



TCP and OSPF Assignment ELC 3080

Name	ID	Sec / BN
مجدي أحمد عباس عبد الحميد الابرق	9210899	3 / 36

Submitted To: Prof. Dr. Rania Osama 3rd Year – Spring 2024

Contents

1.	TCP La	ıb	3
	1.1. E	ffect of TCP Window Size	3
	Comme retransi MATL	Vary iperf3 window size from 1Kbytes to 6 Kbytes in increments of 1 Kbyte, then set it to 12, 16, 24, 33. Plot the average throughput and the average number of retransmissions as function of window size. Ent on the results and explain the zigzg behavior noticed for larger window sizes. Note that number of missions is the 5th column in iperf3 default output & throughput is the 4 th . You can plot data using AB/python/any online graphing tool (Ex. https://chart-studio.plotly.com/create/#/). The window size is the ion in the iperf command (return to INTRO Lab).	e
		Click on any of the TCP a data segment whose source is node n7, dissect the segment by following nt protocol layer headers from TCP->IP->Ethernet identifying how many header bytes are added by each dentify the TCP options used.	5
	1.1.3.	Repeat for an ACK packet sent from node n11	6
	1.2. T	CP short VS long paths	7
	1.2.1. through	Compare the result of throughput with the case when connection was made to node n11. Why nput drops when connecting to n8 although capacities on the two paths are the same.	7
	1.3. H	ligher link Capacity with Drops VS Reliable Lower Capacity	7
	1.3.1.	Compare throughputs in cases a, b, c. Why b is better than c?	9
	1.3.2.	Compare throughputs in cases b, c and d. Which is better? Why?	9
	1.3.3.	Compare throughputs in cases e and f? Which is better? Why?	9
2.	OSPF I	_ab	9
	2.1. O	SPF Link Cost Change	9
	2.1.1. To help	Check what happens to the path between n7 and n11 (as seen after steps 3 and 6)? Explain what happens o visualize the path choose toolbar->widgets-> Throughput	
	connect	Set the cost of eth1 at node n5 back to 10. Establish two iperf3 connections: one from n7 to n11 and the from n11 to n7 both for duration of 500 seconds. Now go to node n4 and set interface cost for interface ting n4 with n5 to 40. What happens in the paths of the two connections? Explain what happens. What do nclude?	
	2.2. O	SPF Database Updates	2
	2.2.1. in the re	Capture and explain the outputs due to execution of step 2. Why some destinations have more than route outing tables?	
	routing	After executing step 3, determine how long it took the network to exchange link state packets and adjust tables. (Hint: you can calculate the required time by observing the time of first OSPF update message and ACK from Wireshark).	d
	2.2.3. updates	After execution of step 4, identify the new routing table and router database at router n2. Explain the in the new routing table and the new database	3

1. TCP Lab

Each of the following sections describes an experiment based on the network topology provided by the Project_OSPF_TCP.imn file. In your report, divide the answers to correspond to one of these sections in exactly the same sequence and same name, for example use Question 2.1, Question 3.2 and so on. The Introductory Lab explained and introduced most of the needed tools like iperf3, vtysh, wireshark, etc. If you have not gone through it, do not start working on the assignment below. It is better to start the network emulation once and run Wireshark captures. You can reset Wireshark captures from experiment to another.

1.1. Effect of TCP Window Size

- Set wireshark filter to display TCP packets only.
- Start an iperf3 server on node n11.
- Start an iperf3 client on node n7 connecting to server on node n11 for a duration of 40 seconds and reporting interval of 10. Note that in iperf, the client is the node sending the traffic and the server simply receives and sends an ACK.
- 1.1.1. Vary iperf3 window size from 1Kbytes to 6 Kbytes in increments of 1 Kbyte, then set it to 12, 16, 24, 32 Kbytes. Plot the average throughput and the average number of retransmissions as function of window size. Comment on the results and explain the zigzag behavior noticed for larger window sizes. Note that number of retransmissions is the 5th column in iperf3 default output & throughput is the 4th. You can plot data using MATLAB/python/any online graphing tool (Ex. https://chart-studio.plotly.com/create/#/). The window size is the last option in the iperf command (return to INTRO Lab).

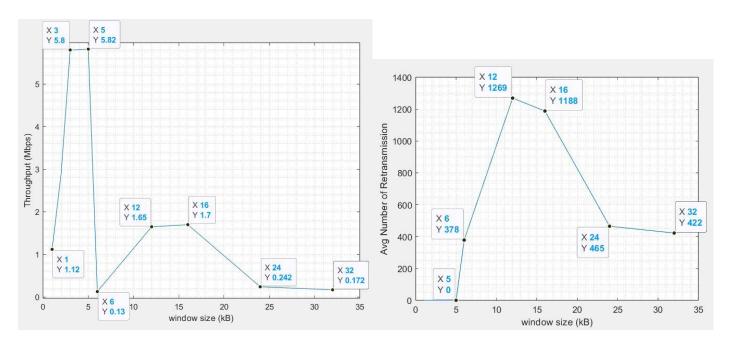


Figure 1: Throughput (Mbps) vs Window Size (kB)

Figure 2: Avg Number of Retransmissions vs Window Size (kB)

Comment on your findings:

