

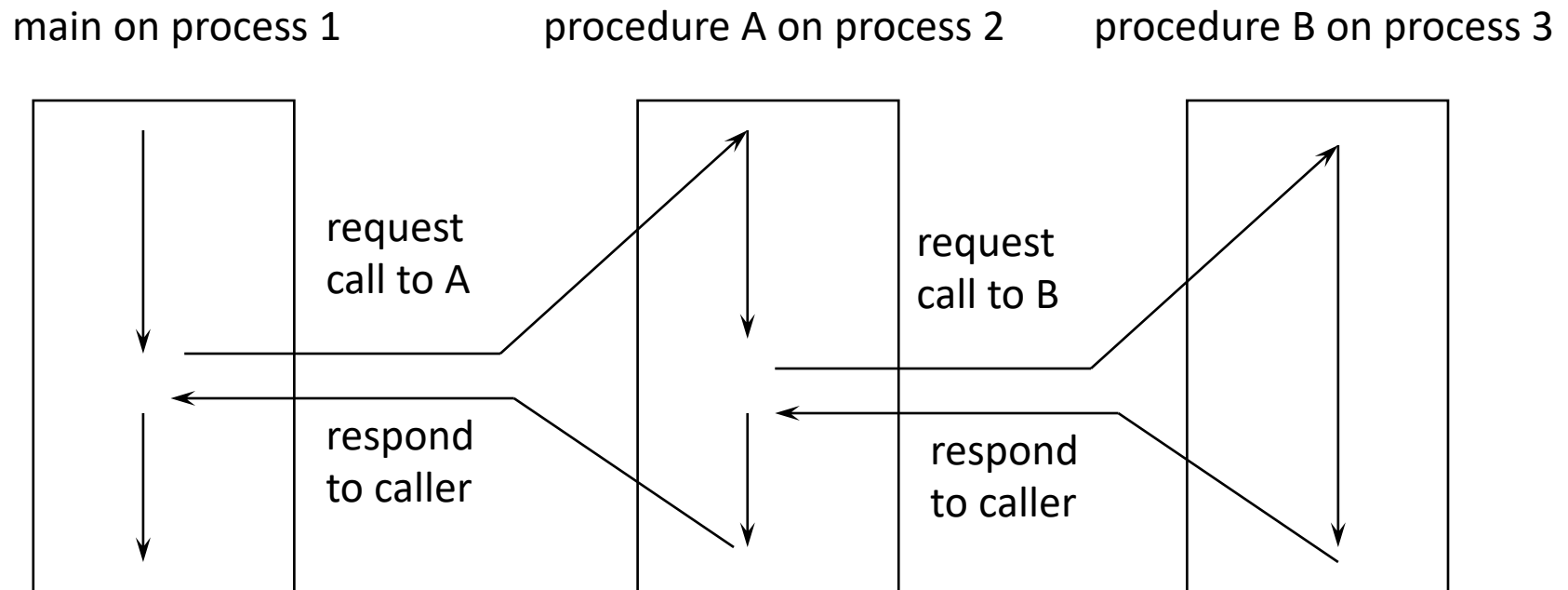
Remote Procedure Call (RPC)

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Reference for study: Van Steen, Tanenbaum,
"Distributed Systems", chapter 4.2

What is Remote Procedure Call?

