

Adjunctions, limits

&

colimits

Def An adjunction $A \xrightleftharpoons[u]{f} B$ between ∞ -categories is comprised of:

- a pair of ∞ -categories A & B

- a pair of ∞ -functors

$$u: A \rightarrow B \quad \& \quad f: B \rightarrow A$$

- a pair of ∞ -natural transformations

$$\eta: id_B \Rightarrow uf \quad \& \quad \varepsilon: fu \Rightarrow id_A$$

So that

$$\begin{array}{ccc}
 \begin{array}{c} B \\ \xrightarrow{\quad u \quad} \\ A \end{array} & = & \begin{array}{c} B \\ \xrightarrow{\quad f \quad} \\ A \end{array} \\
 \begin{array}{ccc} u & \nearrow \varepsilon & \downarrow \eta \\ \downarrow & & \nearrow \\ f & & u \end{array} & = & u \left(\begin{array}{c} B \\ \xrightarrow{\quad = \quad} \\ A \end{array} \right) u \\
 u \xrightarrow{\eta} ufu \xrightarrow{-\varepsilon} u & &
 \end{array}$$

$$\begin{array}{ccc}
 \begin{array}{c} B \\ \xrightarrow{\quad f \quad} \\ A \end{array} & = & \begin{array}{c} B \\ \xrightarrow{\quad = \quad} \\ A \end{array} \\
 \begin{array}{ccc} f & \nearrow \eta & \downarrow \varepsilon \\ \downarrow & & \nearrow \\ u & & f \end{array} & = & f \left(\begin{array}{c} B \\ \xrightarrow{\quad = \quad} \\ A \end{array} \right) f \\
 f \xrightarrow{-\eta} fuf \xrightarrow{\varepsilon -} f & &
 \end{array}$$

Remark. In the the setting of (∞, n) - or (∞, ∞) -categories
this is "pseudo - style" adjunction



It is not most general adjunctions

But : its relationships to the equivalences



to the notions of (co-)limits

Lemma. An adjunction in a 2-category is preserved by any 2-functor

Example. Adjunction between 1-cats

$$\text{Cat} \xrightarrow{\circ} \text{hQCat}$$

$$A \begin{array}{c} \xleftarrow{\quad} \\[-1ex] \perp \\[-1ex] \xrightarrow{\quad} \end{array} B \iff N(A) \begin{array}{c} \xleftarrow{\quad} \\[-1ex] \perp \\[-1ex] \xrightarrow{\quad} \end{array} N(B)$$

regarded as a nerve

Example. Quillen adjunctions

Prop. Given an adjunction $A \begin{array}{c} \xleftarrow{f} \\[-1ex] \perp \\[-1ex] \xrightarrow{u} \end{array} B$ between ∞ -cats

Then

- $\forall \infty\text{-cat } X$

$$\mathrm{Fun}(X, A) \begin{array}{c} \xleftarrow{f_*} \\[-1ex] \perp \\[-1ex] \xrightarrow{u_*} \end{array} \mathrm{Fun}(X, B)$$

- $\forall \infty\text{-cat } X$

$$h\mathrm{Fun}(X, A) \begin{array}{c} \xleftarrow{f_*} \\[-1ex] \perp \\[-1ex] \xrightarrow{u_*} \end{array} h\mathrm{Fun}(X, B)$$

- \forall simplicial set T

$$A^T \begin{array}{c} \xleftarrow{f^T} \\[-1ex] \perp \\[-1ex] \xrightarrow{u^T} \end{array} B^T$$

- If the ambient ∞ -cosmos is cartesian closed, then $\forall \infty\text{-cat } C$

$$A^C \begin{array}{c} \xleftarrow{\quad} \\[-1ex] \perp \\[-1ex] \xrightarrow{\quad} \end{array} B^C$$

Prop. Adjunctions compose

$$\begin{array}{ccc} C & \xrightleftharpoons[f']{\perp} & B \\ u' \swarrow & & \searrow u \\ & \xrightleftharpoons[f]{\perp} & A \end{array} \rightsquigarrow \begin{array}{ccc} C & \xrightleftharpoons[ff']{\perp} & A \\ u'u \swarrow & & \searrow \end{array}$$

Prop. (uniqueness of adjoints)

- If $f \dashv u$ & $f' \dashv u \Rightarrow f \cong f'$
- Conversely, if $f \dashv u$ and $f \cong f' \Rightarrow f' \dashv u$

Lemma. (minimal adjunction data)

$$\begin{array}{ccc} & f & \\ A & \xleftarrow{\perp} & B \\ u & \xrightarrow{\quad} & \end{array} \Leftrightarrow \exists \text{ nat. transf. } id_B \Rightarrow uf \text{ so that } fu \Rightarrow id_A$$

the triangle equality composites are invertible:

$$f \Rightarrow fuf \Rightarrow f \quad \& \quad u \Rightarrow ufu \Rightarrow u$$

