CS 480

Introduction to Artificial Intelligence

September 6, 2022

Announcements / Reminders

- Please follow the Week 02 To Do List instructions
- Quiz #01 will be posted on Friday
 - a separate announcement will be sent to you to let you know that it is available
 - you will have a week or more to complete it
- Written Assignment #01 will be posted this week
- Midterm Exam (consider fixed):
 - October 13th, 2022 during (Thursday) lecture time

Plan for Today

Problem Solving: Searching

Problem-Solving / Planning Agent

- Context / Problem:
 - correct action is NOT immediately obvious
 - a plan (a sequence of actions leading to a goal) may be necessary
- Solution / Agent:
 - come up with a computational process that will search for that plan
- Planning Agent:
 - uses factored or structured representations of states
 - uses searching algorithms

Planning: Environment Assumptions

Works with a "Simple Environment":

- Fully observable
- Single agent (for now -> it can be multiagent)
- Deterministic
- Static
- Episodic
- Discrete
- Known to the agent

Problem-Solving Process

- Goal formulation:
 - adopt a goal (think: desirable state)
 - a concrete goal should help you reduce the amount of searching
- Problem formulation:
 - an abstract representation of states and actions
- Search:
 - search for solutions within the abstract world model
- Execute actions in the solution

Planning: Environment Assumptions

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- Fully observable
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Important and helpful:

Such assumptions **GUARANTEE** a

FIXED sequence of actions as a

solution

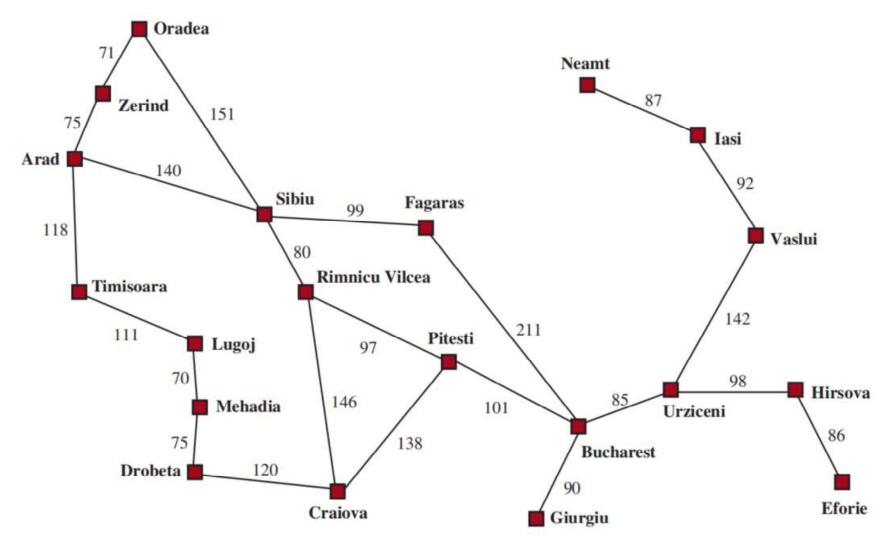
What does it mean?

You can execute the "plan" without worrying about incoming percepts (open-loop control)

Defining Search Problem

- Define a set of possible states: State Space
- Specify Initial State
- Specify Goal State(s) (there can be multiple)
- Define a FINITE set of possible Actions for EACH state in the State Space
- Come up with a Transition Model which describes what each action does
- Specify the Action Cost Function: a function that gives
 the cost of applying action a in state s

Sample Problem: Dracula's Roadtrip



Problem: Get from Arad to Bucharest efficiently (for example: quickly or cheaply).

Search Problem: Dracula's Roadtrip

State Space: a map of Romania

Initial State: Arad

Goal State: Bucharest

Actions:

for example: ACTIONS(Arad) = {ToSibiu,ToTimisoara,ToZerind}

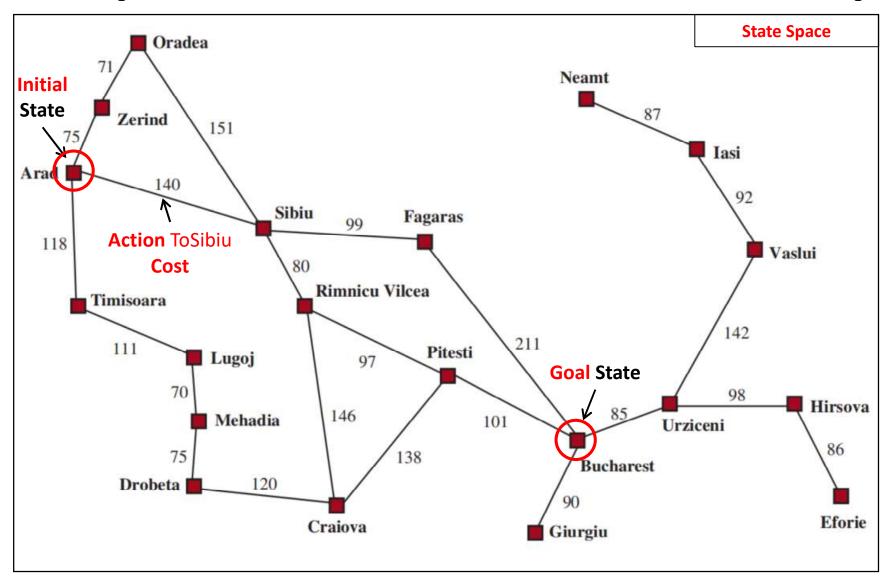
Transition Model:

for example: RESULT(Arad, ToZerind) = Zerind

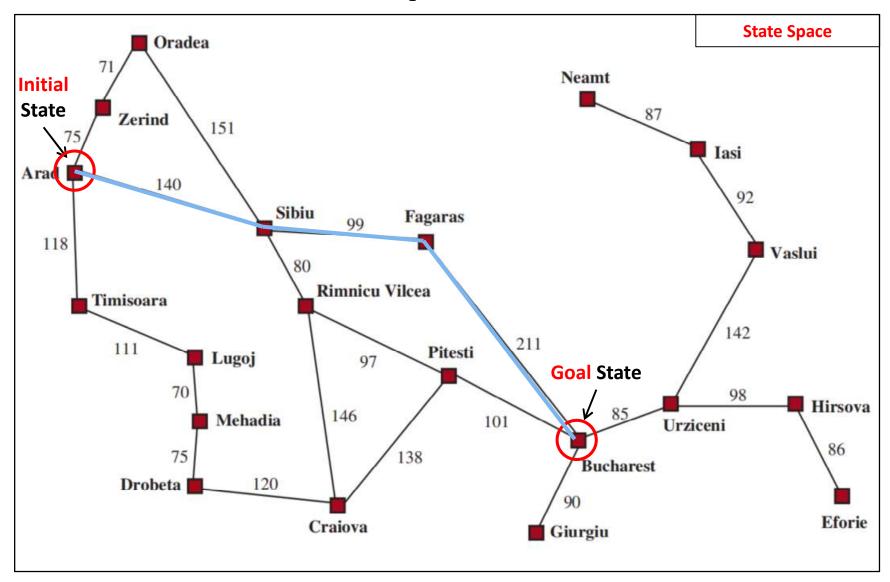
Action Cost Function [ActionCost(S_{current}, a, S_{next})]