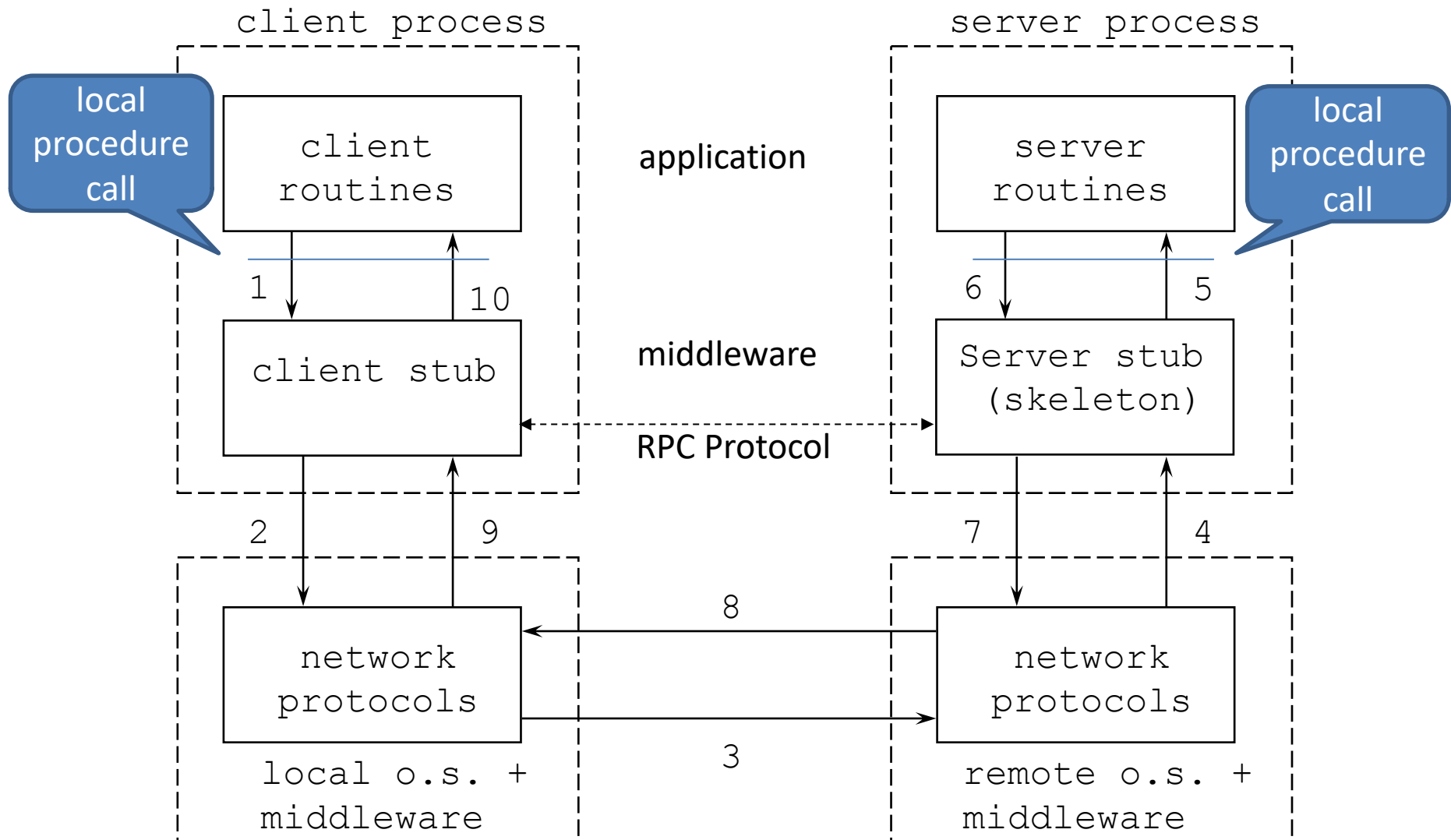
- RPC is the transfer of the (local) procedure call mechanism to a **distributed environment**:



Examples of RPC Implementations

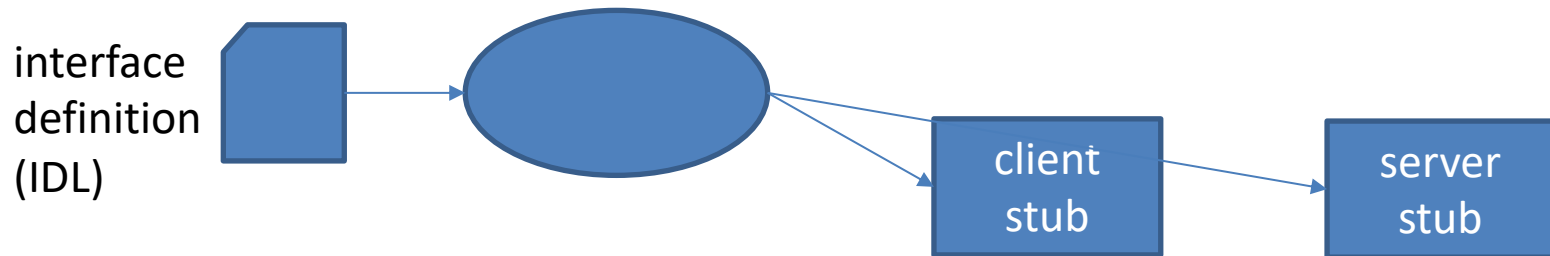
• SUN RPC (ONC)	C Language	 Language independent
• JAVA RMI	Java Language	
• CORBA	Distributed Objects	
• Web services	Distributed Services	
• The Web (REST)	Distributed Resources	
• gRPC		

