OS 161:

Address Space & Memory Management



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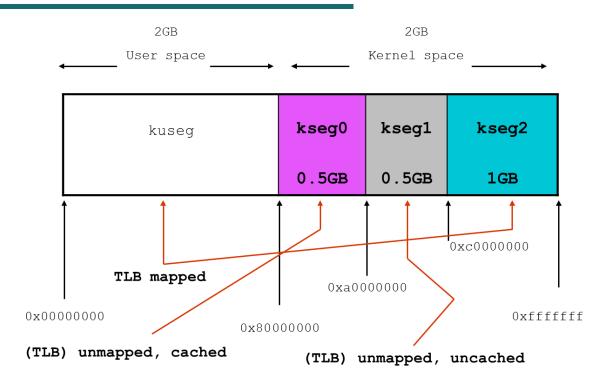
Dumbym and kmalloc

- Memory management of OS161 consists of two elements:
 - Kmalloc allows memory allocation of kernel side
 - Dumbvm memory manager user side
- Memory manager of OS161 use Contiguous allocation
- However, it also uses pagination (minimum amount of available memory is page(s))

Dumbym and kmalloc

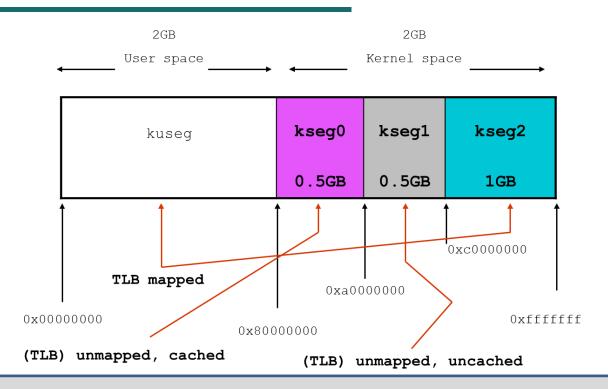
- Allocating by page multiples (4096 byte frame)
- Allocation of memory is done in two levels:
 - getppages (dumbvm.c): calls ram_stealmem (in mutual exclusion) (more generic abstraction)
 - ram_stealmem (ram.c): allocates contiguous RAM starting at firstpaddr, that is increased (containing codes architecture dependent)
- Allocator is common to (for both user and kernel)
 - User memory: as_prepare_load calls getppages for 2 user segments and a stack
 - Dynamic kernel memory: kmalloc is based on alloc_kpages, that calls getppages

MIPS VIRTUAL ADDRESS SPACE



OS161 is running on Mips, a 32 bits processor and addresses at 32 bits. Therefore, logical address space of 4GB.

MIPS VIRTUAL ADDRESS SPACE

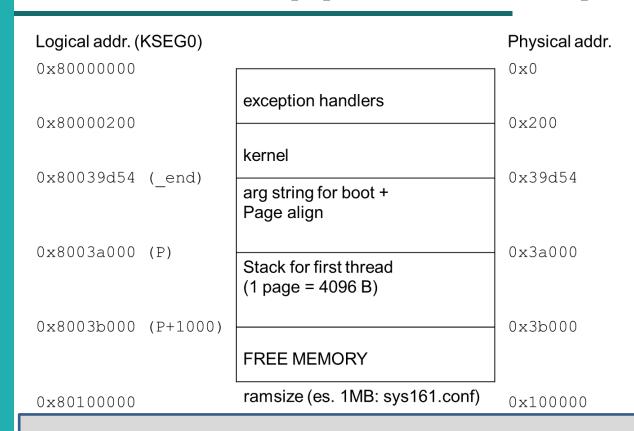


Mips maps the kernel of OS in the logical memory space of the process. So the process sees the user and also kernel in its memory space. (When boot of os161 is done, the kernel is mapped in a physical address that then later, through tlb, it is mapped in this logical memory of the process)

The kuseg is for user memory.

Kseg0 is for kernel, kseg01 I/O devices, and kseg02 is not used.

