Lessons Learned: A Multi-Agent Framework for Code LLMs to Learn and Improve

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Abstract

Recent studies show that LLMs possess different skills and specialize in different tasks. In fact, we observe that their varied performance occur in several levels of granularity. For example, in the code optimization task, code LLMs excel at different optimization categories and no one dominates others. This observation prompts the question of how one leverages multiple LLM agents to solve a coding problem without knowing their complementary strengths a priori. We argue that a team of agents can learn from each other's successes and failures so as to improve their own performance. Thus, a lesson is the knowledge produced by an agent and passed on to other agents in the collective solution process. We propose a lesson-based collaboration framework, design the lesson solicitation–banking–selection mechanism, and demonstrate that a team of small LLMs with lessons learned can outperform a much larger LLM and other multi-LLM collaboration methods.

1 Introduction

Code optimization refers to rewriting the code such that it performs the same task more efficiently [29]. The efficiency is predominantly measured by the runtime, but it could also refer to storage, energy, or other resource consumptions. Code optimization is a less explored coding task in the context of AI compared with the popularly studied code generation task [55], but it is a natural, and sometimes key, step in the software development cycle. Code performance can be improved in many ways, such as using a lower-complexity algorithm, caching and memorization, data alignment, vectorization, and parallelization [44, 4]. Many of these approaches require low-level system knowledge that code LLMs are not explicitly trained with.

We are interested in exploring the use of multiple LLM-powered agents for code optimization. We use multiple LLMs because no one LLM performs the best on all problems, even if they are trained with extensive code data. The common practice of benchmarking uses aggregated performance to rank models [5], but the best-ranked model may not be the best performer for every problem. We take the ParEval benchmark [35], which consists of programming problems in scientific computing, for example. When three open-source, small models and two GPT models are evaluated on this benchmark (see Appendix A for details), GPT-4o [39] is the overall winner. However, on the "geometry" category of problems, Qwen7B [23] outperforms GPT-4o by $2.5\times$ in terms of speedup; and on the "histogram" category, Deepseek7B [18] and Qwen14B [23] outperform GPT-4o by $1.6\times$ (Figure 5). Additionally, among the three open-source models, Qwen7B excels at "search," while

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Original code

```c
for (int i = 0; i < n; ++i)
  for (int j = 0; j < n; ++j)
    for (int k = 0; k < n; ++k)
      C[i][j] += A[i][k] * B[k][j];
```

Naive implementation of matrix multiplication $C = AB$.

Improved code, round 1

```c
for (int i = 0; i < n; ++i)
  for (int k = 0; k < n; ++k)
    for (int j = 0; j < n; ++j)
      C[i][j] += A[i][k] * B[k][j];
```

Lesson: Reordering loops improves cache locality and increases performance. The order of (i,k,j) out of 6 different permutations often performs the best, because of how caches work.

Improved code, round 2

```c
#pragma omp parallel for
for (int i = 0; i < n; ++i)
  for (int k = 0; k < n; ++k)
    for (int j = 0; j < n; ++j)
      C[i][j] += A[i][k] * B[k][j];
```

Lesson: Using OpenMP to parallelize the for-loop further improves performance. Parallelizing only the outermost loop performs the best, due to sufficient parallelism and less parallel scheduling overhead.
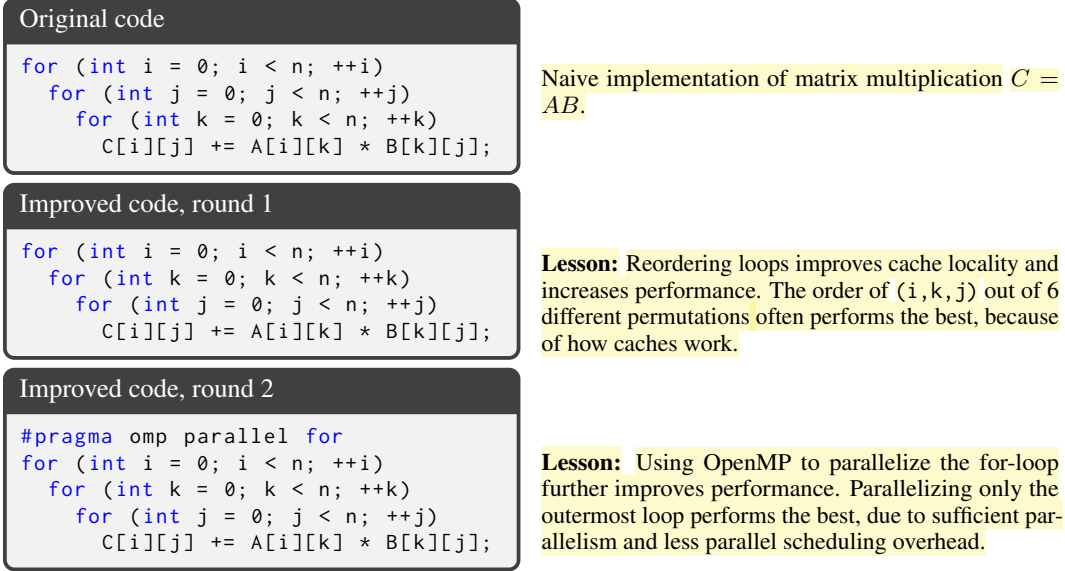
Figure 1: Successively improving matrix-matrix multiplication in C with lessons.

Deepseek7B excels at "scan" and "dense_la" and Qwen14B excels at "sparse_la," "reduce," "fft," and "sort" (Table 4). Such varied performance suggests the opportunity of exploiting the complementary strengths of different LLMs to deliver the best solution.

How does one make use of multiple agents to solve a coding problem? In this work, we advocate the concept of *lessons* inspired by classroom experience. A student's problem-solving skills are not only taught by a teacher at class but also developed through peer learning. For example, after receiving the graded homework, Student A, who cannot complete the solution to a problem may consult with Student B, who earns a perfect score, on tips of the correct steps toward the solution. Meanwhile, Student A can also benefit from learning from Student C, who comes up with a wrong solution, to avoid making similar mistakes. Each student learns from the success and failure lessons of others to improve their own problem-solving skills. LLM agents behave similarly. Pre-training is analogous to classroom teaching, while a pre-trained LLM can improve its skills through prompted with lessons.

A lesson means any information that helps an LLM better solve the problem at hand. For code optimization, such lessons may be optimization strategies applicable to the current code, common pitfalls that programmers are trapped in, or performance feedback from profilers. Note that a code's performance can be improved by a combined use of multiple strategies, step by step. Figure 1 shows the initial steps of a classic example of engineering the performance of matrix-matrix multiplications, which can be improved by 3000x in speedup under extensive optimization [44]. The applied strategies shown are loop reordering and parallelization of for-loops. In a multi-LLM setting, such strategies, or lessons, can be iteratively summarized by one or a few LLMs and learned by others, so that they collectively improve the code performance.

In this work, we propose the framework **LessonL** (pronouced as "lesson-nell") for multiple LLM agents to collaboratively solve problems. Central to the framework is the lesson mechanism for agents to improve their collective intelligence through learning from each other. The main technical innovation is a solicitation–banking–selection framework that generates, deposits, and filters lessons incurred during the collective problem-solving process. Although code optimization is the focal application of LessonL, we demonstrate its use for other tasks (code generation) as well.

LessonL is a novel multi-agent framework that resembles how humans learn to solve problems. Compared with other collaboration frameworks, such as where agents play different roles in the solution process [43, 11, 21, 25, 41], or where agents independently propose solutions that are subsequently communicated and aggregated [46, 13, 26, 33, 54] (see Section 2 for literature review), our framework has a few advantages. First, agents do not need to be distinguished by pre-specified roles, as their complementary strengths for a particular problem may be unknown a priori. Second, communication and prompt contents are economic since lessons are more concise than codes. Third,

more importantly, lessons are interpretable and reusable, allowing the explication of coding knowledge and the creation of educational materials.

Our work contributes the following:

1. a finding that LLMs have complementary strengths even on a fine level of tasks (Appendix A);

2. a novel lesson-based framework for multiple agents to collectively solve problems (Section 3);

3. state-of-the-art performance on code optimization and code generation benchmarks (Section 4.2);

4. empirical evidence that a team of small LLMs can significantly outperform a much larger LLM under similar resource consumptions (Section 4.4);

5. representative code examples and lessons (Appendices J and K).

2 Related Work

Multi-agent collaboration. Recent advances in prompting techniques have enhanced LLMs' reasoning and planning capabilities for complex tasks like mathematics and coding. CoT [48], Self-Consistency [47], ReAct [52], and ToT [51] demonstrate improved problem-solving skills with reasoning while Reflexion [43] leverages self-generated feedback stored in episodic memory to enhance decision-making. As such, LLMs can be used as autonomous agents and recent research efforts propose ways for multiple agents to collaboratively solve problems. Popular multi-agent frameworks either place agents in different roles (such as planner, coder, debugger, reviewer, and tester in a software project) or have agents independently propose solutions, which are collectively refined and consolidated. For role-based methods, see AgentVerse [11], MetaGPT [21], MapCoder [25], ChatDev [41], Self-collaboration [12], SoA [24], and AgentCoder [22]. For individual solution proposals, see MoA [46], LLM-Debate [13, 14], LLM-Blender [26], DyLAN [33], and AgentPrune [54]. Our work LessonL belongs to the latter category. A unique contribution of LessonL is the lesson mechanism for agents to learn from each other and the explication of coding knowledge as a result of the collaborative learning.

Code optimization. It is an important, but under-explored, use case of LLM for code. Prior works focuses on specialized models such as HPC-Coder [36, 37] and HPC-Coder-V2 [8] for high-performance computing, all require curating and/or generating code data and fine-tuning. Besides the usual fine-tuning approaches, PIE [45] proposes additional adaptation techniques, including retrieval-based prompting, performance-conditioning, and self-play. Among the few agentic approaches, SBLLM [16] retrieves optimization examples from an external dataset and Self-Refine [34] iteratively refines the code based on self-generated feedback. In contrast, our multi-agent framework allows the use of more than one agent and it does not rely on an external code dataset.

For extended discussions and more related work, see Appendix B.

3 The LessonL Framework

Our multi-agent collaboration framework is centered around *lessons*, which can be any knowledge or information that helps an agent better solve the problem at hand. When a team of agents participates in the collective problem-solving process, such knowledge is solicited through inspecting each agent's solution and is deposited into a bank for other agents to access. Hence, the solution process becomes iterative, looping between using the existing lessons to update the solutions and using the updated solutions to generate new lessons. For code optimization, such an iterative process can progressively improve the code performance. This process is pictorially illustrated in Figure 2 and sketched in the following pseudocode:

1: Each agent generates an initial solution and the lesson for it. Deposit lessons to the bank.
2: **for** round $t = 1, \dots, T$ **do**
3:     Select k lessons from the bank.
4:     Based on the selected lessons, each agent generates an updated solution.
5:     Each agent generates a new lesson for the updated solution. Deposit lessons to the bank.
6:     Adjust the effectiveness of the k selected lessons.
7: **end for**
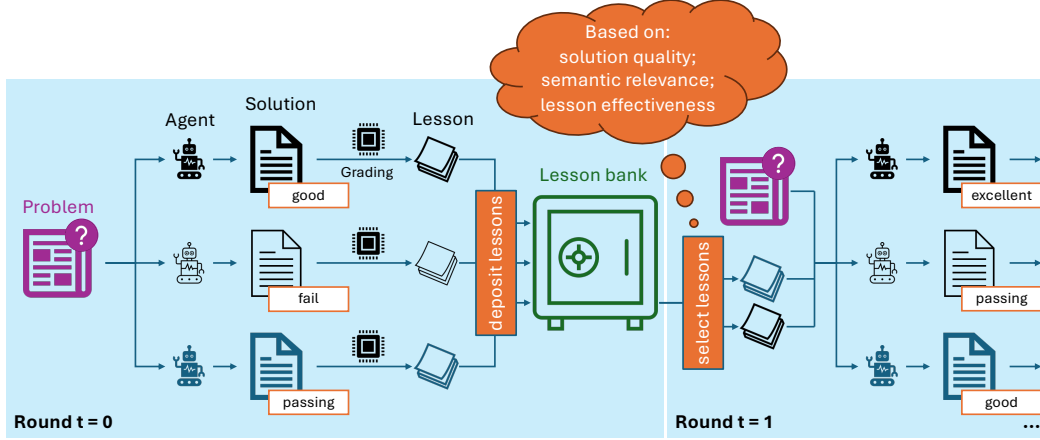8: Return the best solution.

Figure 2: The LessonL framework (which may repeat multiple rounds).

In what follows, we elaborate on a few key components of this framework. The full algorithm is given in Appendix C. We also discuss its extension to other coding tasks, such as code generation.

