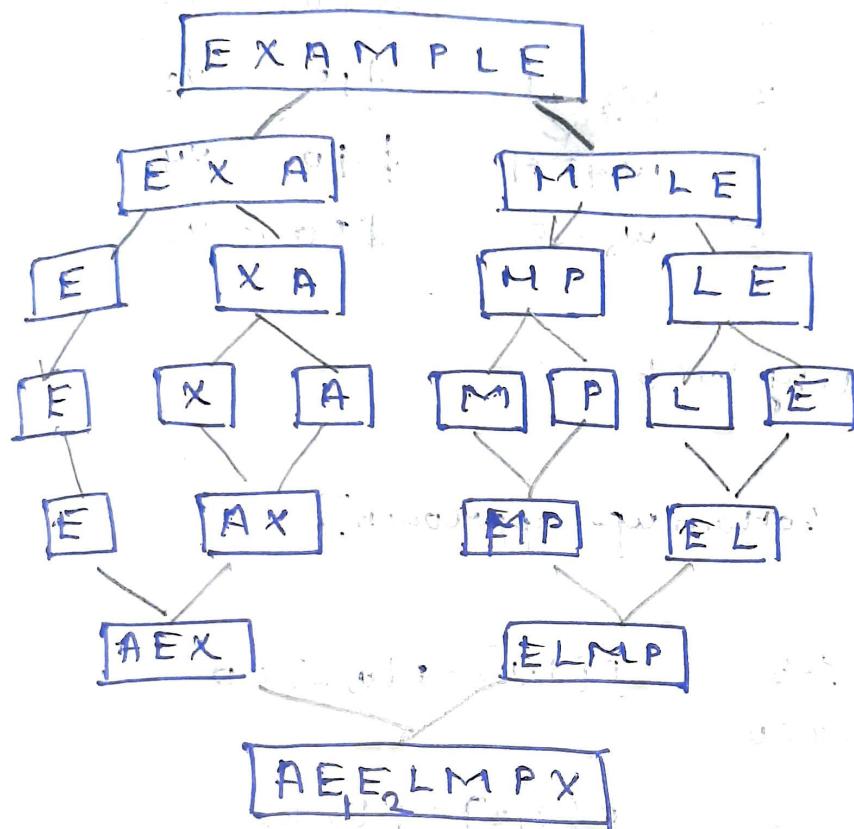


DIVIDE AND CONQUER;

GREEDY APPROACH

- 1) Apply mergesort to sort the list E, X, A, M, P, L, E in alphabetical order.
check if algorithm is stable.



Here, the algorithm is stable because 1st E in EXAMPLE is before 2nd E.

stable algorithm - if 2 objects with equal keys appear in same order in sorted o/p as they appear in i/p array.

2) Apply Quick sort to list A, B, C, D, E
in alphabetical order. check no. of comparisons by choosing

i) First element as pivot

A	B	C	D	E	<u>comparisons</u>
$p = A$	$i = 0$	$j = 4$			
$i = 1$	$B > A$	go to j			1
$j = 4$	$E > A$				2
$j = 3$	$D > A$				3
$j = 2$	$C > A$				4
$j = 1$	$B > A$				5
$j = 0$	$A > A$ (false)				6

since $i > j$, swap $A[j]$ and p

B	C	D	E	
$p = B$	$i = 0$	$j = 3$		
$i = 1$	$C > B$	go to j		7
$j = 3$	$E > B$			8
$j = 2$	$D > B$			9
$j = 1$	$C > B$			10
$j = 0$	$B \leq B$			11

since $i > j$, swap $A[j]$, p

C	D	E	
$p = C$	$i = 0$	$j = 0$	
$i = 1$	$D > C$	go to j	12
$j = 2$	$E > C$		13
$j = 1$	$D > C$		14

$$j=0 \quad C \leq C$$

15

since $i > j$ swap $A[j]$ and P

$$D \leftarrow E$$

$$P = D \quad i=0 \quad j=1$$

$$i=1 \quad E > D \quad \text{go to } j$$

16

$$j=1 \quad E > D \quad \text{so swap}$$

17

$$j=0 \quad D \leq D$$

18

swap $A[j]$ and P

E - No comparison [$\because 18$ element]

\therefore No of comparisons 18

so E is greater than all the elements

ii) Given A, B, C and D. Elements of set E are 1, 2, 3, 4, 5.

