

Non – Preemptive Scheduling



Points To Remember...

- CPU is allocated to the process till it terminates or switches to waiting state
- Execution process is not interrupted even if higher priority ones arrive
- No overhead of switching the process from running to ready state
- Often termed as “rigid”

- ✓ Two processes with same priority are assigned on basis of “arrival time” [first – come – first – serve]
- ✓ Waiting time = Response time
- ✓ Last time instant of gantt chart equals the sum of all the burst time



*Key
Points*

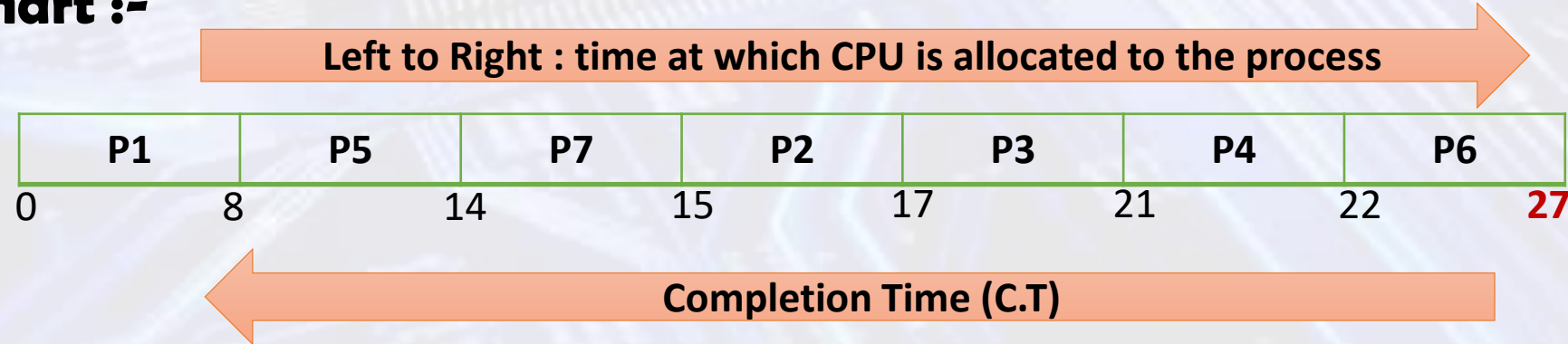
An Example :-

Lesser the
number...
Higher the
priority

	Priority	Arrival Time (A.T)	Burst Time (B.T)	Completion Time (C.T)	Turn-around Time (T.T = C.T – A.T)	Waiting Time (W.T = T.T – B.T)	Response Time (R.T)
P1	3	0	8	8	8-0 = 8	8 – 8 = 0	0
P2	4	1	2	17	17 – 1 = 16	16 – 2 = 14	14
P3	4	3	4	21	21 – 3 = 18	18 – 4 = 14	14
P4	5	4	1	22	22 – 4 = 18	18 – 1 = 17	17
P5	2	5	6	14	14 – 5 = 9	9 – 6 = 3	3
P6	6	6	5	27	27 – 6 = 21	21 – 5 = 16	16
P7	1	10	1	15	15 – 10 = 5	5 – 1 = 4	4

Sum : 27

-: Gantt Chart :-



Advantages

Disadvantages

- Easy to use
- High priority doesn't need to wait for long
- Good mechanism, relative importance of each process is precisely defined

- If a **new high priority process keeps on coming** in the ready queue, then **the process which is in the waiting state may need to wait for long duration of time** [*preemptive scheduling*]
- **Current process is not interrupted**, and hence a **higher priority job may have to wait** [*non-preemptive scheduling*]



We have solution :)

AGING

a technique of **gradually increasing the priority of processes** that wait in the system for a long time.



Starvation
(Indefinite Blocking)