for example: ActionCost(Arad, ToSibiu, Sibiu) = 140

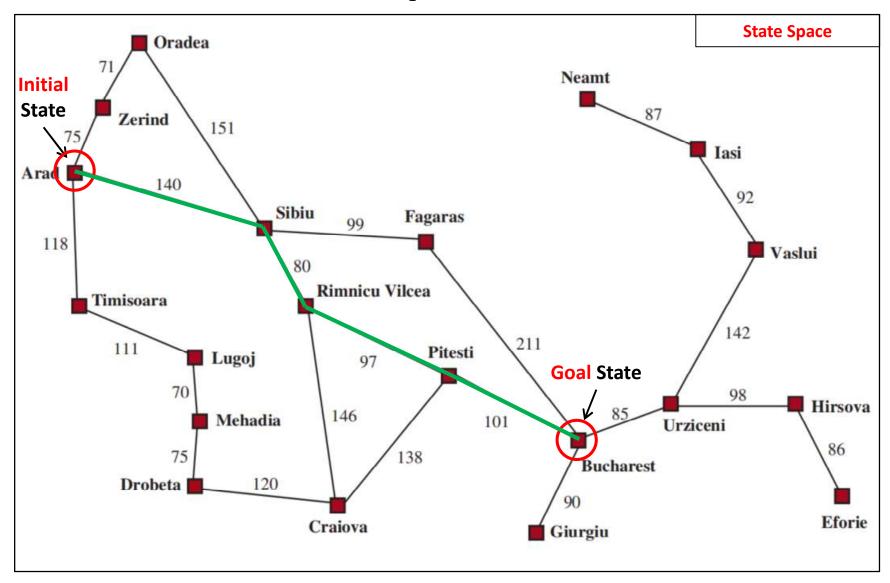
Sample Problem: Dracula's Roadtrip



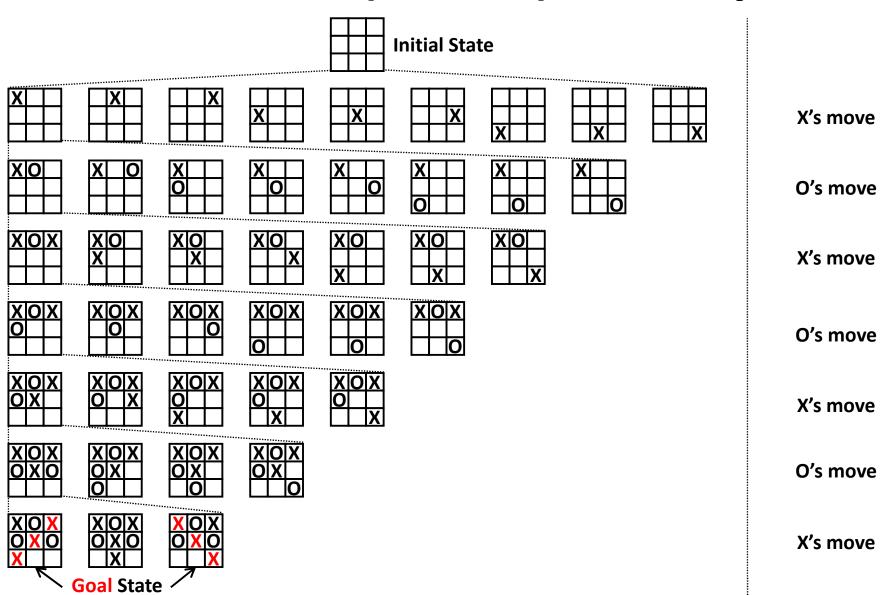
Dracula's Roadtrip: Potential Solution



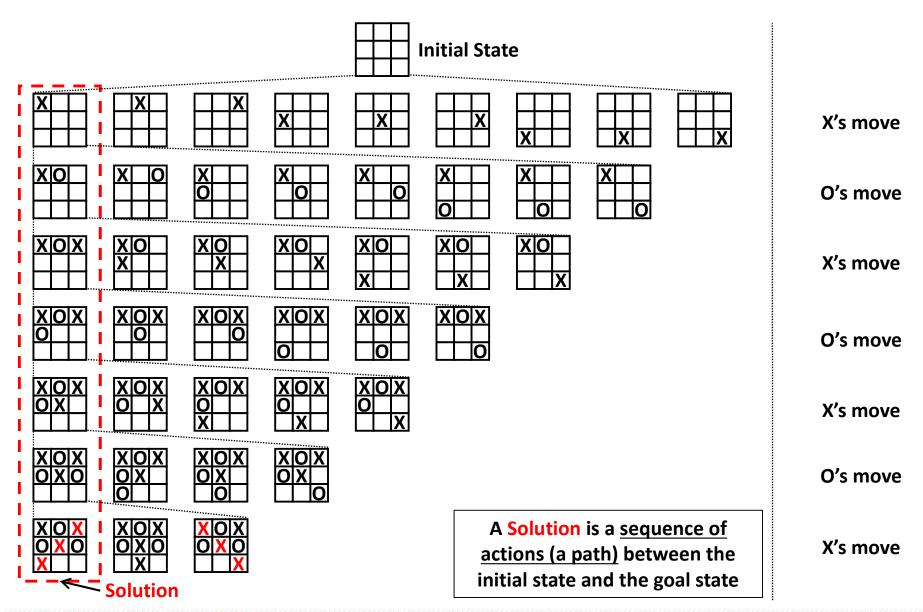
Dracula's Roadtrip: Potential Solution



Tic Tac Toe: (Partial) State Space



Tic Tac Toe: Solution



Chess: (First Move) State Space

Initial
State

AAAAAAA

AAAAA











20 Possible legal <u>first</u> moves: 16 pawn moves 4 knight moves

































Designing the Searching Problem

Analyze and define the Problem / Task

Model and build the State Space