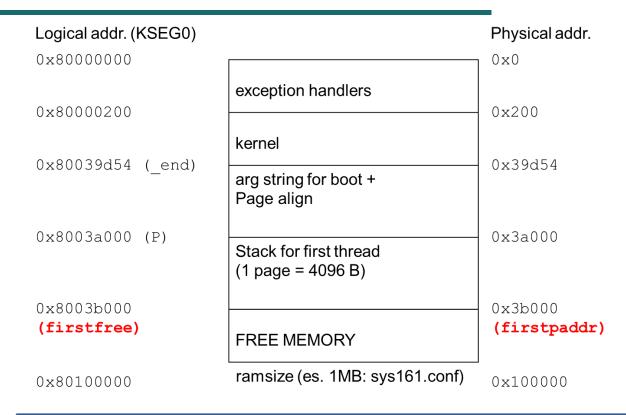
Kernel loader (sys161: start.S)



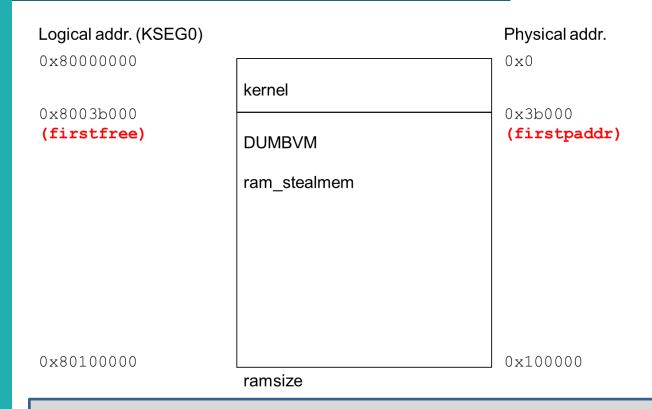
What happens during the boot time of Os161?

Using part of the memory for loading the kernel of OS161, starting from address 0. Starting from exception handlers, kernel, the memory space for the argument of command line of kernel (it is empty right now), stack for the first process that will be created, and free physical memory).

Kernel loader (sys161: start.S)



At the end of the boot phase, OS saves the first available physical address (Free Memory)



Complete layout of physical memory:

Starting from the first free address (first physical address), arriving to the dimension of the memory.

The first free address depends on the size of the kernel while the last depends on the ram size.

ram_bootstrap

```
void
ram bootstrap(void) {
  /* Get size of RAM. */
  size t ramsize = mainbus ramsize();
  if (ramsize > 512*1024*1024) {
    ramsize = 512*1024*1024;
  lastpaddr = ramsize;
  /* Get first free virtual address from where
     start.S saved it. Convert to physical address. */
  firstpaddr = firstfree - MIPS KSEGO;
```

During the boot, the ram_bootstrap is called which is initializing different parameters such as ramsize using function mainbus_ramsize, saving lastpaddr as the ram size, and calculating the first free address with respect to KSEGO.

ram_stealmem (kern/arch/mips/vm/ram.c)

```
paddr_t ram_stealmem(unsigned long npages) {
  paddr_t paddr;
  size_t size = npages * PAGE_SIZE;
  if(firstpaddr + size > lastpaddr) {
    return 0;
  }
  paddr = firstpaddr;
  firstpaddr += size;
  return paddr;
}
```

In **ram.c**, the most important function is **ram_stealmem** that allows to ask for memory.

Npages is the number of pages to ask the OS.

Calculating the size as the multiply of number of pages and page size.

ram_stealmem (kern/arch/mips/vm/ram.c)

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    return 0;
  }
  paddr = firstpaddr;
  firstpaddr += size;
  return paddr;
}
```

Checking whether the first free address + the requested size is more than the last physical address. If yes, retuning 0 which shows failing in memory allocation

ram_stealmem (kern/arch/mips/vm/ram.c)

```
paddr_t ram_stealmem(unsigned long npages) {
  paddr_t paddr;
  size_t size = npages * PAGE_SIZE;
  if(firstpaddr + size > lastpaddr) {
    return 0;
  }
  paddr = firstpaddr;
  firstpaddr += size;
  return paddr;
}
```

Checking whether the first free address + the requested size if more than the last physical address. If not, successful memory allocation, returning the address of the first available physical free memory.

Ram.c is regarding the architecture dependent part. For example, for changing from continues allocation to another one, this module should be changed.

