



# **Finding Maximum Number of Groups With Increasing Length**

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# Introduction to Greedy Technique

## 1 Systematic Approach Approach

The Greedy Technique follows a systematic approach, making locally optimal choices at each stage with the aim of finding a global optimum.

## 2 Immediacy

It focuses on making the best decision at the current moment, without considering the long-term consequences.

## 3 Simple Implementation

The Greedy Technique is relatively straightforward to implement, making it a popular choice for solving complex problems.



DIFFERENT TYPES OF  
GREEDY TECHNIQUE  
REAL TIME EXAMPLES



## Problem Statement:

**Maximum Number of Groups With Increasing Length** You are given a 0-indexed array `usageLimits` of length `n`. Your task is to create groups using numbers from 0 to `n - 1`, ensuring that each number, `i`, is used no more than `usageLimits[i]` times in total across all groups. You must also satisfy the following conditions: Each group must consist of distinct numbers, meaning that no duplicate numbers are allowed within a single group. Each group (except the first one) must have a length strictly greater than the previous group. Return an integer denoting the maximum number of groups you can create while satisfying these conditions. Example 1: Input: `usageLimits = [1,2,5]`

Output: 3 Explanation: In this example, we can use 0 at most once, 1 at most at most twice, and 2 at most five times. One way of creating the maximum maximum number of groups while satisfying the conditions is: Group 1 contains 1 contains the number [2]. Group 2 contains the numbers [1,2]. Group 3 contains the numbers [0,1,2]. It can be shown that the maximum number of number of groups is 3. So, the output is 3..

# Solution:

For usageLimits = [1, 2, 5]:

1. Total Usage Capacity:  $\text{total\_usage} = 1 + 2 + 5 = 8$ .
2. Compute Group Sizes:
  - o For  $k = 1$ ,  $S_1 = 1$ .
  - o For  $k = 2$ ,  $S_2 = 3$ .
  - o For  $k = 3$ ,  $S_3 = 6$ .
  - o For  $k = 4$ ,  $S_4 = 10$ .The total usage capacity of 8 is enough for groups of size up to 3 (since  $S_4 = 10$  exceeds 8).
3. Verify Group Formation:
  - o Verify if you can form 3 groups with sizes 1, 2, and 3, respecting the usage limits. Here, the groups can be formed as follows:
    - Group 1: 1 element.
    - Group 2: 2 elements.
    - Group 3: 3 elements.

This confirms that 3 groups can be formed.



# Code:

```
#include <stdio.h>
```

```
struct Item {
```

```
    int weight;
```

```
    int value;
```

```
};
```

```
int compare(const void a, const void b) {
```

```
    double ratio1 = (double)((((struct Item*)a)->value) / (((struct Item*)a)->weight);
```

```
    double ratio2 = (double)((((struct Item*)b)->value) / (((struct Item*)b)->weight);
```

```
    return ratio2 > ratio1;
```

```
}
```

```
double knapsackGreedy(int capacity, struct Item items[], int n) {
```

```
    qsort(items, n, sizeof(items[0]), compare);
```

```
    int currentWeight = 0;
```

```
    double finalValue = 0.0;
```

```
    for (int i = 0; i < n; i++) {
```

```
        if (currentWeight + items[i].weight <= capacity) {
```

# Output:

```
Output Clear
/tmp/ZG2jqPd9aD.c: In function 'knapsackGreedy':
/tmp/ZG2jqPd9aD.c:12:5: warning: implicit declaration of function 'qsort' [
-Wimplicit-function-declaration]
   12 |     qsort(items, n, sizeof(items[0]), compare);
      |     ^~~~~
/tmp/ZG2jqPd9aD.o
maximum value in Knapsack = 240.00

== Code Execution Successful ==
```

# Advantages of Greedy Technique

## Simplicity

The Greedy Technique is easy to understand and implement, making it a popular choice for problem-solving.