3.1 Lesson Solicitation

Every solution comes with a lesson, either positive or negative. Such lessons explain why the solution is correct or what make it wrong. For code optimization, multiple tools can be used to grade the output code, such as the compiler, the clock, test cases, and profilers. Consider four resulting scenarios:

(a) speed up,   (b) slow down,   (c) functional incorrectness,   (d) syntax error.

We solicit lessons by referring the original code and the modified code, A and B, respectively, and prompting the agent with the resulting scenario and asking for explanations. For (a) and (b), we supplement the scenario description with the measured speedup and for (d), we include the error message reported by the compiler. Such information helps the agent reason with the code pair and articulate precise lessons. For (c), we do not include the test cases because LLMs tend to reason specifically about the test cases, resulting in insufficiently general lessons.

The detailed prompt templates are given in Appendix D.

3.2 Lesson Banking and Selection

In each round, n lessons are generated and deposited to the bank, one from each agent. Then, after several rounds, there accumulate many lessons. It is unwise to feed all of them to the prompt when asking the agent to improve the code, because they may exceed the prompt capacity and also the token consumption can become too costly. Hence, the framework is run with at most k lessons in each round. In practice, it suffices to set k to be, or slightly larger than, n.

A set of lesson selection criteria is in place. First, naturally we want to use lessons about high speedups, because they point to the right directions of optimization. However, positive lessons alone cannot address certain limitations of LLMs, such as the lack of guarantee on code correctness. Hence, negative lessons are still valuable, because they can help the agents avoid similar mistakes. It is, however, challenging to decide which negative lessons are more important than others. Finally, we also consider relevance, in case an agent hallucinates irrelevant lessons. For this, we treat the original code as the query and retrieve semantically relevant lessons from the bank, in a manner similar to retrieval augmented generation [30].

Algorithmically, the selection mechanism goes as follows. If the bank has no more than k lessons, pick them all. Otherwise, sort the lessons according to speedup (more on this in the next subsection) and pick the top $\lceil k/2 \rceil$. Then, sort the remaining lessons according to their cosine similarity with the original code and pick the top $\lfloor k/2 \rfloor$. Any powerful embedding models for code can be used to compute the cosine similarity; for example, CodeBERT [15].

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3.3 Effectiveness Adjustment

The effectiveness of a lesson z is naturally measured by the speedup s — the higher the speedup, the more useful the lesson. Note that s was calculated when z was created together with code y. In other words, s is the speedup of code y over the original code. One's view on the effectiveness of z may change when this lesson is selected to apply later. For example, if the lesson later yields code y' with a worse speedup $s' < s$, should we keep selecting it only because the original s is great?

To more effectively select high-speedup lessons, we need some adjustment to the speedup dynamically. Rather than sorting them according to s, we introduce an adjustment factor f and sort the lessons according to $s \times f$ instead, where f accounts for the performance of the lesson z when actually applied. Specifically, when z is applied in some round t, it incurs speedups $s_j^{(t)}$ for each agent j. We initialize a correction variable c with zero and add $1 + \epsilon$ to it whenever $s_j^{(t)} > s$, or add $1 - \epsilon$ when $s_j^{(t)} < s$, for some $\epsilon \in (0, 1)$. After looping over all n agents, the adjustment factor f is defined as c/n. A value greater than 1 means more output codes enjoy a speedup $> s$ by applying the lesson.

3.4 Extension to Other Coding Tasks

The LessonL framework is general. It can be extended to other coding tasks provided that the key components discussed above are properly adapted. For example, in Python code generation, lessons are needed to be solicited for only one scenario: functional incorrectness, because if the output code passes all test cases, the iteration immediately terminates. We give the prompt templates in Appendix E for completeness. Additionally, we will still perform lesson banking, but k lessons are selected based on the number of test cases passed and the semantic relevance instead of speedup.

4 Experiments

We perform a comprehensive set of experiments to evaluate LessonL. The main finding is that LessonL enables an effective collaboration among LLMs to perform code optimization and other coding tasks. Using our framework, a team of small LLMs that collaborate through the sharing of learned lessons can outperform larger LLMs and multi-LLM collaboration methods.

4.1 Setup

Benchmarks. We use six coding benchmarks to perform the evaluation. (1) ParEval [35] includes 60 coding tasks (per programming mode) related to scientific and parallel computing. We experiment with the serial and OpenMP modes, as they are less demanding on the computational resources required for evaluation. ParEval was originally designed for code generation (i.e., write the code given verbal description). We adapt it to code optimization (i.e., write a faster version of a given source code); see the adaptation details in Appendix F. (2) PolyBench [40] contains 30 numerical tasks with static control flows from domains like linear algebra, image processing, physics, and statistics, adapted for code optimization by us. (3) HumanEval [10] consists of 164 programming problems, assessing language comprehension, algorithms, and basic mathematics. (4) HumanEval+ [32] extends HumanEval with 80x more test cases. (5) MBPP [3] consists of around 1,000 crowd-sourced entry-level Python problems covering fundamentals and standard library use. (6) MBPP+[32] extends MBPP with 35x more test cases. The last four benchmarks evaluate code generation.

LLM agents. We use five models in total. Three are open-source, small models: **Deepseek7B** (deepseek-coder-7b-instruct-v1.5) [18], **Qwen7B** (Qwen2.5-Coder-7B-Instruct) [23], and **Qwen14B** (Qwen2.5-Coder-14B-Instruct) [23]; and two are GPT models: **GPT-4o mini** (gpt-4o-mini) [38] and **GPT-4o** (gpt-4o) [39]. The open-source models are used to evaluate multi-agent collaborations; GPT-4o mini is closed-source but its size is comparable to the open-source models; while GPT-4o is much larger and it sets a strong baseline for single-agent performance.

Baselines. We compare LessonL with three categories of models/methods. (1) **Single-agent standard prompting**. All the above LLMs are experimented with. Among the open-source models, preliminary findings suggest that Qwen14B slightly outperforms the other two (see Table 4 in Appendix A); hence, we omit the results of these two in subsequent tables to save space. (2) **Single-agent reasoning or reflection**. We use Qwen14B as the agent and experiment with CoT [48] and Reflexion [43]. CoT

Table 1: Comparison of model/method performance for two code optimization benchmarks. "Correct" means correctness; "> 2x" means proportion of problems achieving speedup > 2x; "Speedup" is the geometric mean speedup across all problems in the benchmark. The results are reported as the mean and standard deviation over three runs.

ParEval	Serial mode			OpenMP mode		
	Correct	> 2x	Speedup	Correct	> 2x	Speedup
Qwen14B	0.67 ± 0.03	0.14 ± 0.01	1.60 ± 0.03	0.67 ± 0.01	0.48 ± 0.00	2.28 ± 0.02
GPT-4o mini	0.77 ± 0.03	0.14 ± 0.01	1.57 ± 0.12	0.70 ± 0.03	0.55 ± 0.04	2.72 ± 0.09
GPT-4o	0.80 ± 0.00	0.16 ± 0.03	1.72 ± 0.11	0.73 ± 0.05	0.58 ± 0.05	2.93 ± 0.30
CoT	0.61 ± 0.03	0.12 ± 0.03	1.57 ± 0.12	0.65 ± 0.04	0.51 ± 0.03	2.31 ± 0.03
Reflexion	0.67 ± 0.03	$\underline{0.18 \pm 0.00}$	$\underline{1.88 \pm 0.20}$	0.64 ± 0.01	0.46 ± 0.02	2.27 ± 0.06
MapCoder	$\underline{0.88 \pm 0.02}$	0.15 ± 0.02	1.85 ± 0.08	$\underline{0.83 \pm 0.05}$	0.58 ± 0.02	$\underline{3.43 \pm 0.17}$
MoA	0.73 ± 0.04	0.16 ± 0.02	1.76 ± 0.08	0.74 ± 0.04	$\underline{0.59 \pm 0.02}$	2.77 ± 0.11
LessonL	$\mathbf{0.91 \pm 0.02}$	$\mathbf{0.21 \pm 0.01}$	$\mathbf{2.16 \pm 0.11}$	$\mathbf{0.86 \pm 0.01}$	$\mathbf{0.62 \pm 0.02}$	$\mathbf{3.46 \pm 0.03}$

PolyBench	Serial mode			OpenMP mode		
	Correct	> 2x	Speedup	Correct	> 2x	Speedup
Qwen14B	0.46 ± 0.12	0.02 ± 0.02	1.03 ± 0.04	0.49 ± 0.02	0.46 ± 0.02	2.21 ± 0.06
GPT-4o mini	0.41 ± 0.08	0.06 ± 0.02	1.10 ± 0.03	0.55 ± 0.07	0.48 ± 0.05	2.42 ± 0.31
GPT-4o	0.51 ± 0.02	0.06 ± 0.02	1.12 ± 0.03	0.60 ± 0.12	$\underline{0.55 \pm 0.04}$	$\underline{2.73 \pm 0.21}$
CoT	0.51 ± 0.03	0.02 ± 0.04	1.06 ± 0.04	0.39 ± 0.03	0.37 ± 0.04	1.92 ± 0.16
Reflexion	0.41 ± 0.02	0.08 ± 0.07	1.12 ± 0.07	0.56 ± 0.23	0.52 ± 0.04	2.59 ± 0.22
MapCoder	$\underline{0.64 \pm 0.02}$	$\underline{0.12 \pm 0.04}$	$\underline{1.17 \pm 0.03}$	0.55 ± 0.04	0.44 ± 0.05	2.23 ± 0.18
MoA	0.32 ± 0.07	0.08 ± 0.02	1.12 ± 0.04	$\underline{0.61 \pm 0.07}$	0.54 ± 0.05	2.70 ± 0.34
LessonL	$\mathbf{0.77 \pm 0.04}$	$\mathbf{0.13 \pm 0.04}$	$\mathbf{1.32 \pm 0.09}$	$\mathbf{0.71 \pm 0.03}$	$\mathbf{0.62 \pm 0.02}$	$\mathbf{3.40 \pm 0.12}$

applies a chain-style thought process to reason about the steps, while Reflexion iteratively reflects on the task feedback to improve the solution. (3) **Multi-agent collaboration**. We experiment with MapCoder [25] and MoA [46]. MapCoder uses agents for example retrieval, solution planning, coding, and debugging. For our purpose, all agents are Qwen14B. In contrast, each agent in MoA independently codes the solution and refines the aggregated solution. We use the open-source models as agents and use GPT-4o as the final aggregator. Similarly, in our framework LessonL, the agents are open-source models.

For details on baseline implementation, hyperparameters, hardware, and timing, see Appendix G.

4.2 Benchmark Results

We evaluated the performance of LessonL on code optimization and code generation tasks. For code optimization, we studied the ability of LessonL to optimize serial and parallel code drawn from the ParEval and PolyBench benchmarks. For code generation, we investigated LessonL's performance on HumanEval, HumanEval+, MBPP, and MBPP+.

Code optimization task. In the code optimization task, models are given a correct program and are tasked with generating a faster version of that code while maintaining correctness. The **speedup** achieved by a model is measured as the ratio of the runtime of the original and new code. We evaluate the performance of a model by measuring the geometric mean speedup achieved over the set of codes in the benchmark. The geometric mean is preferred over the arithmetic mean because it is more resilient to large outliers that can cause a single code to unduly influence the average. For example, when using the arithmetic mean an algorithmic optimization that improves the asymptotic runtime of a code from $\Theta(n^2)$ to $\Theta(n \log n)$ could result in a $1000\times$ speedup for sufficiently large input, and result in the arithmetic average becoming a poor measure a model's ability to optimize diverse sets of codes. In the case where a new code is incorrect or slower than the original, we consider the speedup to be 1. Similar to [45], which uses the same convention, our rationale is any new code may be discarded in favor of keeping the original code.

Code optimization results. Table 1 presents our experimental results on the code optimization benchmarks. We compared the performance of LessonL with three single-agent models (Qwen14B,

Table 2: Comparison of model/method performance (pass@1) for four code generation benchmarks. GPT-4o and GPT-4o mini results are taken from the EvalPlus leaderboard [32].

	HumanEval	HumanEval+	MBPP	MBPP+
Qwen14B	0.866	0.817	0.831	0.741
GPT-4o mini	0.884	0.835	0.854	0.722
GPT-4o	0.884	<u>0.872</u>	0.876	0.722
CoT	0.884	0.848	0.854	0.741
Reflexion	0.902	0.841	0.889	<u>0.765</u>
MapCoder	0.890	0.823	0.862	0.679
MoA	<u>0.909</u>	<u>0.872</u>	<u>0.897</u>	**0.770**
LessonL	**0.915**	**0.878**	**0.899**	<u>0.765</u>

Table 3: Ablation study. Benchmark: ParEval (serial and OpenMP modes). The results are reported as the mean and standard deviation over three runs.

