

# Binary Search Tree

A binary tree with following property  
(Sometimes a sorted BT)

For each node 'n' with value ( i.e., data present in the node or value of the node or nodes data ) ✓ :

- (i) Value property of all the nodes in the left subtree of node 'n' is smaller than that of value property of node n.
- (ii) Value property of all the nodes in the right subtree of node 'n' is greater than that of value property of node 'n'
- (iii) No duplicates allowed. \*

Operations :-

Traversals :- Like Binary trees

(Inorder, Pre-order, Postorder, etc., traversals)

Search :-

Search a key ( $k$ ) .

(i.e., searching for a node with value property equals to  $k$ )

In the following algorithm

- BSTree represents root
- BSTree . data is representing the data or value property of the root node.
- BSTree . left (BSTree . right) indicates root of left subtree (root of right subtree) of the BSTree

```
Search(BSTree , k )
    if (BSTree is empty)
        return false ;
    endif
    if (BSTree.data == k)
        return true ;
    endif
    if (BSTree.data > k)
        return Search(BSTree.left , k) ;
    else if (BSTree.data < k)
        return Search(BSTree.right , k) ;
    endif
endif .
```

Homework :- Write an iterative algorithm  
for the search operation.

## Searching Minimum.

Minimum ( BSTree )

if ( BSTree.left == NULL )

    return ( BSTree.data );

else

    return Minimum ( BSTree.left );

endif.

end Minimum.

Time Complexity : ( Homework )

( discussed in class )

Searching node with Maximum Value Property.

Maximum ( BSTree )

if ( BSTree . right == NULL )

    return ( BSTree . data )

else

    return Maximum ( BSTree . right )

endif

end Maximum.

Discuss time complexity of above algorithm.

( Discussed in class )

## Adding a new node in Binary Search Tree

New node's  
value properly

```
AddNode ( BSTree , v)
    if ( BSTree == NULL )
        BSTree = CreateTreeNode( v );
        return
    endif
    if ( BSTree.data == v )
        return;
    endif
    if ( BSTree.data > v )
        if ( BSTree.left == NULL )
            BSTree.left = CreateTreeNode( v );
        else
            AddNode ( BSTree.left , v );
        endif
    else if ( BSTree.data < v )
```

```

if (BSTree.right == NULL)
    BSTree.right = CreateTreeNode(v);
else
    AddNode(BSTree.right, v);
endif.
endif.

```

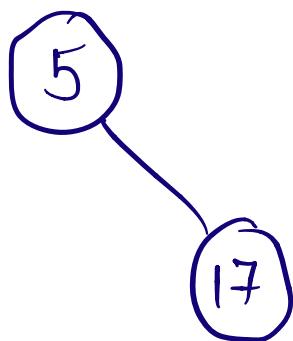
Example :- Assuming Empty tree at the start  
 Construct a Binary Search Tree after the  
 Following set of operations in sequence.

- 1) bst. AddNode( 5 );
- 2) bst. AddNode( 17 );
- 3) bst. AddNode( -1 );
- 4) bst. AddNode( 3 );
- 5) bst. AddNode( 12 );
- 6) bst. AddNode( 5 );

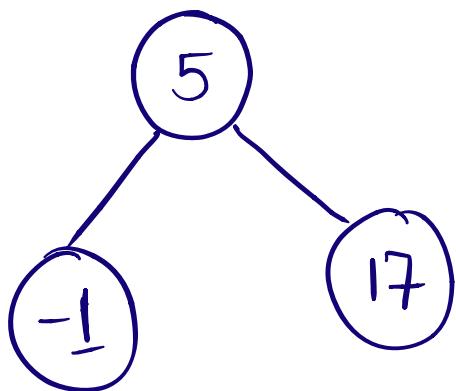
Status After Step 1 .