Proof: \Rightarrow Obvious

$$\Leftarrow \bullet \eta: id_B \Rightarrow uf \quad \& \quad \varepsilon': fu \Rightarrow id_A$$

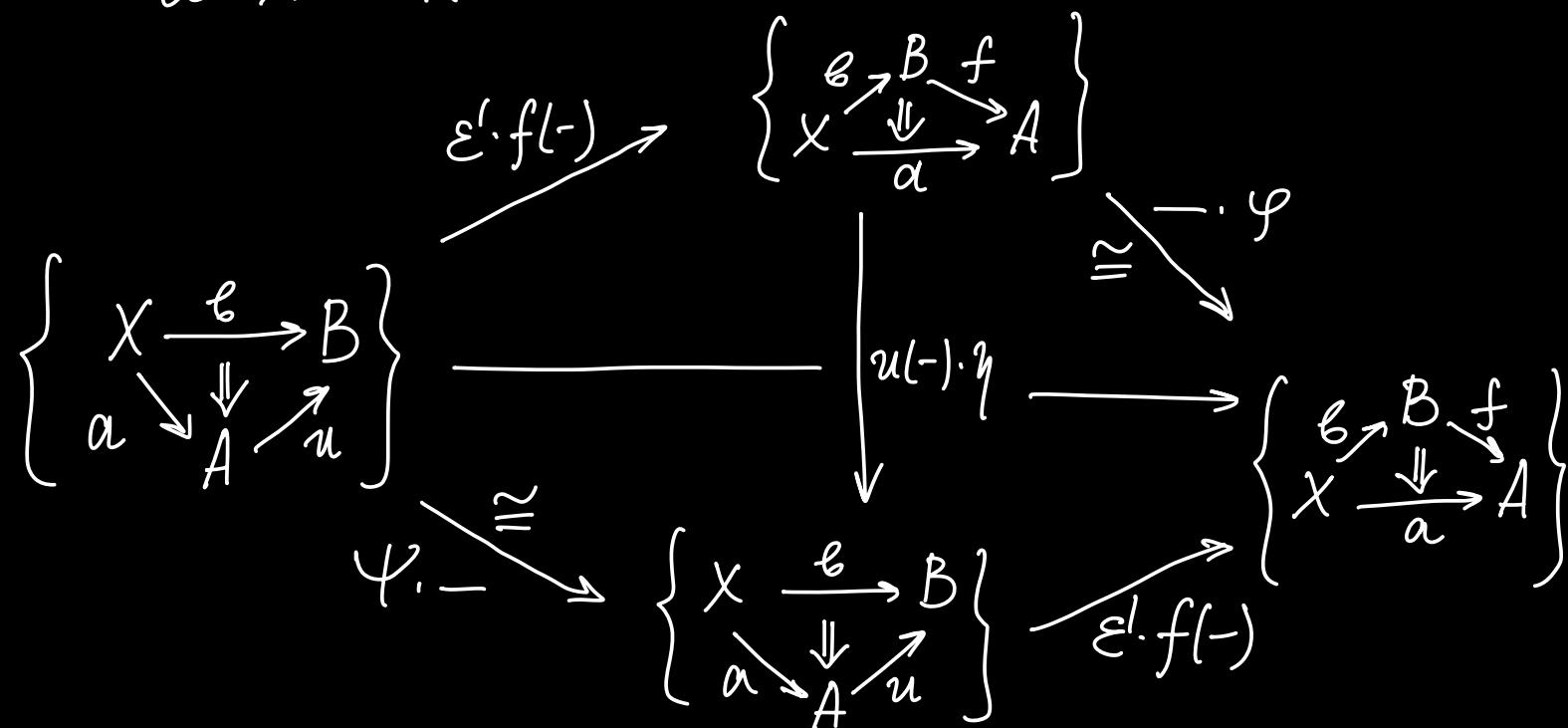
$$\varphi := f \xrightarrow{f\eta} fuf \xrightarrow{\varepsilon'f} f \quad \& \quad \psi := u \xrightarrow{\eta u} ufu \xrightarrow{u\varepsilon'} u$$

are isomorphisms

- Construct an adjunction $f \dashv u$ with unit η and counit ϵ'

- $b \Rightarrow ua \xrightarrow{f \circ -} fb \Rightarrow fua \xrightarrow{\epsilon'} fb \Rightarrow fua \xrightarrow{\epsilon'} a \rightsquigarrow fb \Rightarrow a$

