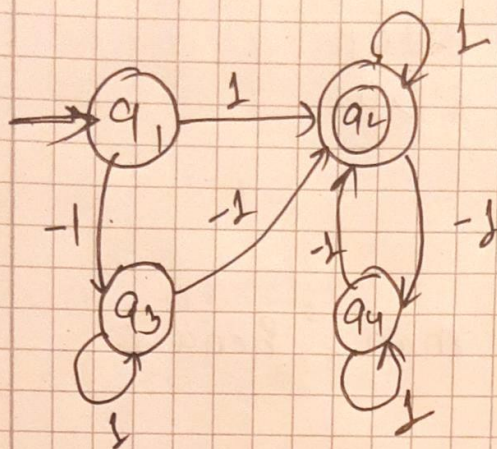


Exercise 1

$$L_1 = \{w \in \{1, -1\}^* : w = w_1 \dots w_n \text{ and } \prod_{i=1}^n w_i = 1\}.$$



The following NFA STD accepts L_1 .
Hence it is a regular language

$$L_2 = \{w \in \{1, -1\}^* : w = w_1 \dots w_n \text{ and } \sum_{i=1}^n w_i = 0\}$$

Assume that L_2 is regular.

By pumping lemma, there exists pumping length p of string s such that,

$$s = xy^iz, \forall i \geq 0$$

$$|xy| \leq p$$

$$|y| > 0$$

Let $s = 1^p (-1)^p$ (This string is contained by L_2)

$$|s| \leq 2p$$

By pumping lemma

$|xy| \leq p$, hence xy contains only 1.

Pumping y , eg. with $i=2$ we get $xgyz$.

In $xgyz$, since $|y| > 0$ and only contains 1, it has more 1's than -1's. $\therefore \sum_{i=1}^n w_i \neq 0$. \square

Ex 2

(a) $L_3 = \{w \in \{0,1\}^* : N(w,0) = 2N(w,1)\}$

Assume that L_3 is regular. By pumping lemma,
 \exists pumping length p , such that, $s = xy^iz \forall i \geq 0$
 $|xy| \leq p$
 $|y| > 0$.

Assume,

$$s = 0^{2p} 1^p$$

This has twice the number of 0^s than 1^s , hence is contained in L_3

$|xy| \leq p$, xy contains only zeros.

Pumping ~~the~~ y by $i=2$, we get $xyyz$.

~~This contains 0^s .~~

The number of ~~zero~~ 0^s in $xyyz$ is greater than $2p$.

Hence, $N(xyyz, 0) \neq 2N(xyyz, 1) \nexists \square$

By contradiction, we prove it is not regular.

(b) $L_4 = \{w \in \{0,1\}^* : N(w,0) = 0, N(w,1) = p, p \text{ is prime}\}$

Assume L_4 is regular. By pumping lemma, \exists pumping length p of s , such that $s = xy^iz \forall i \geq 0$

$$|xy| \leq p$$
$$|y| > 0$$

$$s = 1^p$$

where p is an arbitrary prime number.

$$s = xyz$$

$$|xyz| = p$$

$$\begin{aligned} \text{let } |x| &= a \\ |y| &= b \\ |z| &= c \end{aligned}$$

$$a+b+c = p.$$

Pumping y by $a+c$, i.e. $i = a+c$

$$\begin{aligned} & a + (a+c)b + c \\ &= a + ab + cb + c \\ &= a(1+b) + c(b+1) \\ &= (a+c)(1+b) \end{aligned}$$

Since we have factored the length of the pumped string into $(a+c)(1+b)$, it is not a prime. \square
Hence, we have proved that L_u is not regular by contradiction.

Ex 3 $G_1 = (\{A, B\}, \{0, 1, 2\}, R, A)$

$$\begin{aligned} A &\rightarrow 00A \mid B \\ B &\rightarrow \epsilon \mid 22B \end{aligned}$$

(a) $L(G_1) = \{ 0^{2n} 2^{2m} 1^n \mid m \geq 0, n \geq 0 \}$.