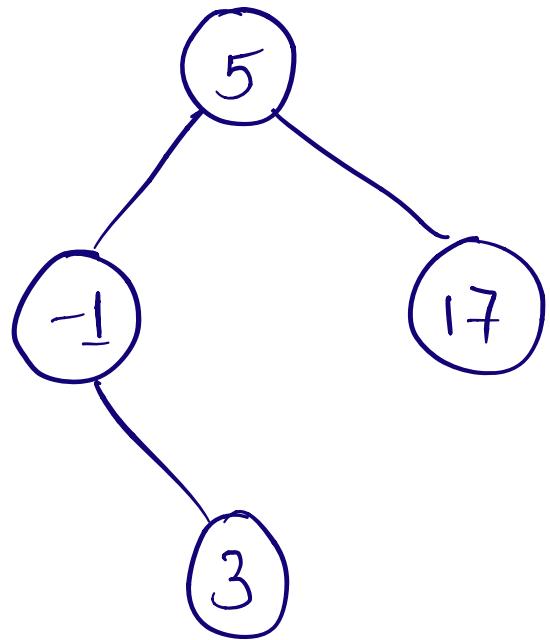
Status After Step 2



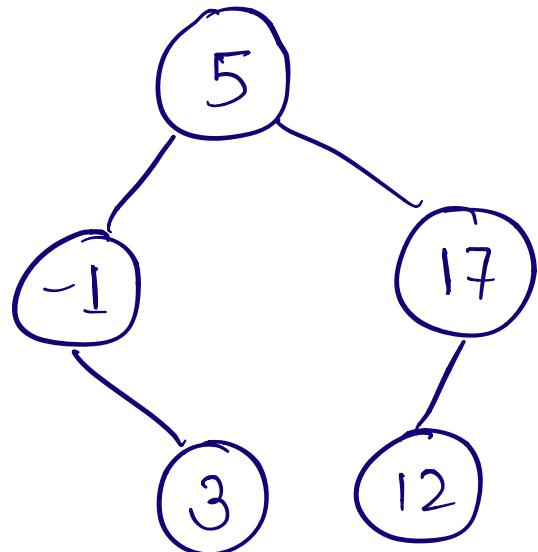
Status After Step 3



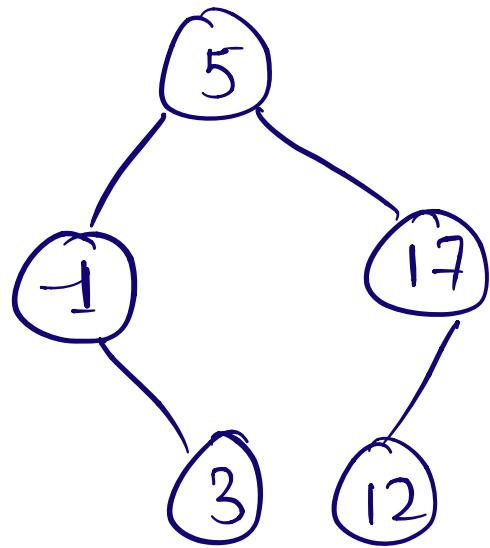
Status After Step 4



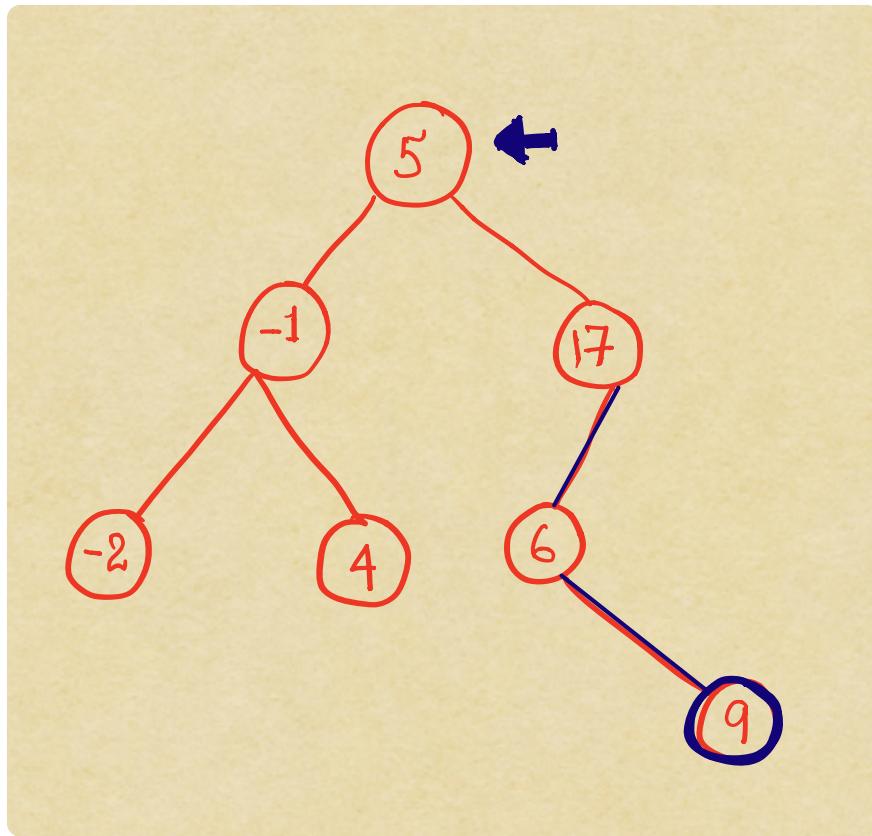
Status After Step 5



Status After Step 6



(Duplicates Not Allowed)



# Succesor

$\text{Succ}(x)$  : Successor of node  $x$

: The node  $y$ , whose data or value property ( $y.data$ ) is the successor of data or value property in sorted order of this tree.

