

Simplicial Categories
&
Homotopy Coherence

How to produce quasi-categories?

- The answer is: by means of the adjunction

$$\mathbb{D} \rightleftarrows \mathbb{R} \quad \text{homotopy coherent nerve}$$

↑
the functor producing
simplicial categories

- The aim of this talk is to introduce such functors

Recall: simplicial functors

- A simplicial functor $F: \underline{\mathcal{C}} \rightarrow \underline{\mathcal{D}}$ consists of functors $F_n: \mathcal{C}_n \rightarrow \mathcal{D}_n$ for each n that commute with the simplicial operator functors
- Example Consider a simplicial functor

$$F: \underline{\mathcal{C}} \longrightarrow \underline{s\text{-}\mathbf{Set}}$$

for each $x \in \underline{\mathcal{C}}$ we specify Fx together with a map

$$\Delta^n \times Fx \rightarrow Fy \quad \forall n\text{-simplex in } \underline{\mathcal{C}}(x, y), \text{ s.t.}$$

the faces and degen. of $\Delta^n \times Fx \rightarrow Fy$ correspond to the ones of the n -sim.

Recall: simplicial natural transformations

- $F, G: \underline{\mathcal{E}} \longrightarrow \underline{\mathcal{D}}$

is given by arrows in the underlying category of $\underline{\mathcal{D}}$
for each object of $\underline{\mathcal{E}}$, s.t. $\alpha_x \in D_0(Fx, Gx) \forall x \in \mathcal{E}$:

- α_x form a natural transformation between F & G .
- $s_0(\alpha_x) \in D_1(Fx, Gx)$ should form $F_1 \Rightarrow G_1$
- $s_0 s_0(\alpha_x) = s_1 s_0(\alpha_x) \in D_2(Fx, Gx) \rightsquigarrow F_2 \Rightarrow G_2$
- the images of the α_x under the unique degeneracy operator $[n] \rightarrow [0]$ form a natural transformation
 $F_n \Rightarrow G_n$

Some natural functors

- Consider functors

$c: \text{Set} \rightarrow \text{sSet}$ — constant functor

$\pi_0: \text{sSet} \rightarrow \text{Set}$ — homotopy functor

$ev_0: \text{sSet} \rightarrow \text{Set}$ — underlying category functor
 $x \xrightarrow{\Delta^0} X$

- These functors are monoidal

- So, they induced the functors c_* , $(\pi_0)_*$ and $(ev_0)_*$ between enriched categories

Some natural functors

- $\mathcal{C} = \mathcal{C}_*: \text{Cat} \longrightarrow \text{Cat}_{\Delta}$
 \uparrow
cat of small cats \nwarrow cat of simplicial cats
- $\pi = (\pi_0)_*: \text{Cat}_{\Delta} \longrightarrow \text{Cat}$
- $u = (ev_0)_*: \text{Cat}_{\Delta} \longrightarrow \text{Cat}$
- We just apply the functors \mathcal{C}, π_0 and ev_0 to
hom-object of the corresponding categories

Weak Equivalences between simplicial cats

Definition

A simplicial functor $F: \underline{\mathcal{E}} \rightarrow \underline{\mathcal{D}}$ between simplicially enriched cats is called a weak equiv. if it induces

- WE on all hom-simplicial sets
(weakly fully faithful)
- an essentially surjective functor

$$\pi(\mathcal{E}) \longrightarrow \pi(\mathcal{D})$$

(weakly essentially surjective)

Simplicial cats and simplicial objects

- Any simplicial cat \underline{E} gives rise to a simplicial object $\underline{E}_\bullet : \Delta^{\text{op}} \rightarrow \text{Cat}$ in Cat :

$$\underline{E}_0 \iff \underline{E}_1 \iff \underline{E}_2 \dots$$

- Each of cat \underline{E}_n has the same objects as \underline{E}
- Define $\underline{E}_n(x, y) := \underline{E}(x, y)_n$,
i.e. arrows in \underline{E}_n are n -simplices in $\underline{E}(x, y)$
- In particular, \underline{E}_0 is the underlying cat of \underline{E}