$$P > E \quad i=0 \quad j=4$$

comparisons

$$j=0 \quad A \leq E$$

1

$$i=1 \quad B \leq E$$

2

$$i=2 \quad C \leq E$$

3

$$i=3 \quad D \leq E$$

4

$$i=4 \quad E \leq E$$

5

i out of bounds

swap $A[j]$ and P

$$A \quad B \quad C \quad D$$

$$P = D \quad i=0 \quad j=3$$

$$i=0 \quad A \leq D$$

6

$$i=1 \quad B \leq D$$

7

$$i=2 \quad C \leq D$$

8

$$i=3 \quad D \leq D$$

9

i out of bounds, swaps $A[j]$, P

A B C

$p = c$	$i=0 + 1, j=2$	$A \leq C$	$B \leq C$	$C \leq C$	20
$i=0$					11
$i=1$					12
$i=2$					13

i out of bounds, swap $A[i][j], p$

A B

$p = B$	$i=0$	$j=1$	$A \leq B$	$B \leq B$	21
					24

i out of bounds, swap $A[i][j], p$

A - single element (no comparison)

No. of comparisons = $\boxed{24}$

3) Apply strassen's algorithm to compute

$$\begin{bmatrix} 2 & 0 & 2 & 1 \\ 4 & 1 & 1 & 0 \\ 0 & 1 & 3 & 0 \\ 5 & 0 & 2 & 1 \end{bmatrix} * \begin{bmatrix} 0 & 1 & 0 & 1 \\ 2 & 1 & 0 & 4 \\ 2 & 0 & 1 & 1 \\ 1 & 3 & 5 & 0 \end{bmatrix}$$

A:

$$A_{00} = \begin{bmatrix} 2 & 0 \\ 4 & 1 \end{bmatrix} \quad A_{01} = \begin{bmatrix} 2 & 1 \\ 2 & 0 \end{bmatrix} \quad A_{10} = \begin{bmatrix} 0 & 1 \\ 5 & 0 \end{bmatrix}$$

$$A_{11} = \begin{bmatrix} 3 & 0 \\ 2 & 1 \end{bmatrix}$$

7. 15 - 9 operations needed for matrix multiplication

$$b_{00} = \begin{bmatrix} 0 & 1 \\ 2 & 1 \end{bmatrix}, b_{01} = \begin{bmatrix} 0 & 1 \\ 0 & 4 \end{bmatrix}, b_{10} = \begin{bmatrix} 2 & 0 \\ 2 & 3 \end{bmatrix}$$

$$b_{11} = \begin{bmatrix} 1 & 1 \\ 5 & 0 \end{bmatrix}$$

$$M_1 = (a_{00} + a_{11}) * (b_{00} + b_{11})$$

$$= \begin{bmatrix} 1 & 0 \\ 4 & 1 \end{bmatrix} + \begin{bmatrix} 3 & 0 \\ 2 & 1 \end{bmatrix} * \begin{bmatrix} 0 & 1 \\ 2 & 1 \end{bmatrix} + \begin{bmatrix} 1 & 1 \\ 5 & 0 \end{bmatrix}$$

$$= \begin{bmatrix} 1 & 0 \\ 6 & 2 \end{bmatrix} * \begin{bmatrix} 1 & 2 \\ 7 & 1 \end{bmatrix}$$

thus for M_1 :

$$M_1 = (a_{00} + a_{11}) * (b_{00} + b_{11}) = 6 * 2 = 12$$

$$m_2 = (a_{10} + a_{11}) * b_{00} = 8$$

$$m_3 = a_{00} * (b_{01} - b_{11}) = 4$$

$$m_4 = a_{11} * (b_{10} - b_{00}) = 12$$

$$m_5 = (a_{00} + a_{01}) * b_{11} = 4 * 1 = 4$$

$$m_6 = (a_{10} - a_{00}) * (b_{00} + b_{01}) = 2 * 3 = 6$$

$$m_7 = (a_{01} - a_{11}) * (b_{10} + b_{11}) = -2 * 8 = -16$$

$$m_1 + m_4 - m_5 + m_7 = 12 + 12 - 4 - 16$$

$$= 4$$

$$m_3 + m_5 = 8$$

$$m_2 + m_4 = 20$$

$$m_1 + m_3 - m_2 + m_6 = 14$$

thus, $M_1 = \begin{bmatrix} 4 & 8 \\ 20 & 14 \end{bmatrix}$

$$p_1 p_2 = (a_{1,0} + a_{1,1}) * b_{0,0}$$

$$= \begin{bmatrix} 0 & 1 \\ 5 & 0 \end{bmatrix} + \begin{bmatrix} 3 & 0 \\ 2 & 1 \end{bmatrix} * \begin{bmatrix} 0 & 1 \\ 2 & 1 \end{bmatrix}$$

$$= \begin{bmatrix} 3 & 1 \\ 7 & 1 \end{bmatrix} * \begin{bmatrix} 0 & 1 \\ 2 & 1 \end{bmatrix}$$