Insider : -2, -1, 4, 5, 6, 9, 17

Pred(4) ← → Succ(5)

Let  $x$  be the node whose successor we are interested in.

Scenario 1:

If  $x$  has right subtree.

$\text{succ}(x) = \text{Minimum } (x \rightarrow \text{right})$

(why ??)

Scenario 2:

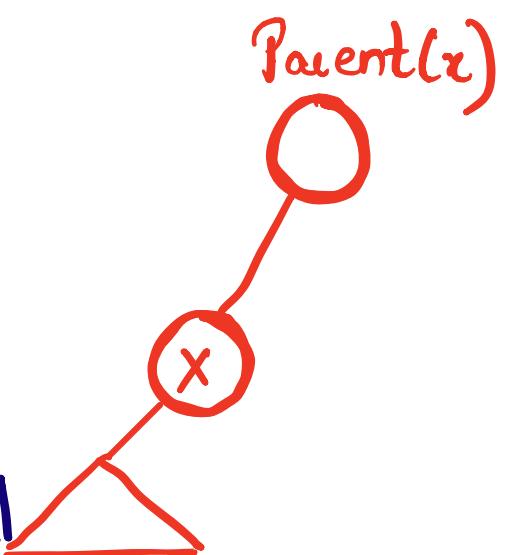
If  $x$  has no right subtree. and  
 $x$  is the left child of the  $\text{parent}(x)$

then

$\text{the succ}(x) = \text{parent}(x)$

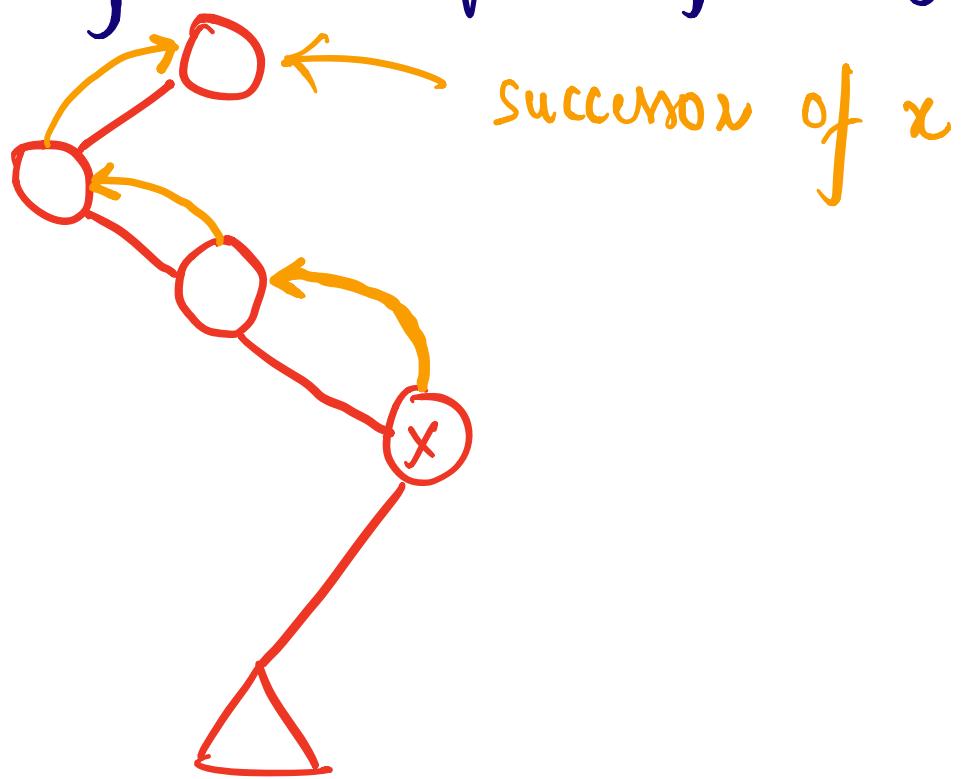
(why ??)