```
static paddr_t
getppages(unsigned long npages) {
  paddr_t addr;
  spinlock_acquire(&stealmem_lock);
  addr = ram_stealmem(npages);
  spinlock_release(&stealmem_lock);
  return addr;
}
```

Getppages is for asking numbers of physical pages, asking for page number and returning pointer to the zone of the memory.

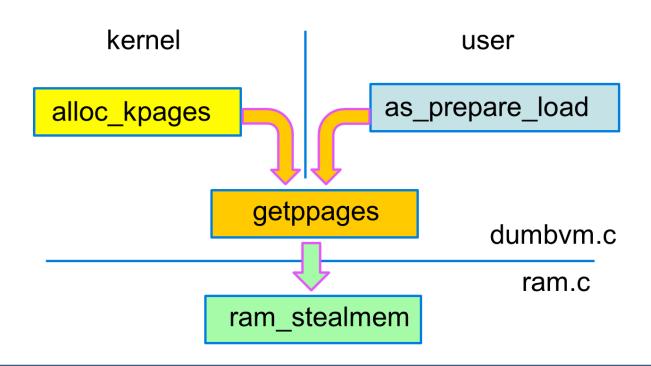
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  spinlock_release(&stealmem_lock);
  return addr;
}
```

Spinlock allows the mutual exclusion in case of multiple process, to avoid corrupting of ram_stealmem function.

```
static paddr t
getppages(unsigned long npages) {
  paddr t addr;
                                 Internal (dumbvm) function
  spinlock acquire(&stealmem lock);
  addr = ram stealmem(npages);
  spinlock release(&stealmem lock);
  return addr;
```

```
static paddr t
                                   Mutual exclusion for
getppages (unsigned long npages)
                                      ram_stealmem
  paddr t addr;
  spinlock acquire(&stealmem lock);
  addr = ram stealmem(npages);
  spinlock release(&stealmem lock);
  return addr;
```

Dumbvm.c (alloc)

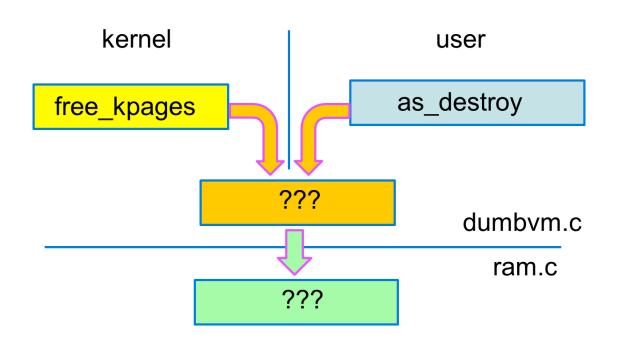


Memory manager architecture of OS161:

ram_stealmem for memory management at low level (architecture dependent) getppages for memory management in dumbvm.c

Allocation function for user side (as_prepare_load) and kernel side (alloc_kpages)

Dumbvm.c – Not implemented >> to do



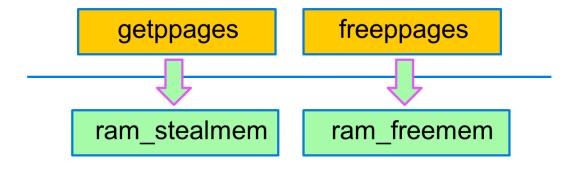
The function for allocation of memory is implemented BUT the functions for releasing memory is not implemented.

free_kpages and **as_destroy** (equivalent of alloc_kpages and as_prepare_load) is defined but EMPTY.

Your Job to solve this issue.

De-alloc (free) in ram.c

- freeppages just an interface to ram freemem
- Data structure (free-list or bitmap) and memory management in ram.c

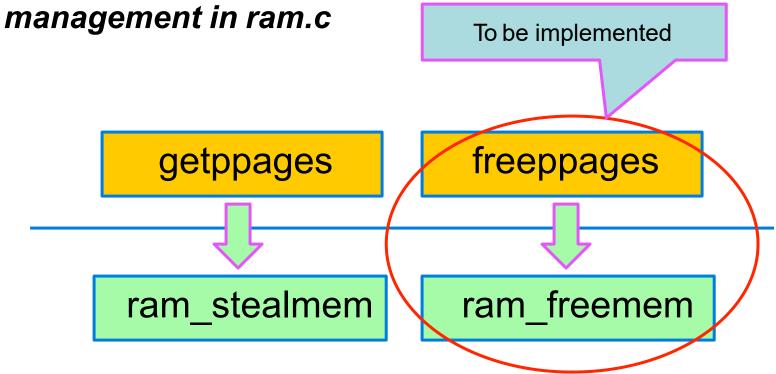


Defining the equivalent system of ram_stealmem and getppages,

Defining a freeppages that we be called by free_kpages and as_destroy (high level)
that will call ram_freemem for freeing memory.

