



BHARATIYA VIDYA BHAVAN'S
SARDAR PATEL INSTITUTE OF TECHNOLOGY
Bhavan's Campus, Munshi Nagar, Andheri (West), Mumbai – 400058-India
Department of Computer Engineering

Name	Balla Mahadev Shrikrishna
UID no.	2023300010
Experiment No.	10

AIM:	To implement a vertex cover approximation algorithm.
Program 1	
PROBLEM STATEMENT :	<p>Details – An Approximate Algorithm is a way of approaching NP-COMPLETENESS for the optimization problem. This technique does not guarantee the best solution. The goal of an approximation algorithm is to come as close as possible to the optimum value in a reasonable amount of time which is at the most polynomial time. Such algorithms are called approximation algorithms or heuristic algorithms. An algorithm for a problem has an approximation ratio of $\rho(n)$ if for any input of size n the cost C of the solution produced by the algorithm is within a factor of $\rho(n)$ of the cost C^* of an optimal solution.</p> $\text{Max}(C/C^*, C^*/C) \leq \rho(n)$ <p>We have an algorithm $\rho(n)$-approximation if it achieves an approximation ratio of $\rho(n)$</p> <p>For maximization problem $0 < C \leq C^*$</p> <p>For minimization problem $0 < C^* \leq C$</p> <p>For the vertex cover problem, the optimization problem is to find the vertex cover with fewest vertices, and the approximation problem is to find the vertex cover with few vertices. Formally, a vertex Cover of a graph G is a set of vertices such that each edge in G is incident to at least one of these vertices. The vertex-cover problem was proven by the NPC. Now, we want to solve the optimal version of the vertex cover problem, i.e., we want to find a minimum size vertex cover of a given graph. We call such vertex cover an optimal vertex cover C^*.</p> <p>Input – Graph $G(V,E)$ of at least 10 vertices.</p> <p>Output – Approximate Vertex Cover for the given graph.</p> <p>Submission –</p> <ol style="list-style-type: none">1) C/C++ source code of implementation2) Verified output for the written source code with multiple inputs3) One page report of Exp. 10



BHARATIYA VIDYA BHAVAN'S
SARDAR PATEL INSTITUTE OF TECHNOLOGY
Bhavan's Campus, Munshi Nagar, Andheri (West), Mumbai – 400058-India
Department of Computer Engineering

ALGORITHM:	<ol style="list-style-type: none">1) Input: An undirected graph with V vertices and E edges.2) Initialization: Mark all vertices as not in cover (inCover array).3) Greedy Selection:<ol style="list-style-type: none">a) While there are uncovered edges, pick an arbitrary edge (u, v) where neither u nor v is in the cover.b) Add both u and v to the cover.4) Termination: When all edges are covered, output the selected vertices.
PROGRAM:	<pre>#include <stdio.h> #include <stdlib.h> #include <stdbool.h> #define MAX_VERTICES 100 int main() { int V, E; int graph[MAX_VERTICES][MAX_VERTICES] = {0}; bool visited[MAX_VERTICES] = {false}; bool inCover[MAX_VERTICES] = {false}; // Get input printf("Enter the number of vertices: "); scanf("%d", &V); printf("Enter the number of edges: "); scanf("%d", &E); printf("Enter %d edges (format: u v):\n", E); for (int i = 0; i < E; i++) { int u, v; scanf("%d %d", &u, &v); // Build undirected graph graph[u][v] = 1; graph[v][u] = 1; } }</pre>



BHARATIYA VIDYA BHAVAN'S
SARDAR PATEL INSTITUTE OF TECHNOLOGY
Bhavan's Campus, Munshi Nagar, Andheri (West), Mumbai – 400058-India
Department of Computer Engineering

```
// Greedy approach: pick an uncovered edge and add both endpoints
to cover
bool edgesRemain = true;
while (edgesRemain) {
    edgesRemain = false;

    // Find an edge that is not covered
    for (int u = 0; u < V && !edgesRemain; u++) {
        for (int v = 0; v < V && !edgesRemain; v++) {
            if (graph[u][v] && !inCover[u] && !inCover[v]) {
                // Add both endpoints to cover
                inCover[u] = true;
                inCover[v] = true;
                edgesRemain = true;
            }
        }
    }
}

// Print result
printf("Marked vertices: ");
for (int i = 0; i < V; i++) {
    if (inCover[i]) {
        printf("%d ", i);
    }
}
printf("\n");

return 0;
}
```



**BHARATIYA VIDYA BHAVAN'S
SARDAR PATEL INSTITUTE OF TECHNOLOGY**

Bhavan's Campus, Munshi Nagar, Andheri (West), Mumbai – 400058-India

Department of Computer Engineering

```
• PS C:\Mahadev\SE\Sem4\DAA\Lab\Lab Sessions\Exp10> gcc vertex_cover.c
• PS C:\Mahadev\SE\Sem4\DAA\Lab\Lab Sessions\Exp10> ./a.exe
Enter the number of vertices: 7
Enter the number of edges: 7
Enter 7 edges (format: u v):
0 1
0 3
1 2
1 4
2 5
4 5
2 6
Marked vertices: 0 1 2 5
• PS C:\Mahadev\SE\Sem4\DAA\Lab\Lab Sessions\Exp10> ./a.exe
Enter the number of vertices: 10
Enter the number of edges: 12
Enter 12 edges (format: u v):
0 1
0 2
3 4
4 5
2 3
3 4
4 6
3 5
2 5
1 4
2 4
5 6
Marked vertices: 0 1 2 3 4 5
• PS C:\Mahadev\SE\Sem4\DAA\Lab\Lab Sessions\Exp10> █
```

RESULT:



BHARATIYA VIDYA BHAVAN'S
SARDAR PATEL INSTITUTE OF TECHNOLOGY
Bhavan's Campus, Munshi Nagar, Andheri (West), Mumbai – 400058-India
Department of Computer Engineering

CONCLUSION:

Name: Balla Mahadev Shrikrishna
UID: 2023300010
Div: A
Batch: A

Exp-10

- * The vertex cover problem involves selecting the smallest set of vertices in a graph such that every edge is incident to at least one vertex in the set.
- * Implementation
 - Graph Representation: Adjacency Matrix
 - Cover Tracking: Boolean array to mark included vertices.
 - Edge Check: Nested loops iterate until no uncovered edges remain.
- * Time Complexity & Space Complexity
 - T.C = $O(V^2)$, i.e. for each edge check all possible vertex pairs.
 - S.C = $O(V^2) + O(V)$ for adj. matrix & visited/cover arrays.
- * Although o/p is not wanted to be optimal, the solⁿ is within a factor of 2 of the minimum, making it practically applicable for real-world problems.