- Let $b: X \rightarrow B$ $a: X \rightarrow A$ be fixed generalized elements



- By 2-of-6 property, we have all six morphisms being bijections

- Define

$$\mathcal{E} := \begin{array}{c} B \\ \nearrow u \quad \downarrow \varepsilon' \\ A \xrightarrow{\quad f \quad} A \\ \searrow \eta \quad \uparrow \psi^{-1} \\ B \end{array}$$

\cong

so that

$$\begin{array}{ccc} B & \xlongequal{\quad} & B \\ f \searrow & \nearrow u & \downarrow \varepsilon \\ A & \xlongequal{\quad} & A \end{array} = \begin{array}{c} B \\ \downarrow \eta \quad \nearrow u \\ A \xrightarrow{\quad f \quad} A \\ \uparrow \psi^{-1} \quad \downarrow \varepsilon' \\ B \end{array} = f \left(\begin{array}{c} B \\ \cong \\ A \end{array} \right) f$$

- Form the pasting equality

$$\begin{array}{ccc}
 \text{Diagram showing the pasting equality } (\psi \circ \eta u) \cdot \psi = \psi & = & \text{Simplified diagram showing } \psi := u \xrightarrow{\eta u} u f u \xrightarrow{u \varepsilon'} u
 \end{array}$$

The left side shows two overlapping regions, A and B. Region A is shaded green and contains a map u from A to B. Region B is shaded blue and contains a map ε' from B back to A. The intersection of A and B contains a map f from A to B. The boundary of A is labeled $u\varepsilon$. The boundary of B is labeled ηu . The overall mapping from A to B is labeled ψ . The right side shows the simplified form of the equality, where the complex overlapping regions are collapsed into a single mapping ψ from A to B.

- But ψ is invertible

$$\text{So, } u\varepsilon \cdot \eta u = \text{id}_u$$



Corollary. (adjoint equivalences) Any equivalence can be promoted to an adjoint equivalence by modifying one of the 2-cells, i.e., the invertible 2-cells in an equivalence can be chosen so as to satisfy the triangle equalities

f, g are inverse equiv. $\Rightarrow f \dashv g \And g \dashv f$

Proof:

$$\begin{array}{ccc}
 A & \begin{array}{c} \xrightarrow{f} \\ \perp \\ \xleftarrow{g} \end{array} & B \\
 A & \xrightarrow{\cong \Downarrow \alpha} & A \\
 B & \xrightarrow{\cong \Downarrow \beta} & B
 \end{array}$$

$$f \xrightarrow{\sim} fg \xrightarrow{\sim} f \And g \xrightarrow{\sim} gfg \xrightarrow{\sim} g$$

So, by the previous lemma f and g fit into an adjunction \square

Prop. Adjunctions are equivalence invariant

$u: A \rightarrow B$ admits a left adjoint $\Leftrightarrow \forall A' \cong A$

$\begin{array}{c} | \\ \text{---} \\ \infty\text{-cats} \end{array}$

$B' \cong B$

the functor

$u': A' \rightarrow B'$

admits a left adjoint

Proof: • $\exists f \dashv u$ with $A' \cong A$, $B' \cong B$ - adjoint equiv.

$$\begin{array}{ccccc} & & f & & \\ A' & \xleftarrow{\sim} & A & \xleftarrow{\perp} & B \\ \perp & \nearrow & \downarrow & \searrow & \perp \\ \sim & & u & & \sim \end{array}$$

• Conversely, if $u': A' \xrightarrow{\sim} A \xrightarrow{u} B \xrightarrow{\sim} B'$ admits a left adj f' then

$$f \Rightarrow fuf \Rightarrow f \quad \& \quad u \Rightarrow ufu \Rightarrow u$$

$$\begin{array}{ccccc} & & f' & & \\ A' & \xleftarrow{\sim} & A & \xleftarrow{\perp} & B \\ \perp & \nearrow & \downarrow & \searrow & \perp \\ \sim & & u & & \sim \end{array}$$



Initial & terminals elements

Def. An initial element in an ∞ -cat A is a left adjoint to the unique functor $! : A \rightarrow \mathbb{1}$
A terminal element is a right adj.

$$\begin{array}{ccc} \mathbb{1} & \begin{array}{c} \xrightarrow{i} \\ \perp \\ \xleftarrow{!} \end{array} & A \\ & \lrcorner & \urcorner \end{array} \qquad \begin{array}{ccc} \mathbb{1} & \begin{array}{c} \xleftarrow{!} \\ \perp \\ \xrightarrow{t} \end{array} & A \\ & \lrcorner & \urcorner \end{array}$$

Lemma. (^{minimal}_{data}) To define an initial element in an ∞ -cat A , it suffices to specify:

- an element $i : \mathbb{1} \rightarrow A$

- a nat. transf. $A \xrightarrow{\begin{array}{c} ! \\ \downarrow \varepsilon \\ i \end{array}} A$ so that $\varepsilon i : i \Rightarrow i$ is huzim.