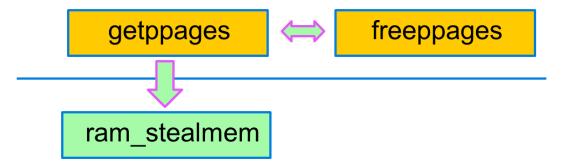
De-alloc (free) in ram.c

- freeppages just an interface to ram_freemem
- Data structure (free-list or bitmap) and memory



De-alloc (free) in ram.c

- Memory not returned to RAM
- Data structure (free-list o bitmap) and memory management in dumbvm.c
- Freeppages coordinates with getppages



Managing the liberation of memory at low level but just at the dumbvm.c level. Eveytime that we are asking for memory allocation, before asking ram_stealmem, looking if the memory page that has been already allocated, is free now. Vantages: working with one system call

Paging in user space

(solution B)

kernel user as_prepare_load alloc_kpages Contiguous alloc for **PAGING** kernel

Paging in user space

(solution B)

kernel user as prepare load alloc kpages Contiguous alloc for **PAGING** kernel Two allocators:

- Paginn for user processes
- Contiguous memory for kernel

Proposed Solution

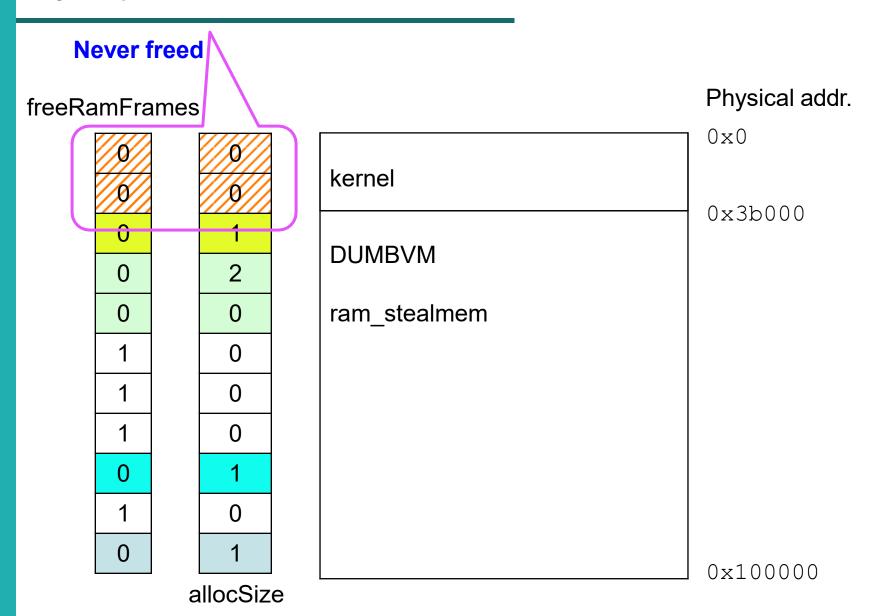
- Contiguous allocation(by pages) common to kernel and user
- Allocator in dumbvm: keep track (using a bitmap) of previously freed pages. In order to alloc
 - When calling getppages, first search among (previously) freed pages (an interval of contiguous free pages)
 - If not found, call ram stealmem
- Bitmap implemented as an array of char (for simiplicity)
 - freeRamFrames[i] = 1/0 (free/alloated): free=FREED! (by freeppages)
- In order to free we need to know
 - Pointer (or index) to first page in interval
 - Size, i.e. number of (contiguous) pages to free
- We need a table to store sizes (number of pages in allocated intervals) for each alloc performad
 - void free kpages (vaddr t addr): table needed as only pointer passed
 - void as_destroy(struct addrspace *as): table not needed ad size is stored in address space
- allocSize[i] = /* numbeer of pages allocated starting at i-th */