Simplicial cats and simplicial objects

- Conversely, any simplicial object $\mathcal{E}_\bullet: \Delta^{\text{op}} \rightarrow \text{Cat}$, s.t.

$$d_i: \mathcal{E}_n \rightarrow \mathcal{E}_{n-1}, \quad s_i: \mathcal{E}_n \rightarrow \mathcal{E}_{n+1}$$

are the identity on objects

- Define

$$\underline{\mathcal{E}}(x, y)_n \text{ to be } \mathcal{E}_n(x, y)$$

- The simplicial action will be specified by the use of functors d_i and s_i

Topological vs. simplicial categories

- We have the following lax monoidal Quillen adjunction

$$\text{sSet} \begin{array}{c} \xrightarrow{\quad I \dashv \quad} \\ \xleftarrow{\quad S \quad} \end{array} \text{Top}$$

- It induces with the lax monoidal localization functors

$$h: \text{sSet} \longrightarrow \text{Ho}(\text{sSet}), \quad h: \text{Top} \longrightarrow \text{Ho}(\text{Top})$$

the change of base adjunction

$$\mathcal{H} := \text{Ho}(\text{sSet})$$

$$\begin{array}{ccccc} \text{Cat}_{\text{sSet}} & \xrightleftharpoons[\quad S_* \quad]{\quad I \dashv \quad} & \text{Cat}_{\text{Top}} \\ h \searrow & & \swarrow h \\ & \text{Cat}_{\mathcal{H}} & \end{array}$$

$\text{Cat}_{\mathcal{H}}$ is the category
of small categories
enriched over homotopy
types

Locally Kan simplicial categories & Bergner's model structure

Definition A simplicial category is locally Kan if each of its hom-space is a Kan complex

- There is a cofibrantly generated model structure on sCat
- We — simplicial functors that descend to H-equivalences
- Fibrant objects — locally Kan simplicial categories
- Generating cofibrations:

$\mathcal{Q}[A]$ is simplicial cat with 0, 1 as objects

Hom-spaces: $\mathcal{Q}[A](0,0) = \mathcal{Q}[A](1,1) = *$, $\mathcal{Q}[A](0,1) = A$
 $\mathcal{Q}[A](1,0) = \emptyset$



Locally Kan simplicial categories & Bergner's model structure

Theorem (Bergner) \exists a cofib. gen. mod. struc on $s\text{Cat}$

whose $W_{\mathcal{E}}$ are $F: \underline{\mathcal{E}} \rightarrow \underline{\mathcal{D}}$, s.t.

$hF: h\underline{\mathcal{E}} \rightarrow h\underline{\mathcal{D}}$ is an H -equiv;

fibrant objects are the locally Kan simpl. catg

and whose cofibs are generated by

$$\{\emptyset \rightarrow *\} \cup \left\{ 2[\partial\Delta^n] \longrightarrow 2[\Delta^n] \right\}_{n \geq 0}$$

Cofibrant simplicial cats & simplicial computads

- An n -arrow $f: a \rightarrow b$ in \underline{b}_n is just an n -simplex in the simplicial set $\underline{\mathcal{E}}(a, b)$ for simplicial cat $\underline{\mathcal{E}}$
- By Eilenberg-Zilber lemma any n -simplex

We say $f = f' \cdot \alpha$
"f has dimension m"
 $\alpha: [n] \rightarrow [m]$ is an epi
non-degen. m-arrow

$$f = f' \cdot \alpha$$

\uparrow $\nwarrow \alpha: [n] \rightarrow [m]$ is an epi
non-degen. m-arrow

Cofibrant simplicial cats & simplicial computads

Definition An arrow in an unenriched cat is atomic if it admits no non-trivial factorizations



Definition A cat is freely generated by a reflexive directed graph if each of its arrows may be uniquely expressed as a composite of atomic arrows