We observe from the throughput versus window size graph that there is a zigzag pattern as the window size increases from 1k to 5k. During this range, the number of retransmissions remains zero because the receiver buffer is not yet full. Consequently, as the window size increases, the throughput also rises until the receiver buffer becomes saturated at a window size of 5k. At the window size of 6k, the graph shows a sudden drop in throughput and an increase in retransmissions due to a significant number of packet losses. When this occurs, the transmitter identifies the issue and adjusts by allocating more space in the buffer. As a result, within a window size range between 6k and 12k, the throughput increases again along with a rise in retransmissions. However, if the window size exceeds 16k, the receiver buffer fills up once more, leading to packet losses and a subsequent decrease in throughput and decrease in number of retransmissions.

```
Figure 4: 2 kB Window Size
                         Figure 3: 1 kB Window Size
                                 MBytes
                                                                                                        0.00-40.00
0.00-40.00
        0.00-40.00 sec
0.00-40.00 sec
                                                                                                                            13.8 MBytes
13.8 MBytes
                                                                                                                                                                                  sender
receiver
                                                                                 receiver
 oot@n7:/tmp/pycore.34469/n7.conf#
                                                                                                 oot@n7:/tmp/pycore.34472/n7.conf# 📗
            Throughput = 1.12 Mbps, Retransmissions = 0
                                                                                                              Throughput = 2.9 Mbps, Retransmissions = 0
                         Figure 5: 3 kB Window Size
                                                                                                                          Figure 6: 4 kB Window Size
                                                 Mbits/sec
                            6.93 MBytes
                                                                                                        nterval Transfer
0.00-40.00 sec 27.7 MBytes
0.00-40.00 sec 27.7 MBytes
                                                                                  receive
                                                                                                perf Done.
coot@n7:/tmp/pycore.34478/n7.conf#
perf Done.
oot@n7:/tmp/pycore.34475/n7.conf#
              Throughput = 5.8 \text{ Mbps}, Retransmissions = 0
                                                                                                             Throughput = 5.81 Mbps, Retransmissions = 0
                         Figure 7: 5 kB Window Size
                                                                                                                          Figure 8: 6 kB Window Size
                                          5.82 Mbits/sec
5.82 Mbits/sec
5.82 Mbits/sec
5.82 Mbits/sec
                                                                                                      Interval
0.00-40.00 sec
0.00-40.00 sec
                                                                                                                                           Bandwidth
        0.00-40.00 sec 27.7 MBytes
0.00-40.00 sec 27.7 MBytes
                                                                                receive
                                                                                                                                                                                receiver
                                                                                                perf Done.
cot@n7:/tmp/pycore.34487/n7.conf#
perf Done.
oot@n7:/tmp/pycore.34481/n7.conf# |
            Throughput = 5.82 Mbps, Retransmissions = 0
                                                                                                            Throughput = 0.13 Mbps, Retransmissions = 378
                        Figure 9: 12 kB Window Size
                                                                                                                        Figure 10: 16 kB Window Size
 ot@n7:/tmp/pycore.34493/n7.conf# iperf3 -c 10.0.13.20 -t
nnecting to host 10.0.13.20, port 5201
-4l local 10.0.12.20 port 38581 connected to 10.0.13.20 p
                                 MBytes
MBytes
                                                                                                                                           1.71 Mbits/sec
1.71 Mbits/sec
1.71 Mbits/sec
1.71 Mbits/sec
                                                                                                                      sec
                                                                                                        nterval Transfer
0.00-40.00 sec 8.12 MBytes
0.00-40.00 sec 8.10 MBytes
        0.00-40.00
                     sec
sec
                           7.86 MBytes
7.84 MBytes
                                                                                  sender
        0,00-40.00
perf Done.
oot@n7:/tmp/pycore.34493/n7.conf# |
                                                                                                perf Done.
oot@n7:/tmp/pycore.34499/n7.conf#
                                                                                                           Throughput = 1.7 Mbps, Retransmissions = 1188
          Throughput = 1.65 Mbps, Retransmissions = 1269
                        Figure 11: 24 kB Window Size
                                                                                                                        Figure 12: 32 kB Window Size
                                                                                                                                  KButes
                                                                                                        0.00-40.00
0.00-40.00
                                                                                                                      sec
sec
                                                                                 sender
                                                                                                                                                                                 sender
receiver
                                                                                 receive
                                                                                                 oot@n7:/tmp/pycore.34505/n7.conf#
 oot@n7:/tmp/pycore.34502/n7.conf# 📗
          Throughput = 0.242 Mbps, Retransmissions = 465
                                                                                                          Throughput = 0.172 Mbps, Retransmissions = 422
```