Für M_2 :

$$m_1 = 4 * 1 = 4 = (a_{0,0} + a_{1,0})(b_{0,0} + b_{1,1})$$

$$m_2 = (a_{1,0} + a_{1,1}) * b_{0,0} = 0$$

$$m_3 = a_{0,0} * (b_{0,1} - b_{1,1}) = 0$$

$$m_4 = a_{1,1} * (b_{1,0} - b_{0,0}) = 2$$

$$m_5 = (a_{0,0} + a_{0,1}) * b_{1,1} = 4$$

$$m_6 = (a_{1,0} - a_{0,0}) + (b_{0,0} + b_{1,1}) = 4$$

$$m_7 = (a_{1,0} - a_{1,1}) * (b_{1,0} + b_{1,1}) = 0$$

$$m_1 + m_4 - m_5 + m_7 = 4 + 2 - 4 + 0 = 2$$

$$m_3 + m_5 = 4$$

$$m_2 + m_4 = 0 + 2 = 2$$

$$\therefore m_1 + m_3 + m_6 - m_2 = 4 + 0 + 0 - 2 = 2$$

$$\boxed{2} = 2$$

$$M_2 = \begin{bmatrix} 2 & 4 \\ 2 & 8 \end{bmatrix}$$

$$\begin{bmatrix} 2 & 1 \\ 2 & 3 \\ 1 & 1 \end{bmatrix}$$

$$M_3 = a_{00} * (b_{01} - b_{11})$$

$$= \begin{bmatrix} 1 & 0 \\ 4 & 1 \end{bmatrix} * \begin{bmatrix} 0 & 1 \\ 0 & 4 \end{bmatrix} - \begin{bmatrix} 1 & 1 \\ 5 & 0 \end{bmatrix}$$

$$= \begin{bmatrix} 1 & 0 \\ 4 & 1 \end{bmatrix} * \begin{bmatrix} -1 & 0 \\ -5 & 4 \end{bmatrix}$$

For M_3 :

$$m_1 = 2 \times 3 = 6$$

$$m_2 = 5 \times -1 = -5$$

$$m_3 = 1 \times -4 = -4$$

$$m_4 = 1 \times -4 = -4$$

$$m_5 = 3 \times -1 = -3$$

$$m_6 = 1 \times 4 = 4$$

$$m_7 = -1 \times -1 = 1$$

$$m_1 + m_4 - m_5 + m_7 = 6 + -4 - -3 + 1 = 6$$

$$m_3 + m_5 = 0$$

$$m_2 + m_4 = -9$$

$$m_1 + m_3 - m_2 + m_6 = 4$$

$$M_3 = \begin{bmatrix} 1 & 0 \\ -9 & 4 \end{bmatrix}$$

$$M_4 = a_{00} * (b_{10} - b_{00})$$

$$= \begin{bmatrix} 3 & 0 \\ 2 & 1 \end{bmatrix} * \begin{bmatrix} 20 \\ 13 \end{bmatrix} - \begin{bmatrix} 0 & 1 \\ 2 & 1 \end{bmatrix} * \begin{bmatrix} 3 & 0 \\ 2 & 1 \end{bmatrix} *$$

$$m_1 = 4 \times 4 = 16$$

$$m_6 = -1 \times 1 = -1$$

$$\begin{bmatrix} 2 & -1 \\ -1 & -2 \end{bmatrix}$$

$$m_2 = 3 \times 2 = 6$$

$$m_7 = -1 \times 1 = -1$$

$$m_3 = 3 \times -3 = -9$$

$$m_1 + m_4 - m_5 + m_7 = 6$$

$$m_4 = 1 \times -3 = -3$$

$$m_3 + m_5 = 3$$

$$m_5 = 3 \times 2 = 6$$

$$m_2 + m_4 = 3$$

$$m_1 + m_3 - m_2 + m_6 = 0$$

$$M_4 = \begin{bmatrix} 6 & -3 \\ 3 & 0 \end{bmatrix}$$

$$M_5 = (a_{00} + a_{01}) * L_{11}$$

$$= \begin{bmatrix} 1 & 0 \\ 4 & 1 \end{bmatrix} + \begin{bmatrix} 2 & 1 \\ 1 & 0 \end{bmatrix} * \begin{bmatrix} 1 & 1 \\ 3 & 0 \end{bmatrix} = \begin{bmatrix} 3 & 1 \\ 3 & 1 \end{bmatrix} + \begin{bmatrix} 1 & 1 \\ 5 & 0 \end{bmatrix}$$

