

Unique Length-3 Palindromic Subsequences

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Introduction

In what follows, several methods are presented. The intention is to demonstrate the versatility of **Rust's** syntax and, in doing so, provide a concise refresher.

Statement: Given a string `s`, return the number of unique palindromes of length three that are subsequences of `s`.

Note that even if there are multiple ways to obtain the same subsequence, it is counted only once.

A palindrome is a string that reads the same forwards and backwards.

A subsequence of a string is a new string generated from the original string by deleting zero or more characters without changing the relative order of the remaining characters.

For example, "ace" is a subsequence of "abcde".

<https://leetcode.com/problems/unique-length-3-palindromic-subsequences/description>

Naive Solution [AC]

The idea is simple, as the problem significantly constrains the solution space by limiting candidates to palindromes of length three. Consequently, the relevant patterns reduce to subsequences of the form "`a_a`".

Interestingly, `String` in **Rust** is UTF-8 encoded and stored contiguously in memory (similar to `Java`). It also exposes a comparable mechanism for traversal via `.chars()`, analogous to `Java`'s `toCharArray()`. The underlying design philosophy shares certain characteristics with that of `Erlang`.

N.B. Every code block has been formatted using `rustfmt`.

```
1 impl Solution {
2     pub fn count_palindromic_subsequence(s: String) -> i32 {
3         let n = s.len();
4         let sc = s.chars().collect::<Vec<char>>();
5         let mut res = 0;
6
7         for left in 0..n {
8             if sc[0..left].contains(&sc[left]) {
9                 continue;
10            }
11
12            let mut right = n - 1;
13            while right > left && sc[right] != sc[left] {
14                right -= 1;
15            }
16
17            let mut seen = vec![0; 26];
18            for i in (left + 1)..right {
```

```

19         let j = (sc[i] as u8 - b'a') as usize;
20         if seen[j] == 0 {
21             res += 1;
22         }
23         seen[j] += 1;
24     }
25 }
26
27     return res;
28 }
29 }
```

In the worst case, this approach runs in $O(n^2)$ time. While the logic is correct, the implementation is not optimal. A more structured version of the same idea, shown below, uses precomputed indices along with a `HashSet` to improve clarity. Its worst-case time complexity is still $O(n^2)$, though membership checks benefit from the expected $O(1)$ performance of hashing.

```

1 impl Solution {
2     pub fn count_palindromic_subsequence(s: String) -> i32 {
3         let sb = s.as_bytes();
4
5         let mut left = vec![-1; 26];
6         let mut right = vec![-1; 26];
7
8         for (i, &ch) in sb.iter().enumerate() {
9             let j = (ch - b'a') as usize;
10            if left[j] == -1 {
11                left[j] = i as i32;
12            }
13            right[j] = i as i32;
14        }
15
16        let mut res = std::collections::HashSet::new();
17
18        for i in 0..26 {
19            if left[i] != -1 && left[i] != right[i] {
20                for mid in (left[i] + 1) as usize..right[i] as usize {
21                    res.insert((i as u8 + b'a', sb[mid]));
22                }
23            }
24        }
25
26        return res.len() as i32;
27    }
28 }
```

A more refined version precomputes frequencies from the right.

```

1 impl Solution {
2     pub fn count_palindromic_subsequence(s: String) -> i32 {
3         let n = s.len(); // safe for lowercase ASCII; otherwise use s.as_bytes().
4         let sb = s.as_bytes();
5
6         let mut right_freq = vec![0; 26];
7         for &b in sb {
8             right_freq[(b - b'a') as usize] += 1;
9         }
10
11         let mut seen = vec![false; 26];
12         let mut res = std::collections::HashSet::new();
13
14         for j in 0..n {
```

```

15     let left = (sb[j] - b'a') as usize;
16     right_freq[left] -= 1;
17
18     for c in 0..26 {
19         if seen[c] && right_freq[c] > 0 {
20             res.insert((c as u8, left as u8));
21         }
22     }
23     seen[left] = true;
24 }
25
26
27     return res.len() as i32;
28 }
29 }
```

Using Bitmask [AC]

Performance can be improved by using a `bitmask` instead of a `HashSet`, thereby avoiding hashing overhead and heap allocations.

```

1 impl Solution {
2     pub fn count_palindromic_subsequence(s: String) -> i32 {
3         let n = s.len();
4         let sb = s.as_bytes();
5
6         let mut left = vec![n; 26];
7         let mut right = vec![0; 26];
8
9         for i in 0..n {
10             let j = (sb[i] - b'a') as usize;
11             left[j] = left[j].min(i);
12             right[j] = i;
13         }
14
15         let mut res = 0;
16
17         for c in 0..26 {
18             if left[c] >= right[c] {
19                 continue;
20             }
21             let mut mask: u32 = 0;
22
23             for i in (left[c] + 1)..right[c] {
24                 let mid = (sb[i] - b'a') as u32;
25                 mask |= 1 << mid;
26             }
27
28             res += mask.count_ones() as i32;
29         }
30
31         return res;
32     }
33 }
```

Using `bitset-style struct` [AC]

This version uses a small bitset-style `struct` instead of a raw `u32` mask. This improves semantic clarity and more closely resembles `std::bitset` in C++.

```

1 #[derive(Clone, Copy)]
2 struct BitSet26 {
3     bits: u32,
4 }
5
6 impl BitSet26 {
7     fn new() -> Self {
8         Self { bits: 0 }
9     }
10
11    fn set(&mut self, idx: usize) {
12        self.bits |= 1 << idx;
13    }
14
15    fn count(&self) -> u32 {
16        self.bits.count_ones()
17    }
18 }
19
20 impl Solution {
21     pub fn count_palindromic_subsequence(s: String) -> i32 {
22         let n = s.as_bytes().len();
23         let bytes = s.as_bytes();
24
25         let mut first = vec![n; 26];
26         let mut last = vec![0; 26];
27
28         for i in 0..n {
29             let j = (bytes[i] - b'a') as usize;
30             first[j] = first[j].min(i);
31             last[j] = i;
32         }
33
34         let mut res = 0;
35
36         for c in 0..26 {
37             if first[c] >= last[c] {
38                 continue;
39             }
40
41             let mut seen = BitSet26::new();
42
43             for i in (first[c] + 1)..last[c] {
44                 let mid = (bytes[i] - b'a') as usize;
45                 seen.set(mid);
46             }
47
48             res += seen.count() as i32;
49         }
50
51         return res;
52     }
53 }

```

Functional Approach [AC]

This version is intentionally written in a functional style and visually resembles solutions commonly written in functional languages such as Erlang.

```

1 impl Solution {
2     pub fn count_palindromic_subsequence(s: String) -> i32 {
3         let bs = s.as_bytes();
4

```

```

5     let (first, last) = s.iter().enumerate().fold(
6         (vec![usize::MAX; 26], vec![0usize; 26]),
7         |(mut left, mut right), (i, &b)| {
8             let j = (b - b'a') as usize;
9             left[j] = left[j].min(i);
10            right[j] = i;
11            (left, right)
12        },
13    );
14
15    (0..26)
16        .filter_map(|ch| {
17            let left = first[ch];
18            let right = last[ch];
19
20            if left < right {
21                Some(
22                    (left + 1..right)
23                        .map(|i| (bs[i] - b'a') as usize)
24                        .collect::<std::collections::HashSet<usize>>()
25                        .len(),
26                )
27            } else {
28                None
29            }
30        })
31        .sum::<usize>() as i32
32    }
33 }

```