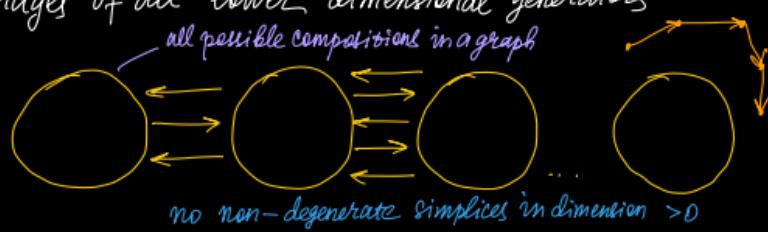
Cofibrant simplicial cats & simplicial computads

Definition A simplicial cat $\mathcal{E}: \Delta^{\text{op}} \rightarrow \text{Cat}$ is a simplicial computad if

- each \mathcal{E}_n is freely generated
- for each surjection $\alpha: [n] \twoheadrightarrow [m]$ and atomic arrow $f \in \mathcal{E}_m$, the arrow $f \cdot \alpha$ is atomic in \mathcal{E}_n

Cofibrant simplicial cats & simplicial computads

- A simplicial computad is a simplicial object in Cat , each of whose categories is freely generated on a set of generating arrows that includes the degenerate images of all lower dimensional generators



Cofibrant simplicial cats & simplicial computads

Lemma

The simplicial computads are the cellular cofibrant objects in $sCat$.

Furthermore, every cofibrant object is cellular, and hence a simplicial computad

Proof

- Prove that any cofibrant simplicial category is a simplicial computad
- Check that a retract of a simplicial computad is a simplicial computad

Cont

- If retract $\mathcal{B} \hookrightarrow \mathcal{G} \rightarrow \mathcal{B}$ in \mathcal{C} of free category
is a free category
- This is it since if $h = fg$ and any two of these are in \mathcal{B} ,
so is the third
- By induction, any arrow in \mathcal{B} is uniquely decomposable
into the shortest composites of atomic arrows of \mathcal{G} that lie in
 \mathcal{B}
- So, at each level a retract of simplicial computed is a
free category.

- Now prove that the degenerate images of atomic arrows in B_n are atomic in B_{n+1}
- This is clear for atomic arrows in B_n that are also atomic in B_n
- Suppose that a degenerate image of some atomic arrow in B_n factors as gf in B_{n+1}
- f is simpl. computad \Rightarrow we have $g'f'$ in B_n with $g' \mapsto g, f' \mapsto f$

- Apply one of the face maps that serves as a retraction of the degeneracy \rightsquigarrow either g or f must map to an identity in B_n
- So, one of g' or f' is an identity \square

A free-forgetful adjunction

$$F : r\text{DirGph} \rightleftarrows \text{Cat} : U$$

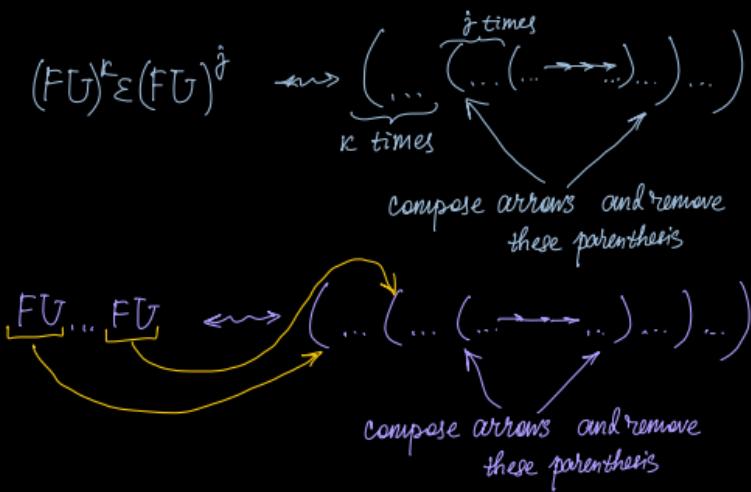
↗
reflexive directed
graphs

- We have the comonad resolution associated to a small cat \mathcal{A}

$$\begin{array}{ccccccc} & & & \xleftarrow{\quad \varepsilon FUFV \quad} & & & \\ & FUF\mathcal{A} & \xleftarrow{\quad F\eta_U \quad} & FUFU\mathcal{A} & \xrightarrow{\quad F\eta_U FV \quad} & FUFUFV\mathcal{A} \dots & \\ & \xleftarrow{\quad FU\varepsilon \quad} & & & \xleftarrow{\quad FU\varepsilon FV \quad} & & \\ & & & & \xrightarrow{\quad FU F\eta_U \quad} & & \\ & & & & \xleftarrow{\quad FU FV \varepsilon \quad} & & \end{array}$$