Recall Inorder Traversal



### Scenario 3:

If  $x$  has no right subtree and  $x$  is not the left child of the parent of  $x$ .



Keep moving up from  $x$ , until you find a parent which branches from the left.

Succesor( node )

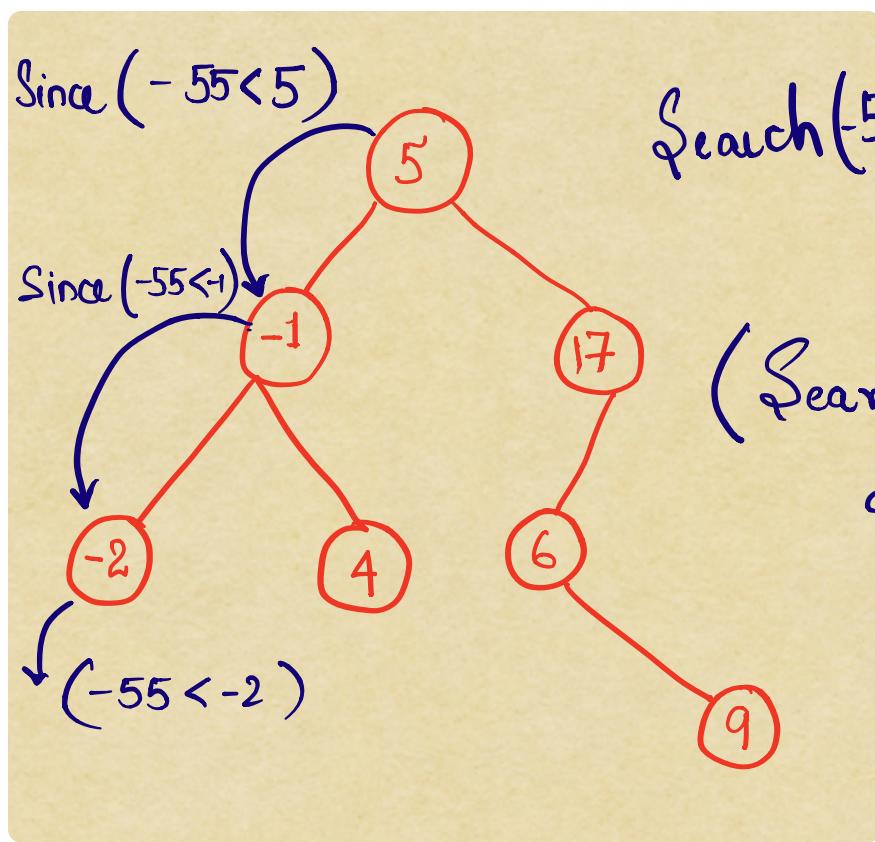
```
if ( node.right != NULL )
    return Minimum( node.right )
endif
temp ← node.parent
while ( temp != NULL &&
        temp.right == node )
    node ← temp ;
    temp ← temp.parent ;
endwhile
return ( temp.data ) ;
```

(Similarly design the algorithm for  
Predecessor)

