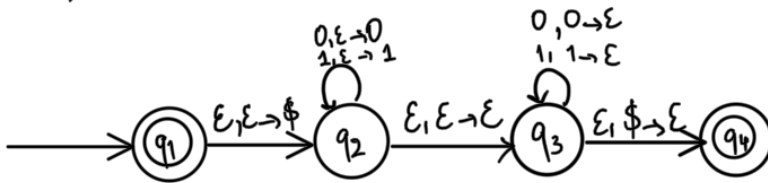


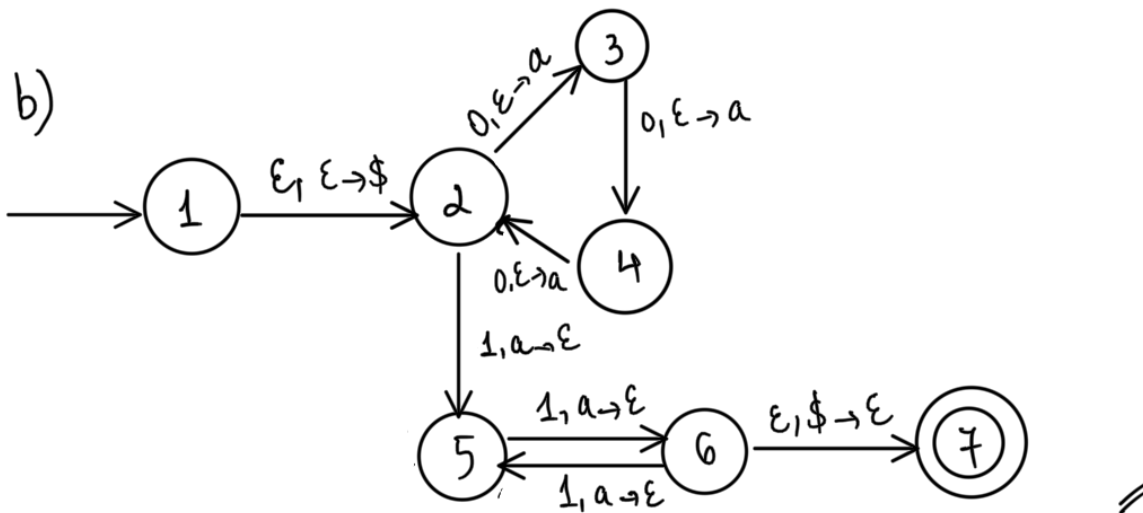
ACC HW 5

Problem 1:

a)



b)



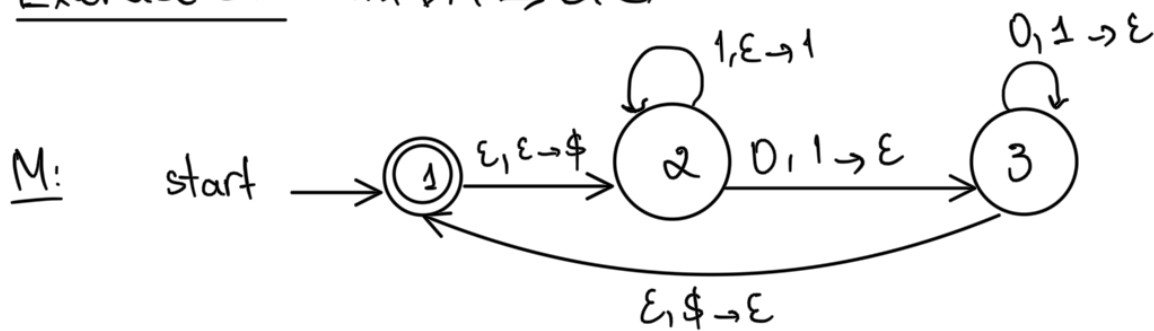
Exercise 2:

$$\begin{aligned}
 L(M_1) &= \{ w = (a)^m (b)^n (c)^{m+n} \mid m, n \in \mathbb{N} \} \\
 &= \{ w : w \text{ has a consecutive sequence of } a, \\
 &\quad \text{then a consecutive sequence of } b, \text{ finally} \\
 &\quad \text{a consecutive sequence of } c : \text{ length of} \\
 &\quad \text{sequence of } c = \text{sum of length of sequences} \\
 &\quad \text{of } a \text{ and } b \text{ (including empty string)} \}
 \end{aligned}$$

$$\begin{aligned}
 L(M_2) &= \{ w : \text{empty string or each prefix of } w \text{ has} \\
 &\quad \text{number of 0s} \geq \text{number of 1s and having the same} \\
 &\quad \text{number of 0 and 1} \}
 \end{aligned}$$



Exercise 3: mPDA \rightarrow CFG



• $\odot = \{A_{11}, A_{12}, A_{13}, A_{21}, A_{22}, A_{23}, A_{31}, A_{32}, A_{33}\}$

• $\odot = \Sigma = \{0, 1\}$

• $\odot = A_{11}$

• \odot (i) $\begin{cases} (r, t) \in \delta(p, a, \epsilon) \\ (q, \epsilon) \in \delta(s, b, t) \end{cases} \parallel \text{From STD, we have:}$