- Note that this comonad resolution is a simplicial
computed $\text{FU}_\bullet \mathcal{A}$
- Hence, it is a cofibrant simplicial category by
the Lemma
above
- $\text{FU}_\bullet \mathcal{A}$ is the free category

$$\begin{array}{c}
 ((\bullet \longrightarrow (\longrightarrow) \dots \longrightarrow)) \quad - \text{a general arrow in } \text{FU}_\bullet \mathcal{A} \\
 ((\longleftarrow (\longrightarrow) \dots \longrightarrow)) \quad - \text{an atomic arrow in } \text{FU}_\bullet \mathcal{A} \\
 \underbrace{\text{FU}_\bullet \text{FU}_\bullet \mathcal{A}}_{\curvearrowright} \longleftrightarrow ((\dots \longrightarrow \dots))
 \end{array}$$



$F(GF)^\kappa \eta (GF)^j G$ ↪ double up the parentheses
that are contained in
exactly k others

$F\dots GF\eta F$ ↪ insert parenthesis around
each individual morphism

- Recall, the main task of this activity is to motivate the correct choice of an adjunction

$$\mathbb{L} : sSet \rightleftarrows sCat$$

A naïve way to construct $\mathbb{L}\Delta$

- A naïve choice for $\mathbb{L}\Delta^\bullet : \Delta \rightarrow \text{sCat}$
 $[n]$ — a discrete simplicial category
But: the right adjoint will be an ordinary nerve
- $\mathbb{L}\Delta^\bullet$ is supposed to be a simplicial category that encodes a "homotopy coherent" diagram of shape $[n]$

A homotopy commutative diagrams

Definition A homotopy commutative diagram of shape \mathcal{A} is a map of reflexive directed graphs

a small cat

$$U\mathcal{A} \longrightarrow U\mathcal{C}$$

that defines a functor

$$\mathcal{A} \longrightarrow h\mathcal{C}$$

- A diagram $F: U\mathcal{A} \rightarrow U\mathcal{C}$ is homotopy commutative if whenever $h = fg$ in \mathcal{A} , Fh and $Fg \cdot Ff$ lie in the same path component of the hom-space

A homotopy commutative diagrams

- When \mathcal{B} is locally Kan, this is the case when
 \exists 1-simplices

$$Fh \xrightarrow{\quad} Fg \cdot Ff$$

&

$$Fg \cdot Ff \xrightarrow{\quad} Fh$$

A homotopy coherent diagram

Definition A homotopy coherent diagram of shape \mathcal{A} is a simplicial functor

$$FU_{\bullet} \mathcal{A} \rightarrow \underline{\mathcal{C}}$$

- By means of the map

$$U\mathcal{A} \xrightarrow{U\eta} UFTU\mathcal{A}$$

One can construct the homotopy commutative diagram

$$U\mathcal{A} \xrightarrow{U\eta} UFTU\mathcal{A} \xleftarrow{U(\mathrm{id}_{U\mathcal{A}})} U\underline{\mathcal{C}}$$

A homotopy coherent natural transform

Definition A hom. coh. nat. tr. is a homotopy coherent diagram of shape

$$\mathcal{A} \times \mathbb{Z},$$

i.e., a simplicial functor

$$FU_{\bullet}(\mathcal{A} \times \mathbb{Z}) \longrightarrow \underline{\mathcal{C}}$$

The main motivation was...

- Given a commutative diagram $F: \mathcal{A} \rightarrow \mathcal{E}$
- Is it possible to form a new diagram in which each object is replaced by a specified homotopy equivalent one?
- Given $\alpha: F \Rightarrow G$, is it possible to replace maps α_n with homotopic ones?
- The answer is no in general, but...