$$m_1 = 4 \times 1 = 4 \quad m_1 + m_4 - m_3 + m_7 = 8$$

$$m_2 = 6 \times 1 = 6 \quad m_3 + m_5 = 3$$

$$m_3 = 3 \times 1 = 3 \quad m_2 + m_4 = 10$$

$$m_4 = 4 \times 1 = 4 \quad m_1 + m_3 + m_2 + m_6 = 5$$

$$m_5 = 4 \times 0 = 0$$

$$m_6 = 2 \times 2 = 4 \quad M_5 = \begin{bmatrix} 1 & 8 \\ 3 & 3 \end{bmatrix}$$

$$m_7 = 0 \times 5 = 0 \quad \begin{bmatrix} 10 & 5 \end{bmatrix}$$

$$M_6 = (a_{10} - a_{00}) + (b_{00} + b_{01})$$

$$= \begin{bmatrix} 0 & 1 \\ 5 & 0 \end{bmatrix} - \begin{bmatrix} 1 & 0 \\ 4 & 1 \end{bmatrix} + \begin{bmatrix} 0 & 1 \\ 2 & 1 \end{bmatrix} + \begin{bmatrix} 0 & 1 \\ 0 & 4 \end{bmatrix}$$

$$= \begin{bmatrix} -1 & 1 \\ 1 & -1 \end{bmatrix} * \begin{bmatrix} 0 & 2 \\ 2 & 5 \end{bmatrix}$$

$$m_1 = -2 \times 5 = -10 \quad m_1 + m_4 - m_3 + m_7 = 2$$

$$m_2 = 0 \times 0 = 0 \quad m_3 + m_5 = 3$$

$$m_3 = -1 \times -3 = 3 \quad m_2 + m_4 = 12$$

$$m_4 = -1 \times 2 = -2$$

$$m_5 = 0 \times 5 = 0 \quad m_1 + m_3 - m_2 + m_6 = -3$$

$$m_6 = 2 \times 2 = 4$$

$$m_7 = 2 \times 7 = 14$$

$$M_6 = \begin{bmatrix} 2 & 3 & 1 \\ -2 & -3 & 1 \end{bmatrix}$$

$$\begin{aligned}
 M_7 &= (a_{01} - a_{11}) * (b_{10} + b_{11}) \\
 &= \begin{bmatrix} 2 & 1 \\ 1 & 0 \end{bmatrix} - \begin{bmatrix} 3 & 0 \\ 2 & 1 \end{bmatrix} \times \begin{bmatrix} 2 & 0 \\ 1 & 3 \end{bmatrix} + \begin{bmatrix} 1 & 1 \\ 5 & 0 \end{bmatrix} \\
 &= \begin{bmatrix} -1 & 1 \\ -1 & -1 \end{bmatrix} \times \begin{bmatrix} 3 & 1 \\ 6 & 3 \end{bmatrix}
 \end{aligned}$$

$$m_1 = -2 \times 6 = -12$$

$$m_2 + m_3 + 2 \times 3 = -6 \quad m_1 + m_4 - m_5 + m_7 = 3$$

$$m_3 = -1 \times -2 = 2$$

$$m_4 = -1 \times 3 = -3$$

$$m_5 = 0 \times 3 = 0$$

$$m_6 = 0 \times 4 = 0$$

$$m_7 = 2 \times 9 = 18$$

$$m_1 + m_3 + m_2 + m_6 = -4$$

$$M_2 = \begin{bmatrix} 3 & 2 \\ -9 & -4 \end{bmatrix}$$

$$\therefore M_1 + M_4 + M_5 + M_7 = \begin{bmatrix} 4 & 8 \\ 20 & 4 \end{bmatrix} + \begin{bmatrix} 6 & -3 \\ 3 & 0 \end{bmatrix}$$

$$= \begin{bmatrix} 8 & 3 \\ 20 & 5 \end{bmatrix} + \begin{bmatrix} 3 & 2 \\ -9 & -4 \end{bmatrix}$$

$$= \begin{bmatrix} 8 & 4 \\ 4 & 5 \end{bmatrix}$$

$$M_3 + M_5 = \begin{bmatrix} -1 & 0 \\ -9 & 4 \end{bmatrix} + \begin{bmatrix} 8 & 5 \\ 10 & 5 \end{bmatrix} = \begin{bmatrix} 7 & 5 \\ 2 & 9 \end{bmatrix}$$