(b)
$$\begin{aligned} X &\rightarrow XZ \mid Ym \\ Y &\rightarrow YY \mid m \mid n \\ Z &\rightarrow n \mid mm \end{aligned}$$

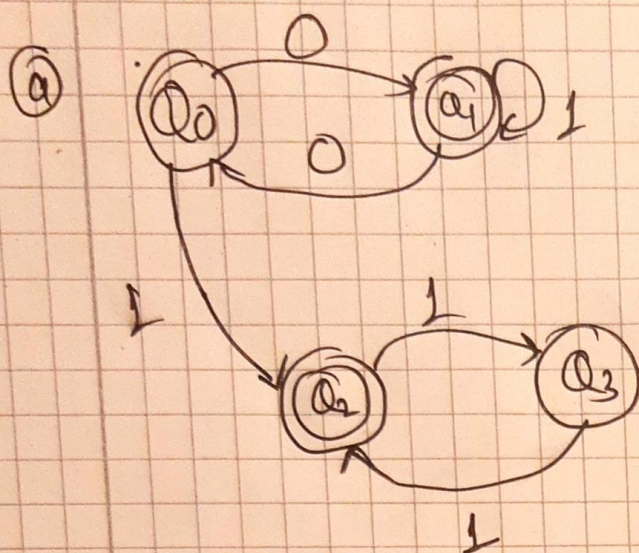
$$L(G_2) = \{ w \text{ over } \{0, 1\} \mid$$

the first part of w has at least two symbols ending with m , concatenated with second part that has a string with zero m s or even number of m 's.

$$[mn]^+ m n^* (mm)^*$$

Ex 2

Ex 4



$$G = (Q, \Sigma, R, Q_0)$$

$$Q = (\{Q_0, Q_1, Q_2, Q_3\}, \{0, 1\}, R, Q_0)$$

$$R :- Q_0 \rightarrow 0Q_1 \mid 1Q_2$$

$$Q_1 \rightarrow 1Q_2 \mid 0Q_0 \mid \epsilon$$

$$Q_2 \rightarrow 1Q_3 \mid \epsilon$$

$$Q_3 \rightarrow 1Q_2$$

6) Convert $G_3 = (\{X, Y, Z\}, \{m, n\}, R, X)$

$$X \rightarrow XZ | Ym$$

$$Y \rightarrow YY | nn$$

$$Z \rightarrow \epsilon | mm$$

into Chomsky normal form

Step 1.

Add a starting S'

$$S' \rightarrow X$$

$$X \rightarrow XZ | Ym$$

$$Y \rightarrow YY | nn$$

$$Z \rightarrow \epsilon | mm$$

Step 2

Remove null productions.

Null production: $Z \rightarrow \epsilon$

$$X \rightarrow XZ | Ym | X$$

$$Y \rightarrow YY | nn$$

$$Z \rightarrow mm$$

$$S' \rightarrow X$$

Step 3: Remove unit production.

$$X \rightarrow X, S' \rightarrow X$$

Removing $X \rightarrow X$: $X \rightarrow XZ | Ym, Y \rightarrow YY | nn, Z \rightarrow mm, S' \rightarrow X$

Removing $S' \rightarrow X$: $X \rightarrow XZ | Ym, Y \rightarrow YY | nn, Z \rightarrow mm, S' \rightarrow XZ | Ym$

Step 4: fix $S' \rightarrow Ym, X \rightarrow Ym$.

$$R^c: X \rightarrow XZ | YL$$

$$Y \rightarrow YY | nn$$

$$Z \rightarrow mm$$

$$S' \rightarrow XZ | YL$$

$$L \rightarrow m$$

Chomsky form: $G_3 = (\{X, Y, Z, L\}, \{m, n\}, R^c, S')$