		Serial mode			OpenMP mode		
		Correct	> 2x	Speedup	Correct	> 2x	Speedup
(0)	LessonL	0.91 ± 0.02	$\mathbf{0.21 \pm 0.01}$	$\mathbf{2.16 \pm 0.11}$	$\mathbf{0.86 \pm 0.01}$	$\mathbf{0.62 \pm 0.02}$	3.46 ± 0.03
(1)	k high-speedup only	$\mathbf{0.92 \pm 0.02}$	$\mathbf{0.21 \pm 0.01}$	2.08 ± 0.04	0.83 ± 0.02	0.58 ± 0.00	3.20 ± 0.04
(2)	k high-relevance only	0.91 ± 0.02	0.18 ± 0.02	1.96 ± 0.08	0.84 ± 0.02	0.61 ± 0.01	3.40 ± 0.16
(3)	No speedup adjustment	$\mathbf{0.92 \pm 0.03}$	$\mathbf{0.21 \pm 0.03}$	2.05 ± 0.05	$\mathbf{0.86 \pm 0.02}$	0.61 ± 0.02	3.28 ± 0.10
(4)	Random k lessons	0.91 ± 0.02	0.19 ± 0.02	2.03 ± 0.05	0.82 ± 0.02	0.61 ± 0.11	$\mathbf{3.47 \pm 0.28}$
(5)	No lessons	0.89 ± 0.01	0.20 ± 0.00	2.05 ± 0.01	0.82 ± 0.01	0.56 ± 0.01	3.01 ± 0.16

GPT-4o-mini, and GPT-4o) and four prompting / agentic methods (CoT, Reflexion, MapCoder, and MoA). We report each method's average speedup, proportion of correct code, and the proportion of codes that achieve a speedup greater than $2\times$.

A few observations follow. First, LessonL achieves the best results across metrics and benchmarks This highlights the attractiveness of its novel collaborative lesson mechanism, which enables agents to learn from each other and iteratively improve solutions. Second, multi-agent collaborations generally surpass single-agent approaches. Following LessonL, MapCoder and MoA often yield the next best results with role specialization or solution aggregation. Third, achieving large speedups is more challenging for serial code than for OpenMP code. In serial mode, most methods yield an speedup below $2\times$, with less than 20% of solutions surpassing $2\times$ speedup. Conversely, in OpenMP mode, most methods achieve more than $2\times$ speedup, with more than 40% of solutions surpassing $2\times$. This is expected, as many scientific computing and image processing tasks are inherently parallelizable. For instance, LessonL achieves an average $3.4\times$ speedup with 8 threads across both benchmarks.

Code generation. While code optimization is the main testbed of LessonL, we also experiment with a widely studied coding task—code generation. Table 2 compares the results for four commonly used benchmarks by using the correctness (pass@1) metric. We see that LessonL is the best in three out of four benchmarks, with MoA coming next. In the remaining benchmark, the ranking positions of these two frameworks are swapped. This observation echos that of code optimization, corroborating the usefulness of multi-agent collaboration for solving coding tasks and highlighting the effectiveness of our lesson framework.

4.3 Ablation Study

We conducted a study of LessonL to investigate the impact of its design components and the number of collaboration rounds.

Ablating components of the framework. We analyzed the selection mechanisms for lessons in our ablation study. The results are illustrated in Table 3 and compare five variations of the LessonL framework: (0) the full version of LessonL; (1) lessons are selected based only on speedup; (2) lessons selected based only on relevance; (3) lessons selected based on speedup and relevance, but without a speedup adjustment; (4) random selection of lessons; (5) no lessons used.

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Table 3 reports the results. We see that most ablated variants suffer a decrease of speedup. In serial mode, variants without high-speedup lesson prioritization (2,4) suffer most, while in OpenMP mode, variants (1,3) show larger performance drops. This trend also applies to the proportion of problems with >2x speedup. This highlights different optimal strategies for each mode. Serial mode benefits from high-speedup lesson selection with dynamic adjustments, while OpenMP mode favors high-relevance lessons and speedup adjustments. Interestingly, in serial mode, the correctness metric sometimes benefits from ablations, even though the relationship between lessons and correctness is indecisive. While variant (4) marginally outperforms LessonL in OpenMP speedup (+0.01), LessonL demonstrates significantly better stability (0.03 vs 0.28 standard deviation). Finally, The consistent underperformance of variant (5) across all metrics confirms that lessons are fundamental to LessonL.

Varying the number of iterations. Figure 3 plots the performance (speedup) achieved by LessonL and alternative methods when varying the number of collaboration rounds. The concept of a "round" is different across certain achitectures: for MoA a round is a "layer" in the MoA architecture, and for MapCoder it is a "debugging round". The performance of GPT-4o is also included to provide a baseline for comparison. We see that while LessonL and Reflexion consistently benefit from using more rounds, the same is not true for the MoA and MapCoder frameworks. The performance of MoA actually decreases when increasing the number of rounds, and there is no clear trend when using more than 2 rounds in MapCoder. Even when using a much smaller model, Reflexion surpasses GPT-4o and continues to improve when further increasing the number

Figure 3: Performance over rounds (or called "layers"). Benchmark: ParEval (serial mode).

of rounds. In contrast, MoA's performance drops below GPT-4o when using more than 3 rounds, even though MoA actually uses the GPT-4o model as an aggregator.

We studied the outputs of each method and provide examples in Appendix H. For a program that finds the kth smallest element of an array, we observe that MoA implements the divide-and-conquer algorithm in early rounds, achieving high speedup. In subsequent rounds, however, MoA introduces overheads and eventually resorts to a simple solution that calls std::nth_element, which results in slower code. In contrast, in the example of performing convolutions on an image, LessonL increases the speedup step by step by removing redundant zero-additions, separating boundary and interior computations, and avoiding extra data copies. These examples illustrate how the lesson mechanism can selectively inject useful information to aid improvements in coding solutions.

4.4 Cost Analysis

Code performance is not the only dimension that determines the effectiveness of a code model, considering the complex interplay of monetary costs and time costs in single-LLM and multi-LLM methods. We investigate where LessonL stands for the code optimization task by considering budget constraints with respect to money and time. For this, we follow [46] and extract the pricing information of LLMs from API providers or guesstimate the price based on similar models. We also use the number of floating point operations (FLOPS) as a proxy of latency when considering the time cost. See Appendix I for the detailed estimation. The estimation may be time sensitive (for example, price varies over time and across providers) and may include educated guesses (for example, the sizes of closed-source models), but it paints a reasonable picture of the performance landscape.

Figure 4 shows two speedup-versus-cost plots for ParEval, with the Pareto front indicated by the dashed line. Besides the afore-mentioned models/methods, we add Qwen14B(20) and CoT(20), which prompt the respective model/method 20 times and select the fastest code (an inference-time scaling approach [6]). We see a few clusters from the plots. First, LessonL is Pareto optimal because of its superior speedup compared with competitors. Second, Qwen14B, GPT-4o mini, and CoT are also Pareto optimal, or nearly optimal, because of the low costs in both money and time. Third, GPT-4o is on the Pareto front regarding time and barely misses the front regarding money (its dollar cost is slightly higher than that of LessonL). We consider all these methods to be cost-effective. For the remaining methods, inference-time scaling (CoT(20) and Qwen14B(20)) does not help and neither do Reflexion or multi-agent methods. In fact, MapCoder and MoA appear to be the least
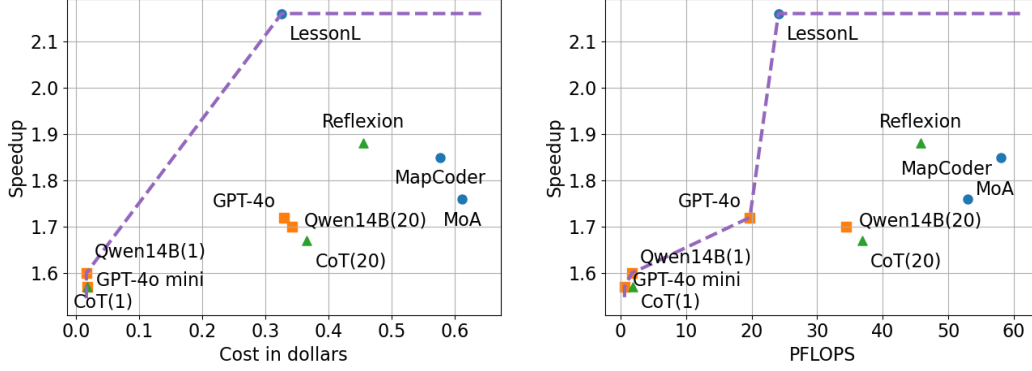
Figure 4: Performance versus costs and latency. Benchmark: ParEval (serial mode). The dashed line is the Pareto front.

cost-effective, because they not only iterate multiple times, but also use multiple agents. This signifies the challenge of multi-agent methods, which are generally the most costly, and in contrast reflects the attractiveness of LessonL. Finally, compared with the much larger model GPT-4o, LessonL with small models yields significantly better speedups by consuming similar resources.

4.5 Case Study

A few examples reveal the interesting lessons learned by the LLMs in our framework; see details in Appendices J (Geometry) and K (DFT). In the Geometry example, the original code finds the smallest distance among n points in two dimensions by using the straightforward $O(n^2)$ implementation. The final optimized code uses a divide-and-conquer algorithm that reduces the complexity to $O(n \log n)$ and achieves a 74.31x speedup. This optimization is a complete change of algorithm and it is nontrivial even for human programmers. Several of the lessons affirm the benefit of divide-and-conquer while others point out syntactic errors and discourage the agents to repeat them.

More interesting is the DFT example, where the original code implements the straightforward algorithm in $O(N^2)$ cost. It is well known that DFT can be implemented by using the FFT algorithm in $O(N \log N)$ cost and some agents try to implement it, but they fail with subtle bugs. The lessons alert these failures and eventually the agents choose to optimize the code by precomputing the repeatedly used exponentials, yielding a considerable 10.83x speedup. Moreover, one lesson stands out by stating a syntactic error caused by using Intel intrinsic instructions. The evaluation pipeline and hardware do not support particular intrinsics and the involved agent later refrains from using them in the implementation.

5 Conclusions

We have presented a lesson-based framework LessonL for multiple agents to collaboratively solve coding problems. This framework allows each agent to learn from other agents' successes and failures to improve its own solution. Empirical evaluations indicate that with lessons learned over the course, our framework outperforms other methods to achieve state-of-the-art performance, and a team of small LLMs can significantly outperform a much larger LLM under a similar resource budget. Case studies show that the agents perform many skillful optimizations, such as performing divide-and-conquer and precomputation. An avenue of future research is to improve the autonomy of agents, including their decision-making ability in lesson selection.

Limitations. Even though experiment results demonstrate competitive latency of LessonL compared with large LLMs, learning from lessons will defer the time to first token, because lesson solicitation and subsequent trials are overheads. This may affect user experience. However, the negative impact may be offset by the rich and educational information—lessons—delivered together with the solution.

Broader impacts. Code optimization is an important step of the software development cycle, especially for domains such as high performance computing, real-time animation, and transaction

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systems. This work develops a cost-effective approach to improving programmer productivity, by using LLMs to suggest optimized code. The interpretability of lessons lowers the barrier of mastering the low-level system knowledge required for optimization. Nevertheless, LLM outputs are not without flaws and practitioners should verify the lessons if they plan to apply the knowledge elsewhere.

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Contents

A   Complementary Strengths of LLM Agents

Section 1 motivates the use of multiple LLM agents to solve coding tasks, because it is unknown a priori which coding agent performs the best for each problem. We supplement the discussions with benchmarking results on the adapted ParEval benchmark [35] here (see Appendix F for the adaptation details). Figure 5 compares the overall performance of five LLMs, three open-source and two closed-source. Specifically, the models are:

- deepseek-coder-7b-instruct-v1.5, abbreviated as Deepseek7B;
- Qwen2.5-Coder-7B-Instruct, abbreviated as Qwen7B;
- Qwen2.5-Coder-14B-Instruct, abbreviated as Qwen14B;
- gpt-4o-mini, abbreviated as GPT-4o mini;
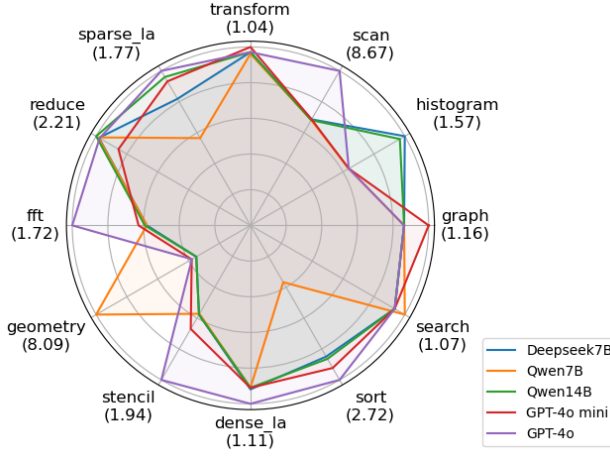- gpt-4o, abbreviated as GPT-4o.

Figure 5: Comparison of models on their coding abilities for each category of the ParEval benchmark (serial mode). The coding ability is measured by the speedup when asked to optimize a code. Geometric mean of the speedups is taken over all problems in the same category. The number in the parenthesis is the maximum speedup for the corresponding category over all models. Performance for each category is normalized by the maximum speedup.