1.1.2. Click on any of the TCP a data segment whose source is node n7, dissect the segment by following different protocol layer headers from TCP->IP->Ethernet identifying how many header bytes are added by each layer, identify the TCP options used.

```
🔣 798 65.274124 10.0.12.20 10.0.13.20 TCP 626 [TCP Fast Retransmission] 38599 > targus-...data1 [PSH, ACK] Seq=19568 Ack=1 Win=1116 Len=558 TSval=1137763 TSecr=1137762
     Arrival Time: May 17, 2024 16:21:54.058059000 EEST
Epoch Time: 1715952114.058059000 seconds
     [Time delta from previous captured frame: 0.000481000 seconds]
[Time delta from previous displayed frame: 0.000481000 seconds
     [Time since reference or first frame: 65.274124000 seconds]
     Frame Number: 798
Frame Length: 626 bytes (5008 bits)
     Capture Length: 626 bytes (5008 bits)
     [Frame is marked: False]
[Frame is ignored: False]
     [Protocols in frame: sll:ip:tcp:data]
[Coloring Rule Name: Bad TCP]
[Coloring Rule String: tcp.analysis.flags]
 Linux cooked capture
    Packet type: Sent by us (4)
Link-layer address type: 1
Link-layer address length: 6
     Source: 00:00:00_aa:00:18 (00:00:00:aa:00:18)
Protocol: IP (0x0800)
 Internet Protocol Version 4, Src: 10.0.12.20 (10.0.12.20), Dst: 10.0.13.20 (10.0.13.20)
  ▽ Differentiated Services Field: 0x00 (DSCP 0x00: Default; ECN: 0x00: Not-ECT (Not ECN-Capable Transport))
     0000 00. = Differentiated Services Codepoint: Default (0x00)
......00 = Explicit Congestion Notification: Not-ECT (Not ECN-Capable Transport) (0x00)
Total Length: 610
     Identification: 0x89e0 (35296)
 Identification: 0x89e0 (35296)

Flags: 0x02 (Don't Fragment)
0..... = Reserved bit: Not set
.1.... = Don't fragment: Set
.0. ... = More fragments: Not set
Fragment offset: 0
  Time to live: 64
Protocol: TCP (6)
Header checksum: 0x818e [correct]
       [Good: True]
     [Bad: False]
Source: 10.0.12.20 (10.0.12.20)
     Destination: 10.0.13.20 (10.0.13.20)
 Transmission Control Protocol, Src Port: 38599 (38599), Dst Port: targus-getdatal (5201), Seq: 19568, Ack: 1, Len: 558 Source port: 38599 (38599)
     Destination port: targus-getdatal (5201)
     [Stream index: 3]
Sequence number: 19568
                                 (relative sequence number)
    [Next sequence number: 20126 (relative sequence
Acknowledgement number: 1 (relative ack number)
                                      (relative sequence number)]
    Acknowledgement number: 1
Header length: 32 bytes
  ♥ Flags: 0x018 (PSH, ACK)
       000. . . . = Reserved: Not set
. . . . . . = Nonce: Not set
. . . . . . = Congestion Window Reduced (CWR): Not set
                                                                                                                        Note:
                                                                                                                        The timestamp option includes a timestamp
       .... .0.. ... = ECN-Echo: Not set .... .0. ... = Urgent: Not set
                                                                                                                        from the sender, which is then echoed by the
       .... = Acknowledgement: Set
       .... 1... = Push: Set
.... 0.. = Reset: Not set
                                                                                                                        receiver in each packet once it's enabled
       .... .... ..0. = Syn: Not set
.... .... 0 = Fin: Not set
                                                                                                                        during connection setup. This helps calculate
                                                                                                                        round-trip time samples to estimate packet
     Window size value: 1116
     [Calculated window size: 1116]
                                                                                                                        loss and extends the 32-bit sequence number.
     [Window size scaling factor: 1]

    □ Checksum: 0x2f7c [validation disabled]

                                                                                                                        On fast connections, sequence numbers can
        [Good Checksum: False]
        [Bad Checksum: False]
                                                                                                                        reset quickly, causing potential confusion
  ▽ Options: (12 bytes)
       No-Operation (NOP)
                                                                                                                        between old and new data. The PAWS
        No-Operation (NOP)

▼ Timestamps: TSval 1137763. TSecr 1137762

                                                                                                                        (Protection Against Wrapped Sequence
          Kind: Timestamp (8)
           Length: 10
                                                                                                                        numbers) system resolves this by discarding
          Timestamp value: 1137763
                                                                                                                        segments with outdated timestamps.
           Timestamp echo reply: 1137762
  ▽ [TCP Analysis Flags]
        ▽ [This frame is a (suspected) fast retransmission]
           ▽ [Expert Info (Warn/Sequence): Fast retransmission (suspected)]
                [Message: Fast retransmission (suspected)]
[Severity level: Warn]
                [Group: Sequence]
          [This frame is a (suspected) retransmission]
```

Figure 13: TCP Packet

Comment on your findings here:

- Frame Length: \rightarrow 626 bytes.
- \triangleright IP Header: → 20 bytes
- ➤ TCP Header: \rightarrow 32 bytes (12 Option + 20 Without Option)
- ightharpoonup TCP Options: ightharpoonup 12 bytes (2 NOP + 10 Timestamp)
- \rightarrow TCP Options Timestamp: \rightarrow 10 bytes (4T_x + 4R_x + 1 Length + 1 Kind)
- \triangleright Ethernet Header: = Frame Length IP Header TCP Header = 626 20 32 = 574 bytes
- As Shown in Fig. 13: Packet Sent from n7 to n11 has large Size (626 B) as it contains Data & Headers.

1.1.3. Repeat for an ACK packet sent from node n11.

```
790 65.272130 10.0.13.20 10.0.12.20 TCP 68 [TCP Dup ACK 789#1] targus-getdatal > 38599 [ACK] Seq=1 Ack=19568 Win=1079 Len=0 TSval=1137762 TSecr=1137762
 Frame 790: 68 bytes on wire (544 bits), 68 bytes captured (544 bits)
    Arrival Time: May 17, 2024 16:21:54.056065000 EEST
    Epoch Time: 1715952114.056065000 seconds
    [Time delta from previous captured frame: 0.000496000 seconds]
[Time delta from previous displayed frame: 0.000496000 seconds]
    [Time since reference or first frame: 65.272130000 seconds] Frame Number: 790
    Frame Length: 68 bytes (544 bits)
    Capture Length: 68 bytes (544 bits)
[Frame is marked: False]
    [Frame is ignored: False]
[Protocols in frame: sll:ip:tcp]
    [Coloring Rule Name: Bad TCP]
[Coloring Rule String: tcp.analysis.flags]
   Packet type: Sent by us (4)
Link-layer address type: 1
Link-layer address length: 6
    Source: 00:00:00_aa:00:1a (00:00:00:aa:00:1a)
Protocol: IP (0x0800)
Internet Protocol Version 4, Src: 10.0.13.20 (10.0.13.20), Dst: 10.0.12.20 (10.0.12.20)
    Header length: 20 bytes
▼ Differentiated Services Field: 0x00 (DSCP 0x00: Default; ECN: 0x00: Not-ECT (Not ECN-Capable Transport))
0000 00.. = Differentiated Services Codepoint: Default (0x00)
.....00 = Explicit Congestion Notification: Not-ECT (Not ECN-Capable Transport) (0x00)
    Total Length: 52
    Identification: 0x13ed (5101)
Flags: 0x02 (Don't Fragment)

0.... = Reserved bit: Not set

1..... = Don't fragment: Set

.... = More fragments: Not se
                  = More fragments: Not set
    Time to live: 64
Protocol: TCP (6)

∀ Header checksum: 0xf9af [correct]

      [Bad: False]
    Source: 10.0.13.20 (10.0.13.20)
Destination: 10.0.12.20 (10.0.12.20)
Transmission Control Protocol, Src Port: targus-getdatal (5201), Dst Port: 38599 (38599), Seq: 1, Ack: 19568, Len: 0
     Source port: targus-getdatal (5201)
    Destination port: 38599 (38599)
    [Stream index: 3]
Sequence number: 1
                              (relative sequence number)
    Acknowledgement number: 19568
                                           (relative ack number)
    Header length: 32 bytes

▽ Flags: 0x010 (ACK)

      000. .... = Reserved: Not set
       ...0 .... = Nonce: Not set
       .... 0... = Congestion Window Reduced (CWR): Not set
       .... .0.. .... = ECN-Echo: Not set
       .... ..0. .... = Urgent: Not set
      .... ...1 .... = Acknowledgement: Set
.... 0... = Push: Not set
       .... .0.. = Reset: Not set
       .... .... ..0. = Syn: Not set
.... .... 0 = Fin: Not set
    Window size value: 1079
    [Calculated window size: 1079]
    [Window size scaling factor: 1]

    ∨ Checksum: 0x2d4e [validation disabled]

       [Good Checksum: False]
       [Bad Checksum: False]
 ▽ Options: (12 bytes)
      No-Operation (NOP)
       No-Operation (NOP)
    ▽ Timestamps: TSval 1137762, TSecr 1137762
         Kind: Timestamp (8)
          Length: 10
          Timestamp value: 1137762
          Timestamp echo reply: 1137762
 ▽ [SEQ/ACK analysis]
      [TCP Analysis Flags]
       [This is a TCP duplicate ack]
[Duplicate ACK #: 1]
    ▽ [Duplicate to the ACK in frame: 789]
      ▼ [Expert Info (Note/Sequence): Duplicate ACK (#1)]

[Message: Duplicate ACK (#1)]
             [Severity level: Note]
             [Group: Sequence]
```