$$M_2 + M_4 = \begin{bmatrix} 2 & 4 \\ 2 & 8 \end{bmatrix} + \begin{bmatrix} 6 & -3 \\ 3 & 0 \end{bmatrix} = \begin{bmatrix} 8 & 1 \\ 5 & 8 \end{bmatrix}$$

$$M_1 + M_3 - M_2 + M_6 = + 7 \quad \text{Ans}$$

$$\begin{bmatrix} 4 & 8 \\ 20 & 4 \end{bmatrix} + \begin{bmatrix} -1 & 0 \\ -9 & 4 \end{bmatrix} - \begin{bmatrix} 2 & 4 \\ 2 & 8 \end{bmatrix} + \begin{bmatrix} 2 & 3 \\ -2 & -3 \end{bmatrix}$$

$$= \begin{bmatrix} 3 & 7 \\ 7 & 7 \end{bmatrix}$$

Ans: $\begin{bmatrix} 5 & 4 & 7 & 3 \\ 4 & 5 & 2 & 9 \\ 8 & 1 & 3 & 7 \\ 5 & 8 & 7 & 7 \end{bmatrix}$

4 Find minimal spanning tree of graph given below using

i) Prim's Algorithm

$$a(-,-) = b(a,3)$$

$$c(a,5)$$

$$d(a,4)$$

$$e(-,\infty)$$

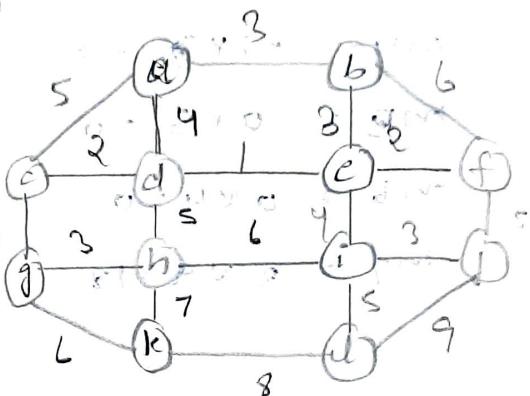
$$f(-,\infty)$$

$$g(-,\infty)$$

$$h(-,\infty) = k(-,\infty)$$

$$i(-,\infty) = l(-,\infty)$$

$$j(-,\infty)$$



2) $b(a,3) = c(a,5) = d(a,4) = e(b,3) =$

$$f(b,6) = g(-,\infty) = h(-,\infty)$$

$$i(-,\infty) = j(-,\infty) = k(-,\infty)$$

$$l(-,\infty)$$

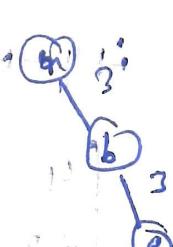


3) $e(b,3) = c(a,5) = d(e,2) = f(2,2)$

$$g(-,\infty) = h(-,\infty) = i(e,4)$$

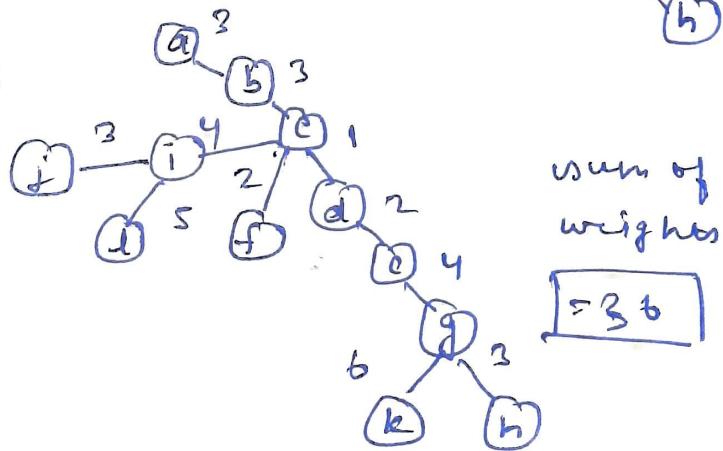
$$j(-,\infty)$$

$$k(-,\infty) = l(-,\infty)$$



- 4) $d(e, 2) = c(d, 2) \cdot f(e, 2) \cdot g(-, \infty) \cdot h(d, 5)$