Moreover, in Table 4, we break down their performance into each category of the benchmark. We also highlight the best one among the open-source models. One sees that each model excels in some of the categories and no one dominates others.

Table 4: Speedup achieved by a model for each category of the ParEval benchmark (serial mode). The speedup is computed as the geometric mean. Three open-source models (i.e., excluding GPT-4o mini and GPT-4o) are compared with the best **boldfaced**.

	Deepseek7B	Qwen7B	Qwen14B	GPT-4o mini	GPT-4o
graph	**1.00**	**1.00**	**1.00**	1.16	1.00
histogram	**1.57**	1.00	1.52	1.00	1.00
scan	**5.95**	5.91	5.91	5.95	8.67
transform	**1.01**	1.00	**1.01**	1.04	1.01
sparse_la	1.45	1.00	**1.70**	1.65	1.77
reduce	2.17	2.18	**2.21**	1.89	2.16
fft	1.00	1.00	**1.02**	1.08	1.72
geometry	2.84	**8.09**	2.84	3.08	3.07
stencil	**1.12**	1.11	**1.12**	1.30	1.94
dense_la	**1.02**	1.01	1.01	1.01	1.11
sort	2.31	1.00	**2.35**	2.51	2.72
search	1.00	**1.07**	1.00	1.00	1.00
Overall	1.57	1.50	**1.59**	1.58	1.78

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B More Related Work

An LLM can be used as a problem solver in many ways. LATM [7] proposes a tooling framework wherein an LLM can create tools on its own and reuse them later. LATS [57] shows that LLMs can be integrated with Monte Carlo tree search for more effective planning. PHP [56] is an iterative framework where solutions are refined based on hints. L2MAC [20] introduces a system that mirrors the von Neumann architecture of a computer, utilizing a prompt program and file storage to generate long and coherent outputs. SFS [31] frames code generation as a black-box optimization problem, employing search techniques to enhance solution diversity and efficiency.

Moreover, multiple LLMs can collaborate for problem solving. MoA [46] orchestrates multiple LLM agents in a layered structure to aggregate and synthesize the solution. LLM-Debate [13, 14] enables structured inter-agent debates to reach consensus. LLM-Blender [26] ranks, filters and fuses top k responses from multiple LLMs, improving the performance on instruction following. DyLAN [33] employs dynamic team adjustment for different tasks. AgentPrune [54] eliminates redundant communication through graph pruning. GPTSwarm [58] optimizes collaboration through graph-based interactions. ChatLLM [19] facilitates interactions among multiple dialogue-based agents, enabling them to provide feedback and collectively improve the decision-making process. EoT [53] promotes cross-model communication among LLMs, integrating various communication paradigms and a confidence evaluation mechanism to enhance reasoning accuracy. AutoGen [50] automates the development of agent-based applications, streamlining processes such as data collection, model training, and evaluation to efficiently adapt pre-trained models to specialized tasks. FrugalGPT [9] introduces strategies to reduce the computational and financial costs associated with deploying LLMs, focusing on optimizing model architectures and inference techniques to achieve cost-effective performance without significantly compromising accuracy. MacNet [42] utilizes directed acyclic graphs to organize LLM agents, streamlining their interactive reasoning via topological ordering.

C Algorithm of LessonL

The LessonL framework presented in Section 3 is summarized in Algorithm 1, together with subroutines Algorithms 2–4.

Algorithm 1 The LessonL Framework

Input: code x; number of agents, n; number of rounds, T; number of selected lessons, k
1: $Z_{all} = \{\}$ ▷ initialize the lesson bank
2: **for** agent $j = 1, \ldots, n$ **do**   ▷ initial round
3:      Get new code $y_j^{(0)} = \text{LLM}_j(x)$
4:      Get speedup $s_j^{(0)} = \text{Exec}(y_j^{(0)})$ and lesson $z_j^{(0)} = \text{LLM}_j(x, y_j^{(0)}, s_j^{(0)})$
5:      Initialize factor $f_j^{(0)} = 1$
6:      $Z_{all}.\text{add}((z_j^{(0)}, s_j^{(0)}, f_j^{(0)}))$ ▷ insert lesson with speedup and factor into the lesson bank
7: **end for**
8: **for** round $t = 1, \ldots, T$ **do**
9:      **if** $tn \leq k$, let $Z^{(t)} = Z_{all}$; **otherwise then**   ▷ select lessons $Z^{(t)}$ used in the next round
10:         $Z_s, Z_{remain} = \text{SELECTHIGHSPEEDUP}(Z_{all}, \text{cnt} = \lceil k/2 \rceil)$ ▷ select half of the lessons
11:         $Z_q = \text{SELECTHIGHRELEVANCE}(Z_{remain}, \text{cnt} = \lfloor k/2 \rfloor)$ ▷ select another half
12:         $Z^{(t)} = Z_s \cup Z_q$ ▷ all selected lessons
13:     **end if**
14:     **for** agent $j = 1, \ldots, n$ **do**   ▷ next round
15:         Get new code $y_j^{(t)} = \text{LLM}_j(x, Z^{(t)})$ ▷ use selected lessons $Z^{(t)}$ to prompt
16:         Get speedup $s_j^{(t)} = \text{Exec}(y_j^{(t)})$ and lesson $z_j^{(t)} = \text{LLM}_j(x, y_j^{(t)}, s_j^{(t)})$
17:         Initialize factor $f_j^{(t)} = 1$
18:         $Z_{all}.\text{add}((z_j^{(t)}, s_j^{(t)}, f_j^{(t)}))$ ▷ insert lesson with speedup and factor into the lesson bank
19:     **end for**
20:     $\text{ADJUSTFACTOR}(Z^{(t)}, \{s_j^{(t)}\}_j)$ ▷ adjust factors for selected lessons based on performance
21: **end for**
22: **return** best among $\{y_1^{(0)}, \ldots, y_n^{(0)}, \ldots, y_1^{(T)}, \ldots, y_n^{(T)}\}$

Algorithm 2 SELECTHIGHSPEEDUP(Z_{all}, cnt)

1: Sort lessons in $Z_{all} = \{(z, s, f)\}$ in descending order of $s \times f$ ▷ speedup × factor
2: Let Z_s include top cnt lessons whose $s \times f \geq$ threshold
3: $Z_{remain} = Z_{all} \backslash Z_s$
4: **return** Z_s, Z_{remain}

Algorithm 3 SELECTHIGHRELEVANCE(Z_{remain}, cnt)

1: Sort lessons in $Z_{remain} = \{(z, s, f)\}$ in descending order of $\cos(z, x)$ ▷ $\cos(\text{lesson}, \text{code})$
2: Let Z_q include top cnt lessons
3: **return** Z_q

Algorithm 4 ADJUSTFACTOR($Z^{(t)}, \{s_j^{(t)}\}_j$)

1: **for** each lesson (z, s, f) in $Z^{(t)}$ **do**
2:      Initialize correction $c \leftarrow 0$
3:      **for** agent $j = 1, \ldots, n$ **do**
4:          **if** $s < s_j^{(t)}$ **then**   ▷ original speedup $s <$ yielded speedup when applying the lesson
5:              Update correction $c \leftarrow c + (1 + \epsilon)$
6:          **else**
7:              Update correction $c \leftarrow c + (1 - \epsilon)$
8:          **end if**
9:      **end for**
10:     Update factor $f \leftarrow c/n$
11: **end for**

D   Prompt Templates for Code Optimization

This section provides the prompt templates for code optimization. The initial prompt differs from subsequent prompts in that it starts with no lessons. The lesson solicitation prompt has several variants depending on the quality of the output code (speedup, slowdown, incorrect output, compilation error).

D.1   Initial Prompt

> You are given a piece of code written in {language}. Your task is to rewrite it in the same language to improve its performance (i.e., execution time). {use_openmp} Do not change the input/output behaviors of the code. Include the generated code between "'{language} and "'.
>
> // Code:
> {source_code}

In our experiments, the variable language is "C++". The variable use_openmp is set to "You should use OpenMP to parallelize the code." for OpenMP prompts; otherwise it is empty.

D.2   Lesson Solicitation

> The following are two functionally equivalent codes. They are compiled by using the same compiler and executed in the same environment. Code B runs {faster_or_slower} than Code A with a speedup {speedup}x. Explain why Code B is {faster_or_slower}. Be brief in the explanations. Use only one or two sentences.
>
> // Code A:

{source_code}

// Code B:
{target_code}

The following two codes are not functionally equivalent; that is, given the same input, they produce different outputs. Explain the reasons that make Code B nonequivalent to Code A. Be brief in the explanations. Use only one or two sentences.

// Code A:
{source_code}

// Code B:
{target_code}

The following are two codes. Code B attempts to improve the performance of Code A, but it has syntactic errors. Explain why Code B cannot be compiled. You may get hints from the compiler output provided after Code B. Be brief in the explanations. Use only one or two sentences.

// Code A:
{source_code}

// Code B:
{target_code}

Compiler output:
{compiler_output}

D.3 Subsequent Prompt With Lessons

You are given a piece of code written in {language}. Your task is to rewrite it in the same language to improve its performance (i.e., execution time). {use_openmp} Do not change the input/output behaviors of the code. Some lessons regarding correctness and performance are provided to help you rewrite the code. Include the generated code between "'{language} and "'.

// Code:
{source_code}

While you rewrite the code, consider the following lessons. If Code A and Code B appear in the lessons, Code A refers to the given code and Code B refers to an attempted rewrite. Code B may not be optimal and it could be even worse than Code A.

{lesson[0]}

{lesson[1]}

{lesson[2]}

{lesson[3]}

Besides the above lessons, consider other optimization strategies that can more significantly improve the performance of the given code.

Each item in the lesson list can be one of the following:

• "Lesson {idx} slightly improves the code performance. {lesson_content} However, despite the code performance improvement, the speedup is only marginal."

• "Lesson {idx} significantly improves the code performance. {lesson_content}"

• "Lesson {idx} degrades the code performance. {lesson_content}"

• "Lesson {idx} compromises code equivalence. {lesson_content}"

• "Lesson {idx} produces non-compilable code. {lesson_content}"

E   Prompt Templates for Code Generation

This section provides the prompt templates for code generation. The initial prompt differs from subsequent prompts in that it starts with no lessons. Following common practice, we include an example in the prompt, demonstrating the function signature, function description, and function completion.

E.1   Initial Prompt

> You are given a function signature in {language} together with a docstring that explains what the function does. Your task is to implement the function according to the docstring. You should restate the function signature and docstring. Include the generated code between "'{language} and "'. For example, given the function signature and docstring
>
> {example_function_signature_and_docstring}
>
> You should respond with
>
> {example_generation}
>
> ### Here is the function to implement:
> {function_signature_and_docstring}

In our experiments, the variable language is "Python".

The variable example_function_signature_and_docstring is

```
def sum(a: float, b: float) -> float:
    """ Return the sum of two floats a and b """
```

The variable example_generation is

```
'''Python
def sum(a: float, b: float) -> float:
    """ Return the sum of two floats a and b """
    return a + b
'''
```

> The following completed code is incorrect; i.e., it does not exactly reflect the description in the docstring. The code passes only {num_pass_cases} test cases out of {num_total_cases}, leaving {num_fail_cases} failed. Explain why the code is incorrect (that is, why it fails some test cases). Be brief in the explanations. Use only one or two sentences.
>
> ### Completed code:
> {completed_code}

E.3 Subsequent Prompt With Lessons

> You are given a function signature in {language} together with a docstring that explains what the function does. Your task is to implement the function according to the docstring. You should restate the function signature and docstring. Some lessons are provided to help you implement the function. Include the generated code between "'{language} and "'. For example, given the function signature and docstring
>
> {example_function_signature_and_docstring}
>
> You should respond with
>
> {example_generation}
>
> ### Here is the function to implement:
> {function_signature_and_docstring}
>
> While you implement the function, consider the following lessons.
>
> {lesson[0]}
>
> {lesson[1]}
>
> {lesson[2]}
>
> {lesson[3]}

Each item in the lesson list is:

• "Lesson {idx} reasons why the code does not pass all test cases. {lesson_content}"

F ParEval Modification for Code Optimization

We run experiments on a code optimization benchmark adapted from ParEval [35], which is an existing benchmark for LLM code generation. At a high level, our code optimization benchmark tasks an LLM with optimizing a piece of source code. The LLM's response is evaluated based on its correctness as well as the speedup it can achieve over the source code. In this section, we discuss our adaptation and improvement of ParEval.

F.1 Adaptation from ParEval

In the original ParEval benchmark, the LLM is tasked with completing a function definition. To test correctness, the LLM generated code is run on custom test cases and it is compared to a provided baseline.

In our code optimization benchmark, we use the baseline as the source code to be optimized. Some changes are made to the source code for particular problems, such as adding definitions of custom types, so that the source code can run on its own.