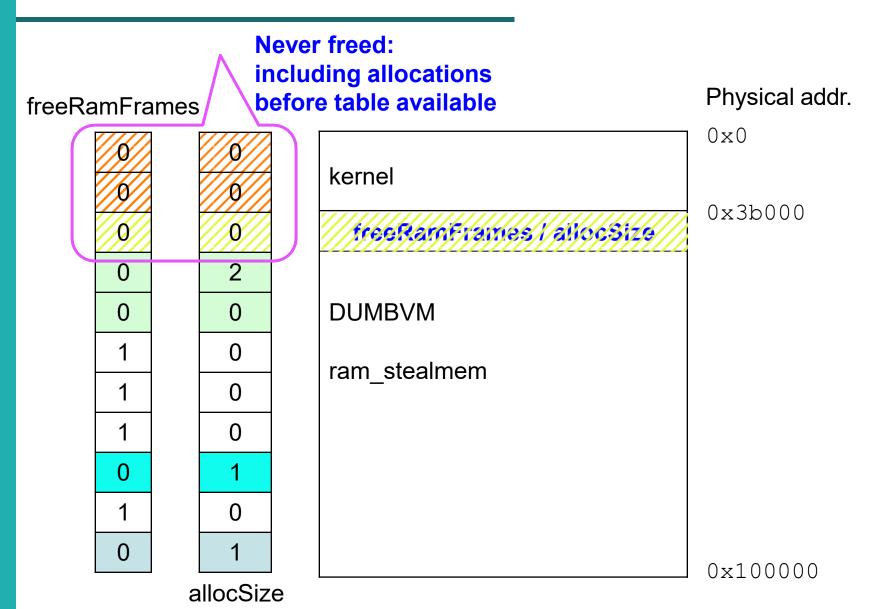
Global Variables and Test Function

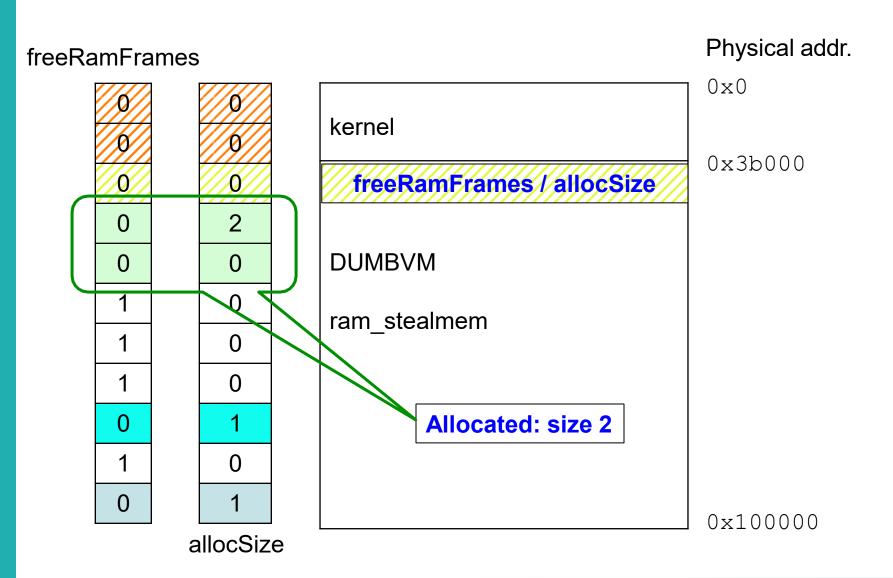
```
static struct spinlock freemem lock = SPINLOCK INITIALIZER;
static unsigned char *freeRamFrames = NULL;
static unsigned long *allocSize = NULL;
static int nRamFrames= 0;
static int allocTableActive = 0;
static int isTableActive () {
  int active;
  spinlock acquire(&freemem lock);
  active = allocTableActive;
  spinlock release(&freemem lock);
  return active;
```