Figure 14: ACK Packet

Comment on Your findings:

- Frame Length: \rightarrow 68 bytes.
- ightharpoonup IP Header: ightharpoonup 20 bytes
- ➤ TCP Header: \rightarrow 32 bytes (12 Option + 20 Without Option)
- \triangleright TCP Options: → 12 bytes (2 NOP + 10 Timestamp)
- ightharpoonup TCP Options Timestamp: \rightarrow 10 bytes (4T_x + 4R_x + 1 Length + 1 Kind)
- \triangleright Ethernet Header: = Frame Length IP Header TCP Header = 68 20 32 = 16 bytes.
- ➤ As Shown in Fig. 14: Packet Sent from n11 to n7 has Small Size (68 B) as it contains Headers only.

1.2. TCP short VS long paths

- Run an iperf3 server on node n8.
- Start an iperf3 client on node n7 connecting to server on node n8 for a duration of 40 seconds and reporting interval of 10 with window size 4K.
- 1.2.1. Compare the result of throughput with the case when connection was made to node n11. Why throughput drops when connecting to n8 although capacities on the two paths are the same.

Figure 15: Short Path (n7 to n11) Throughput = 5.81 Mbps

Figure 16: Long Path (n7 to n8) Throughput = 3.78 Mbps

Comment in your findings here:

- \triangleright Throughput from n7 to n11 (Short Path): = 5.81 Mbps
- \triangleright Throughput from n7 to n8 (Long Path): = 3.78 Mbps
- ➤ We observe that when n8 starts serving instead of n11, the throughput decreases because the path from n7 to n8 is longer than the path from n7 to n11, resulting in delay increase (RTT and Time out is Increased).
- ➤ By increasing Time out with the same window size, the network utilization is decreased. So, throughput decreases as shown in the Fig. 15 & Fig. 16.

1.3. Higher link Capacity with Drops VS Reliable Lower Capacity

This questions for this part are based on the path between n7 and n11.

For each case of the following, run iperf3 client from n7 to n11 with window size 4K.

- A Select link between n4 and n5, configure it to have capacity of 10 Mbps with zero loss in both directions.
- B Select link between n4 and n5, configure it to have capacity of 3 Mbps with zero loss in both directions.
- C Select link between n4 and n5, configure it to have capacity of 10 Mbps with 5% loss in both directions.
- D Select link between n4 and n5, configure it to have capacity of 100 Mbps with 10% loss in both directions.
- E Select link between n4 and n5, configure it to have capacity of 10 Mbps with 1% loss in direction from n4 to n5 and 0% loss in the other direction.
- F Select link between n4 and n5, configure it to have capacity of 10 Mbps with 0% loss in direction from n4 to n5 and 1% loss in the other direction.