F.2 Benchmarking

ParEval performs benchmarking by taking the average runtime of the baseline code and the LLM generated code over a fixed number of runs. This approach can produce noisy results if the execution time of the baseline code or LLM generated code is very short. To address this issue, we use Nanobench, a lightweight benchmarking library [2]. Nanobench varies the number of benchmarking runs based on the runtime of the code as well as its variance, thereby producing stable results on codes with shorter runtimes. Overall, Nanobench produces consistent and stable results, which ensures that LLMs are accurately evaluated on the benchmark.

F.3 Input Generation

We ensure that the randomness used when generating inputs for the baseline code and the LLM generated code are the same. Fixing randomness among the two is necessary when benchmarking as the runtime of certain problems can vary greatly depending on the specific inputs provided. For example, consider the task of finding a value in a list. If the value is the first element in the list, then the runtime will be much faster compared to when the value is not in the list. We ensure all inputs are generated from a seeded random generator, and ensure that the seed is the same when running the baseline and when running the LLM generated code.

We also improve the random input generation for some problems to ensure that all cases are considered and thoroughly tested. First, we look at improving the input generation for graph problems, which includes finding the number of connected components in a graph, finding the size of the largest connected component in the graph, and finding the shortest path length in a graph. Previously, ParEval generated random graphs using the $G(n, p)$ Erdos-Renyi model. Each edge in the n-vertex graph would be independently added with probability p. The issue with this type of generation using a fixed probability p for large graphs is that the resulting graph is almost surely connected, which means the number of connected components is 1, the size of the largest connected component is n, and there always exists a path between any two vertices. Therefore, some incorrect codes can pass the provided test cases as the test cases are not varied enough. Therefore, to mitigate this issue, we vary the distribution of generated graphs by implementing a simple random graph algorithm, which first randomly selects the number of edges in the graph, and then for each edge randomly selects the two vertices incident to the edge. We limit the number of edges in the graphs to generate graphs with a wide range of connected components for the purposes of these problems. Alternative random graph generation algorithms can be used, but we find this simple algorithm works for the specific graph problems in the benchmark.

Next, we look at improving the input generation for a search problem which tasks the user with finding a target value in a list. The current code from ParEval generates random inputs from a narrow range of values, which means that when benchmarking on large input sizes, all values in the range will be generated, which has the effect that any target value will almost surely be in the list. This allows incorrect codes, such as a function that always returns true, to potentially pass the provided test cases on large input sizes.

F.4 Other Fixes

We also fix several bugs in the existing ParEval benchmark when adapting it to the code optimization setting. Several problems encounter overflow issues when scaling up the input size, as the generated input values are quite large. We fix this behavior by changing the range of random values generated for these specific problems. In addition, we fix bugs in the baseline codes of several problems that produce incorrect results, such as dense and sparse matrix LU decomposition, finding the convex hull and computing its perimeter, and computing the maximum contiguous subarray sum.

G Additional Information of Evaluation Setup

G.1 Baseline Implementation for Code Optimization

CoT: We prompt the agent by adding "Let's think step by step" at the end to solicit reasoning.

Reflexion: Reflexion was originally used for code generation, wherein each iteration, synthetic test cases are constructed and verbal feedback of the generated code is produced based on the code's correctness on these test cases. The iterations will terminate when the generated code is correct. For code optimization, we replace the testing of correctness with the measurement of speedup. We have Reflexion run ten iterations and choose the code with the best speedup over these iterations as the final solution.

MapCoder: MapCoder was originally used for code generation. Given a coding task, MapCoder first retrieves similar tasks as examples, then proposes a few plans to solve the current task, and executes the plans one-by-one in the decreasing order of its confidence. For every plan, a few iterations of coding–debugging are performed. MapCoder will terminate when a plan leads to correct code. For code optimization, we replace the testing of correctness in the debugging step with the measurement of speedup. We have MapCoder produce three plans and run five iterations for each plan. We choose the code with the best speedup over all plans and all iterations as the final solution.

MoA: We follow the default configuration, using three layers and having GPT-4o as the aggregator.

G.2 Baseline Implementation for Code Generation

Different from ParEval, the benchmarks for code generation offer a fixed set of test cases. HumanEval/MBPP offer very few test cases (which is a motivation of developing HumanEval+/MBPP+). Hence, in some of the baselines, there is a choice between using the given set of test cases or using syntheticly generated test cases in each iteration. The drawback of using the given set is that the generated code may overfit the test cases when the set is small, whereas the drawback of using a synthetic set, which is generated by an LLM, is that the test cases may be incorrect. We consider overfit doing more harm and use syntheticly generated test cases instead. Following Reflexion [43], each time, six test cases are generated by Qwen14B.

For all four benchmarks, we use EvalPlus [32] to evaluate the generation.

LessonL: As opposed to the code optimization case, for code generation, LessonL terminates the iterations whenever correct code, as measured by the synthetic test cases, is produced. The pass@1 score is reported by using the given test cases.

CoT: Same as the configuration in the code optimization case.

Reflexion: We use a maximum of ten iterations.

MapCoder: We use three plans and a maximum of five iterations each.

MoA: Same as the configuration in the code optimization case.

G.3 Hyperparameters, Hardware, and Timing

Hyperparameters. The hyperparameters used by LessonL are summarized in Table 5, unless otherwise stated (e.g., in ablation studies). We barely tune them. Additionally, LLM parameters for all models are also included in Table 5. Both code optimization and code generate tasks share the same set of hyperparameters, except that code generation doesn't have THRESHOLD in SELECTHIGH-SPEEDUP.

Table 5: Hyperparameters.

Number of agents, n	3
Number of rounds, T	4
Number of selected lessons, k	4
Threshold in SELECTHIGHSPEEDUP	1.1
ϵ in ADJUSTFACTOR	0.1
LLM temperature	0.2
LLM frequency penalty	0.5

Hardware. For code optimization benchmarks (ParEval and PolyBench), the LLMs are hosted on an Amazon EC2 g6e.12xlarge instance through vLLM [28]. This instance contains 4x Nvidia L40S

GPUs (192 GB). Additionally, we use c5.4xlarge instances, each with 16x Intel Xeon Platinum vCPUs (3.4 GHz), to ensure resource isolation from the g6e.12xlarge instance when evaluating code correctness and speedup. All baselines and methods use one single c5.4xlarge exclusively for performance evalution to ensure consistent hardware resources. For code generation (HumanEval, HumanEval+, MBPP, MBPP+), we conduct experiments on a server with 8 A6000 GPUs. The LLMs are hosted on A6000 GPUs.

Timing. A full run of LessonL on ParEval takes a few hours. The runtime for other methods and/or other benchmarks falls in a similar scale.

H Supplementary Information for Figure 3

Figure 3 in the main text shows the performance variation as the number of rounds increases for methods that do allow iterations (MoA, MapCoder, and Reflexion). In MoA, they are called "layers." While prior work suggests that these methods' performance may improve with more iterations, we find that it does not necessarily happen for code optimization. In what follows, for each method (including LessonL), we show an example of the output code at each round and discuss the speedup variation. For readability, the codes were edited by removing nonessential comments (such as problem description) and white spaces (including indentation).

H.1 LessonL

Below are the outputs of LessonL. The speedup generally improves over rounds. The code is to perform convolution on an image. In round 0, the code is improved by removing the redundant zero-additions. The code in round 1 is identical to that in round 0. In round 2, the code handles the interior region and the boundary separately, yielding further improvement. Round 3 achieves a significant further improvement by removing the temporary storage and avoiding extra data copies. The code in round 4 is identical to that in round 3.

Original code

```cpp
#pragma once
#include <vector>

const int edgeKernel[3][3] = {{-1, -1, -1}, {-1, 8, -1}, {-1, -1, -1}};

/* Convolve the edge kernel with a grayscale image. Each pixel will be
replaced with
   the dot product of itself and its neighbors with the edge kernel.
   Use a value of 0 for pixels outside the image's boundaries and clip
   outputs between 0 and 255.
   imageIn and imageOut are NxN grayscale images stored in row-major.
   Store the output of the computation in imageOut.
   Example:
   input: [[112, 118, 141, 152],
   [93, 101, 119, 203],
   [45, 17, 16, 232],
   [82, 31, 49, 101]]
   output: [[255, 255, 255, 255],
   [255, 147, 0, 255],
   [36, 0, 0, 255],
   [255, 39, 0, 255]]
*/
void convolveKernel(std::vector<int> const& imageIn, std::vector<int>
&imageOut, size_t N) {
  for (size_t i = 0; i < N; i++) {
    for (size_t j = 0; j < N; j++){
      int sum = 0;
      for (int k = -1; k < 2; k++) {
```

```cpp
        for (int l = -1; l < 2; l++){
          int x = i + k;
          int y = j + l;
          if ((x < 0) || (x >= N) || (y < 0) || (y >= N)) {
            sum += 0;
          } else {
            sum += imageIn[x * N + y] * edgeKernel[k + 1][l + 1];
          }
        }
      }
      if (sum < 0) {
        imageOut[i * N + j] = 0;
      } else if (sum > 255) {
        imageOut[i * N + j] = 255;
      } else {
        imageOut[i * N + j] = sum;
      }
    }
  }
}
```

LessonL, round 0, speedup 1.71x

```cpp
#pragma once
#include <vector>
#include <algorithm>

const int edgeKernel[3][3] = {{-1, -1, -1}, {-1, 8, -1}, {-1, -1, -1}};

void convolveKernel(std::vector<int> const& imageIn, std::vector<int>
&imageOut, size_t N) {
  for (size_t i = 0; i < N; i++) {
    for (size_t j = 0; j < N; j++) {
      int sum = 0;
      for (int k = -1; k < 2; k++) {
        for (int l = -1; l < 2; l++) {
          int x = i + k;
          int y = j + l;
          if (x >= 0 && x < N && y >= 0 && y < N) {
            sum += imageIn[x * N + y] * edgeKernel[k + 1][l + 1];
          }
        }
      }
      imageOut[i * N + j] = std::clamp(sum, 0, 255);
    }
  }
}
```

LessonL, round 1, speedup 1.71x

```cpp
#pragma once
#include <vector>
#include <algorithm> // For std::clamp

const int edgeKernel[3][3] = {{-1, -1, -1}, {-1, 8, -1}, {-1, -1, -1}};

void convolveKernel(std::vector<int> const& imageIn, std::vector<int>
&imageOut, size_t N) {
  for (size_t i = 0; i < N; i++) {
```

```
    for (size_t j = 0; j < N; j++) {
      int sum = 0;
      for (int k = -1; k < 2; k++) {
        for (int l = -1; l < 2; l++) {
          int x = i + k;
          int y = j + l;
          if (x >= 0 && x < N && y >= 0 && y < N) {
            sum += imageIn[x * N + y] * edgeKernel[k + 1][l + 1];
          }
        }
      }
      imageOut[i * N + j] = std::clamp(sum, 0, 255);
    }
  }
}
```

```cpp
#pragma once
#include <vector>
#include <algorithm> // For std::clamp

const int edgeKernel[3][3] = {{-1, -1, -1}, {-1, 8, -1}, {-1, -1, -1}};

void convolveKernel(std::vector<int> const& imageIn, std::vector<int>
&imageOut, size_t N) {
  // Temporary vector to store intermediate results
  std::vector<int> temp(N * N, 0);

  // Process the interior pixels
  for (size_t i = 1; i < N - 1; i++) {
    for (size_t j = 1; j < N - 1; j++) {
      int sum = 0;
      for (int k = -1; k < 2; k++) {
        for (int l = -1; l < 2; l++) {
          sum += imageIn[(i + k) * N + (j + l)] * edgeKernel[k + 1][l
          + 1];
        }
      }
      temp[i * N + j] = sum;
    }
  }

  // Process the boundary pixels
  for (size_t i = 0; i < N; i++) {
    for (size_t j = 0; j < N; j++) {
      if (i == 0 || i == N - 1 || j == 0 || j == N - 1) {
        int sum = 0;
        for (int k = -1; k < 2; k++) {
          for (int l = -1; l < 2; l++) {
            int x = i + k;
            int y = j + l;
            if (x >= 0 && x < N && y >= 0 && y < N) {
              sum += imageIn[x * N + y] * edgeKernel[k + 1][l + 1];
            }
          }
        }
        temp[i * N + j] = sum;
      }
    }
  }
```

```
  // Apply clipping to the results
  for (size_t i = 0; i < N * N; i++) {
    imageOut[i] = std::clamp(temp[i], 0, 255);
  }
}
```

