EXP NO: 08 DATE:

GENERATE THREE ADDRESS CODE FOR A SIMPLE PROGRAM USING LEX AND YACC

AIM:

To design and implement a LEX and YACC program that generates three-address code (TAC) for a simple arithmetic expression or program. The program will:

- Recognize expressions like addition, subtraction, multiplication, and division.
- Generate three-address code that represents the operations in a way that could be directly translated into assembly code or intermediate code for a compiler.

ALGORITHM:

1. Lexical Analysis (LEX) Phase:

Input: A string containing an arithmetic expression (e.g., a = b + c * d;).

Output: A stream of tokens such as identifiers (variables), numbers (constants), operators, and special characters (like =, ;, (), etc.).

- 1. Define the Token Patterns:
 - o ID: Identifiers (variables) are strings starting with a letter and followed by letters or digits (e.g., a, b, result). o NUMBER: Constants (e.g., 1, 5, 100).
 - o OPERATOR: Arithmetic operators (+, -, *, /). o ASSIGNMENT: Assignment operator (=). o PARENTHESIS: Parentheses for grouping ((and)).
 - o WHITESPACE: Spaces, tabs, and newline characters (which should be ignored).
- 2. Write Regular Expressions for the Tokens:
 - o ID -> [a-zA-Z_][a-zA-Z0-9_]* o NUMBER -> [0-9]+ o OPERATOR -> [\+\-*/] o ASSIGN -> "=" o PAREN -> [\(\))]
 - o WHITESPACE -> [\t\n]+ (skip whitespace)
- 3. Action on Tokens:
 - o When a token is matched, pass it to YACC using yylval to store the token values.
- 2. Syntax Analysis and TAC Generation (YACC) Phase:

Input: Tokens provided by the LEX lexical analyzer.

Output: Three-address code for the given arithmetic expression.

- 1. Define Grammar Rules:
 - o Assignment:

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statement: ID '=' expr

This means an expression is assigned to a variable.

o Expressions:

bash CopyEdit expr: expr

OPERATOR expr

An expression can be another expression with an operator (+, -, *, /).

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expr:

NUMBER

expr: ID

expr: '(' expr ')'

- 2. Three-Address Code Generation:
 - o For every arithmetic operation, generate a temporary variable (e.g., t1, t2, etc.) to hold intermediate results.
 - o For a = b + c, generate:

ini

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t t1 = b +

ca = t1

o For a = b * c + d, generate:

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t1 = b * c

t2 = t1 + d

a = t2

3. Temporary Variable Management:

 Keep a counter (temp_count) for generating unique temporary variable names

(t0, t1, t2, ...).

o Each time a new operation is encountered, increment the temp_count to generate a new temporary variable.

4. Rule Actions:

o When a rule is matched (e.g., expr OPERATOR expr), generate the TAC and assign temporary variables for intermediate results.

Detailed Algorithm:

- 1. Initialize Lexical Analyzer:
 - o Define the token patterns for ID, NUMBER, OPERATOR, ASSIGN, PAREN, and WHITESPACE.
- 2. Define the Syntax Grammar:
 - o Define grammar rules for:
 - \Box Assignments: ID = expr
 - □ Expressions: expr -> expr OPERATOR expr, expr -> NUMBER, expr -> ID, expr -> (expr)
- 3. Token Matching:
 - LEX: Match input characters against the defined regular expressions for tokens.
 - o YACC: Use the tokens to parse and apply grammar rules.
- 4. TAC Generation:
 - o For Assignment:
 - ☐ Upon parsing ID = expr, generate a temporary variable for the result of expr and assign it to the variable ID. o For Arithmetic Operations:
 - ☐ For each operator (e.g., +, -, *, /), generate temporary variables for intermediate calculations.
- 5. Output TAC:
 - o Print the generated three-address code, with each expression and its intermediate results represented by temporary variables.

PROGRAM:

3address.1

```
%{
#include "3address.tab.h"
#include <string.h>
#include <stdlib.h>
%}
```

ID [a-zA-Z][a-zA-Z0-9]*

```
NUM [0-9]+
%%
{ID} { yylval.str = strdup(yytext); return ID; }
{NUM} { yylval.str = strdup(yytext); return NUM; }
     { return '='; }
     { return ';'; }
"("
     { return '('; }
     { return ')'; }
"+"
    { return '+'; }
"_"
      { return '-'; }
!!*!!
      { return '*'; }
     { return '/'; }
[\t\n]; // skip whitespace
%%
int yywrap() {
return 1; }
3address.y
%{
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
int tempCount = 0;
char* createTemp() {     char* temp =
(char*)malloc(10); sprintf(temp,
"t%d", tempCount++); return
temp;
}
void yyerror(const char* s);
int yylex();
%}
%union {
  char* str;
%token <str> ID NUM
```

```
%type <str> expr
%left '+' '-'
%left '*' '/'
%%
stmt:
   ID '=' expr ';' {
printf("\%s = \%s\n", \$1, \$3);
expr:
   expr'+' expr {
      char* temp = createTemp();
printf("%s = %s + %s\n", temp, $1, $3);
      $ = temp;
  expr'-' expr {
                        char* temp =
                    printf("\%s = \%s -
createTemp();
%s\n", temp, $1, $3);
      $ = temp;
   }
  expr'*' expr {
                        char* temp =
                    printf("%s = %s *
createTemp();
%s\n", temp, $1, $3);
      $ = temp;
  expr'/'expr {
                       char* temp =
                    printf("\%s = \%s /
createTemp();
%s\n", temp, $1, $3);
      $ = temp;
  | '(' expr ')' {
$$ = $2;
   }
  | ID {
      $\$ = strdup(\$1);
  | NUM {
      $\$ = strdup(\$1);
```

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```
%%

void yyerror(const char* s) {
  printf("Syntax Error: %s\n", s);
}

int main() {    printf("Enter an arithmetic expression :\n");    yyparse();    return 0;
}
```

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OUTPUT:

```
yacc -d expr.y lex expr.l gcc
y.tab.c lex.yy.c -o expr_parser
./expr_parse
r a = b * c +
d; t0 = b * c
t1 = t0 + d a
= t1
```

Implementation	
Output/Signature	

RESULT:

Thus the process effectively tokenizes the input, parses it according to defined grammar rules, and generates the corresponding Three-Address Code, facilitating further compilation or interpretation stages.