Case	Throughput	Figures
A	5.81 Mbps	root@n7:/tmp/pycore.34478/n7.conf# iperf3 -c 10.0.13.20 -t 40 -i 10 -w 4k Connecting to host 10.0.13.20, port 5201 [4] local 10.0.12.20 port 38586 connected to 10.0.13.20 port 5201 [ID] Interval
В	2.85 Mbps	root@n7:/tmp/pycore.34523/n7.conf# iperf3 -c 10.0.13.20 -t 40 -i 10 -w 4k Connecting to host 10.0.13.20, port 5201 [4] local 10.0.12.20 port 38611 connected to 10.0.13.20 port 5201 [ID] Interval Transfer Bandwidth Retr Cwnd [4] 0.00-10.00 sec 3.33 MBytes 2.85 Mbits/sec 0 14.1 KBytes [4] 10.00-20.00 sec 3.41 MBytes 2.86 Mbits/sec 0 14.1 KBytes [4] 20.00-30.00 sec 3.39 MBytes 2.84 Mbits/sec 0 14.1 KBytes [4] 30.00-40.00 sec 3.41 MBytes 2.86 Mbits/sec 0 14.1 KBytes [4] 30.00-40.00 sec 3.41 MBytes 2.86 Mbits/sec 0 14.1 KBytes [1D] Interval Transfer Bandwidth Retr [4] 0.00-40.00 sec 13.6 MBytes 2.85 Mbits/sec 0 sender [4] 0.00-40.00 sec 13.6 MBytes 2.85 Mbits/sec receiver iperf Bone, root@n7:/tmp/pycore.34523/n7.conf#
C	493 kbps	root@n7:/tmp/pycore.34526/n7.conf# iperf3 -c 10.0.13.20 -t 40 -i 10 -w 4k Connecting to host 10.0.13.20, port 5201 [4] local 10.0.12.20 port 38614 connected to 10.0.13.20 port 5201 [ID] Interval Transfer Bandwidth Retr Cund [4] 0.00-10.00 sec 737 KBytes 603 Kbits/sec 39 4.24 KBytes [4] 10.00-20.00 sec 663 KBytes 570 Kbits/sec 38 4.24 KBytes [4] 20.00-30.00 sec 696 KBytes 570 Kbits/sec 38 4.24 KBytes [4] 30.00-40.00 sec 305 KBytes 250 Kbits/sec 22 2.83 KBytes [4] 0.00-40.00 sec 2.35 MBytes 493 Kbits/sec 137 sender [4] 0.00-40.00 sec 2.34 MBytes 492 Kbits/sec 137 receiver iperf Done. root@n7:/tmp/pycore.34526/n7.conf#
D	132 kbps	root@n7:/tmp/pycore.34530/n7.conf# iperf3 -c 10.0.13.20 -t 40 -i 10 -w 4k Connecting to host 10.0.13.20, port 5201 [4] local 10.0.12.20 port 38618 connected to 10.0.13.20 port 5201 [ID] Interval Transfer Bandwidth Retr Cund [4] 0.00-10.00 sec 233 KBytes 191 Kbits/sec 25 4.24 KBytes [4] 10.00-20.00 sec 33.9 KBytes 27.8 Kbits/sec 9 4.24 KBytes [4] 20.00-30.00 sec 120 KBytes 98.5 Kbits/sec 10 4.24 KBytes [4] 30.00-40.00 sec 263 KBytes 215 Kbits/sec 25 4.24 KBytes [4] 30.00-40.00 sec 560 KBytes 133 Kbits/sec 69 sender [1D] Interval Transfer Bandwidth Retr [4] 0.00-40.00 sec 650 KBytes 133 Kbits/sec 69 sender [4] 0.00-40.00 sec 645 KBytes 132 Kbits/sec receiver iperf Done. root@n7:/tmp/pycore.34530/n7.conf#
E	3.56 Mbps	root@n7:/tmp/pycore.34533/n7.conf# iperf3 -c 10.0.13.20 -t 40 -i 10 -w 4k Connecting to host 10.0.13.20, port 5201 [4] local 10.0.12.20 port 38621 connected to 10.0.13.20 port 5201 [ID] Interval
F	2.8 Mbps	root@n7:/tmp/pycore.34537/n7.conf# iperf3 -c 10.0.13.20 -t 40 -i 10 -w 4k Connecting to host 10.0.13.20, port 5201 [4] local 10.0.12.20 port 38625 connected to 10.0.13.20 port 5201 [ID] Interval Transfer Bandwidth Retr Cund [4] 0.00-10.00 sec 3.82 MBytes 3.21 Mbits/sec 25 5.66 KBytes [4] 10.00-20.00 sec 3.34 MBytes 2.80 Mbits/sec 25 8.48 KBytes [4] 20.00-30.00 sec 2.78 MBytes 2.33 Mbits/sec 28 4.24 KBytes [4] 30.00-40.00 sec 3.41 MBytes 2.86 Mbits/sec 24 4.24 KBytes [4] 0.00-40.00 sec 13.3 MBytes 2.80 Mbits/sec 100 sender [4] 0.00-40.00 sec 13.3 MBytes 2.80 Mbits/sec receiver iperf Bone. root@n7:/tmp/pycore.34537/n7.conf#

1.3.1. Compare throughputs in cases a, b, c. Why b is better than c? Comment in your findings here:

We observe that the throughput in case (A) is higher than in case (B), and the throughput in case (B) is higher than in case (C). This indicates that case (B) performs better than case (C), Although the link capacity in case (C) '10 Mbps' is greater than in case (B) '3Mbps' but in case (C) there are higher losses (5%) that necessitate retransmissions, unlike in case (B) that has loss rate is 0%. In addition to that, the throughput in case (A) is greater than that of case (B) due to the greater link capacity in case (A) compared to case (B) as both case (A) and case (B) have equal losses.

1.3.2. Compare throughputs in cases b, c and d. Which is better? Why? Comment in your findings here:

Case (B) is better than cases (C) and (D) [Throughput_B > Throughput_C > Throughput_D] because the loss in case (B) is zero in both directions, eliminating the need for the client or server to retransmit any packets. Although the capacity in case (B) is lower than in cases (C) and (D), packet retransmissions are necessary in cases (C) and (D) due to non-zero loss (for case (C) = 5% loss & for case (D) = 10% loss).

1.3.3. Compare throughputs in cases e and f? Which is better? Why? Comment in your findings here:

Throughput for Case (E) is better than case (F) because large packets are often sent from n5 to n4, while the packets sent from n4 to n5 are usually smaller. Consequently, losses will affect the (n5 to n4) link more than (n4 to n5) link.

2. OSPF Lab

Each of the following sections describes an experiment based on the network topology provided by the **Project OSPF TCP.imn** file.

The <u>Introductory Lab</u> explained and introduced most of the needed tools like iperf3, vtysh, wireshark, etc. If you have not gone through it, do not start working on the assignment below.

It is better to start the network emulation once and run Wireshark captures. You can reset Wireshark captures from experiment to another.

2.1. OSPF Link Cost Change

- 1. Stop any running iperf3 clients.
- 2. Set all links to have zero loss in the two directions with 100 Mbps speed.
- 3. Run iperf3 between n7 and n11 for a duration of 500 seconds or longer. Identify the path between n7 and n11.
- 4. Open vtysh on node n5 by opening a bash terminal and typing vtysh. Now we can configure the router and link costs.
- 5. Type the following in vtysh: show Ip ospf interface eth1 (To know the ethernet number of each interface you can choose toolbar->view->show->interface names) This displays information about interface eth1. Note its ospf cost.
- 6. Type the following configure terminal.
 - interface eth1 (the interface for the link between n5 and n4)
 - ospf cost 40 (set cost to 40)

2.1.1. Check what happens to the path between n7 and n11 (as seen after steps 3 and 6)? Explain what happens. To help visualize the path choose toolbar->widgets-> Throughput.

```
Hello, this is Quagga (version 0.99.20.1).
Copyright 1996-2005 Kunihiro Ishiguro, et al.

n5# show ip ospf interface eth1
eth1 is up
ifindex 116, MTU 1500 bytes, BW 0 Kbit <UP,BROADCAST,RUNNING,MULTICAST>
Internet Address 10.0.4.2/24, Area 0.0.0.0
MTU mismatch detection:enabled
Router ID 10.0.5.1, Network Type BROADCAST, Cost: 10
Transmit Delay is 1 sec, State DR, Priority 1
Designated Router (ID) 10.0.5.1, Interface Address 10.0.4.2
Backup Designated Router (ID) 10.0.3.1, Interface Address 10.0.4.1
Multicast group memberships: OSPFAllRouters OSPFDesignatedRouters
Timer intervals configured, Hello 10s, Dead 40s, Wait 40s, Retransmit 5
Hello due in 7.095s
Neighbor Count is 1, Adjacent neighbor count is 1
n5#
```