 $i(e, 4) \cdot j(-, \infty) \cdot k(-, \infty) \cdot l(-, \infty)$
- 5) $c(d, 2) = f(e, 2) \cdot g(e, 4) \cdot h(d, e) \cdot i(e, 4)$
 $j(-, \infty) \cdot k(-, \infty) \cdot l(-, \infty)$
- 6) $f(e, 2) = g(e, 4) \cdot h(d, e) \cdot i(e, 4)$
 $j(f, 5) \cdot k(-, \infty) \cdot l(-, \infty)$
- 7) $g(e, 4) \cdot h(g, 3) \cdot i(e, 4) \cdot j(f, 5) \cdot k(g, 6)$
 $l(-, \infty)$
- 8) $h(g, 3) \cdot i(e, 4) \cdot j(f, 5) \cdot k(g, 6)$
 $l(-, \infty)$
- 9) $i(e, 4) \cdot j(2, 3) \cdot k(g, 6)$
 $l(i, s)$
- 10) $j(2, 3) \cdot k(g, 6) \cdot l(i, s)$
- 11) $l(i, s) \cdot k(g, 6)$
- 12) $k(g, 6)$ - empty

final trace:



ii) Kruskal's Algorithm

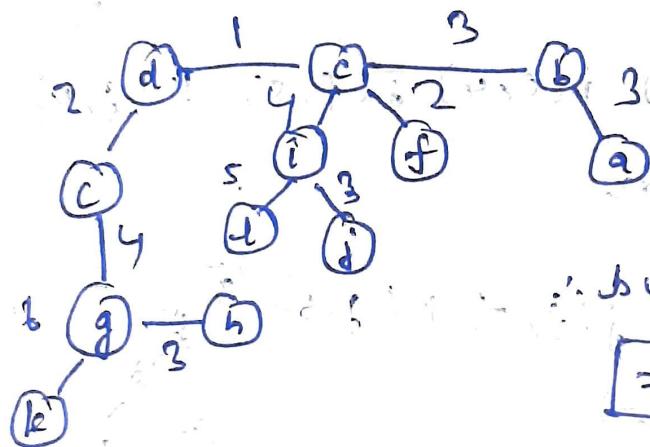
1 2 2 3 3
de cd cf ab eb

3 3 4 4 4
gh ij ad eg ci

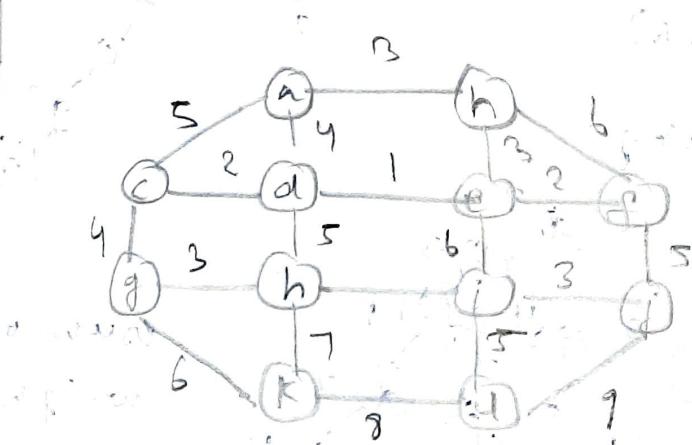
5 5 5 6 6
ac ah fj li bt hi gk

7 8 9

hk rl mf



- 5) Find shortest path in the graph given below considering source vertex as 'a' using dijkstra's algorithm.



$$a(-, 0) \quad b(a, 3) \quad c(a, \epsilon) \quad d(a, 4)$$

$$c(-,\omega) \ f(-,\omega) \ g(-,\omega)$$

$$h(n, \omega) \in (-\infty, \omega] \cup [\omega, \infty) \quad (a)$$

$$k(-y\omega) \cup (-,\omega)$$

$$5(2, 3)$$

$$c(4,5) d(9,4) \in (b, b)$$

$$f(b, q) g(-, \infty) h(-, \infty)$$

i (-, ∞) j (-, ∞) k (-, ∞)