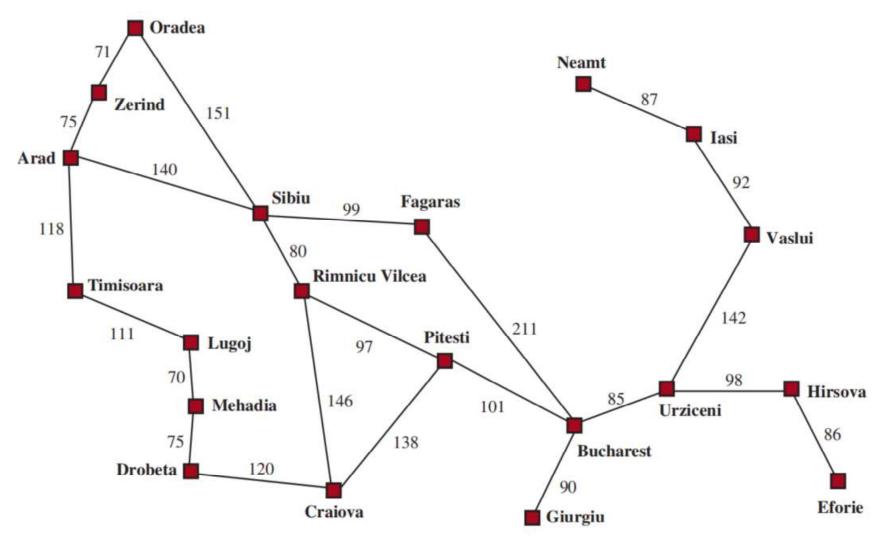
Select searching algorithm

Search

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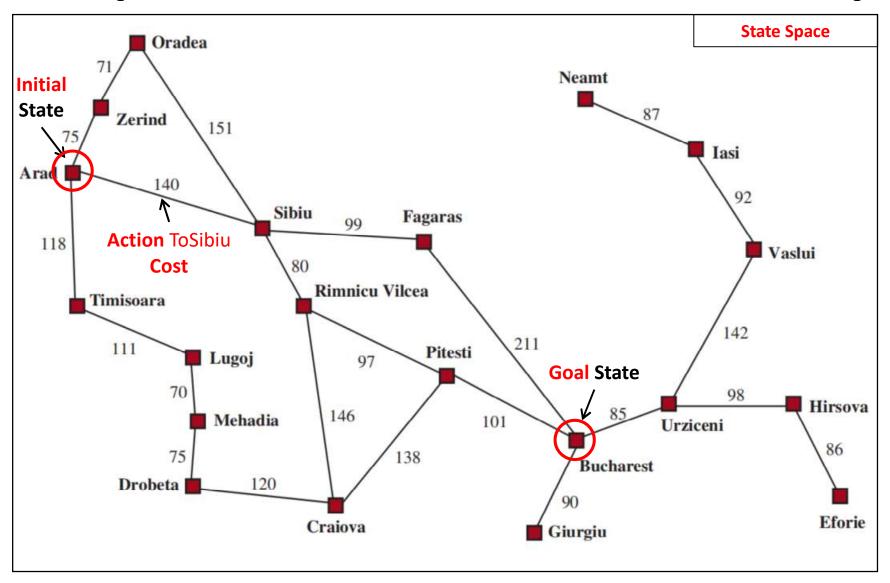
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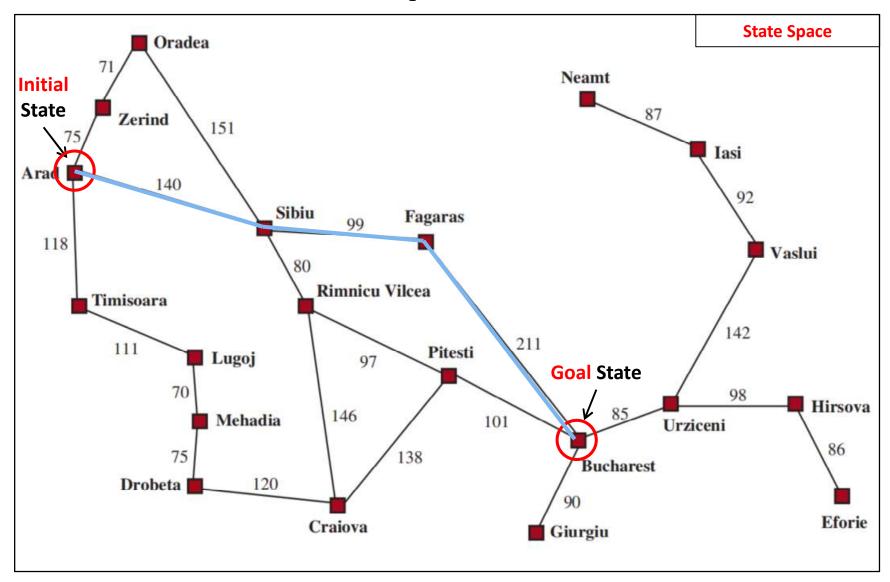
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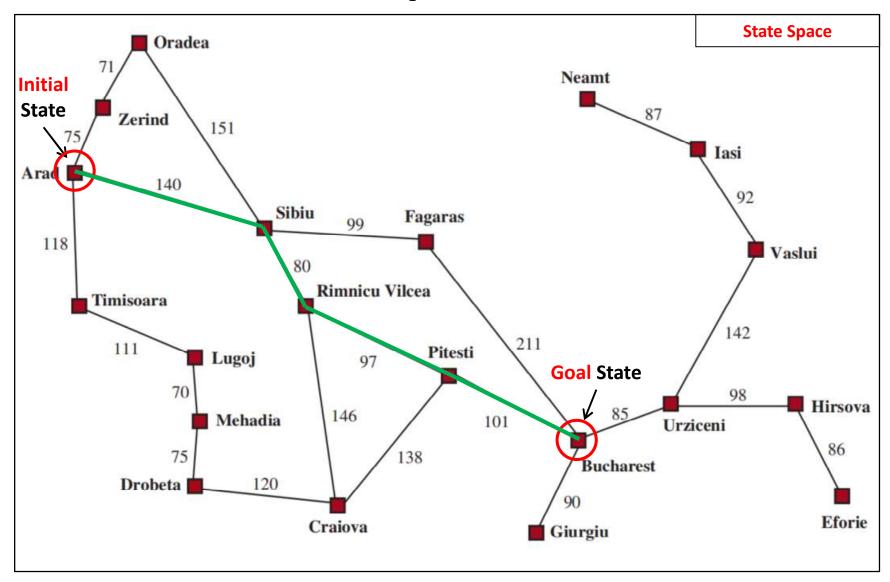
Sample Problem: Dracula's Roadtrip



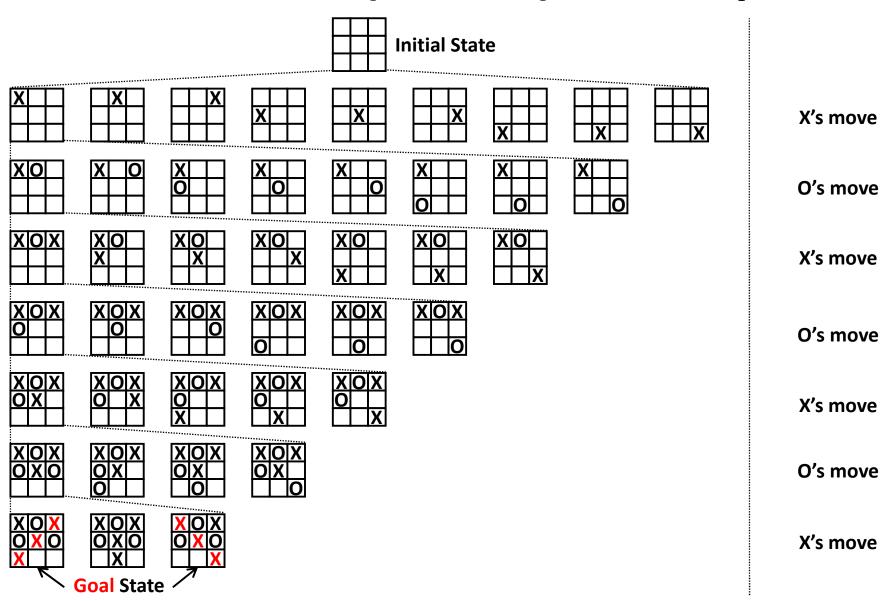
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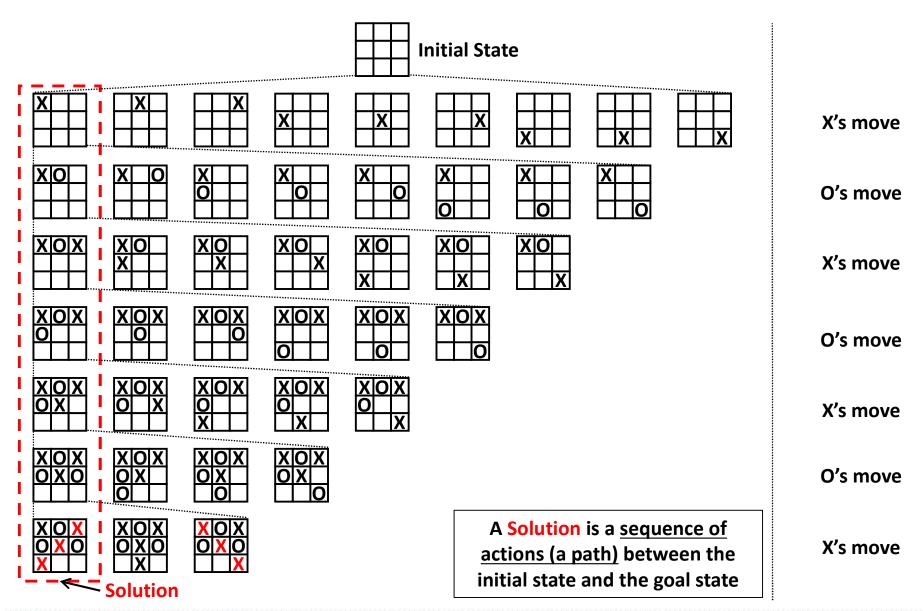
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Tic Tac Toe: (Partial) State Space