The RPC model



The RPC Middleware

- Language-based support
 - RPC is offered as a feature of a programming language (e.g. Java RMI), exploiting language-specific serialization
 - high transparency, but limited to single-language apps
- Stub Generation



- stubs can be generated for multiple languages

Issues to be handled by the RPC Middleware

- Caller and callee are **heterogeneous** (different languages, hw/sw platforms, etc.)
 - procedure arguments and return values must be converted (*marshaling/unmarshalling*)
- Caller and callee run on **different** processes and hosts
 - the server must be *located* before issuing the call (*dynamic linking*)
 - multiple incoming calls must be *synchronized* on the server
 - if and how arguments can be passed by reference must be defined (*address spaces are disjoint*)
 - partial failure is possible and must be managed

Managing Partial Failure (Remote Call Semantics)

- Remote call semantics differ from local call semantics:
 - A remote call may **fail** (server crashes, network outages)
 - A remote call may be executed **more than once** (message duplications)
- RPC middleware protocols can handle such problems only **partially**
 - ⇒ programmers must be aware of them
 - ⇒ Normally, two kinds of return are possible:
 - successful call
 - unsuccessful call

Remote call semantics

- According to what RPC protocol is used, different call semantics are obtained.
- Call Semantics Definitions:
 - **at least once**
 - **at most once**
 - **exactly once**

Typical RPC Implementation with At Least Once Semantics Calls

- Upon successful return
=> at least once semantics
- Upon unsuccessful return
=> no guarantee given (the procedure may have been executed 0 or more times)
- Procedures must be **idempotent** (i.e. not sensitive to number of calls).
 - Often procedures can be made idempotent. Example:
append X to file Y => write X at offset k in file Y

- Transport: UDP
- RPC protocol includes a way to associate response to request (and manages retransmissions)

Typical RPC Implementations with At Most Once Semantics Calls

- The client is guaranteed that, in any case, the remote procedure will be executed at most once:
 - Upon successful return
 - exactly once semantics
 - Upon unsuccessful return
 - at most once semantics

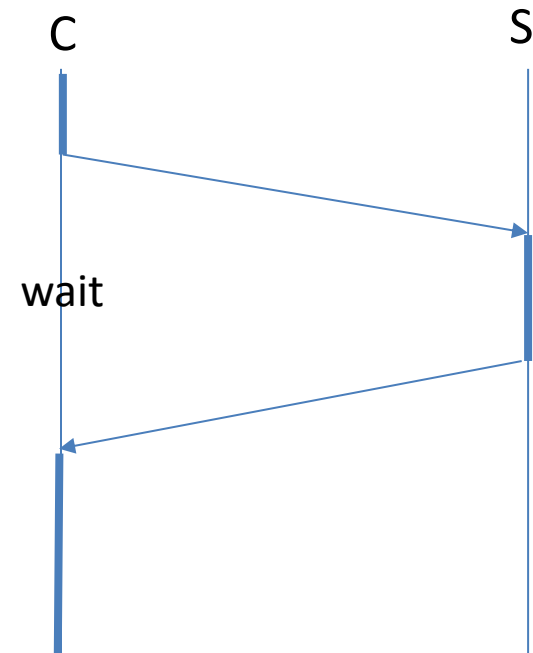
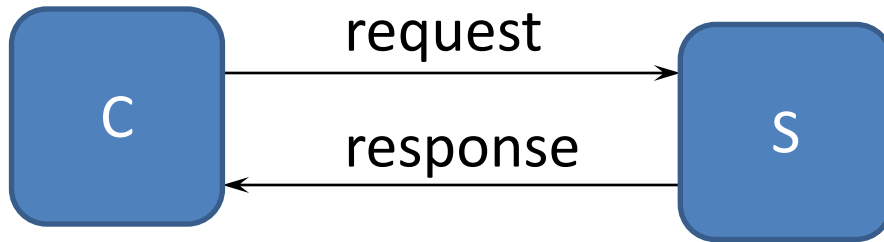
- Transport: TCP
- RPC protocol associates response to request via the connection

