (P, a) to (a, P), i.e phides of. i. we need to perform test whether gican be exan-converted before P. Test 1, 2 and 5 do not need to be repeated, but new version of test 3 & 4 are used, with order (9,P) 3. Is a entirely on the apposite side of P's Plane from the viewpoint? = z V test 3' pass .. we move of to the end of there is no order the list & it becomes the in which p and of new P. can be excan conjusted correctly. Either for 4. Is pentirely on the same of must be oblit. side of 9's plane as the viewpt? The original unsplit folygon is discarded, its pieces are inserted in Proper z-order & the algorithm proceeds from the start. 

To artoid looping, we modify our approach by marking each polygon that is moved to the end of the list. Then, whenever the first 1111 fire test fails & the current Polygon of is TIL marked, we do not tests 3' and 4". Instead, we split at either for 9 & reinsert the Pieces Worst-case scenario can two polygons fails all the tests even when they are already ordered correctly? IIII only the z-coordinates of P 0.5 0.5 each vertex is shown. The algorithm scan convert f first. "it is on back of Q. Rotate of colochuise until it 107 begin to obscure P, do not interest Zestente, eso they must be compared. P& O. Since, P& O, have overlapping All the test 1,2,3,4,5,3 R4 fails. . the Polygon will be explit, even though Pican be escan converted before a. Disadvantage of Depth-Sort Algo (). Polygons have to be cort first. D. we need to explit polygons to volve cyclic & intersecting objects.