Tic Tac Toe: Solution



Chess: (First Move) State Space

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4 knight moves

































Designing the Searching Problem

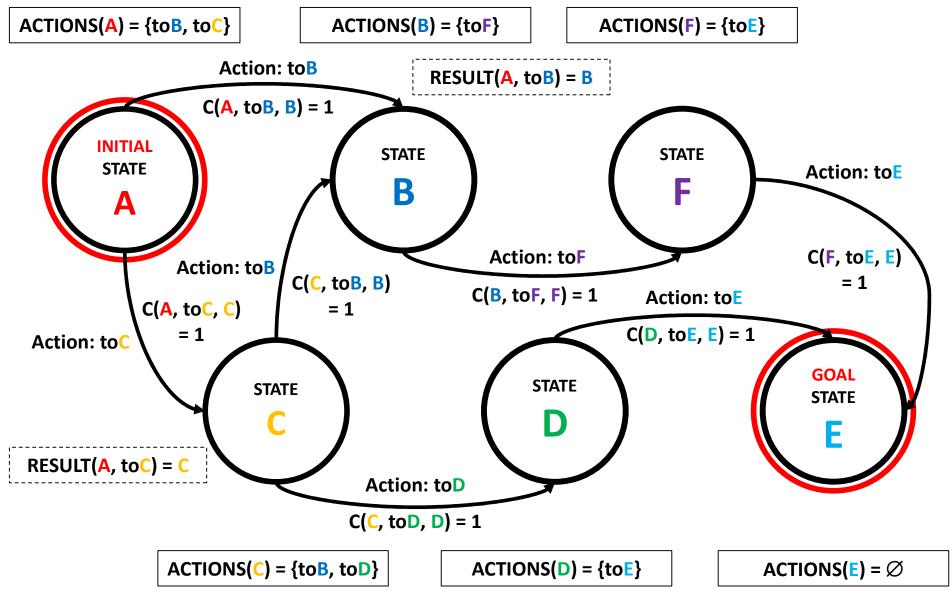
Analyze and define the Problem / Task

Model and buid the State Space

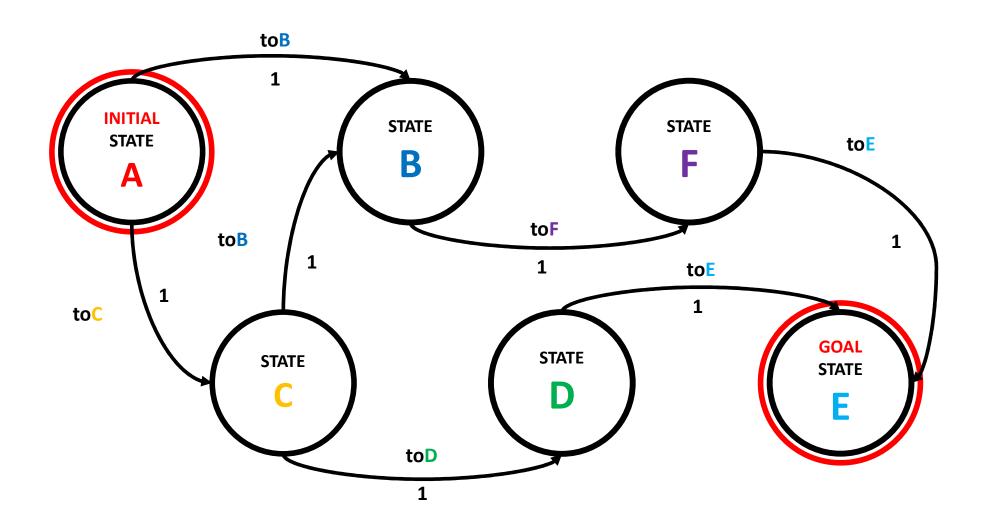
Select searching algorithm

Search

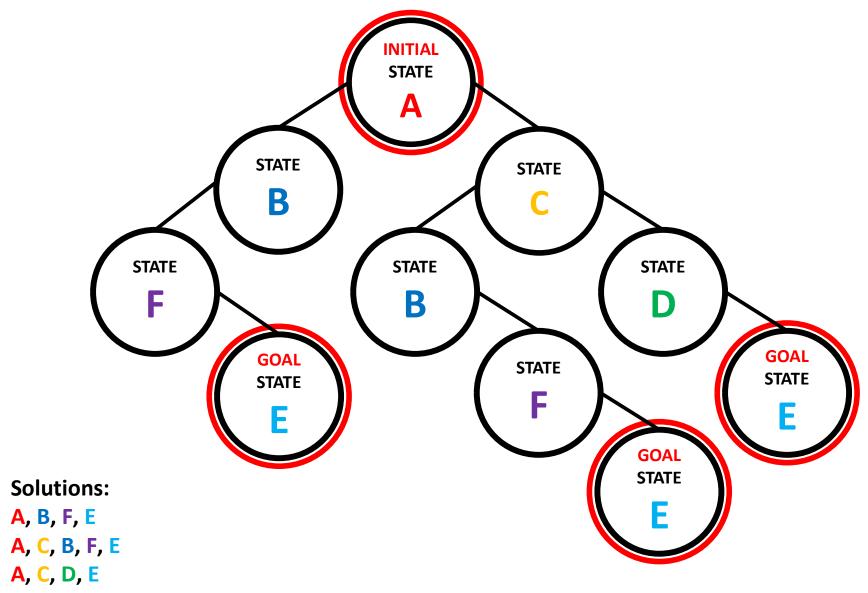
State Space Model: A Graph



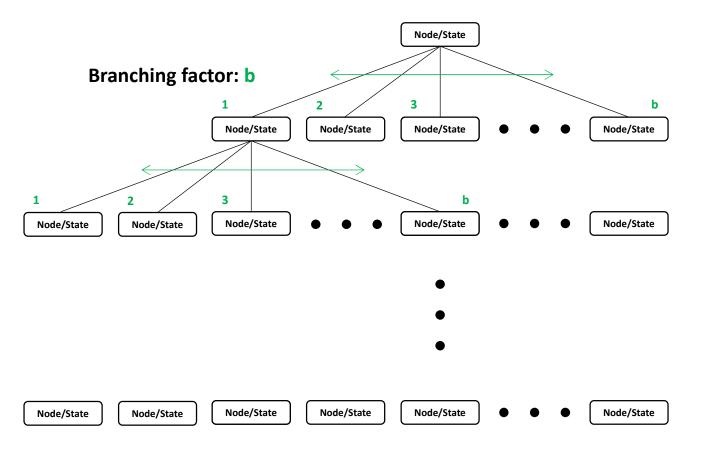
State Space Model: A Graph



Searching State Space: Search Tree



Search Tree Challenges: Size



Total number of nodes / states: $1 + b + b^2 + b^3 + ... + b^d \rightarrow O(b^d)$

Quickly becomes unmanageable and impossible to search with brute force!

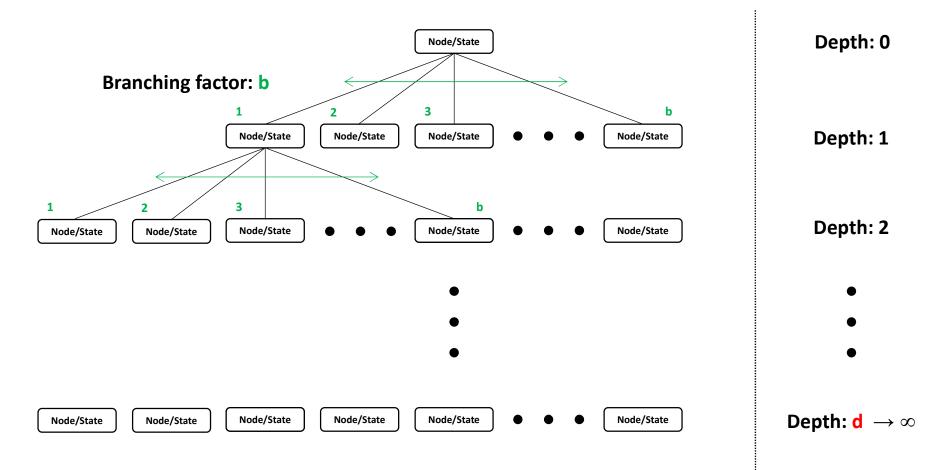
Depth:
$$0 | N_0 = 1$$

Depth:
$$1 | N_1 = b$$

Depth: 2 |
$$N_2 = b^2$$

Depth:
$$d \mid N_d = b^d$$

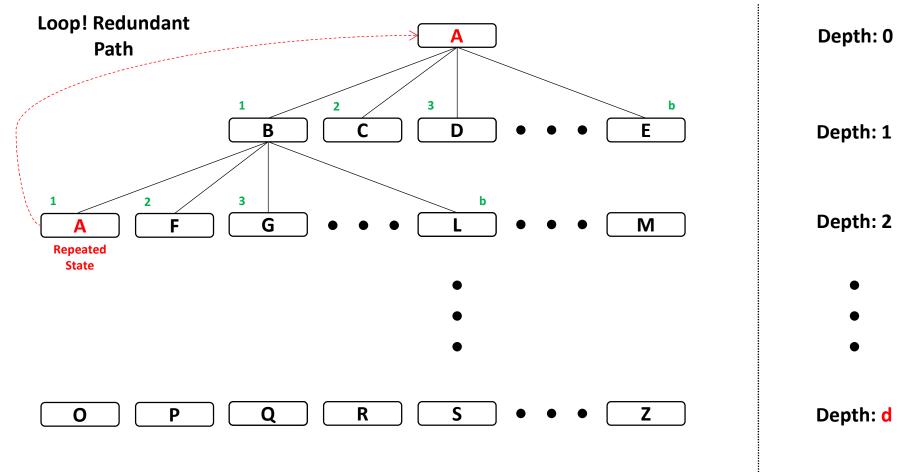
Search Tree Challenges: Infiniteness



Unmanageable and impossible to search with brute force!

Memory and time use grows quickly!