## Deletion of node with value property v.

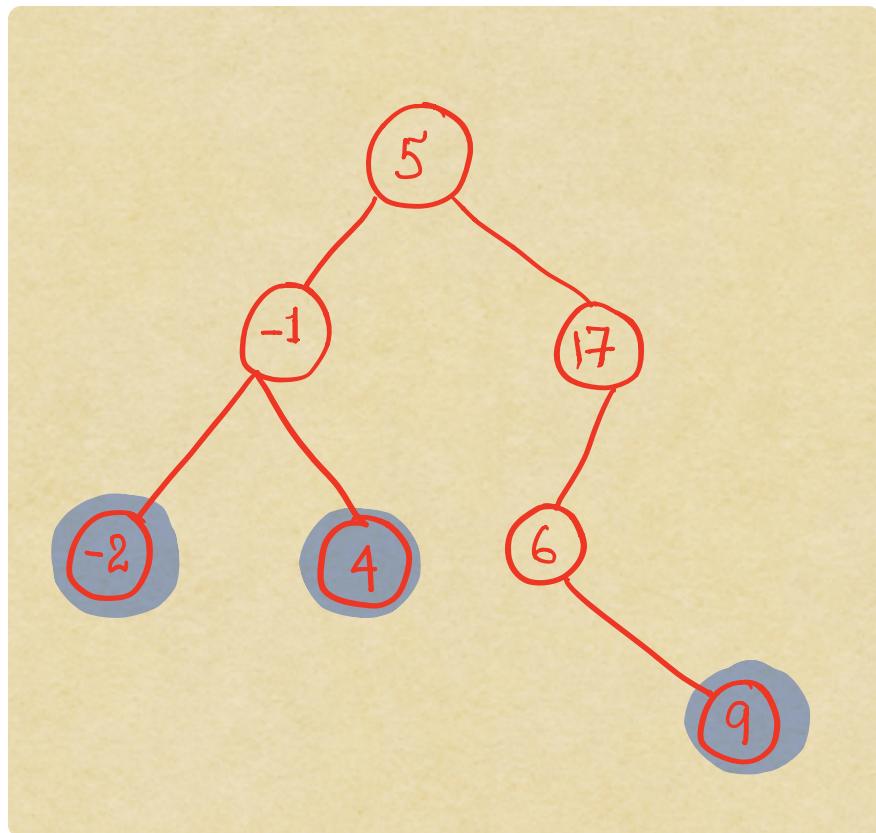
Deletion of a node with value property v without disturbing the value and structure property of binary search tree.

Consider the following tree for demonstration



$\text{delete}(-55)$ : No node with value property  $-55$  is there so nothing will be deleted

"If the node to be deleted is a leaf node"

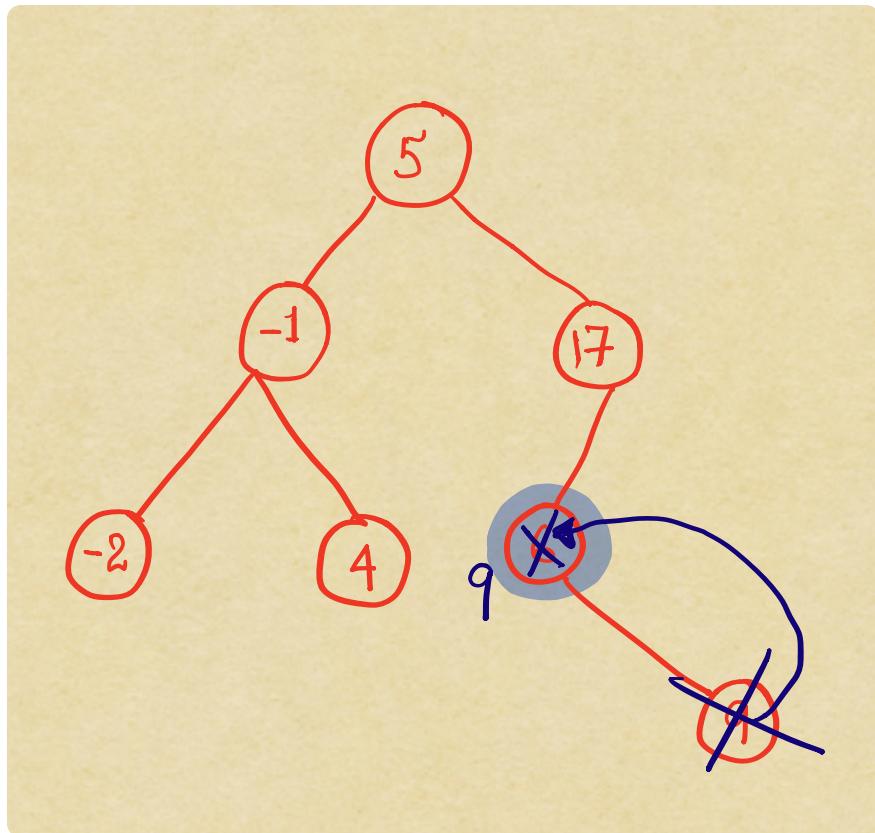


Locate or search the node with the value  
properly, if it is a leaf node (i.e.,  
node's left and right child is NULL)

then

Just Remove the Node from the BST.