$$d(-, \omega)$$

$$d(a, y)$$

$$c(a,s) \in d(s) f(b,a)$$

$$g(-\infty) \text{ and } h(-\infty)$$

$$f(-,\infty) \models (-,\infty) I (-,\infty)$$

$c(a, s)$

$$e(d,s) \cdot f(b,g) \cdot g(c,s) \cdot h(d,g)$$

$$i(-,\infty) \cup (-,\infty) k(-,\infty) \cup (-,\infty)$$

$$e(d,s)$$

$f(e, \tau) g(l, \tau) h(d, \tau)$ if e, d

$$j(-,\omega) k(-,\omega) + l(-,\omega)$$

$f(e, 7) \quad g(c, 9) \quad h(d, 9) \quad i(c, 9) \quad j(f, 22)$

$$g(e, \alpha) = h(d, \alpha) + f(f_1, \alpha) k(k_1, \alpha)$$

$$h(\alpha, \beta) = i(e, \beta) j(f, \beta)$$

$$k(g, \omega) \perp (-, \omega)$$

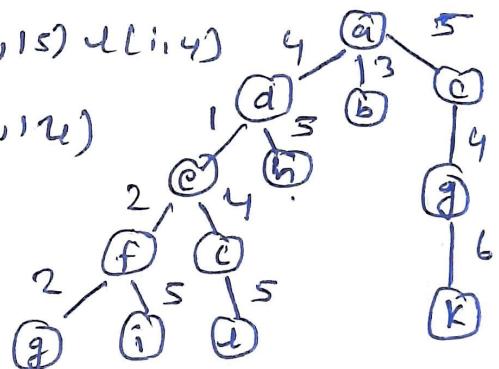
$i(e, q)$

$f(4,12) \neq (9,15)$.

j[f,12]

K(9,15)

$k(g, 15)$



DYNAMIC PROGRAMMING

SOLVING 0/1 KNAPSACK PROBLEM

1)

item	Weight	value	
1	$w_1 = 3$	\$25	v1
2	$w_2 = 2$	\$20	v2
3	$w_3 = 1$	\$15	v3
4	$w_4 = 4$	\$40	v4
5	$w_5 = 5$	\$50	v5

capacity $W = 6$

Soln:

[use bottom up approach]

$$j-w_i = 1-3 \\ = -2 < 0$$

$$F(1,2) = F(0,2) = 0$$

$$j-w_i = 2-3 \\ = -1 < 0$$

$$F(2,2) = F(0,2) = 0$$

$$j-w_i = 3-3 = 0 \\ F(3,3) = \max [F(0,3), F(0,0)] \\ + 25$$

$$j-w_i = 4-3 = 1 \\ F(2,4) = \max [F(0,4), F(0,2) + 25] \\ = 1 > 0$$

$$\text{So ans} = \max [0, 25] = 25$$

$$j-w_i = 5-3 \\ = 2 > 0$$

$$F(1,5) = \max [F(0,5),$$

$$F(0,2) + 25]$$

$$= 25$$

	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	0	0	25	25	25	25
2	0	0	20	25	25	45	45
3	0	15	20	35	40	45	60
4	0	15	20	35	40	55	60
5	0	15	20	35	40	55	65

$$j-w_i = 6-3 \geq 0 \quad F(2,6) = \max[F(0,6), F(0,3)+25]$$

similarly;

$$j-w_i = 1-3 = -2 < 0 \quad F(2,1) = F(0,1) = 0$$

$$j-w_i = 2-2 = 0 \quad F(3,2) = \max[0, 20] = 20$$

$$j-w_i = 3-2 = 1 > 0 \quad F(2,3) = \max[25, 20] = 25$$

$$j-w_i = 4-2 = 2 > 0 \quad F(2,4) = \max[25, 20] = 25$$

$$j-w_i = 5-2 = 3 > 0 \quad F(2,5) = \max[25, 45] = 45$$

$$j-w_i = 6-2 = 4 > 0 \quad F(2,6) = \max[25, 45] = 45$$

$$j-w_i = 1-1 = 0 \quad F(3,1) = \max[0, 15] = 15$$

$$j-w_i = 2-1 = 1 > 0 \quad F(3,2) = \max[20, 15] = 20$$

$$j-w_i = 3-1 = 2 > 0 \quad F(3,3) = \max[25, 35] = 35$$

$$j-w_i = 4-1 = 3 > 0 \quad F(3,4) = \max[25, 40] = 40$$

$$j-w_i = 5-1 = 4 > 0 \quad F(3,5) = \max[45, 40] = 45$$

$$j-w_i = 6-1 = 5 > 0 \quad F(3,6) = \max[45, 60] = 60$$

$$j-w_i = 1-4 = -3 \leq 0 \quad F(4,1) = F(3,2) = 15$$

$$j-w_i = 2-4 = -2 \leq 0 \quad F(4,2) = F(3,2) = 20$$

$$j-w_i = 3-4 = -1 \leq 0 \quad F(4, 3) = F(3, 3) = 35$$

$$j-w_i = 4-4 = 0 \quad F(4, 4) = \max[40, 40] = 40$$

$$j-w_i = 5-4 = 1 \quad F(4, 5) = \max[45, 55] = 55$$

$$j-w_i = 6-4 = 2 \geq 0 \quad F(4, 6) = \max[60, 60] = 60$$

$$j-w_i = -4 \quad F(5, 1) = F(4, 1) = 15$$

$$j-w_i = -3 \quad F(5, 2) = F(4, 2) = 20$$

$$j-w_i = -2 \quad F(5, 3) = F(4, 3) = 35$$

$$j-w_i = -1 \quad F(5, 4) = F(4, 4) = 40$$

$$j-w_i = 0 \quad F(5, 5) = \max[55, 50] = 55$$

$$j-w_i = 1 \quad F(5, 6) = \max[60, 65] = 65$$

Now, find what items are included in knapsack.