LessonL, round 3, speedup 9.45x

```
#pragma once
#include <vector>
#include <algorithm> // For std::clamp

const int edgeKernel[3][3] = {{-1, -1, -1}, {-1, 8, -1}, {-1, -1, -1}};

void convolveKernel(std::vector<int> const& imageIn, std::vector<int>
&imageOut, size_t N) {
  // Process interior pixels
  for (size_t i = 1; i < N - 1; i++) {
    for (size_t j = 1; j < N - 1; j++) {
      int sum = 0;
      for (int k = -1; k < 2; k++) {
        for (int l = -1; l < 2; l++) {
          sum += imageIn[(i + k) * N + (j + l)] * edgeKernel[k + 1][l
          + 1];
        }
      }
      imageOut[i * N + j] = std::clamp(sum, 0, 255);
    }
  }

  // Process boundary pixels
  for (size_t i = 0; i < N; i++) {
    for (size_t j = 0; j < N; j++) {
      if (i == 0 || i == N - 1 || j == 0 || j == N - 1) {
        int sum = 0;
        for (int k = -1; k < 2; k++) {
          for (int l = -1; l < 2; l++) {
            int x = i + k;
            int y = j + l;
            if ((x < 0) || (x >= N) || (y < 0) || (y >= N)) {
              sum += 0;
            } else {
              sum += imageIn[x * N + y] * edgeKernel[k + 1][l + 1];
            }
          }
        }
        imageOut[i * N + j] = std::clamp(sum, 0, 255);
      }
    }
  }
}
```

LessonL, round 4, speedup 9.50x

```
#pragma once
#include <vector>
#include <algorithm> // For std::clamp
```

```cpp
const int edgeKernel[3][3] = {{-1, -1, -1}, {-1, 8, -1}, {-1, -1, -1}};

void convolveKernel(std::vector<int> const& imageIn, std::vector<int>
&imageOut, size_t N) {
  // Process interior pixels
  for (size_t i = 1; i < N - 1; i++) {
    for (size_t j = 1; j < N - 1; j++) {
      int sum = 0;
      for (int k = -1; k < 2; k++) {
        for (int l = -1; l < 2; l++) {
          sum += imageIn[(i + k) * N + (j + l)] * edgeKernel[k + 1][l
          + 1];
        }
      }
      imageOut[i * N + j] = std::clamp(sum, 0, 255);
    }
  }

  // Process boundary pixels
  for (size_t i = 0; i < N; i++) {
    for (size_t j = 0; j < N; j++) {
      if (i == 0 || i == N - 1 || j == 0 || j == N - 1) {
        int sum = 0;
        for (int k = -1; k < 2; k++) {
          for (int l = -1; l < 2; l++) {
            int x = i + k;
            int y = j + l;
            if ((x < 0) || (x >= N) || (y < 0) || (y >= N)) {
              sum += 0;
            } else {
              sum += imageIn[x * N + y] * edgeKernel[k + 1][l + 1];
            }
          }
        }
        imageOut[i * N + j] = std::clamp(sum, 0, 255);
      }
    }
  }
}
```

H.2 MoA

Below are the outputs of MoA. The speedup deteriorates over rounds/layers. The code finds the kth smallest element of an array. The code in round 0 achieves the highest speedup by implementing the divide-and-conquer algorithm. In round 1, the speedup reduces because of the type cast from time_t to unsigned int. In round 2, the speedup further reduces because the implementation is changed from divide-and-conquer to using std::nth_element. The codes in round 3, 4, and 5 are identical to that in round 2.

Original code

```cpp
#pragma once
#include <algorithm>
#include <vector>

/* Find the k-th smallest element of the vector x.
   Example:

   input: x=[1, 7, 6, 0, 2, 2, 10, 6], k=4
```

```
   output: 6
*/
int findKthSmallest(std::vector<int> const& x, int k) {
  std::vector<int> x_copy = x;
  std::sort(x_copy.begin(), x_copy.end());
  return x_copy[k-1];
}
```

MoA, round 0, speedup 12.36x

```
#pragma once
#include <vector>
#include <cstdlib> // For std::rand() and std::srand()

int partition(std::vector<int>& x, int left, int right, int
pivotIndex) {
  int pivotValue = x[pivotIndex];
  std::swap(x[pivotIndex], x[right]); // Move pivot to end
  int storeIndex = left;
  for (int i = left; i < right; ++i) {
    if (x[i] < pivotValue) {
      std::swap(x[storeIndex], x[i]);
      ++storeIndex;
    }
  }
  std::swap(x[right], x[storeIndex]); // Move pivot to its final place
  return storeIndex;
}

int quickselect(std::vector<int>& x, int left, int right, int k) {
  if (left == right) { // If the list contains only one element,
    return x[left];  // return that element
  }

  // Select a random pivotIndex
  int pivotIndex = left + std::rand() % (right - left + 1);

  // Find the pivot position in a sorted list
  pivotIndex = partition(x, left, right, pivotIndex);

  // The pivot is in its final sorted position
  if (k == pivotIndex) {
    return x[k];
  } else if (k < pivotIndex) {
    return quickselect(x, left, pivotIndex - 1, k);
  } else {
    return quickselect(x, pivotIndex + 1, right, k);
  }
}

int findKthSmallest(std::vector<int> const& x, int k) {
  std::vector<int> x_copy = x;
  std::srand(time(0)); // Seed the random number generator
  return quickselect(x_copy, 0, x_copy.size() - 1, k - 1);
}
```

MoA, round 1, speedup 8.10x

```
#pragma once
```

```cpp
#include <vector>
#include <cstdlib> // For std::rand() and std::srand()
#include <ctime>  // For std::time()

int partition(std::vector<int>& x, int left, int right, int
pivotIndex) {
  int pivotValue = x[pivotIndex];
  std::swap(x[pivotIndex], x[right]); // Move pivot to end
  int storeIndex = left;
  for (int i = left; i < right; ++i) {
    if (x[i] < pivotValue) {
      std::swap(x[storeIndex], x[i]);
      ++storeIndex;
    }
  }
  std::swap(x[right], x[storeIndex]); // Move pivot to its final place
  return storeIndex;
}

int quickselect(std::vector<int>& x, int left, int right, int k) {
  if (left == right) { // If the list contains only one element,
    return x[left];  // return that element
  }

  // Select a random pivotIndex
  int pivotIndex = left + std::rand() % (right - left + 1);

  // Find the pivot position in a sorted list
  pivotIndex = partition(x, left, right, pivotIndex);

  // The pivot is in its final sorted position
  if (k == pivotIndex) {
    return x[k];
  } else if (k < pivotIndex) {
    return quickselect(x, left, pivotIndex - 1, k);
  } else {
    return quickselect(x, pivotIndex + 1, right, k);
  }
}

int findKthSmallest(std::vector<int> const& x, int k) {
  std::vector<int> x_copy = x;
  std::srand(static_cast<unsigned int>(std::time(0))); // Seed the
  random number generator
  return quickselect(x_copy, 0, x_copy.size() - 1, k - 1);
}
```

MoA, round 2, speedup 6.71x

```cpp
#pragma once
#include <vector>
#include <algorithm>

int findKthSmallest(std::vector<int> const& x, int k) {
  std::vector<int> x_copy = x;
  std::nth_element(x_copy.begin(), x_copy.begin() + k - 1,
  x_copy.end());
  return x_copy[k - 1];
}
```

30

MoA, round 3, speedup 6.72x

```cpp
#pragma once
#include <vector>
#include <algorithm>

int findKthSmallest(std::vector<int> const& x, int k) {
  std::vector<int> x_copy = x;
  std::nth_element(x_copy.begin(), x_copy.begin() + k - 1,
  x_copy.end());
  return x_copy[k - 1];
}
```

MoA, round 4, speedup 6.74x

```cpp
#pragma once
#include <algorithm>
#include <vector>

int findKthSmallest(std::vector<int> const& x, int k) {
  std::vector<int> x_copy = x;
  std::nth_element(x_copy.begin(), x_copy.begin() + k - 1,
  x_copy.end());
  return x_copy[k - 1];
}
```

MoA, round 5, speedup 6.72x

```cpp
#pragma once
#include <algorithm>
#include <vector>

int findKthSmallest(std::vector<int> const& x, int k) {
  std::vector<int> x_copy = x;
  std::nth_element(x_copy.begin(), x_copy.begin() + k - 1,
  x_copy.end());
  return x_copy[k - 1];
}
```

H.3 MapCoder

Below are the outputs of MapCoder. The speedup does not follow a monotonic trend over rounds. MapCoder works by using one agent to generate a few plans and evaluate each plan with a confidence rating. The plans are executed in the decreasing order of confidence, but it does not necessarily mean that the earlier executed plans will lead to better performing codes. In this example, the original code is to compute the highest node degree in an undirected graph. This amounts to computing the sum of each row of the adjacency and comparing different rows. All the plans perform some sort of loop unrolling, the idea of which is gaining less attraction nowadays because loop unrolling should be delegated to the compiler for better instruction parallelism and fewer comparisons. As a result, none of the output codes runs faster than the original code.

Oiginal code

```cpp
#pragma once
#include <algorithm>
#include <vector>
#include <queue>
#include <stack>
#include <deque>
#include <limits>

/* Compute the highest node degree in the undirected graph. The graph
is defined in the adjacency matrix A.
   A is an NxN adjacency matrix stored in row-major. A is undirected.
   Example:
   input: [[0, 0, 0, 1], [0, 0, 0, 1], [0, 0, 0, 1], [1, 1, 1, 0]]
   output: 3
*/
int maxDegree(std::vector<int> const& A, size_t N) {
  int highestDegree = 0;
  for (int i = 0; i < N; i += 1) {
    int degree = 0;
    for (int j = 0; j < N; j += 1) {
      degree += A[i * N + j];
    }
    highestDegree = std::max(highestDegree, degree);
  }
  return highestDegree;
}
```

MapCoder, round 0, speedup 1.00x

```cpp
#pragma once
#include <algorithm>
#include <vector>
#include <numeric> // For std::accumulate

int maxDegree(const std::vector<int>& A, size_t N) {
  int highestDegree = 0;
  for (size_t i = 0; i < N; i += 2) { // Loop Unrolling: Process two
  rows at a time
    int degree1 = std::accumulate(A.begin() + i * N, A.begin() + (i +
    1) * N, 0);
    int degree2 = (i + 1 < N) ? std::accumulate(A.begin() + (i + 1) *
    N, A.begin() + (i + 2) * N, 0) : 0;
    highestDegree = std::max(highestDegree, std::max(degree1,
    degree2));
  }
  return highestDegree;
}
```

MapCoder, round 1, speedup 0.96x

```cpp
#pragma once
#include <algorithm>
#include <vector>
#include <queue>
#include <stack>
#include <deque>
#include <limits>

inline int maxDegree(const std::vector<int>& A, size_t N) {
```

```
    int highestDegree = 0;
    for (size_t i = 0; i < N; i += 4) { // Loop unrolling by a factor of
    4
      int degree0 = 0, degree1 = 0, degree2 = 0, degree3 = 0;
      if (i + 0 < N) {
        for (size_t j = 0; j < N; j += 4) { // Inner loop unrolling by a
        factor of 4
          degree0 += A[(i + 0) * N + j + 0];
          degree0 += A[(i + 0) * N + j + 1];
          degree0 += A[(i + 0) * N + j + 2];
          degree0 += A[(i + 0) * N + j + 3];
          if (i + 1 < N) {
            degree1 += A[(i + 1) * N + j + 0];
            degree1 += A[(i + 1) * N + j + 1];
            degree1 += A[(i + 1) * N + j + 2];
            degree1 += A[(i + 1) * N + j + 3];
          }
          if (i + 2 < N) {
            degree2 += A[(i + 2) * N + j + 0];
            degree2 += A[(i + 2) * N + j + 1];
            degree2 += A[(i + 2) * N + j + 2];
            degree2 += A[(i + 2) * N + j + 3];
          }
          if (i + 3 < N) {
            degree3 += A[(i + 3) * N + j + 0];
            degree3 += A[(i + 3) * N + j + 1];
            degree3 += A[(i + 3) * N + j + 2];
            degree3 += A[(i + 3) * N + j + 3];
          }
        }
      }
      highestDegree = std::max({highestDegree, degree0, degree1,
      degree2, degree3});
    }
    return highestDegree;
}
```

MapCoder, round 2, speedup 0.65x

```cpp
#pragma once
#include <algorithm>
#include <vector>
#include <queue>
#include <stack>
#include <deque>
#include <limits>

int maxDegree(const std::vector<int>& A, size_t N) {
  int highestDegree = 0;
  for (size_t i = 0; i < N; i += 1) {
    int degree = 0;
    // Unrolling the inner loop
    for (size_t j = 0; j < N; j += 4) {
      if (j + 0 < N) degree += A[i * N + j + 0];
      if (j + 1 < N) degree += A[i * N + j + 1];
      if (j + 2 < N) degree += A[i * N + j + 2];
      if (j + 3 < N) degree += A[i * N + j + 3];
    }
    highestDegree = std::max(highestDegree, degree);
  }
  return highestDegree;
```

```
}
```

MapCoder, round 3, speedup 0.74x

```cpp
#pragma once
#include <algorithm>
#include <vector>
#include <queue>
#include <stack>
#include <deque>
#include <limits>

int maxDegree(std::vector<int> const& A, size_t N) {
  int highestDegree = 0;
  for (int i = 0; i < N; i += 1) {
    int degree = 0;
    for (int j = 0; j < N; j += 4) { // Unroll the loop by a factor of
    4
      degree += A[i * N + j];
      if (j + 1 < N) degree += A[i * N + j + 1];
      if (j + 2 < N) degree += A[i * N + j + 2];
      if (j + 3 < N) degree += A[i * N + j + 3];
    }
    highestDegree = std::max(highestDegree, degree);
  }
  return highestDegree;
}
```

MapCoder, round 4, speedup 1.00x

```cpp
#pragma once
#include <algorithm>
#include <vector>
#include <queue>
#include <stack>
#include <deque>
#include <limits>

int maxDegree(std::vector<int> const& A, size_t N) {
  int highestDegree = 0;
  for (int i = 0; i < N; i += 1) {
    int degree = 0;
    int j = 0;
    // Unroll the inner loop by a factor of 4
    for (; j <= N - 4; j += 4) {
      degree += A[i * N + j];
      degree += A[i * N + j + 1];
      degree += A[i * N + j + 2];
      degree += A[i * N + j + 3];
    }
    // Handle the remaining elements
    for (; j < N; j += 1) {
      degree += A[i * N + j];
    }
    highestDegree = std::max(highestDegree, degree);
  }
  return highestDegree;
}
```

H.4 Reflexion

Below are the outputs of Reflexion. The speedup generally improves over rounds. The code is to perform cell XOR. The code in round 0 fails to compile. The code in round 1 is identical to the original code. The code in round 2 is marginally slower than the original code. This is likely because of the timing noise; in fact, this code precomputes the common expression 'i * N + j', which can slightly improves the runtime. In round 3, the code performance drastically improves, because half of the if-tests are eliminated. The code in round 4 is identical to that of round 3.