If the node to be deleted has exactly one child.



Search the node with the specified value property.  
if the node to be deleted has exactly one child  
then promote the child

Eg :- node to be deleted is 6

`deleteNode ( BSTree , v )`

`if ( BSTree == NULL )`

`return`

`endif.`

`if ( BSTree.data > v )`

`if ( BSTree.left != NULL )`

`deleteNode ( BSTree.left , v );`

`endif`

`return`

`endif`

`if ( BSTree.data < v )`

`if ( BSTree.right != NULL )`

`deleteNode ( BSTree.right , v );`

`endif`

`return`

`endif`

  
searching  
or  
locating  
the node.

```

if ( BSTree.parent == NULL )
    BSTree = NULL ;
    return
endif

if ( BSTree.left == NULL &&
     BSTree.right == NULL )
    if ( BSTree.parent.left == BSTree )
        BSTree.parent.left = NULL
    else
        BSTree.parent.right = NULL
    endif
    return
endif

```

If node to delete is a root node.

If the node to delete is a leaf node

```
if ( BSTree.left != NULL &&  
    BSTree.right == NULL )
```

```
    BSTree.left.parent = BSTree.parent  
    if ( BSTree.parent.left == BSTree )
```

```
        BSTree.parent.left = BSTree.left
```

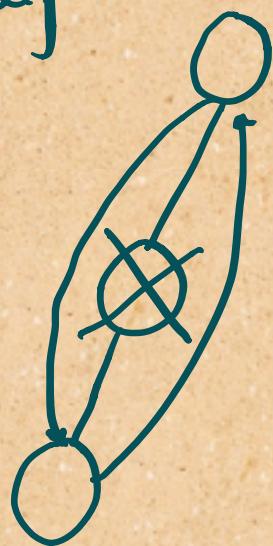
```
    else
```

```
        BSTree.parent.right = BSTree.left
```

```
    endif
```

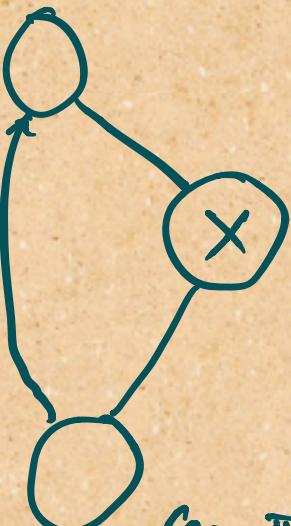
```
    return
```

```
endif
```



(Case I)

If the node  
to delete has one left  
child only



(Case II)

if ( BSTree.left == NULL ||  
    BSTree.right != NULL)

    BSTree.right.parent = BSTree.parent  
    if ( BSTree.parent.left == BSTree )

        BSTree.parent.left = BSTree.right

    else

        BSTree.parent.right = BSTree.right

    endif

    return

endif

If node to delete  
has one right child  
only.

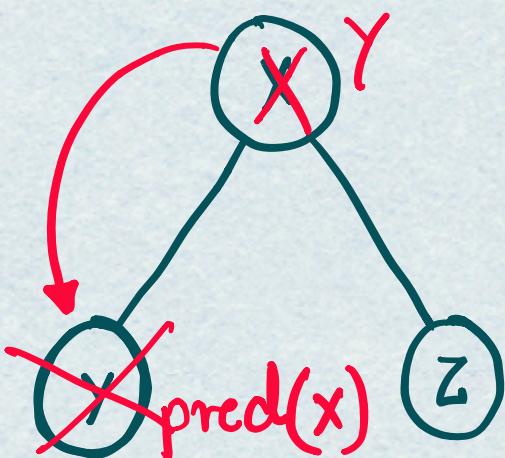
```
if ( BSTree.left != NULL    ||
      BSTree.right != NULL)
  pdata <- pred(BSTree) ; (pred. of the
                           node to be
                           deleted)
```

BSTree.data  $\leftarrow$  pdata

delete ( BSTree.left , pdata )

return

endif.



If node to delte  
has both left and  
right child.

# Complexities of

- Search
- Adding A node
- Deleting a node
- Maximum
- Minimum



Covered in  
class