Proof: • ∞ -cat $1 \in K$ is 2-terminal in the homotopy 2-cat hK

- $i: 1 \rightarrow A$ is a section of $!: A \rightarrow 1$

So, $!i = \text{id}_1$ & $\eta: 1 \xrightarrow{\sim} !i$ is an iso

We have a triangle

$$\begin{array}{ccc} & A & = A \\ i \nearrow & \downarrow \varepsilon \nearrow & \uparrow i \\ 1 & \xlongequal{\quad} & 1 \end{array} = i \left(\begin{array}{c} A \\ = \\ 1 \end{array} \right) i$$

$$\begin{array}{ccc} A & = A & = ! \left(\begin{array}{c} A \\ = \\ 1 \end{array} \right) ! \\ \downarrow \varepsilon / ! & \downarrow \eta \searrow & \downarrow \\ 1 & \xlongequal{\quad} & 1 \end{array}$$

$$i \xrightarrow{i\eta} i !i \xrightarrow{\varepsilon i} i \qquad ! \xrightarrow{\eta !} !i ! \xrightarrow{! \varepsilon} !$$

$\varepsilon: i! \rightarrow \text{id}$ this does not

$\eta: \text{id} \Rightarrow !i$ We want:

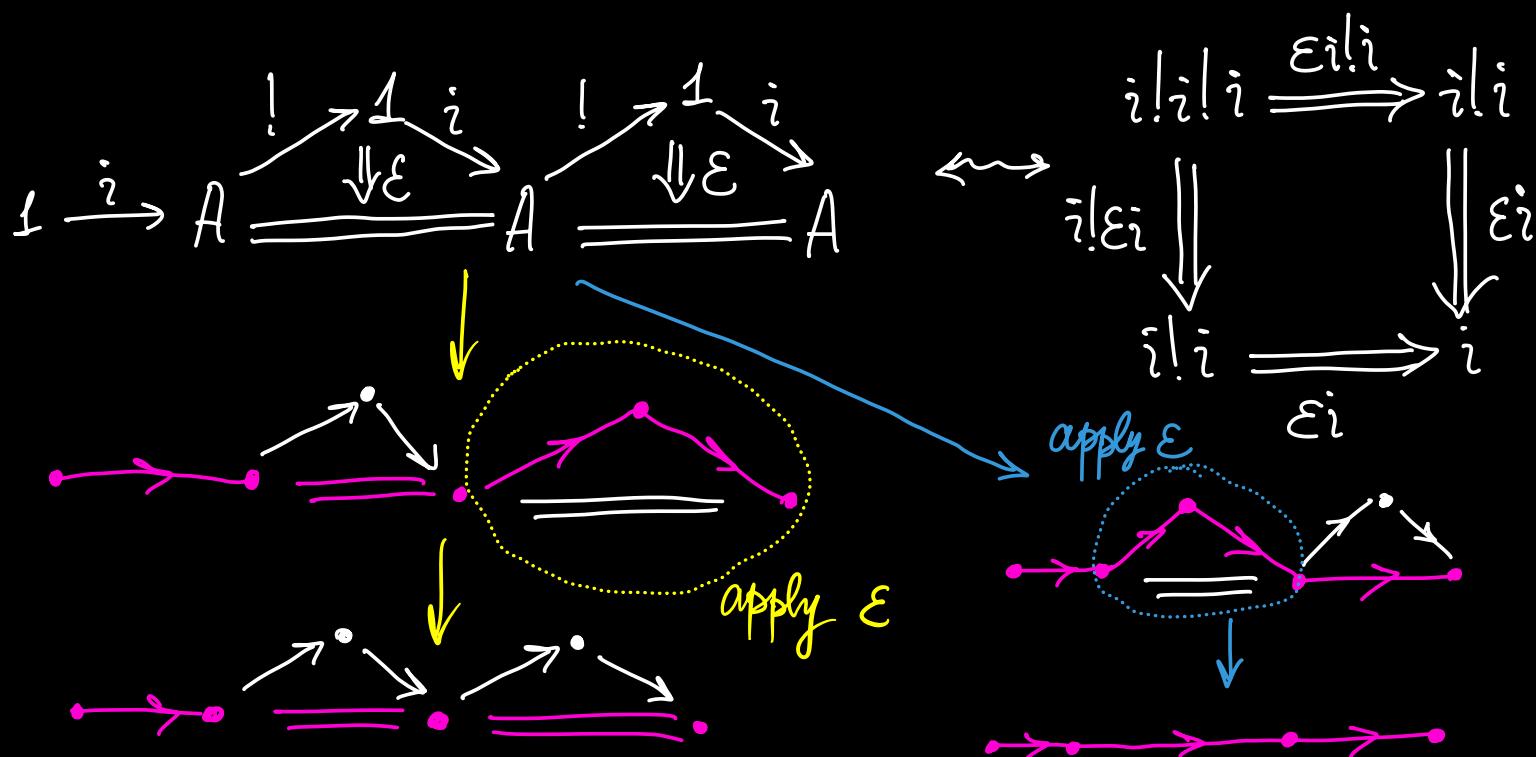
$$\varepsilon i = \text{id}_i$$

this holds automatically

$$! \varepsilon(a) = !(a)$$

- To do this, it suffices to require an isomorphism
 $\varepsilon_i : \tilde{i} \cong i$ (See the prop. above)

- Consider the horizontal composite and represent it as a vertical composite in two ways:



- So, $\varepsilon_i \cdot \varepsilon_i = \varepsilon_i \Rightarrow \varepsilon_i = \text{id}_i$ since ε_i is an isomorphism by the assumption

Lemma. (uniqueness) Any two initial elements in an ∞ -cat A are isomorphic in hA and if $\alpha \cong_i$ in hA $\Rightarrow \alpha$ is initial too

Proof:

- \forall left adjoints i and i' to $! : A \rightarrow 1$ are nat. isomorphic
- $\forall \alpha : 1 \rightarrow A$ that is isomorphic to left adj to $! : A \rightarrow 1$ is itself a left adjoint
- A functor isomorphism $i \cong i'$ gives $i \cong i'$ in hA

Lemma. An element $i: 1 \rightarrow A$ is initial $\Leftrightarrow \forall f: X \rightarrow A$
 \exists a unique 2-cell with boundary

$$\begin{array}{ccc} & 1 & \\ i \swarrow & \Downarrow \exists! & \searrow i \\ X & \xrightarrow{f} & A \end{array}$$

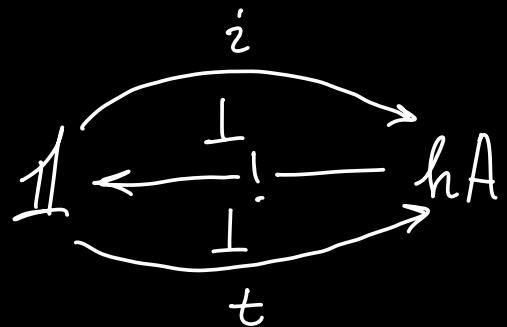
Proof.

- $i \dashv ! \rightsquigarrow 1 \cong h\text{Fun}(X, 1)$

$$h\text{Fun}(X, 1) \xrightleftharpoons[i_*]{!} \perp \xrightleftharpoons{i^*} h\text{Fun}(X, A)$$

So, the constant functor $i!: X \rightarrow A$ is initial in $h\text{Fun}(X, A)$

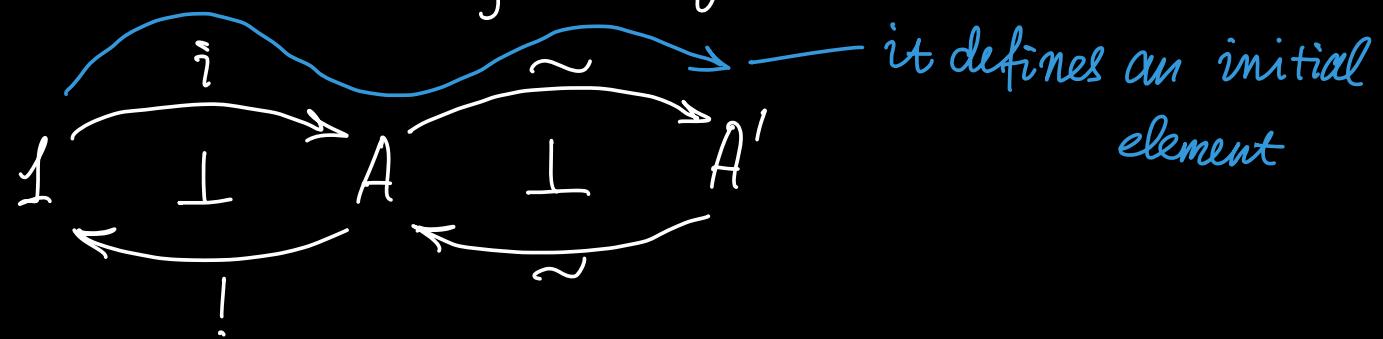
- Conversely, if $i: 1 \rightarrow A$ satisfies univ. property of the statement
 Then apply it to the $\text{id}_A: A \rightarrow A$ any by Lemma above
 We are done □