Original code

```cpp
#pragma once
#include <vector>

/* Set every cell's value to 1 if it has exactly one neighbor that's a
1. Otherwise set it to 0.
   Note that we only consider neighbors and not input_{i,j} when
   computing output_{i,j}.
   input and output are NxN grids of ints in row-major.
   Example:
   input: [[0, 1, 1, 0],
   [1, 0, 0, 0],
   [0, 0, 0, 0],
   [0, 1, 0, 0]
   output: [[0, 0, 1, 1],
   [1, 0, 0, 1],
   [0, 0, 1, 0],
   [1, 0, 1, 0]]
*/
void cellsXOR(std::vector<int> const& input, std::vector<int> &output,
size_t N) {
  for (size_t i = 0; i < N; ++i) {
    for (size_t j = 0; j < N; ++j) {
      int count = 0;
      if (i > 0 && input[(i-1)*N + j] == 1) ++count;
      if (i < N-1 && input[(i+1)*N + j] == 1) ++count;
      if (j > 0 && input[i*N + j-1] == 1) ++count;
      if (j < N-1 && input[i*N + j+1] == 1) ++count;
      output[i*N + j] = (count == 1) ? 1 : 0;
    }
  }
}
```

Reflexion, round 0, compilation failed

```cpp
#pragma once
#include <vector>

// Lookup table for the count of neighbors
const int neighborCountLookup[9] = {
  0, 0, 0, // 0 neighbors
  1, 0, 0, // 1 neighbor
  0, 0, 0, // 2 neighbors
  0, 0, 0, // 3 neighbors
  0, 0, 0, // 4 neighbors
  0, 0, 0, // 5 neighbors
  0, 0, 0, // 6 neighbors
  0, 0, 0, // 7 neighbors
  0, 0, 0  // 8 neighbors
};
```

```
void cellsXOR(std::vector<int> const& input, std::vector<int> &output,
size_t N) {
  for (size_t i = 0; i < N; ++i) {
    for (size_t j = 0; j < N; ++j) {
      int count = 0;
      if (i > 0) {
        if (j > 0) count += input[(i-1)*N + j-1];
        count += input[(i-1)*N + j];
        if (j < N-1) count += input[(i-1)*N + j+1];
      }
      if (j > 0) count += input[i*N + j-1];
      if (j < N-1) count += input[i*N + j+1];
      if (i < N-1) {
        if (j > 0) count += input[(i+1)*N + j-1];
        count += input[(i+1)*N + j];
        if (j < N-1) count += input[(i+1)*N + j+1];
      }
      output[i*N + j] = neighborCountLookup[count];
    }
  }
}
```

Reflexion, round 1, speedup 1.00x

```
#pragma once
#include <vector>

void cellsXOR(std::vector<int> const& input, std::vector<int> &output,
size_t N) {
  for (size_t i = 0; i < N; ++i) {
    for (size_t j = 0; j < N; ++j) {
      int count = 0;
      if (i > 0 && input[(i-1)*N + j] == 1) ++count;
      if (i < N-1 && input[(i+1)*N + j] == 1) ++count;
      if (j > 0 && input[i*N + j-1] == 1) ++count;
      if (j < N-1 && input[i*N + j+1] == 1) ++count;
      output[i*N + j] = (count == 1) ? 1 : 0;
    }
  }
}
```

Reflexion, round 2, speedup 0.95x

```
#pragma once
#include <vector>

void cellsXOR(std::vector<int> const& input, std::vector<int> &output,
size_t N) {
  for (size_t i = 0; i < N; ++i) {
    for (size_t j = 0; j < N; ++j) {
      int count = 0;
      size_t idx = i * N + j;
      if (i > 0) {
        if (input[idx - N] == 1) ++count;
      }
      if (i < N - 1) {
        if (input[idx + N] == 1) ++count;
      }
```

```
      if (j > 0) {
        if (input[idx - 1] == 1) ++count;
      }
      if (j < N - 1) {
        if (input[idx + 1] == 1) ++count;
      }
      output[idx] = (count == 1) ? 1 : 0;
    }
  }
}
```

Reflexion, round 3, speedup 10.81x

```
#pragma once
#include <vector>

void cellsXOR(std::vector<int> const& input, std::vector<int> &output,
size_t N) {
  for (size_t i = 0; i < N; ++i) {
    for (size_t j = 0; j < N; ++j) {
      int count = 0;
      size_t idx = i * N + j;
      if (i > 0) {
        count += input[idx - N];
      }
      if (i < N - 1) {
        count += input[idx + N];
      }
      if (j > 0) {
        count += input[idx - 1];
      }
      if (j < N - 1) {
        count += input[idx + 1];
      }
      output[idx] = (count == 1) ? 1 : 0;
    }
  }
}
```

Reflexion, round 4, speedup 10.71x

```
#pragma once
#include <vector>

void cellsXOR(std::vector<int> const& input, std::vector<int> &output,
size_t N) {
  for (size_t i = 0; i < N; ++i) {
    for (size_t j = 0; j < N; ++j) {
      int count = 0;
      size_t idx = i * N + j;
      if (i > 0) {
        count += input[idx - N];
      }
      if (i < N - 1) {
        count += input[idx + N];
      }
      if (j > 0) {
        count += input[idx - 1];
      }
```

```
    if (j < N - 1) {
      count += input[idx + 1];
    }
    output[idx] = (count == 1) ? 1 : 0;
  }
 }
}
```

I Supplementary Information for Figure 4

Figure 4 in the main text shows the monetary cost and FLOPS for various models and methods. For this, several pieces of information need to be estimated: the unit cost, FLOPS per token, and the number of tokens.

The pricing information for the open-source models is based on `https://api.together.ai/models` and that for the GPT models is from `https://openai.com/api/pricing`. Specifically, we estimate the unit cost of Deepseek7B, Qwen7B, and Qwen14B to be $0.20, $0.20, and $0.30, respectively, per 1M tokens. Meanwhile, GPT-4o mini costs $0.15/$0.60 per 1M input/output tokens, while GPT-4o costs $2.50/$10.00 per 1M input/output tokens. The pricing information is retrieved in April 2025 and it may be time-sensitive. The number of tokens is estimated based on the CodeLlama-7b-Instruct-hf tokenizer from Huggingface [49]. For multi-agent methods, we assume each agent consumes/generates the same amount of tokens. Table 6 summarizes the information for one experiment on the ParEval benchmark.

Table 6: (Estimated) Costs for running a ParEval experiment. For multi-agent methods, the listed number of tokens is the sum over all agents. For MoA, the first row shows the number of tokens for the open-source agents, while the second row for the GPT-4o aggregator (separately priced). For CoT, Reflexion, and MapCoder, all agents are Qwen14B.

	Speedup	#Input tokens	#Output tokens	#Total tokens	Cost ($)
Qwen14B(1)	1.60	23421	33550	56971	0.017
Qwen14B(20)	1.70	468420	671001	1139421	0.342
GPT-4o mini	1.57	20815	26399	47214	0.019
GPT-4o	1.72	20815	27800	48615	0.330
CoT(1)	1.57	23681	37119	60800	0.018
CoT(20)	1.67	477220	742380	1219600	0.366
Reflexion	1.88	1126390	389520	1515910	0.455
MapCoder	1.85	1210143	713383	1923526	0.577
MoA	1.76	344316	172973	517289	0.612
		79005	29441	108446	
LessonL	2.16	1009359	387994	1397353	0.326

For FLOPS per token, we take the approach of [27] and estimate it to be $24n_{\text{layer}}d_{\text{model}}^2$, where n_{layer} is the number of layers and d_{model} is the hidden dimension. For the open-source models Deepseek and Qwen, these two hyperparameters come from [18, 23]. On the other hand, the information of the GPT models is unknown. We treat GPT-4o mini an 8B model and GPT-4o a 200B model, following [1]. Then, we estimate n_{layer} and d_{model} by borrowing the information from the LLaMA models of similar sizes [17]. Specifically, we reuse the configuration of LLaMA 8B for GPT-4o mini and that of LLaMA 405B for GPT-4o, except that for the latter case n_{layer} is cut by half. The estimated FLOPS are given in Table 7.

J Case Study: Geometry

Let us study a case about finding the smallest distance among n points in two-dimensional space. The original code performs a naive computation by enumerating all possible pairs and finding the minimum distance, which takes $O(n^2)$ time. The final optimized code reduces the complexity from

Table 7: (Estimated) Model configurations and FLOPS per token.

Model	n_{layer}	d_{model}	FLOPS
DeepSeek7B	32	4096	12884901888
Qwen7B	28	3584	8631877632
Qwen14B	48	5120	30198988800
GPT-4o mini	32	4096	12884901888
GPT-4o	63	16384	405874409472

$O(n^2)$ to $O(n \log n)$ and achieves an outstanding speedup of 74.31x. The algorithm starts by sorting the points based on their x-coordinates and then performs divide-and-conquer. The divide part separates a point set in two equal halves. The conquer part (recursively) calculates the minimum distance for each of the two halves; let us call the smaller of the two distances, d. Then, the combine part filters out points with an x-distance from the middle point greater than d, sorts the remaining on their y-coordinates, ignores point pairs with a y-distance greater than d, and returns the smallest distance among the remaining points.

Qwen14B summarizes the reason of the gain as the following:

> *"Code B runs significantly faster than Code A because it sorts the points by their x-coordinate before comparing them pairwise. This sorting operation reduces the number of comparisons made in the nested loop, which is a time-consuming operation in Code A. By sorting the points first, we can eliminate unnecessary comparisons and thus speed up the execution time."*

We see that it attributes the initial sorting operation to be the main reason of the outstanding result. This reason is partially correct, but sorting is only the precursor step of divide-and-conquer that leads to a reduced $O(n \log n)$ cost. We proceed to investigate all the generated lessons over the course and find that indeed some lessons highlight the importance of divide-and-conquer.

Sorting is mentioned in lessons 1, 2, and 4, while divide-and-conquer is mentioned in lessons 1, 4, and 5. The remaining lessons are all about compilation failures, including forgetting '#pragma once' (lesson 3), no implementation of 'std::set' for 'Point' (lesson 6), mismatch of argument types (lesson 7), undefined functions (lesson 8), and use of function before definition (lesson 9). These lessons reflect the complex implementation of a divide-and-conquer algorithm, which can cause many errors (even for human programmers). Overall, the lessons are helpful to guide an LLM to generate correct and high-performing codes.

The original code, the final code, and the full text of the lessons are provided below. For readability, the codes were edited by removing nonessential comments (such as problem description) and white spaces (including indentation).