The main motivation was...

Preposition (Cisner-Petit) Given a homotopy coherent diagram $F: \mathcal{A} \rightarrow \underline{\mathcal{E}}$ in a locally Kan simplicial category and a family of homotopy equivalences

$$F_a \longrightarrow G_a$$

This data extends to a homotopy coherent diagram

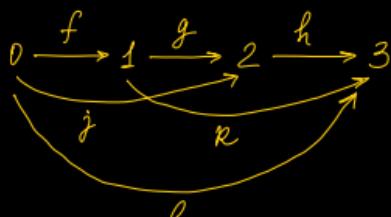
$$G: \mathcal{A} \rightarrow \underline{\mathcal{E}}$$

and homotopy coherent map

$$F \Rightarrow G$$

A fruitful example

- Consider the category $[3]$



Describe the
hom-space
in $F[3](0, 3)$

- The vertices of $F[3](0, 3)$ are the paths of edges:
 ℓ , Rf , hj and hgk (*)
- To obtain the 1-simplices one should write one pair of parenthesis for each in (*)

A fruitful example

- The non-degenerate simplices fit into the diagram

$$\begin{array}{ccc} \ell & \xrightarrow{(hf)} & hf \\ (hj) \downarrow & \swarrow (hg\,f) & \downarrow (hg)\,(f) \\ hj & \xrightarrow{(h)\,(gf)} & hg\,f \end{array}$$

Question

What does $FU_{[3]}(0,3)$ look like?

The desired definition of \mathbb{L}

Definition $\mathbb{L}[\Delta^n] := \text{FT}_\bullet[n]$

That is, $\mathbb{L}[\Delta^n]$ is a homotopy coherent diagram of shape $[n]$

Lemma The cat Cat_Δ admits all small colimits. Hence, there exists a unique colimit-preserving functor

$\mathbb{L}[-]: \text{Set} \rightarrow \text{Cat}_\Delta$
which sends Δ^n to $\mathbb{L}[\Delta^n]$ already defined

Consider an appropriate left Kan extension

Another definition of \mathcal{L}

Definition Let \mathcal{I} be a finite non-empty linearly ordered set, and let i, j be elements of this set.

We let

$$\mathcal{P}_{i,j} := \left\{ I \subseteq \mathcal{I} \mid i, j \in I \text{ and } k \in I \Rightarrow i < k < j \right\},$$

i.e., $\mathcal{P}_{i,j}$ consists of all subsets of $[i, j] \subseteq \mathcal{I}$ which contain i and j

- $\mathcal{P}_{i,j}$ is a poset

Another definition of \mathbb{L}

- For $i \leq j \leq k$ in \mathbb{J} , there is a canonical map of posets

$$P_{i,j} \times P_{j,k} \longrightarrow P_{i,k}$$

$$(I, I') \longmapsto I \cup I'$$

Definition Objects of $\mathbb{L}[\Delta^{\mathbb{J}}]$ — the elements of \mathbb{J}

$$\text{Hom}_{\mathbb{L}[\Delta^{\mathbb{J}}]}(i, j) := \begin{cases} \emptyset, & \text{if } i > j \\ N(P_{i,j}), & \text{if } i \leq j \end{cases}$$

Composition is defined via the previous observation

Some simple properties

- $\forall n \geq 1 \quad \mathcal{N}(P_{0,n}) \cong (\Delta^1)^{n-1}, \quad P_{i,j} \cong P_{0,j-i}$
- For $i \leq j$, the simplicial set $\text{Hom}_{\mathbb{L}[\Delta^n]}(i, j)$ is contractible
- There is a unique isomorphism
$$\pi(\mathbb{L}[\Delta^n]) \cong [n]$$
which is identity on objects
- By adjunction $\mathbb{L}[\Delta^n] \xrightarrow{\cong} c[n]$ a we of simplicial cans

Non - properties of \mathbb{L}

- \mathbb{L} does not preserve products
- The space $\mathbb{L}(x,y)$ are essentially never Kan complexes or even quasi-categories

The homotopy coherent nerve

Definition Let \mathcal{C} be a simplicially enriched category.