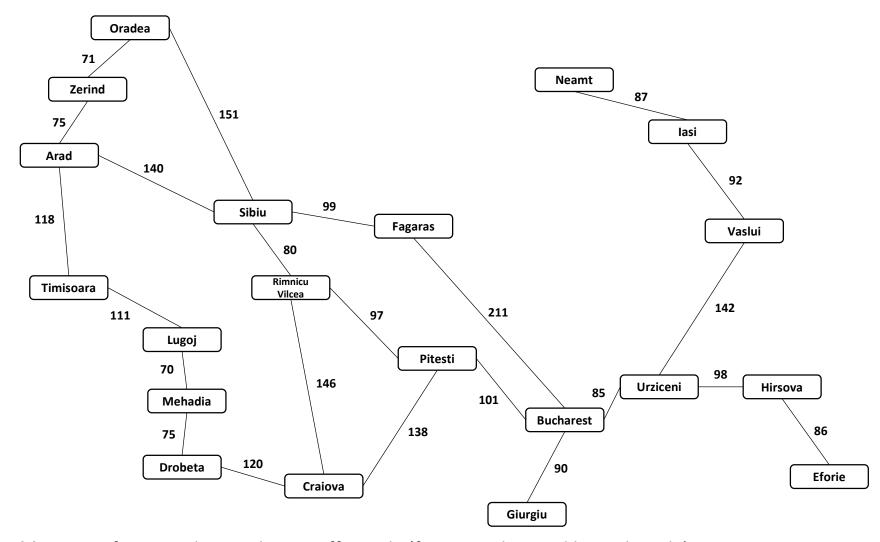
Search Tree Challenges: Loops



This would lead to an infinite state sequence repetition if not handled!

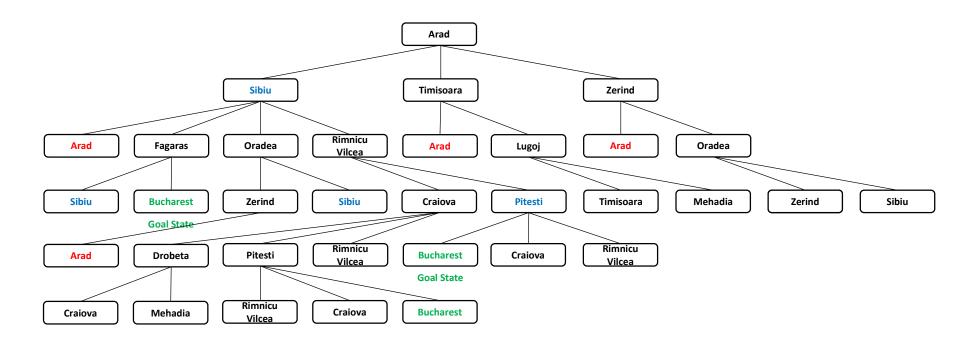
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Sample Problem: Dracula's Roadtrip



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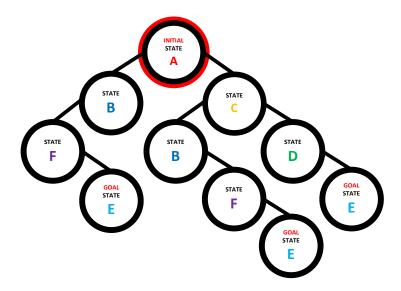
Dracula's Roadtrip as a Tree



INCOMPLETE! I need to redraw it in smarter way

Search Tree: Implementations

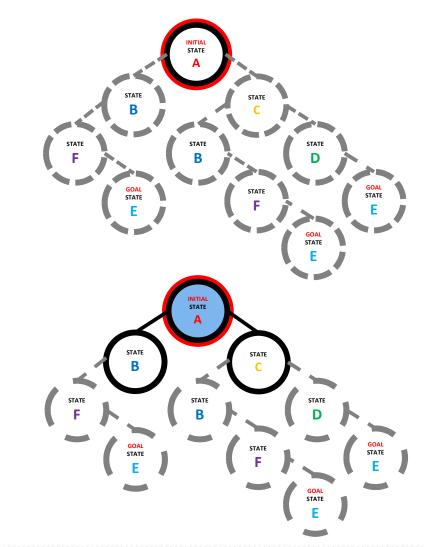
Build entire search tree



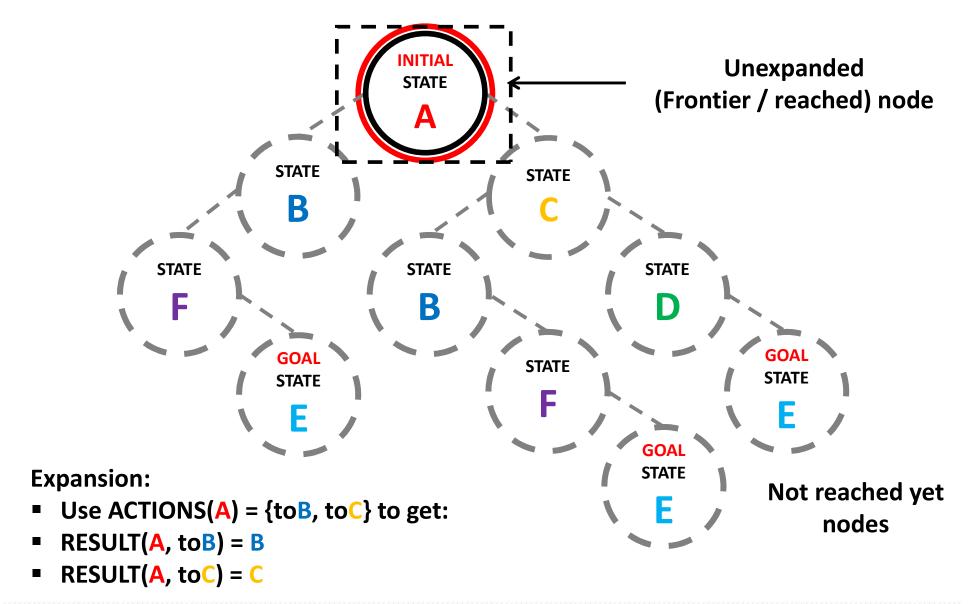
Challenges:

- memory requirements
- impossible for infinite number of states

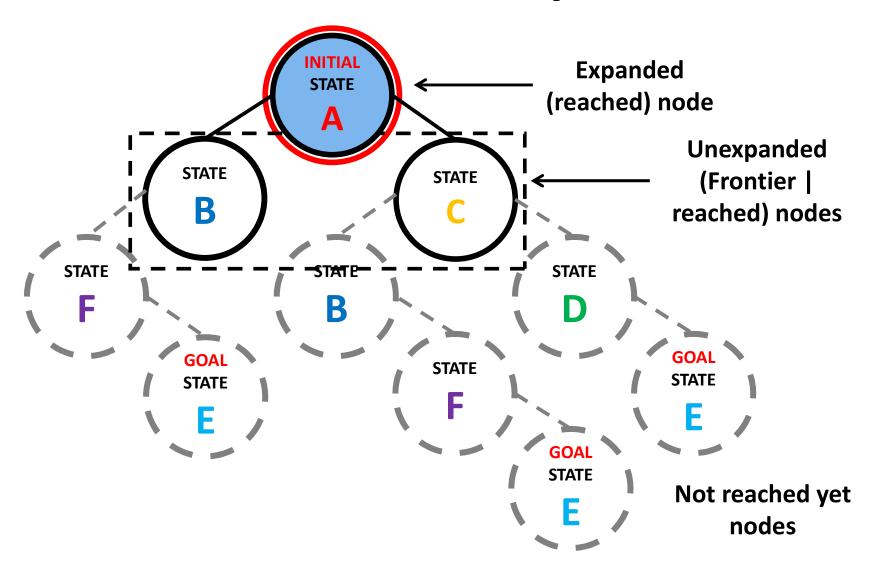
Expand/generate nodes as you go



Search Tree: Node Expansion



Search Tree: Node Expansion



Chess: State Node Expansion

Use game rules to generate subsequent possible game tree states / nodes!