Original code

```cpp
#pragma once
#include <vector>
#include <limits>

struct Point {
  double x, y;
};

double distance(Point const& p1, Point const& p2) {
  return std::sqrt(std::pow(p2.x-p1.x, 2) + std::pow(p2.y-p1.y, 2));
}

/* Return the distance between the closest two points in the vector
points.
```

```
    Example:
    input: [{2, 3}, {12, 30}, {40, 50}, {5, 1}, {12, 10}, {3, 4}]
    output: 1.41421
*/
double closestPair(std::vector<Point> const& points) {
  // The polygon needs to have at least two points
  if (points.size() < 2)    {
    return 0;
  }
  auto getDist = [](Point const& a, Point const& b) {
    return std::sqrt(std::pow(b.x-a.x, 2) + std::pow(b.y-a.y, 2));
  };
  double minDist = std::numeric_limits<double>::max();
  for (size_t i = 0; i < points.size() - 1; i++) {
    for (size_t j = i + 1; j < points.size(); j++) {
      const double dist = getDist(points[i], points[j]);
      if (dist < minDist) {
        minDist = dist;
      }
    }
  }
  return minDist;
}
```

Optimized code, speedup 74.31x

```
#include <vector>
#include <algorithm>
#include <cmath>
#include <limits>

struct Point {
  double x, y;
};

double distance(const Point& p1, const Point& p2) {
  return std::sqrt(std::pow(p2.x - p1.x, 2) + std::pow(p2.y - p1.y,
  2));
}

double closestPairHelper(const std::vector<Point>& points, size_t
start, size_t end) {
  if (end - start <= 3) {
    double minDist = std::numeric_limits<double>::max();
    for (size_t i = start; i < end; ++i) {
      for (size_t j = i + 1; j < end; ++j) {
        minDist = std::min(minDist, distance(points[i], points[j]));
      }
    }
    return minDist;
  }

  size_t mid = (start + end) / 2;
  double leftMin = closestPairHelper(points, start, mid);
  double rightMin = closestPairHelper(points, mid, end);
  double minDist = std::min(leftMin, rightMin);

  std::vector<Point> strip;
  for (size_t i = start; i < end; ++i) {
    if (std::abs(points[i].x - points[mid].x) < minDist) {
      strip.push_back(points[i]);
```

```
    }
  }

  std::sort(strip.begin(), strip.end(), [](const Point& a, const
  Point& b) {
    return a.y < b.y;
  });

  for (size_t i = 0; i < strip.size(); ++i) {
    for (size_t j = i + 1; j < strip.size() && strip[j].y - strip[i].y
    < minDist; ++j) {
      minDist = std::min(minDist, distance(strip[i], strip[j]));
    }
  }

  return minDist;
}


double closestPair(const std::vector<Point>& points) {
  if (points.size() < 2) {
    return 0;
  }
  std::vector<Point> sortedPoints = points;
  std::sort(sortedPoints.begin(), sortedPoints.end(), [](const Point&
  a, const Point& b) {
    return a.x < b.x;
  });
  return closestPairHelper(sortedPoints, 0, sortedPoints.size());
}
```

Lesson 1 (selected)

Code B is significantly faster because it sorts the points by their x-coordinate before comparing
each pair of points. This sorting step reduces the number of comparisons made in the nested
loop, which is a time-consuming operation. By sorting the points first, we ensure that
similar pairs of points (those with similar x-coordinates) are closer together in the sorted
array, reducing the number of unnecessary distance calculations. This results in a significant
speedup because it eliminates redundant computations.

Lesson 2 (selected)

Code B is significantly faster because it implements a divide-and-conquer strategy, also
known as the "plane sweep" algorithm, which is more efficient for this type of problem. It
sorts the points based on their x-coordinates and then uses a strip-based approach to find the
closest pair of points. This approach reduces the time complexity from $O(n^2)$ to $O(n \log n)$,
making it much faster for large datasets.

In Code A, the algorithm compares all pairs of points, leading to a time complexity of $O(n^2)$.
This makes it less efficient for larger datasets compared to Code B, which has a better time
complexity.

Additionally, Code B optimizes by only considering points within a certain distance (the
minimum distance found so far) in the strip, reducing unnecessary comparisons. This
optimization is not present in Code A and contributes to its slower performance.

Lesson 3 (selected)

Code B cannot be compiled because it is missing the '#pragma once' directive at the beginning of the file, which is used to prevent multiple inclusions of the same header file. Additionally, there are no syntactic errors in Code B that would cause it to fail compilation based on the provided compiler output.

Lesson 4 (selected)

Code B is significantly faster because it implements a more efficient algorithm for finding the closest pair of points. Specifically, it uses a divide-and-conquer strategy combined with the concept of a "strip" to reduce the time complexity from $O(n^2)$ to $O(n \log n)$.

In Code A, the algorithm compares every pair of points, which leads to a time complexity of $O(n^2)$. This is because there are $n * (n-1)/2$ comparisons in the worst case.

On the other hand, Code B uses a divide-and-conquer approach. It first sorts the points by their x-coordinate and then recursively finds the minimum distance in each half. This step has a time complexity of $O(n \log n)$.

Additionally, Code B also considers pairs of points that lie within a "strip" (a range of y-coordinates) around the vertical line that divides the two halves. This strip is sorted by y-coordinate and then checked for pairs that are closer than the current minimum distance. The stripClosest function has a time complexity of $O(n)$, which is used when there are more than 3 points in the strip.

Lesson 5

Code B is significantly faster because it implements the divide-and-conquer algorithm for finding the closest pair of points, which reduces the time complexity from $O(n^2)$ in Code A to approximately $O(n \log n)$.

Lesson 6

Code B cannot be compiled because it attempts to use a 'std::set' to store 'Point' objects, but 'Point' does not have a defined comparison operator. The compiler needs a way to compare 'Point' objects to insert them into the set.

Lesson 7

Code B cannot be compiled because the 'stripClosest' function is called with a 'std::vector<Point>' and a 'double', but the function expects a 'std::vector<Point>' and an 'int'. The compiler output would indicate a type mismatch error at the call to 'stripClosest'.

Lesson 8

Code B cannot be compiled because the function 'closestPairUtil' is declared but not defined, leading to a linker error. The compiler output would indicate an unresolved external symbol for 'closestPairUtil'.

Lesson 9

Code B cannot be compiled because the 'stripClosest' function is defined after it is used in the 'closestPairUtil' function. In C++, functions must be declared before they are used, or defined inline within the scope where they are called. The compiler error indicates that it does not recognize 'stripClosest' when it is first referenced in 'closestPairUtil'.

K   Case Study: DFT

Let us study a case about performing the discrete Fourier transform (DFT). Mathematically, the DFT of a length-N array x is an array X of the same length, where

$$X_k = \sum_{n=0}^{N-1} x_n \cdot e^{-i2\pi nk/N} = \sum_{n=0}^{N-1} x_n[\cos(-2\pi nk/N) + \sin(-2\pi nk/N)], \quad k = 0, 1, \ldots, N-1. \tag{1}$$

Because of the periodicity of trigonometric functions, many nk values are congruent modulo N and thus many complex exponentials are repeatedly computed. A natural idea of improving the runtime of the DFT is to precompute these exponentials and perform a table lookup whenever needed. Indeed, the optimized code adopts this idea, reflected by lessons 1, 2, 5, 6, and 8, and agreed with by Qwen14B's assessment:

> *"Code B is significantly faster than Code A because it precomputes the complex exponential factors, reducing redundant calculations in the inner loop. This optimization minimizes the computational overhead and accelerates the DFT process."*

It is interesting to note that the above mathematical definition of DFT implies an $O(N^2)$ computational cost, while it is well known that a divide-and-conquer algorithm, called the fast Fourier transform (FFT), can bring down the cost to $O(N \log N)$. Some agents in some rounds do try FFT, as lessons 4, 7, and 9 suggest, but they fail. Two of these lessons say that the optimized code "is nonequivalent to" the original code, while another lesson says there is a logical error. We investigate into the resulting codes and find that they make an error by flipping the sign (note the negative sign inside the trigonometric functions in (1)). In other words, the incorrectness of the codes is caused by a sign error, rather than the nonequivalence between DFT and FFT. While the lessons are not fully correct, they do discourage the agents to try the FFT idea, as the agents are probably not capable of precisely implementing FFT in context.

Another lesson (lesson 3) is incurred from compilation errors. The code attempts to use Intel intrinsic instructions but it fails compilation. Such attempts requires CPU support and the use of the corresponding compiler flags, which are not supported by our evaluation pipeline. This lesson discourages the use of Intel intrinsics in the benchmarking.

The original code, the final code, and the full text of the lessons are provided below. For readability, the codes were edited by removing nonessential comments (such as problem description) and white spaces (including indentation).

Original code

```cpp
#pragma once
#include <vector>
#include <complex>
#include <cmath>

#if !defined(M_PI)
#define M_PI 3.14159265358979323846
#endif
/* Compute the discrete fourier transform of x. Store the result in
output.
   Example:
   input: [1, 4, 9, 16]
   output: [30+0i, -8-12i, -10-0i, -8+12i]
*/
void dft(std::vector<double> const& x,
std::vector<std::complex<double>> &output) {
  int N = x.size();
  output.resize(N, std::complex<double>(0, 0)); // Resize the output
  vector and initialize with 0
```

```
  for (int k = 0; k < N; k++) { // For each output element
    std::complex<double> sum(0, 0);
    for (int n = 0; n < N; n++) { // For each input element
      double angle = 2 * M_PI * n * k / N;
      std::complex<double> c(std::cos(angle), -std::sin(angle)); //
      Euler's formula
      sum += x[n] * c;
    }
    output[k] = sum;
  }
}
```

Optimized code, speedup 10.83x

```
#pragma once
#include <vector>
#include <complex>
#include <cmath>

#if !defined(M_PI)
#define M_PI 3.14159265358979323846
#endif

void dft(std::vector<double> const& x,
std::vector<std::complex<double>> &output) {
  int N = x.size();
  output.resize(N, std::complex<double>(0, 0)); // Resize the output
  vector and initialize with 0

  // Precompute the complex exponential factors
  std::vector<std::complex<double>> W(N);
  for (int k = 0; k < N; k++) {
    double angle = -2 * M_PI * k / N;
    W[k] = std::complex<double>(std::cos(angle), std::sin(angle));
  }

  for (int k = 0; k < N; k++) { // For each output element
    std::complex<double> sum(0, 0);
    for (int n = 0; n < N; n++) { // For each input element
      sum += x[n] * W[(n * k) % N];
    }
    output[k] = sum;
  }
}
```

Lesson 1 (selected)

Code B is significantly faster because it precomputes the complex exponential factors and stores them in a vector, W. This avoids the need to compute the same factor multiple times within the nested loop, which can lead to significant performance improvements due to reduced computation time and increased cache efficiency. In Code B, each factor is computed only once and then reused in the inner loop, while in Code A, each factor is recomputed for every input element.

Lesson 2 (selected)

Code B is significantly faster than Code A because it precomputes the complex exponential factors, reducing redundant calculations in the inner loop. This optimization minimizes the overhead associated with trigonometric function calls and repeated angle calculations, leading to a substantial performance improvement.

Lesson 3 (selected)

The compiler output is not provided, but based on the code, the issue likely lies in the use of '_mm256_set_pd' and '_mm256_mul_pd' with 'double' values. These intrinsics expect '__m256d' (256-bit vector of 4 double-precision floating-point numbers), but the code is passing individual 'double' values. Additionally, the horizontal add operations are incorrectly used to sum up the results.

Lesson 4 (selected)

Code B is nonequivalent to Code A because it uses the Cooley-Tukey FFT algorithm, which is a fast Fourier transform (FFT) method, while Code A implements the direct computation of the Discrete Fourier Transform (DFT). The FFT algorithm in Code B is more efficient for large input sizes, but it does not handle real-valued inputs directly; instead, it processes complex numbers. Code A correctly handles real-valued inputs by converting them to complex numbers with zero imaginary parts before computing the DFT.

Lesson 5

Code B is faster because it precomputes the complex exponential factors (W[k]) and reuses them in the inner loop, avoiding redundant calculations. In Code A, the complex exponential factor is calculated inside the inner loop for each element of x, which results in a higher number of operations and thus slower execution time.

Lesson 6

Code B is significantly faster because it precomputes the complex exponential factors (W[n]) and reuses them in the inner loop, which reduces the number of trigonometric function calls and thus improves performance. In Code A, each iteration of the inner loop computes a new complex exponential, while in Code B, these values are precomputed and stored in the W vector, leading to a more efficient use of computational resources.

Lesson 7

Code B is not equivalent to Code A because it uses the Fast Fourier Transform (FFT) algorithm instead of the direct computation of the Discrete Fourier Transform (DFT). The FFT is an efficient algorithm for computing the DFT, but it requires the input size to be a power of two. If the input size is not a power of two, Code B will still attempt to compute the FFT, which can lead to incorrect results or errors. In contrast, Code A directly computes the DFT using nested loops, which does not have this restriction but may be less efficient for large input sizes.

Lesson 8

Code B is significantly faster than Code A because it precomputes the complex exponential factors, reducing redundant calculations in the inner loop. This optimization minimizes the overhead of trigonometric function calls and repeated angle calculations, leading to a substantial performance improvement.

Lesson 9

Code B contains a logical error in the recursive call to 'dft' within the Cooley-Tukey FFT algorithm. The output vector is being modified in-place, and the recursive calls are overwriting parts of the output vector that are still needed for the computation. This leads to incorrect results and potential crashes.