## Speed

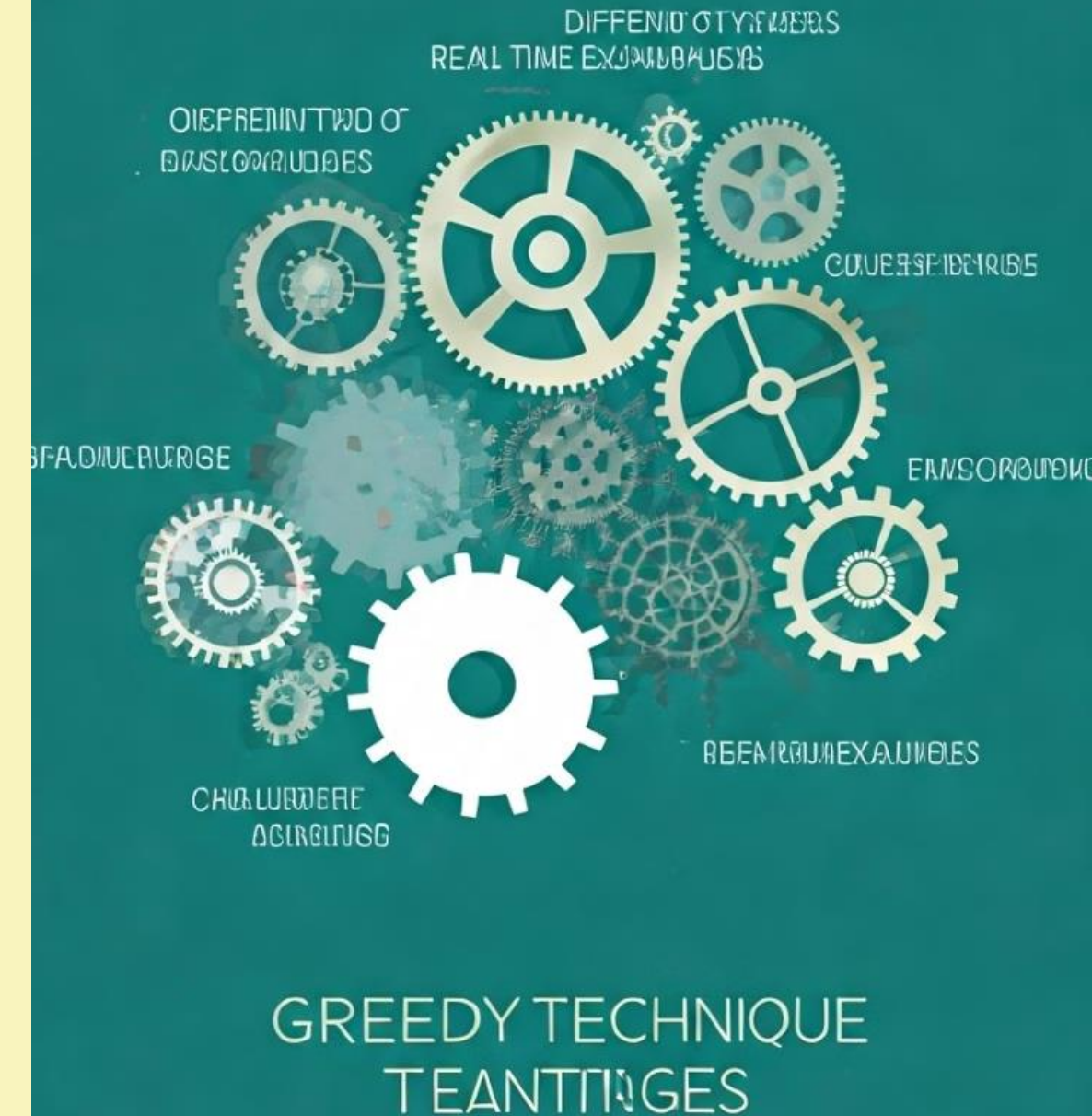
It can quickly find a solution, as it does not require extensive analysis or backtracking.

## Memory Efficiency

The Greedy Technique does not require storing a large amount of data, making it memory-efficient.

## Applicability

It can be applied to a wide range of problems, from scheduling to optimization tasks.



# Disadvantages of Greedy Technique

## No Guarantee of Optimality

The Greedy Technique does not always always find the globally optimal solution, solution, as it focuses on local optimality. optimality.

## Lack of Flexibility

It cannot backtrack or reconsider previous decisions, limiting its ability to ability to adapt to changing situations. situations.

## Potential for Suboptimal Solutions

In some cases, the Greedy Technique Technique may lead to suboptimal solutions that are not the best overall. overall.

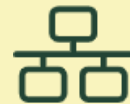


# Applications of Greedy Technique



## Scheduling

Allocating resources and tasks tasks based on immediate needs.



## Network Optimization Optimization

Improving network efficiency efficiency by prioritizing the the most critical connections. connections.



## Resource Allocation

Distributing limited resources resources to maximize immediate benefits.



## Graph Algorithms

Traversing and exploring graphs using the Greedy Technique.



# Greedy Technique Algorithms

1

## Kruskal's Algorithm

Finds the minimum spanning tree of a weighted graph by adding the cheapest available edge at each step.

2

## Dijkstra's Algorithm

Computes the shortest path between nodes in a graph by repeatedly choosing the node with the smallest distance.

3

## Huffman Coding

Constructs an optimal prefix code by repeatedly combining the combining the two least frequent symbols into a new node.



# Implementing Greedy Technique in Projects



1

## Problem Analysis

Identify the problem and its characteristics to determine if the Greedy Technique is an appropriate approach.

2

## Algorithm Selection

Choose the right Greedy Technique algorithm based on the problem requirements and constraints.

3

## Implementation

Carefully implement the selected algorithm, ensuring it meets the project's objectives.

4

## Testing and Optimization

Thoroughly test the implementation and optimize it for efficiency, if necessary.

# Conclusion and Key Takeaways

## 1 Powerful Approach

The Greedy Technique is a powerful problem-solving approach that can be applied to a wide range of problems.

## 2 Balancing Advantages and Disadvantages

Understanding the strengths and limitations of the Greedy Technique is crucial for effective implementation.

## 3 Continuous Learning

Staying up-to-date with the latest Greedy Technique algorithms and their applications is essential for staying competitive.