20 Possible legal <u>first</u> moves: 16 pawn moves 4 knight moves

































Designing the Searching Problem

Analyze and define the Problem / Task

Model and buid the State Space

Select searching algorithm

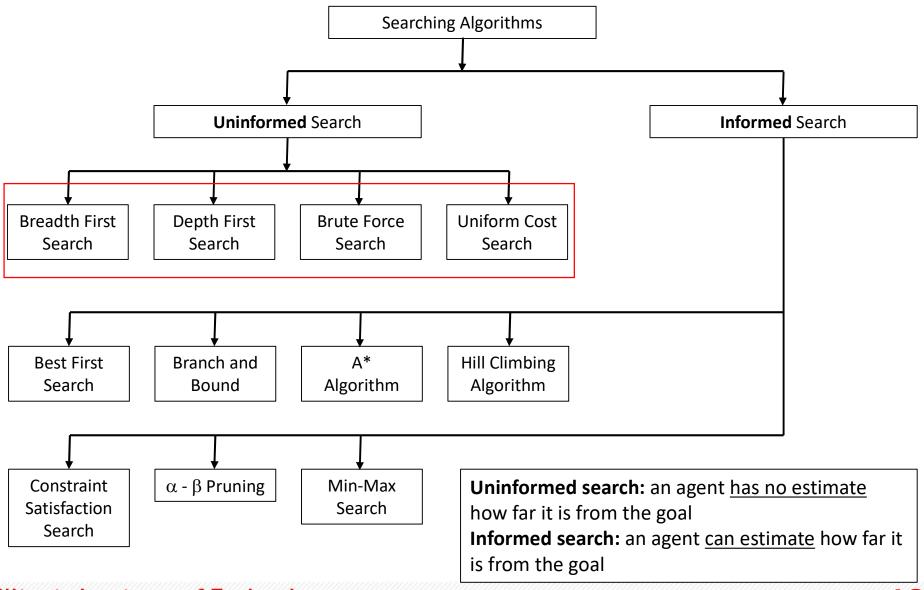
Search

Measuring Searching Performance

Search algorithms can be evaluated in four ways:

- Completeness: Is the algorithm guaranteed to find a solution when there is one, and to correctly report failure when there is not?
- Cost optimality: Does it find a solution with the lowest path cost of all solutions?
- Time complexity: How long does it take to find a solution? (in seconds, actions, states, etc.)
- Space complexity: How much memory is needed to perform the search?

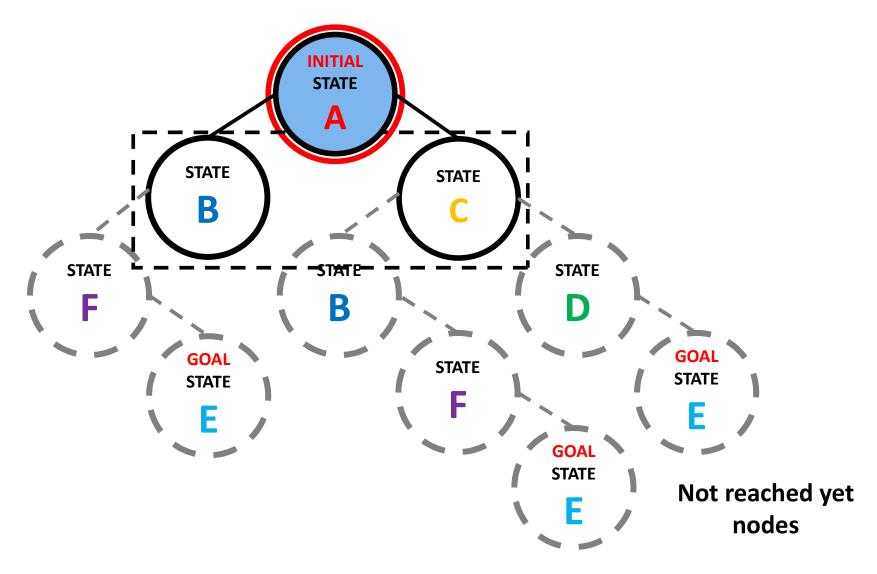
Selected Searching Algorithms



Uninformed Searching

- Breadth First Search (BFS):
 - Will find a solution with a minimal number of actions
 - Large memory requirement
 - Only relatively small problem instances are tractable
- Depth First Search:
 - May NOT find a solution with a minimal number of actions
 - Requires less memory than BFS (for tree search
 - Backtracking (one child / successor generated at a time)
- Brute Force Search: depends on the approach -> bad
- Uniform Cost Search: minimize solution / path cost

Expansion: Which Node to Expand?



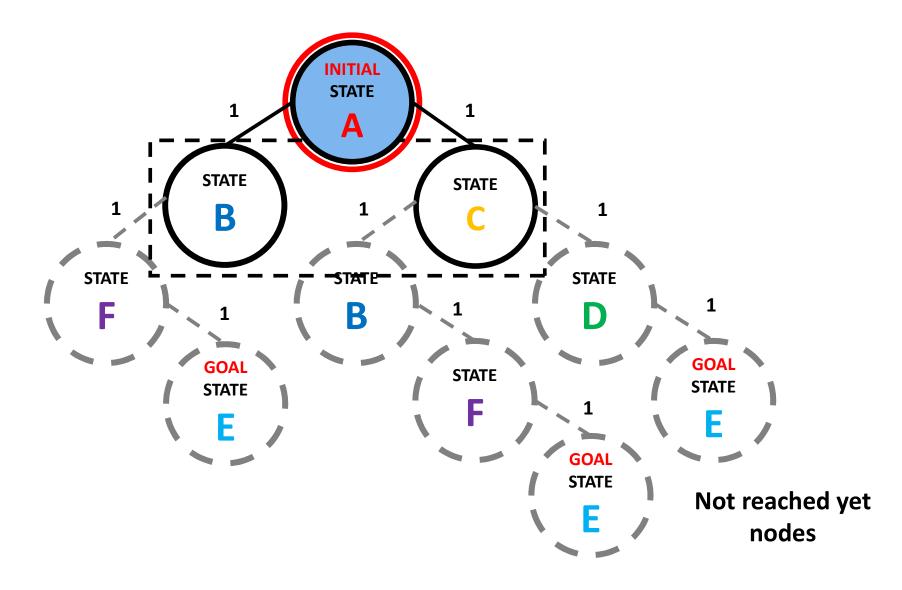
Evaluation function

Calculate / obtain:

f(n) = f(State n)
f(n) = f(relevant information about State n)

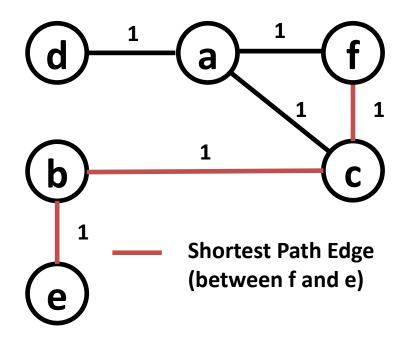
A state n with minimum f(n) should be chosen for expansion
What about ties?

Search Tree: Uniform Action Cost



Uniform Cost Search | Dijkstra's Algo

Weighted Graph G



Popular algorithms:

Dijkstra's algorithm

Shortest Path Problem

Shortest path problem:

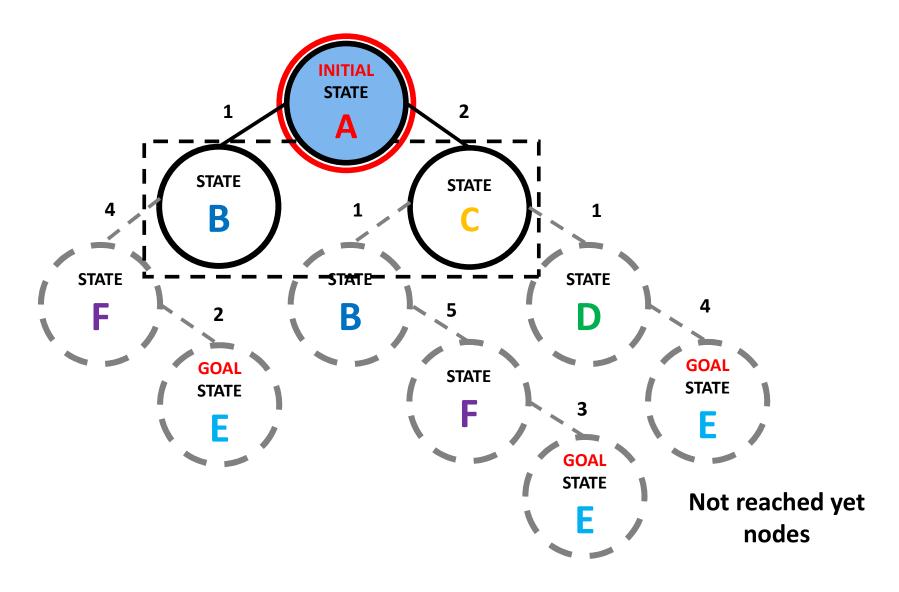
Given a weighted graph G(V, E, w) and two vertices a, b in V, find the shortest path between vertices a and b (all edge weights are equal).