- Transport: UDP
- RPC protocol includes
 - response-request association
 - retransmission management
 - duplicate checking at server

Other RPC issues: Security and Performance

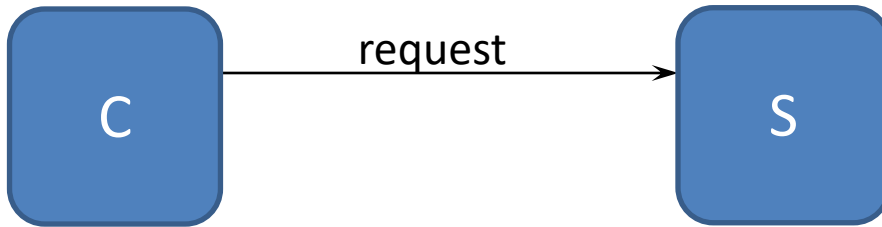
- Main *security requirements*:
 - access control
 - authentication
 - confidentiality
 - integrity
- Main *performance issues*:
 - a remote call is **orders of magnitude** slower than a local call

Classical (Synchronous) RPC

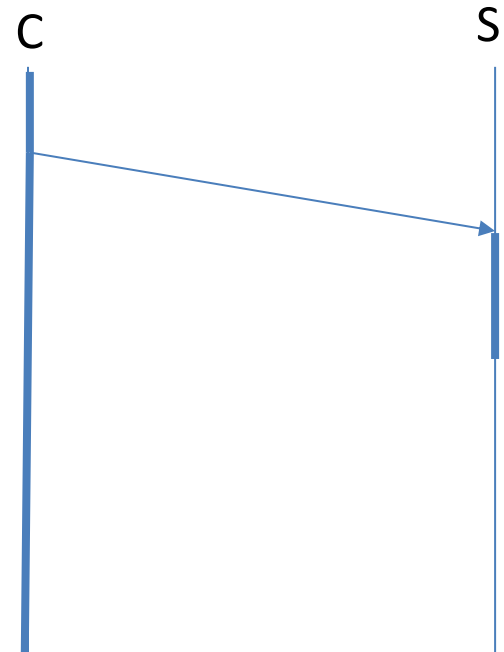


- Request:
 - Procedure identification
 - Input parameters
- Response:
 - Success/Failure indication
 - Return value and output parameters (or exceptions)

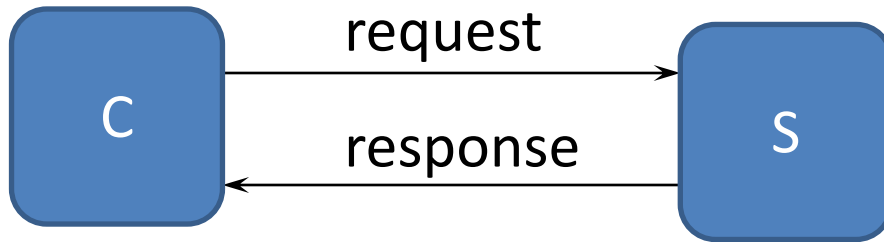
One-Way (Asynchronous) RPC



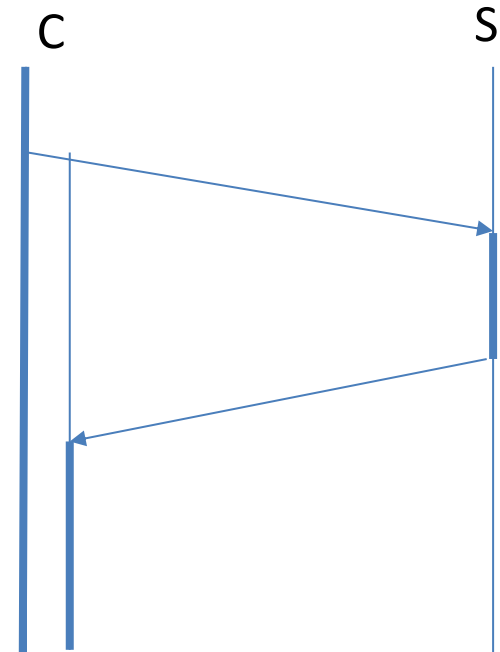
- No output arguments/return value
- No exceptions