+ For $p=1 \rightarrow a=\epsilon \Rightarrow r=2, t=\$ \Rightarrow s=3, b=\epsilon, q=1$
 $\Rightarrow \text{Add } \boxed{A_{11} \rightarrow A_{23}}$

+ For $p=2 \rightarrow a=1, r=2, t=1 \Rightarrow \begin{cases} s=2, b=0, q=3 \\ s=3, b=0, q=3 \end{cases}$

$\Rightarrow \text{Add } \boxed{A_{23} \rightarrow 1A_{22}0 \mid 1A_{23}0}$

+ For $p=3 \Rightarrow$ No satisfying edge for (i)

All together R:

+ $A_{11} \rightarrow \epsilon \mid A_{11}A_{11} \mid A_{12}A_{21} \mid A_{13}A_{31} \mid A_{23}$

+ $A_{22} \rightarrow \epsilon \mid A_{21}A_{12} \mid A_{22}A_{22} \mid A_{23}A_{32}$

+ $A_{33} \rightarrow \epsilon \mid A_{31}A_{13} \mid A_{32}A_{23} \mid A_{33}A_{33}$

+ $A_{12} \rightarrow A_{11}A_{12} \mid A_{12}A_{22} \mid A_{13}A_{32}$

+ $A_{13} \rightarrow A_{11}A_{13} \mid A_{12}A_{23} \mid A_{13}A_{33}$

+ $A_{21} \rightarrow A_{21}A_{11} \mid A_{22}A_{21} \mid A_{23}A_{31}$

+ $\underline{A_{23}} \rightarrow A_{21}A_{13} \mid A_{22}A_{23} \mid A_{23}A_{33} \mid 1A_{22}0 \mid 1A_{23}0$

+ $A_{31} \rightarrow A_{31}A_{11} \mid A_{32}A_{21} \mid A_{33}A_{31}$

+ $A_{32} \rightarrow A_{31}A_{12} \mid A_{32}A_{22} \mid A_{33}A_{32}$

b) 1100 111000

$$\begin{aligned} A_{11} &\rightarrow A_{11}A_{11} \rightarrow A_{23}A_{23} \rightarrow 1A_{23}0 \ 1A_{23}0 \\ &\rightarrow 11A_{22}00 \ 11A_{23}00 \rightarrow 11\varepsilon 00 \ 111A_{22}000 \\ &\rightarrow 1100 \ 111\varepsilon 000 \rightarrow 1100 \ 111000 \end{aligned}$$

□

Exercise 4:

$$\begin{aligned} \Phi_{\Sigma}: \Sigma^* &\rightarrow \mathbb{N}^n \\ \omega &\rightarrow (|\omega|_{\sigma_1}, \dots, |\omega|_{\sigma_n}) \end{aligned}$$

Lemma 1: Any context free language over $\Sigma = \{\sigma\}$ is regular

Proof: Let $\Sigma = \{\sigma\}$: fixed alphabet & L : any context free language over Σ .

- a) Pumping lemma $\Rightarrow \exists q: |s| \geq q$, then
 $\omega = uvxyz$ satisfying the conditions ($u, v, x, y, z \in \{\sigma\}^*$)
- (i) $\forall i \geq 0, uv^ixy^iz \in L$
 - (ii) $|vy| > 0 \Rightarrow$ at least one of them is not empty
 - (iii) $|vxy| \leq q$

Choose $p = q, r = |uxz|, s = |vy| \Rightarrow r \geq 0, s > 0$ (by ii)

$|\omega| = r + s$ and since $u, v, x, y, z \in \{\sigma\}^*$

$$\Rightarrow \omega = \sigma^{r+s} = \sigma^r \cdot \sigma^s \quad (\text{satisfy (a)}) \quad \checkmark$$

Choose $t = |x| \Rightarrow t \leq r$ ($|uxz| = r$).

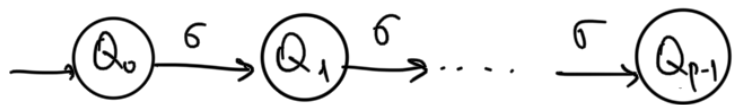
$$\text{By (iii)} \Rightarrow 0 < t + s \leq q = p \quad (\text{satisfy (b)}) \quad \checkmark$$

$$\text{By (i)} \quad uv^ixy^iz = \sigma^r \cdot \sigma^{si} \in L \quad (\text{satisfy (c)}) \quad \checkmark$$

$$b) L_p = \{\omega \in L: |\omega| \geq p\} \subseteq L.$$

$L'_p = L \setminus L_p$: regular & express L : terms of L_p & L'_p

Consider the NFA N



$\forall w \in L'_p \Rightarrow$ we will make each state $Q_{|w|}$ accepted

Hence $L'_p = L(N) \Rightarrow L'_p$: regular.

Moreover $L'_p = L \setminus L_p \Rightarrow L = L_p \cup L'_p$