Figure 17: n5 interface 1

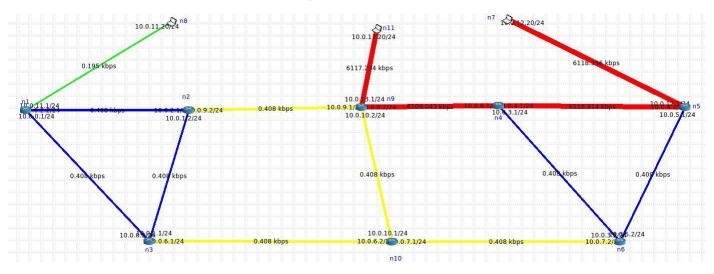


Figure 18: Path From n7 to n11 before changing n5 cost.

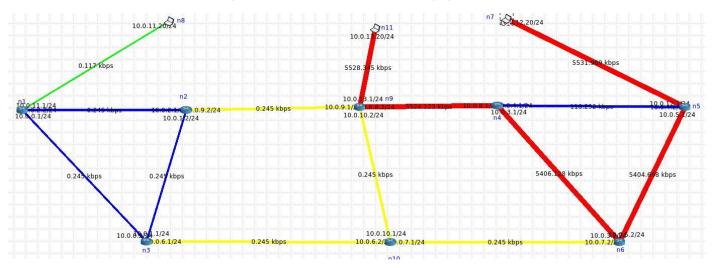


Figure 19: Path From n7 to n11 after changing n5 cost to 40

Comment on your findings:

As illustrated in Fig. 18 & 19, increasing the cost of the interface between n4 and n5 causes the data packets to take an alternative route to achieve a lower overall path cost, with the help of Dijkstra's Algorithm.

The routers begin to recalculate the best route to determine the shortest path between n7 and n11 as shown in Fig. 19

2.1.2. Set the cost of eth1 at node n5 back to 10. Establish two iperf3 connections: one from n7 to n11 and the second from n11 to n7 both for duration of 500 seconds. Now go to node n4 and set interface cost for interface connecting n4 with n5 to 40. What happens in the paths of the two connections? Explain what happens. What do you conclude?

root@n4:/tmp/pycore.59083/n4.conf# vtysh

Hello, this is Quagga (version 0.99,20.1).
Copyright 1996-2005 Kunihiro Ishiguro, et al.

n4# show ip ospf interface eth1
eth1 is up
ifindex 361, MTU 1500 bytes, BW 0 Kbit <UP,BROADCAST,RUNNING,MULTICAST>
Internet Address 10.0.4.1/24, Area 0.0.0.0

MTU mismatch detection:enabled
Router ID 10.0.3.1, Network Type BROADCAST, Cost: 10
Transmit Delay is 1 sec, State Backup, Priority 1
Designated Router (ID) 10.0.5.1, Interface Address 10.0.4.2
Backup Designated Router (ID) 10.0.3.1, Interface Address 10.0.4.1
Multicast group memberships: OSPFAllRouters OSPFDesignatedRouters
Timer intervals configured, Hello 10s, Dead 40s, Wait 40s, Retransmit 5
Hello due in 3,903s
Neighbor Count is 1, Adjacent neighbor count is 1
n4#

Figure 20: n4 interface 1

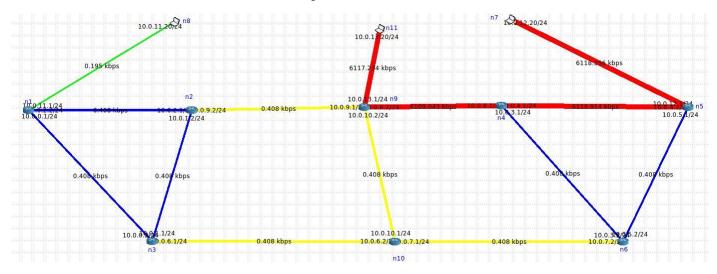


Figure 21: Path from n7 to n11 and vice versa before changing n4 cost.

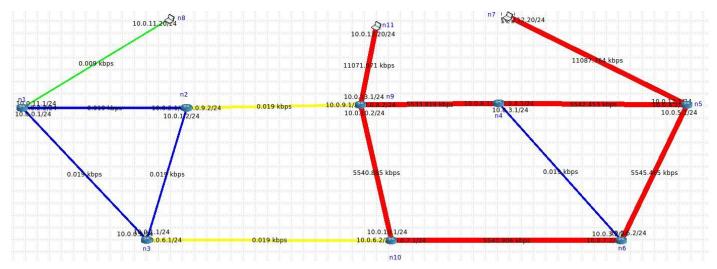


Figure 22: Path from n7 to n11 and vice versa after changing n4 cost to 40

Comment on your findings:

For the same reason mentioned above, the path becomes much longer. The increasing cost of a link forces the router to choose an alternative path, even if it is longer, to reduce congestion.

As OSPF could handle such situations.by increasing the cost of $n4 \rightarrow n5$ and leaving $n5 \rightarrow n4$ the same. The shortest path from n7 to n11 is not the same as the shortest path from n11 to n7.