BFS and UCS: Pseudocode

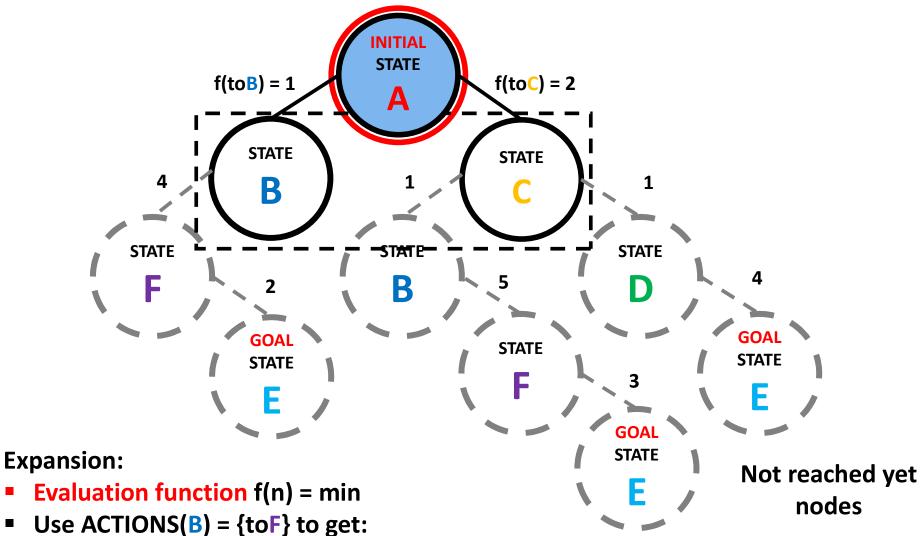
```
function Breadth-First-Search(problem) returns a solution node or failure
  node \leftarrow \text{NODE}(problem.\text{INITIAL})
  if problem.Is-GOAL(node.STATE) then return node
  frontier \leftarrow a FIFO queue, with node as an element
  reached \leftarrow \{problem.INITIAL\}
   while not IS-EMPTY(frontier) do
     node \leftarrow Pop(frontier)
     for each child in EXPAND(problem, node) do
       s \leftarrow child.STATE
       if problem.IS-GOAL(s) then return child
       if s is not in reached then
          add s to reached
          add child to frontier
  return failure
```

function UNIFORM-COST-SEARCH(*problem*) **returns** a solution node, or *failure* **return** BEST-FIRST-SEARCH(*problem*, PATH-COST)

Search Tree: Variable Action Cost

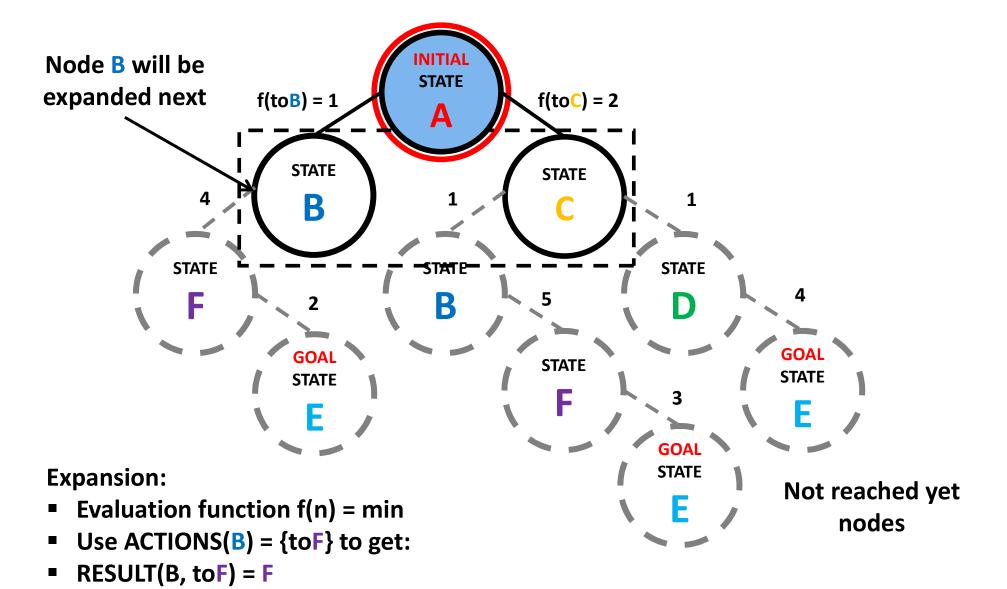


Search Tree: Variable Action Cost



- RESULT(B, toF) = F

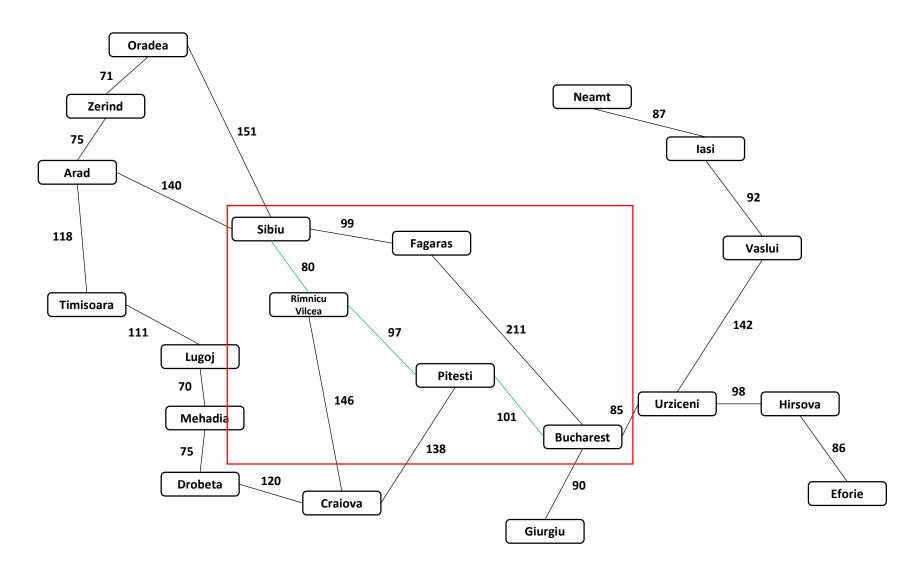
Search Tree: Best-First Search



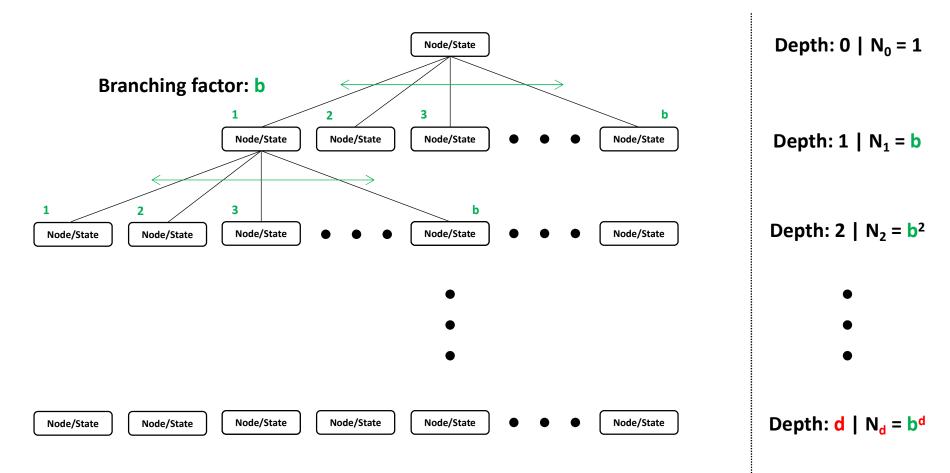
Best-First Search: Pseudocode

```
function BEST-FIRST-SEARCH(problem, f) returns a solution node or failure
  node \leftarrow Node(State=problem.INITIAL)
  frontier \leftarrow a priority queue ordered by f, with node as an element
  reached \leftarrow a lookup table, with one entry with key problem. INITIAL and value node
  while not IS-EMPTY(frontier) do
     node \leftarrow Pop(frontier)
    if problem.Is-GOAL(node.STATE) then return node
    for each child in EXPAND(problem, node) do
       s \leftarrow child.STATE
       if s is not in reached or child. PATH-COST < reached[s]. PATH-COST then
          reached[s] \leftarrow child
         add child to frontier
  return failure
function EXPAND(problem, node) yields nodes
  s \leftarrow node.STATE
  for each action in problem. ACTIONS(s) do
    s' \leftarrow problem.RESULT(s, action)
     cost \leftarrow node.PATH-COST + problem.ACTION-COST(s, action, s')
    yield Node(State=s', Parent=node, Action=action, Path-Cost=cost)
```