- Remark.
- Being homotopy initial is weaker than being initial in the ∞ -cat
 - But a homotopy initial element in a complete $(\infty, 1)$ -cat defines an initial element

Lemma. If A has an initial element and $A \cong A'$ then A' has an initial element and they are respected by the equivalence up to isomorphism

Proof: $A \cong A' \rightsquigarrow$ an adjoint equivalence



By the uniqueness of initial elements, the equivalence $A' \xrightarrow{\sim} A$ preserves initial elements \square

Limits & colimits

(co-)limits should be interpreted as homotopy ones in ∞ -cat

(Co-)limits of a diagram valued inside an ∞ -cat A in some ∞ -cosmos

→ indexed by a simpl. set J in an ∞ -cat A in a generic ∞ -cosmos

→ indexed by an ∞ -cat J and valued in an ∞ -cat A in a cartesian closed ∞ -cosmos

Def. (diagram ∞ -cat)

For $A \in K$, $J \in sSet$

or

$J \in K_{\text{cart. closed}}$

define a diagram of shape J in A as an element

$$d: 1 \rightarrow A^J$$

$$\begin{aligned} sSet^{\text{op}} \times \mathcal{K} &\rightarrow \mathcal{K} \\ (\mathcal{I}, A) &\mapsto A^{\mathcal{I}} \end{aligned}$$

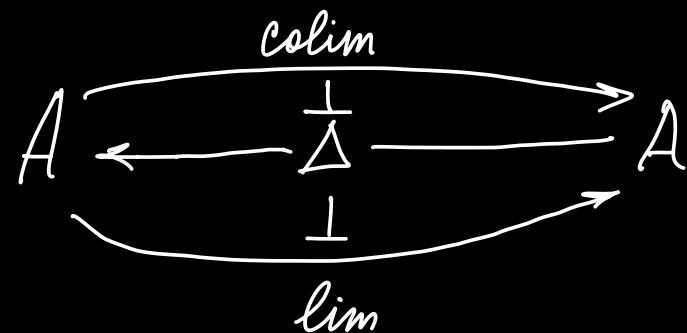
$$\begin{aligned} \mathcal{K}^{\text{op}} \times \mathcal{K} &\longrightarrow \mathcal{K} \\ (\mathcal{I}, A) &\longmapsto A^{\mathcal{I}} \end{aligned}$$

- Bifunctors

Apply Bifunctor to $! : \mathcal{I} \rightarrow \mathbf{1}$

$\Delta : A \rightarrow A^{\mathcal{I}}$ - constant diagram functor

Def. (limit & colimit functor)



Lemma. Products or coproducts in an ∞ -cat A also define define ones in its homotopy cat hA

Proof: . If \mathcal{I} is a set $\Rightarrow A^{\mathcal{I}} \cong \prod_{\mathcal{I}} A$

$$\begin{array}{ccc} h\mathcal{K} & \xrightarrow{h\text{Fun}(1,-)} & \text{Cat} \\ A & \longmapsto hA & \leftarrow \text{preserves products} \end{array}$$

$$h(A^{\mathcal{I}}) \cong h\left(\prod_{\mathcal{I}} A\right) \cong \prod_{\mathcal{I}} hA \cong (hA)^{\mathcal{I}}$$

\mathcal{I} is a set

So,

$$(hA)^{\mathcal{I}} \cong h(A^{\mathcal{I}})$$

Put $\mathcal{I} = \emptyset \triangleleft$

The definition is insufficiently general!

Def. (absolute lifting diagram)

Given a cospan $C \xrightarrow{g} A \xleftarrow{f} B$ in a 2-cat

An absolute left lifting of g through f is given by

a 1-cell ℓ & 2-cell λ :

$$\begin{array}{ccc} & B & \\ \ell \nearrow & \Downarrow \lambda & \downarrow f \\ C & \xrightarrow{g} & A \end{array} \quad \text{s.t.}$$

$$\begin{array}{ccc} X & \xrightarrow{\ell} & B \\ c \downarrow & \Updownarrow \chi & \downarrow f \\ C & \xrightarrow{g} & A \end{array} = \begin{array}{ccc} X & \xrightarrow{\exists! \uparrow \sum} & B \\ \downarrow & \ell \nearrow & \Downarrow \lambda \downarrow \\ C & \xrightarrow{g} & A \end{array} \quad \begin{array}{l} \text{Any 2-cell } \chi \\ \text{factors through} \\ (\ell, \lambda) \end{array}$$

an absolute right lifting of g through f :

$$\begin{array}{ccc} & B & \\ \lrcorner \nearrow \Downarrow \rho & \downarrow f & \\ C & \xrightarrow{g} & A \end{array}$$

$$\begin{array}{ccc} X & \xrightarrow{\ell} & B \\ c \downarrow & \Downarrow \chi & \downarrow f \\ C & \xrightarrow{g} & A \end{array} = \begin{array}{ccc} X & \xrightarrow{\ell} & B \\ c \downarrow & \exists! \Downarrow \zeta & \downarrow f \\ C & \xrightarrow{g} & A \end{array}$$