Then its simplicial nerve (or the homotopy-coherent nerve) is the simplicial set

$$\mathcal{K}(\mathcal{C})_n := \text{Hom}_{\text{Cat}_{\Delta}}(\mathbb{I}^{\wedge}[\Delta^n], \mathcal{C})$$

- If \mathcal{C} is an ordinary cat

$$\mathcal{N}(\mathcal{C}) \cong \mathcal{K}(c\mathcal{C})$$

a constant enrichment

$$\begin{aligned} \triangleright \quad \mathcal{K}(c\mathcal{C})_n &= \text{Hom}_{\text{Cat}_{\Delta}}(\mathbb{I}^{\wedge}[\Delta^n], c\mathcal{C}) \cong \text{Hom}_{\text{Cat}}(\pi(\mathbb{I}^{\wedge}[\Delta^n]), \mathcal{C}) \\ &\cong \text{Hom}_{\text{Cat}}([n], \mathcal{C}) = \mathcal{N}(\mathcal{C})_n \end{aligned}$$

△

An adjunction and a Quillen equivalence

Theorem The pair $\mathcal{L} \dashv \mathcal{R}$ forms a Quillen equivalence between Joyal's model structure for quasi-categories and Bergner's model structure for simplicial categories

- So, the homotopy coherent nerve of a locally Kan simplicial category is a quasi-category

↗ a source of examples

An adjunction and a Quillen equivalence

Lurie's definition of a $\mathcal{W}\mathcal{E}$ in the Joyal's model
Corollary The map $X \rightarrow Y$ of simplicial sets ^{structure} is a categorical equivalence \Leftrightarrow the induced functor

$$h\mathbb{I}\mathbb{C} X \longrightarrow h\mathbb{I}\mathbb{C} Y$$

is an equivalence of \mathcal{H} -categories

Some fruitful results on $\mathbb{L} \dashv \mathbb{R}$

Theorem Let \mathcal{A} be any small cat. Then

$$\mathbb{L}\mathcal{N}\mathcal{A} \cong \text{FU}_{\bullet}\mathcal{A}$$

Theorem If \mathcal{E} is a locally Kan simplicial cat, then

$\mathbb{R}\mathcal{E}$ is a quasi-cat

Example. For a simplicial modal cat \mathcal{M} we have a right Quillen bifunctor $\mathcal{M}^{\text{op}} \times \mathcal{M} \rightarrow \text{sSet}$

- The category \mathcal{M}_{QR} is locally Kan

- So, $\mathbb{R}\mathcal{M}_{QR}$ is the quasi-cat, and $\text{hRM}_{QR} \cong \text{hM}$

The original motivation by Boardman & Vogt

- Homotopy coherent diagrams and natural transformations assemble into a quasi-cat
- This was an example that motivated the original motivation by Boardman & Vogt:

\forall small cat \mathcal{A} and \forall locally Kan simplicial cat $\underline{\mathcal{E}}$

$$\mathbb{D}^N(\mathcal{A}x[n]) \rightarrow \underline{\mathcal{E}} \leftarrow \text{an } n\text{-simplex}$$

Vertices — homotopy coherent diagrams $\mathcal{A} \rightarrow \underline{\mathcal{E}}$ of some simpl. set

Edges — homotopy coherent natural transf. between $\mathcal{A} \rightarrow \underline{\mathcal{E}}$

The original motivation by Boardman & Vogt

- By adjunction, n -simplices are

$$\begin{array}{c} \mathcal{N}(A \times [n]) \rightarrow \underline{\mathcal{R}\mathcal{E}} \\ \parallel \\ N^A \times \Delta^n \end{array}$$

- Once again, by adjunction the n -simpl. are iso to

$$\Delta^n \rightarrow (\underline{\mathcal{R}\mathcal{E}})^{N^A}$$

- So, by Yoneda, our simplicial set is isomorphic to

$$(\underline{\mathcal{R}\mathcal{E}})^{N^A} \text{ a quasi-cat} \Rightarrow \text{it is a quasi-cat}$$

Thank you!