capacity = 6 - 5 = 1

$$\text{i)} F(5, 6) > F(4, 6) \Rightarrow 65 > 60 \Rightarrow \text{item 5}$$

Hence, 5th item included.

$$\text{ii)} \text{capacity} = 6-5 = 1$$

$$F(5, 1) = F(4, 1) = 15$$

$$F(4, 1) = F(3, 1) = 15$$

So consider $F(3, 1) > F(2, 2) \Rightarrow 15 > 0$

Hence, 3rd item is included.

Now capacity is 1 - 1 = 0 (knapsack is full)

Therefore, 3rd & 5th items are used
to fill knapsack and most optimal
solution is $F(5, b)$ with value of 65.

2)

To find Transitive closure (Warshall's Alg)

$$R_0 = A = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$R_1 = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$R_{\text{ref}} = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$R_3 = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$R_4 = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

3) To solve : All pairs shortest path (Floyd's Algorithm)

$$D_0 = \begin{bmatrix} 0 & 2 & \infty & 1 & 8 \\ 6 & 0 & 3 & 2 & \infty \\ \infty & \infty & 0 & 4 & \infty \\ \infty & \infty & 2 & 0 & 3 \\ 3 & \infty & \infty & \infty & 0 \end{bmatrix}$$

$$D_0 = \begin{bmatrix} 0 & 2 & \infty & 1 & 8 \\ 6 & 0 & 3 & 2 & \infty \\ \infty & \infty & 0 & 4 & \infty \\ \infty & \infty & 2 & 0 & 3 \\ 3 & \infty & \infty & \infty & 0 \end{bmatrix}$$

$$D_{1,1} = \begin{bmatrix} 0 & 2 & \infty & 1 & 8 \\ 6 & 0 & 3 & 2 & 14 \\ \infty & \infty & 0 & 4 & \infty \\ \infty & \infty & 2 & 0 & 3 \\ 3 & 5 & \infty & 4 & 0 \end{bmatrix}$$

$$D_{2,2} = \begin{bmatrix} 0 & 2 & 5 & 1 & 8 \\ 6 & 0 & 3 & 2 & 14 \\ \infty & \infty & 0 & 4 & \infty \\ \infty & \infty & 2 & 0 & 3 \\ 3 & 5 & 1 & 4 & 0 \end{bmatrix}$$

$$D_3 = \begin{bmatrix} 0 & 2 & 5 & 1 & 8 \\ 6 & 0 & 3 & 2 & 14 \\ \infty & \infty & 0 & 4 & \infty \\ \infty & \infty & 2 & 0 & 3 \\ 3 & 5 & 8 & 4 & 0 \end{bmatrix}$$

$$D_{4,2} = \begin{bmatrix} 0 & 2 & 3 & 2 & 4 \\ 6 & 0 & 3 & 2 & 5 \\ \infty & \infty & 0 & 4 & 7 \\ \infty & \infty & 2 & 0 & 3 \\ 3 & 5 & 6 & 4 & 0 \end{bmatrix}$$

$$D_S = \begin{bmatrix} a & b & c & d & e \\ a & 0 & 2 & 3 & 1 & 4 \\ b & 6 & 0 & 3 & 2 & 5 \\ c & 10 & 12 & 0 & 4 & 7 \\ d & 6 & 3 & 2 & 0 & 3 \\ e & 3 & 5 & 6 & 4 & 6 \end{bmatrix}$$

Step 1: Find the transpose of matrix D_S

$$D_S^T = \begin{bmatrix} a & b & c & d & e \\ a & 0 & 6 & 10 & 3 \\ b & 2 & 0 & 12 & 3 \\ c & 3 & 3 & 0 & 6 \\ d & 3 & 2 & 4 & 4 \\ e & 4 & 5 & 7 & 6 \end{bmatrix}$$

Step 2: Find the transpose of D_S^T

$$D_S^{TT} = \begin{bmatrix} a & b & c & d & e \\ a & 0 & 2 & 3 & 1 & 4 \\ b & 6 & 0 & 3 & 2 & 5 \\ c & 10 & 12 & 0 & 4 & 7 \\ d & 6 & 3 & 2 & 0 & 3 \\ e & 3 & 5 & 6 & 4 & 6 \end{bmatrix}$$