"Absolute" means

Lemma. Left (right) lifting diagrams are stable under restriction of their domain object:

If (ℓ, λ) - abs. left lifting of g through, then $\forall c : X \rightarrow C$ (ℓ_c, λ_c) defines an absolute left lifting of g_c through f

$$\begin{array}{ccc} X & \xrightarrow{c} & C \\ & \nearrow \ell & \nearrow \lambda \\ & \xrightarrow{g} & A \end{array}$$

Example: $\text{id}_B \Rightarrow u f$ is the unit of $f \dashv u$



(f, η) defines an absolute left lifting diagram

$$\begin{array}{ccc} & f & \nearrow \\ B & \xrightarrow{\quad \eta \quad} & A \\ & \text{id} & \downarrow \end{array}$$

Dually for $\varepsilon: f u \Rightarrow \text{id}_A$

$$\begin{array}{ccc} & u & \nearrow \\ A & \xrightarrow{\quad \varepsilon \quad} & B \\ & \text{id} & \downarrow \end{array}$$

Proof:

$$\begin{array}{ccc} X & \longrightarrow & B \\ \downarrow & \Downarrow \alpha & \downarrow \\ A & \xlongequal{\quad} & A \end{array} = \begin{array}{ccc} X & \xrightarrow{\exists! \beta} & B \\ \downarrow & u & \downarrow \varepsilon \\ A & \xlongequal{\quad} & A \end{array}$$

$$\alpha: f b \Rightarrow a \rightsquigarrow \beta: b \Rightarrow ua$$

$$\begin{array}{ccc}
 & \xleftarrow{\beta} & \xrightarrow{\alpha} \\
 \Downarrow & f_* & \Uparrow \\
 h\text{Fun}(X, B) & \xrightleftharpoons[\perp]{\quad} & h\text{Fun}(X, A) \\
 \Downarrow u_* & &
 \end{array}$$

If $f \dashv u$ then it is an adjunction and (u, ϵ) defines an absolute right lifting of id_A through f

Conversely, the unit & triangle identities of an adj. can be extracted from the univ. prop. of the right lifting diagram ▲

In particular, $\text{colim} \dashv \Delta \dashv \lim$ define

$$\begin{array}{ccc}
 \text{colim} & \nearrow A & \\
 A^J & \Downarrow \eta & \downarrow \Delta \\
 & = & \\
 & A^J &
 \end{array}
 \quad \& \quad
 \begin{array}{ccc}
 & \nearrow A & \\
 A^J & \Downarrow \epsilon & \downarrow \\
 & = & \\
 & A^J &
 \end{array}$$

These univ. properties are preserved under restriction. It motivates

Def. (limit & colimit) A colimit of a family of diagrams

$d: D \rightarrow A^{\mathcal{I}}$ of shape \mathcal{I} in an ∞ -cat A is given by
an absolute left lifting diagram

$$\begin{array}{ccc} & \text{colim } d & \nearrow A \\ D & \xrightarrow{\eta} & \downarrow \Delta \\ & \searrow A & \end{array}$$

comprised of a generalized element $\text{colim } d: D \rightarrow A$ &
a colimit cone $\eta: d \Rightarrow \Delta \text{colim } d$

Dually,

$$\begin{array}{ccc} \lim d & \nearrow A \\ D & \xrightarrow{\varepsilon} & \downarrow \Delta \\ & \searrow A & \end{array}$$

$\varepsilon: \Delta \lim d \Rightarrow d$ — a limit cone

a limit of a family of
diagrams $d: D \rightarrow A^{\mathcal{I}}$ of
shape \mathcal{I} in an ∞ -cat A

Example An initial element $i: 1 \rightarrow A$ can be regarded as a colimit of the empty diagram

The constant diagram functor:

$$!: A \rightarrow A^{\mathbb{J}} = A^{\emptyset} = 1$$

$$\begin{array}{ccc} 1 & \xrightarrow{i} & A \\ & \parallel & \downarrow ! \\ & = & 1 \end{array} \quad \begin{array}{c} X \xrightarrow{f} A \\ \downarrow ! \qquad \uparrow \chi \qquad \downarrow \\ 1 = 1 \end{array} = \begin{array}{c} X \xrightarrow{f} A \\ \downarrow ! \qquad \nearrow ? \qquad \downarrow ! \\ 1 \xrightarrow{i} \parallel 1 \end{array}$$