Paths are $(n7 \rightarrow n5 \rightarrow n4 \rightarrow n9 \rightarrow n11)$ & $(n11 \rightarrow n9 \rightarrow n10 \rightarrow n6 \rightarrow n5 \rightarrow n7)$

2.2. OSPF Database Updates

- 1. Start wireshark and have its filter to capture only OSPF related packets.
- 2. Go to another router say router n2, open vtysh and issue the commands:
 - config terminal
 - show in ospf database (check this link to understand more)
 - show in ospf route.
- 3. On router n4, set link cost of eth1 to 20. Capture the link state packets advertised in Wireshark.
- 4. Go to router n4, go out of vtysh, or open new bash terminal and issue the following commands to disconnect router n4 from the network:
 - ifconfig eth0 down
 - ifconfig eth1 down
 - ifconfig eth2 down

2.2.1. Capture and explain the outputs due to execution of step 2. Why some destinations have more than route in the routing tables?

```
Hello, this is Quagga (version 0.93,20,1).
Copyright 1996-2005 Kunihiro Ishiguro, et al.

n2# show ip ospf database

OSPF Router with ID (10.0.1.2)

Router Link States (Area 0.0.0.0)

Link ID ADV Router Age Seq# CkSum Link count 10.0.0.1 10.0.0.1 623 0x80000000 0xf2bd 3 10.0.0.2 10.0.0.2 623 0x80000000 0x9809 3 10.0.1.2 10.0.3.1 10.0.3.1 615 0x80000000 0x8d0a 3 10.0.3.1 10.0.3.1 615 0x80000000 0xdde 3 10.0.5.1 10.0.5.1 621 0x80000000 0xdde 3 10.0.5.2 10.0.5.2 620 0x80000000 0xdde 3 10.0.5.2 10.0.5.2 620 0x80000000 0xdde 3 10.0.6.2 10.0.8.2 619 0x80000000 0xdde 3 10.0.5.2 10.0.5.2 620 0x80000000 0xdde 3 10.0.5.2 10.0.5.1 621 0x80000000 0xdde 3 10.0.5.2 10.0.5.2 620 0x80000000 0xdde 3 10.0.5.2 10.0.3.2 620 0x80000000 0xdde 3 10.0.0.2 10.0.3.2 624 0x80000000 0xdde 3 10.0.0.2 624 0x80000000 0xdde 3 10.0.0.2 624 0x80000000 0x5fcb 10.0.0.2 624 0x80000000 0x5fcb 10.0.0.2 624 0x80000001 0x5fcb 10.0.3.2 10.0.3.2 623 0x80000001 0x5fcb 10.0.3.2 10.0.3.2 623 0x80000001 0x5fcb 10.0.4.2 10.0.5.1 621 0x80000001 0x6bb3 10.0.4.2 10.0.5.1 621 0x80000001 0x6bb3 10.0.6.2 625 0x80000001 0x6bb3 10.0.6.2 625 0x80000001 0x5bb 10.0.8.2 10.0.8.2 625 0x80000001 0x5bb 10.0.8.2 10.0.8.2 629 0x80000001 0x5bb 10.0.9.1 10.0.8.2 629 0x80
```

```
via 10,0,2.2, eth1 [10] area: 0,0,0.0 directly attached to eth0 [10] area: 0,0,0.0 directly attached to eth1 [30] area: 0,0,0,0 via 10,0,9.1, eth2 [30] area: 0,0,0.0 via 10,0,9.1, eth2 [40] area: 0,0,0.0 via 10,0,9.1, eth2
  10.0.1.0/24
  10,0,2,0/24
  10.0.3.0/24
  10,0,4,0/24
  10.0.5.0/24
  10,0,6,0/24
  10.0.7.0/24
  10.0.8.0/24
  10.0.9.0/24
  10,0,10,0/24
  10.0.11.0/24
  10,0,12,0/24
  10.0.13.0/24
     ====== OSPF router routing table ==
======= OSPF external routing table ======
```

Figure 23: n2 Database

Figure 24: n2 Routing table.

Comment on your findings:

From the above routing table there are multiple paths available that have the same cost. This means that the network switch can choose any one of these paths to send data. Additionally, if one of these paths becomes unavailable (goes down), there are still other paths that can be used, ensuring continued network connectivity.

The database for router n2 is illustrated in the figures, and it consists of two segments:

Router (n2) Link State: This describes the network's routers, links, and topology, constructed at each router after exchanging hello packets with other routers.

Designated Router Link State: Similar to the router link state but generated by the designated router, which is elected to manage communication on a multi-access network segment, reducing overhead and ensuring efficient data exchange.

Figures 23 and 24 show the n2 database and routing table, respectively. Both tables have fields including:

Link ID: Identifies the router or network for which the LSA is originated.

ADV Router: Specifies the originating router of the LSA.

Age: Indicates the time since the LSA was last updated.

SEQ#: The sequence number of the LSA, used to compare and accept new information.

Checksum: Validates the LSA packet.

Link Count: Number of links advertised within the LSA.

The routing table in figure 24 includes destination addresses and their corresponding routes. OSPF supports Equal Cost Multipath (ECMP), meaning multiple paths can be assigned to the same destination if they have equal cost, balancing load and increasing network fault tolerance. Some destinations have multiple equal-cost paths, while others have a single path.

2.2.2. After executing step 3, determine how long it took the network to exchange link state packets and adjust routing tables. (Hint: you can calculate the required time by observing the time of first OSPF update message and the last ACK from Wireshark).

No.	Time	Source	Destination	Protocol Lengt	n Info		
10322	1380.229640	10.0.6.2	224.0.0.5	OSPF 8	4 Hello Packet		_
10323	1380.229645	10.0.7.1	224.0.0.5	OSPF 8	4 Hello Packet	No.	Ti
10324	1380.229647	10.0.10.1	224.0.0.5	OSPF 8	4 Hello Packet	10389	13
10325	1388.370918	10.0.4.1	224.0.0.5	OSPF 20	8 LS Update	10390	13
10326	1388.370924	10.0.4.1	224.0.0.5	OSPF 20	8 LS Update	10391	13
10327	1388.370933	10.0.8.1	224.0.0.5	OSPF 20	8 LS Update	10392	13
10328	1388.370936	10.0.8.1	224.0.0.5	OSPF 20	8 LS Update	10393	13
10329	1388.371087	10.0.8.1	224.0.0.5	OSPF 20	8 LS Update	10394	13
10330	1388.371090	10.0.4.1	224.0.0.5	OSPF 20	8 LS Update	10395	13
10331	1388.371166	10.0.9.1	224.0.0.5	OSPF 11	2 LS Update	10396	13
10332	1388.371170	10.0.9.1	224.0.0.5	OSPF 11	2 LS Update	10397	13
10333	1388.371179	10.0.10.2	224.0.0.5	OSPF 11	2 LS Update	10398	13
10334	1388.371182	10.0.10.2	224.0.0.5	OSPF 11	2 LS Update	10399	13
10335	1388.371224	10.0.5.1	224.0.0.5	OSPF 11	2 LS Update	10400	
10336	1388.371226	10.0.5.1	224.0.0.5	0SPF 11	2 LS Update	10401	
10337	1388.371335	10.0.9.1	224.0.0.5	OSPF 11	2 LS Update	10402	