Two-way Asynchronous RPC

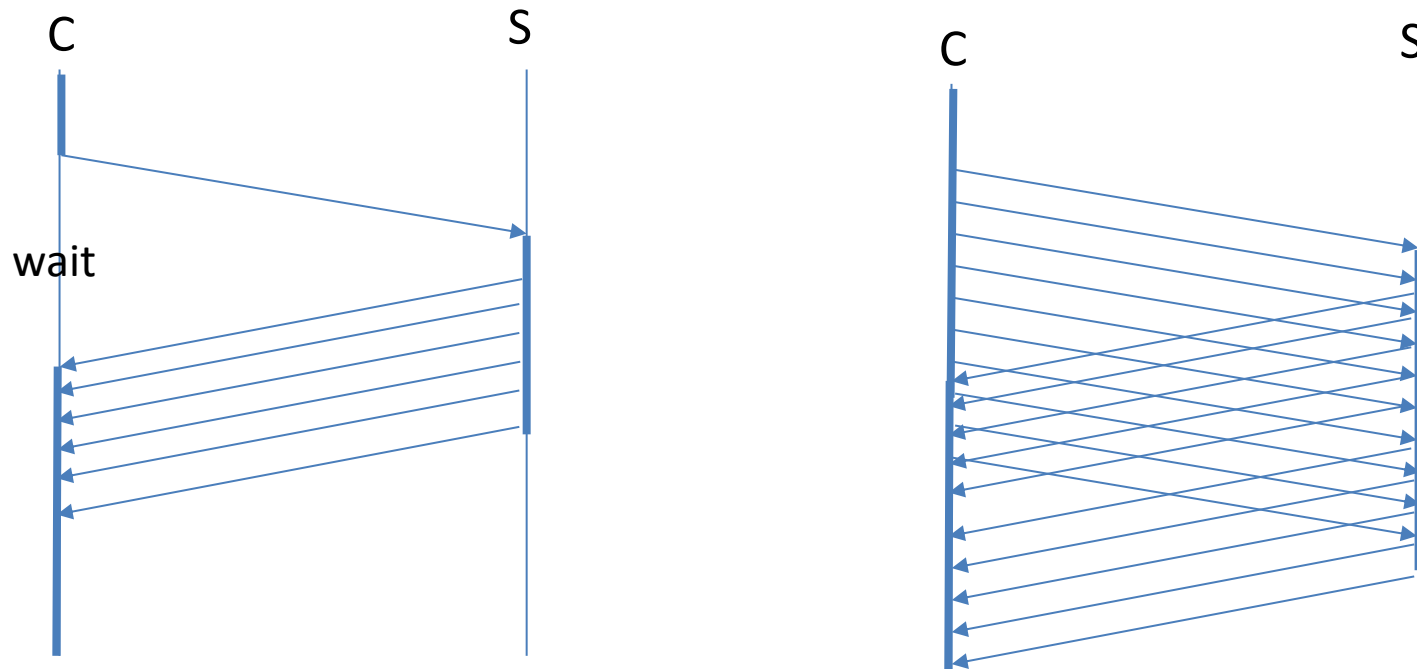


- Like Synchronous RPC
- Same protocol, but:
 - Client does not wait
 - Callback manages the response



RPC with Streaming

- All flavors of RPC can exploit **streaming** for request and/or response:
 - a request or response can be a *stream* of messages rather than a *single* message



Distributed Applications Developed on top of Socket APIs

- Sockets only provide communication services
- The implementation of the interaction protocol is up to the programmer:
 - data encoding for exchanged data
 - interaction procedures, management of partial failures, etc.
- Emphasis is more on communication than on application logic
- Normally used to implement
 - middleware
 - custom application-level protocols

Developing Applications on top of RPC

- The application can be developed as normal centralized code (but having in mind that interaction will be distributed!)
- Application modules are distributed on multiple hosts, and the application is tested in the distributed environment.
- The programmer can focus more on application logic than on communication

Developing Services on top of RPC

- Applications are conceived as combinations of services
- Services are designed and developed as reusable components that interact via RPC
- Advantages:
 - services are simple reusable components
 - scaling out can be achieved by creating multiple service instances (based on request volume)