Best First Search: Issue



Let's Go Back to Depth First Search

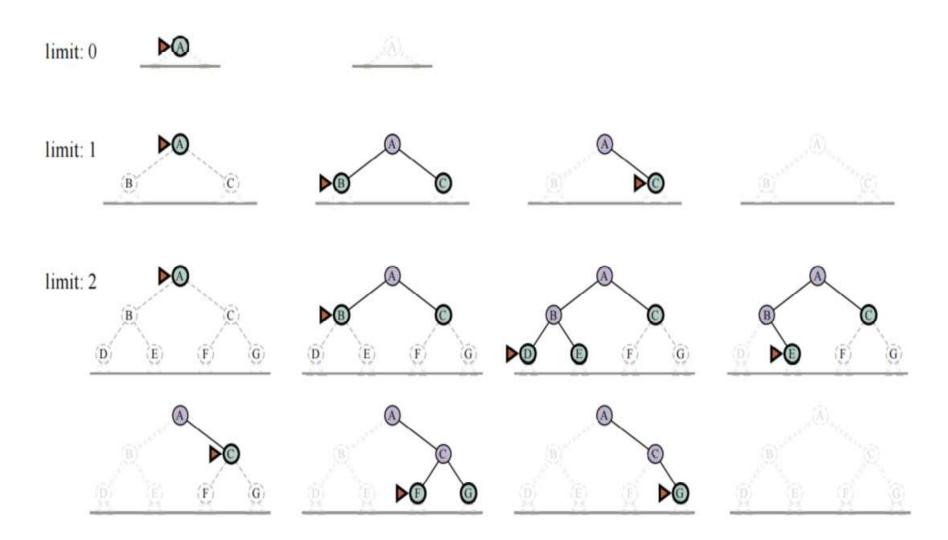


Tree depth is an issue!

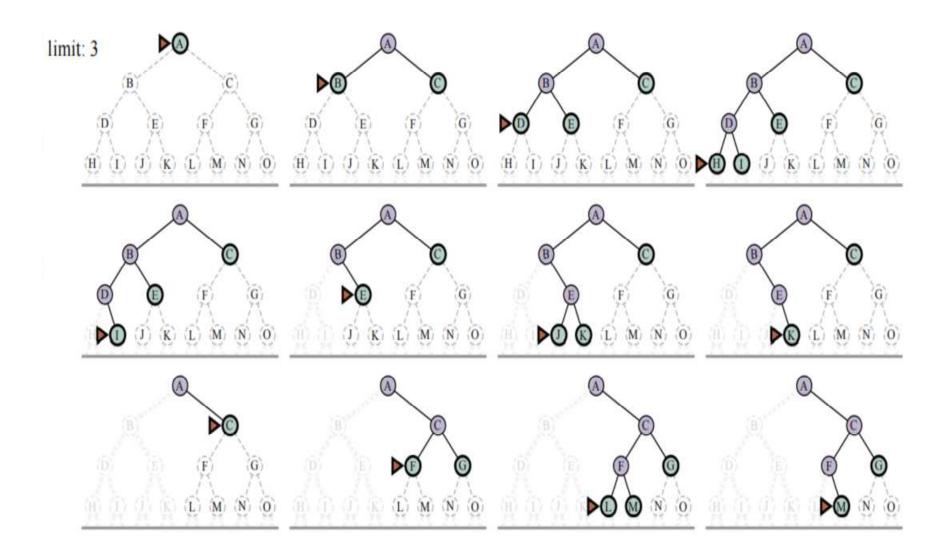
"Controlled" DFS: Pseudocode

```
function ITERATIVE-DEEPENING-SEARCH(problem) returns a solution node or failure
  for depth = 0 to \infty do
     result \leftarrow DEPTH-LIMITED-SEARCH(problem, depth)
    if result \neq cutoff then return result
function DEPTH-LIMITED-SEARCH(problem, \ell) returns a node or failure or cutoff
  frontier \leftarrow a LIFO queue (stack) with NODE(problem.INITIAL) as an element
  result \leftarrow failure
  while not IS-EMPTY(frontier) do
     node \leftarrow Pop(frontier)
    if problem.Is-GOAL(node.STATE) then return node
    if DEPTH(node) > \ell then
       result \leftarrow cutoff
    else if not Is-CYCLE(node) do
       for each child in EXPAND(problem, node) do
         add child to frontier
  return result
```

Iterative Deepening DFS: Illustration



Iterative Deepening DFS: Illustration

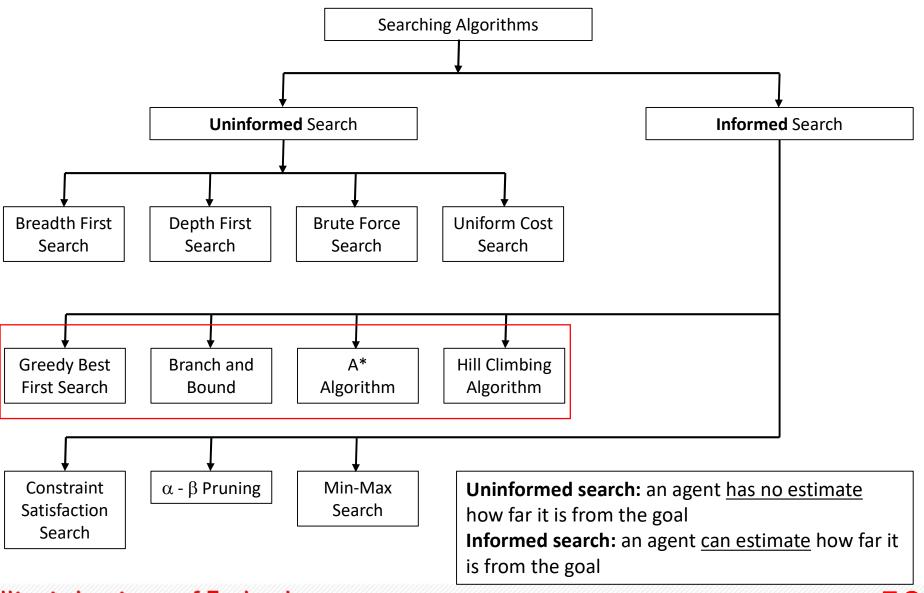


Uninformed Search Algorithms

Criterion	Breadth- First	Uniform- Cost	Depth- First	Depth- Limited	Iterative Deepening	Bidirectional (if applicable)
Complete? Optimal cost? Time Space	Yes^1 Yes^3 $O(b^d)$ $O(b^d)$	$\operatorname{Yes}^{1,2}$ Yes $O(b^{1+\lfloor C^*/\epsilon \rfloor})$ $O(b^{1+\lfloor C^*/\epsilon \rfloor})$	No No $O(b^m)$ $O(bm)$	No No $O(b^\ell)$ $O(b\ell)$	${ m Yes^1} \ { m Yes^3} \ O(b^d) \ O(bd)$	${ m Yes^{1,4}} \ { m Yes^{3,4}} \ O(b^{d/2}) \ O(b^{d/2})$

Figure 3.15 Evaluation of search algorithms. b is the branching factor; m is the maximum depth of the search tree; d is the depth of the shallowest solution, or is m when there is no solution; ℓ is the depth limit. Superscript caveats are as follows: 1 complete if b is finite, and the state space either has a solution or is finite. 2 complete if all action costs are $\geq \epsilon > 0$; 3 cost-optimal if action costs are all identical; 4 if both directions are breadth-first or uniform-cost.

Selected Searching Algorithms



Informed Search and Heuristics

Informed search relies on domain-specific knowledge / hints that help locate the goal state.

h(n): heuristic function - estimated cost of the cheapest path from State n to the goal state