10389 1388.730011 10.0.2.2 224.0.0.5 OSPF 80 LS Acknowledge 10390 1388.730014 10.0.2.2 224.0.0.5 OSPF 80 LS Acknowledge 10391 1388.730439 10.0.0.1 224.0.0.5 OSPF 80 LS Acknowledge 10392 1388.730443 10.0.2.2 224.0.0.5 OSPF 80 LS Acknowledge 10393 1388.884656 10.0.10.1 224.0.0.5 OSPF 80 LS Acknowledge 10394 1380.884671 10.0.10.1 224.0.0.5 OSPF 80 LS Acknowledge 10395 1388.885669 10.0.10.1 224.0.0.5 OSPF 80 LS Acknowledge	No.	Time	Source	Destination	Protocol Leng	th Info
10391 1388.730439 10.0.0.1 224.0.0.5 OSPF 80 LS Acknowledge 10392 1388.730443 10.0.2.2 224.0.0.5 OSPF 80 LS Acknowledge 10393 1388.884656 10.0.10.1 224.0.0.5 OSPF 80 LS Acknowledge 10394 1388.884657 10.0.10.1 224.0.0.5 OSPF 80 LS Acknowledge	1038	39 1388.730011	10.0.2.2	224.0.0.5	OSPF	80 LS Acknowledge
10392 1388.730443 10.0.2.2 224.0.0.5 OSPF 80 LS Acknowledge 10393 1388.884656 10.0.10.1 224.0.0.5 OSPF 80 LS Acknowledge 10394 1388.884671 10.0.10.1 224.0.0.5 OSPF 80 LS Acknowledge	1039	00 1388.730014	10.0.2.2	224.0.0.5	OSPF	80 LS Acknowledge
10393 1388.884656 10.0.10.1 224.0.0.5 OSPF 80 LS Acknowledge 10394 1388.884671 10.0.10.1 224.0.0.5 OSPF 80 LS Acknowledge	1039	1 1388.730439	10.0.0.1	224.0.0.5	0SPF	80 LS Acknowledge
10394 1388.884671 10.0.10.1 224.0.0.5 OSPF 80 LS Acknowledge	1039	2 1388.730443	10.0.2.2	224.0.0.5	OSPF	80 LS Acknowledge
	1039	3 1388.884656	10.0.10.1	224.0.0.5	OSPF	80 LS Acknowledge
10395 1388.885069 10.0.10.1 224.0.0.5 OSPF 80 LS Acknowledge	1039	94 1388.884671	10.0.10.1	224.0.0.5	OSPF	80 LS Acknowledge
	1039	95 1388.885069	10.0.10.1	224.0.0.5	OSPF	80 LS Acknowledge
10396 1390.108222 10.0.3.2 224.0.0.5 OSPF 84 Hello Packet	1039	96 1390.108222	10.0.3.2	224.0.0.5	0SPF	84 Hello Packet
10397 1390.108222 10.0.1.2 224.0.0.5 OSPF 84 Hello Packet	1039	7 1390.108222	10.0.1.2	224.0.0.5	0SPF	84 Hello Packet
10398 1390.108231 10.0.3.2 224.0.0.5 OSPF 84 Hello Packet	1039	98 1390.108231	10.0.3.2	224.0.0.5	0SPF	84 Hello Packet
10399 1390.108237 10.0.1.2 224.0.0.5 OSPF 84 Hello Packet	1039	99 1390.108237	10.0.1.2	224.0.0.5	0SPF	84 Hello Packet
10400 1390.108246 10.0.5.2 224.0.0.5 OSPF 84 Hello Packet	1046	00 1390.108246	10.0.5.2	224.0.0.5	0SPF	84 Hello Packet
10401 1390.108252 10.0.2.1 224.0.0.5 OSPF 84 Hello Packet	1046	1 1390.108252	10.0.2.1	224.0.0.5	0SPF	84 Hello Packet
10402 1390.108253 10.0.5.2 224.0.0.5 OSPF 84 Hello Packet	1046	2 1390.108253	10.0.5.2	224.0.0.5	0SPF	84 Hello Packet

Figure 25: n2 First Update Packet

Figure 26: n2 Last ACK Packet

Comment on your findings:

From the above graphs we conclude that the time taken by the Network to exchange link state packets and adjust routing tables \rightarrow Time taken = 1388.885069 - 1388.370918 = 0.514151 sec

2.2.3. After execution of step 4, identify the new routing table and router database at router n2. Explain the updates in the new routing table and the new database.

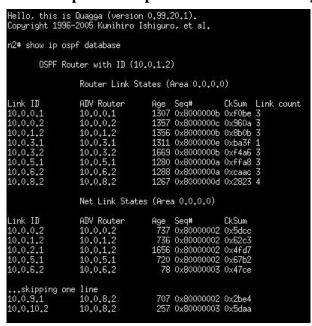


Figure 27: n2 Database after Shutting down n4

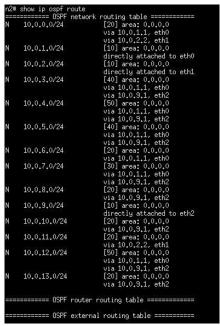


Figure 28: n2 Routing Table after Shutting down n4

Comment on your findings:

If the node (n4) is disconnected, another node will take its place which is (n2). Although n2 has a higher cost than n4, the network will update its database to reflect this change. This ensures that data can still be routed effectively, even though the overall path cost will be higher.