We want: an initial element defines an absolute left lifting diagram whose 2-cell is identity

But the existence & uniqueness of $\nearrow ?$ follow from initiality of $i: 1 \rightarrow A$ among all generalized elements $f: X \rightarrow A$

Example (exercise) In a cartesian closed ∞ -cosmos, $i: 1 \rightarrow A$ can be regarded as a limit of $\text{id}_A: A \rightarrow A$

Theorem Right adjoints preserve limits

Proof: • $\exists A$ admits limits of $d: 1 \rightarrow A^J$

$$\begin{array}{ccc} & \lim_d \nearrow & A \\ 1 & \Downarrow \lambda \downarrow & \downarrow \Delta \\ & d \searrow & A^J \end{array}$$

- $f \dashv u \rightsquigarrow f^J \dashv u^J$

- Show that

$$\begin{array}{ccccc} & \lim_d \nearrow & A & \xrightarrow{u} & B \\ 1 & \Downarrow \lambda \downarrow & \downarrow \Delta & & \downarrow \Delta \\ & d \searrow & A^J & \longrightarrow & B^J \end{array} \quad \text{is again an absolute right lifting diagram}$$

- Consider

$$\begin{array}{ccc}
 X & \xrightarrow{f} & B \\
 ! \downarrow \chi & \Downarrow \chi & \downarrow \Delta \\
 1 & \xrightarrow[d]{\quad} & A^J \xrightarrow{u^J} B^J
 \end{array}$$

- Add a square to the right hand side, composing with f

$$\begin{array}{ccccc}
 X & \xrightarrow{f} & B & \xrightarrow{f} & A \\
 ! \downarrow \chi & \Downarrow \chi & \downarrow \Delta & & \downarrow \Delta \\
 1 & \xrightarrow[d]{\quad} & A^J & \xrightarrow{u^J} & B^J \xrightarrow{f^J} A^J \\
 & & & \Downarrow \varepsilon^J & \\
 & & & \curvearrowright & \\
 & & & &
 \end{array}$$

\Updownarrow just compose
& factorize

- Add a square to the right hand side, comparing with α

$$\begin{array}{ccc}
 \begin{array}{c}
 \begin{array}{ccccc}
 X & \xrightarrow{\epsilon} & B & \xrightarrow{f} & A \\
 \downarrow ! & \nearrow \exists! \Downarrow \zeta & \downarrow \lambda & \downarrow \Delta & \downarrow \Delta \\
 1 & \xrightarrow{d} & A^J & \xrightarrow{u^J} & B^J
 \end{array} & = &
 \begin{array}{ccccc}
 X & \xrightarrow{\epsilon} & B & \xrightarrow{f} & A \\
 \downarrow ! & \Downarrow \chi & \downarrow \Delta & \downarrow \Delta & \downarrow \Delta \\
 1 & \xrightarrow{d} & A^J & \xrightarrow{u^J} & B^J \\
 & & \downarrow \varepsilon^J & \xrightarrow{f^J} & A^J \xrightarrow{u^J} B^J
 \end{array}
 \end{array}
 \end{array}$$

$$\begin{array}{ccc}
 & = &
 \begin{array}{ccccc}
 X & \xrightarrow{\epsilon} & B & = & B \\
 \downarrow ! & & \downarrow \Delta & & \downarrow \Delta \\
 1 & \xrightarrow{d} & A^J & \xrightarrow{u^J} & B^J \\
 & & \downarrow \varepsilon^J & \xrightarrow{f^J} & A^J \xrightarrow{u^J} B^J
 \end{array}
 \end{array}$$

apply a triangle identity of $f^J \dashv u^J$

So, χ factors through the composite of ζ and η

Uniqueness
is left as an exercise

Corollary Equivalences preserve limits

Moreover,

Prop. An equivalence $f: A \xrightarrow{\sim} B$ preserves, reflects and creates limits and colimits

Proof: Prove that f reflects limits and colimits

- Consider a family of diagrams $d: K \rightarrow A^J$ that admits limits in B

$$\begin{array}{ccccc} & \text{lim} & & & \\ K & \xrightarrow{d} & A^J & \xrightarrow{\sim} & B^J \\ & & \Downarrow \lambda & & \downarrow \Delta \end{array}$$

after comp.
with $A \xrightarrow{\sim} B$

- The composite 2-cell

$$\begin{array}{ccccc}
 & \text{lim} & & & \\
 K & \xrightarrow{d} & A^J & \xrightarrow{\sim} & B \\
 & & \downarrow \lambda & & \downarrow \Delta \\
 & & B^J & \xrightarrow{\sim} & A^J \\
 & & \downarrow \cong & &
 \end{array}$$

again defines
an absolute
right lifting
diagram

The univ. prop. can be verified



Thank you!