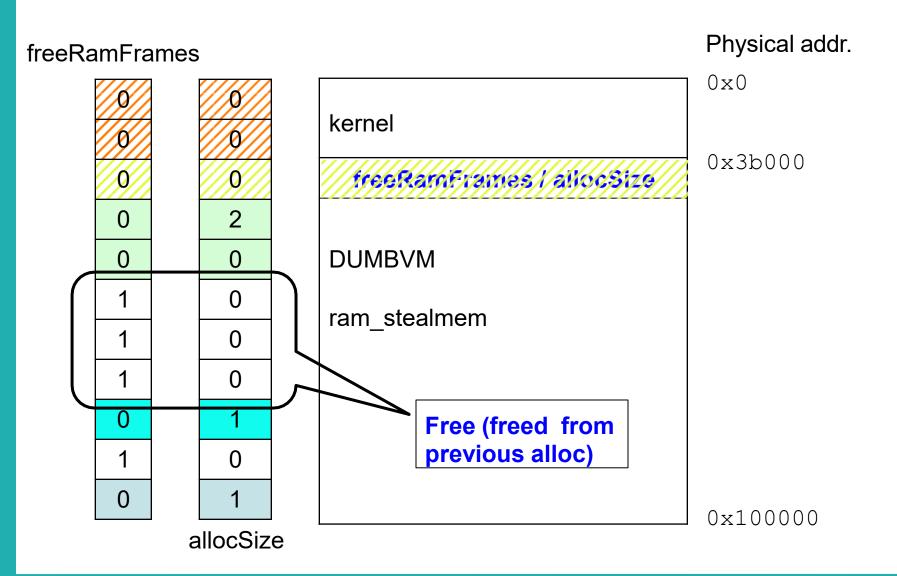
Global Variables and Test Function

```
static struct spinlock freemem lock = SPINLOCK INITIALIZER;
static unsigned char *freeRamFrames = NULL;
static unsigned long *allocSize = NULL;
static int nRamFrames= 0;
static int allocTable
static int
  int active D, ynamic arrays as RAM size known
           at Boot (depends on da sys161.conf)
  spinlock
  active = Alternative: over-dimensioned static arrays!
  spinlock release(&freemem lock);
  return active;
```









Free_kpages & as_destroy

```
void free kpages(vaddr t addr) {
  if (isTableActive()) {
    paddr t paddr = addr - MIPS KSEG0;
    long first = paddr/PAGE SIZE;
    KASSERT (nRamFrames>first);
    freeppages(paddr, allocSize[first]);
void as destroy(struct addrspace *as) {
  dumbvm can sleep();
  freeppages (as->as pbase1, as->as npages1);
  freeppages (as->as pbase2, as->as npages2);
  freeppages(as->as stackpbase, DUMBVM STACKPAGES);
  kfree (as);
```

Initialization

```
void vm bootstrap(void) {
  int i;
  nRamFrames = ((int)ram getsize())/PAGE SIZE;
  /* alloc freeRamFrame and allocSize */
  freeRamFrames = kmalloc(sizeof(unsigned char)*nRamFrames);
  allocSize = kmalloc(sizeof(unsigned long)*nRamFrames);
  if (freeRamFrames==NULL | | allocSize==NULL) {
    /* reset to disable this vm management */
    freeRamFrames = allocSize = NULL; return;
  for (i=0; i<nRamFrames; i++) {</pre>
    freeRamFrames[i] = (unsigned char)0; allocSize[i] = 0;
  spinlock acquire(&freemem lock);
  allocTableActive = 1;
  spinlock release(&freemem lock);
```

getppages

```
static paddr t getppages (unsigned long npages) {
 paddr t addr;
  /* try freed pages first */
  addr = getfreeppages(npages);
  if(addr== 0) {/* call stealmem */
    spinlock acquire(&stealmem lock);
    addr = ram stealmem(npages);
    spinlock release(&stealmem lock);
  if(addr != 0 && isTableActive()) {
    spinlock acquire (&freemem lock);
    allocSize[addr/PAGE SIZE] = npages;
    spinlock release(&freemem lock);
  return addr;
```

getfreeppages

```
static paddr t getfreeppages (unsigned long npages) {
 paddr t addr;
  long i, first, found, np = (long)npages;
  if (!isTableActive()) return 0;
  spinlock acquire(&freemem lock);
  // Linear search of free interval
  for (i=0, first=found=-1; i<nRamFrames; i++) {</pre>
    if (freeRamFrames[i]) {
      if (i==0 | !freeRamFrames[i-1])
        first = i; /* set first free in an interval */
      if (i-first+1 >= np)
        found = first;
```

getfreeppages

```
if (found \ge 0) {
  for (i=found; i<found+np; i++) {</pre>
    freeRamFrames[i] = (unsigned char)0;
  allocSize[found] = np;
  addr = (paddr t) found*PAGE SIZE;
else {
  addr = 0;
spinlock release(&freemem lock);
return addr;
```

getfreeppages

```
static int freeppages (paddr t addr, unsigned long npages) {
  long i, first, np=(long)npages;
  if (!isTableActive()) return 0;
  first = addr/PAGE SIZE;
  KASSERT(allocSize!=NULL);
  KASSERT (nRamFrames>first);
  spinlock acquire(&freemem lock);
  for (i=first; i<first+np; i++) {</pre>
    freeRamFrames[i] = (unsigned char)1;
  spinlock release(&freemem lock);
  return 1;
```