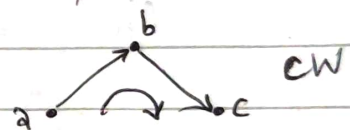


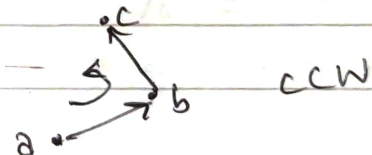
* Orientation

Three kinds of orientation for three pts.:

- Clockwise (CW) : right turn



- Counter CW : left turn



- Collinear (COLL) : no turn



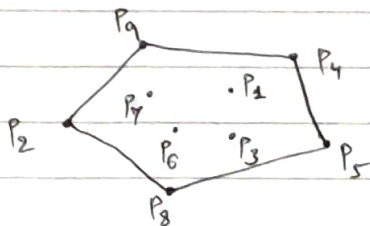
Orientatⁿ can be characterized by sign of $\Delta(a, b, c)$

$$\begin{aligned} \text{If } \Delta(a, b, c) < 0 &\Rightarrow \text{CW} \\ \Delta(a, b, c) = 0 &\Rightarrow \text{COLL} \\ \Delta(a, b, c) \geq 0 &\Rightarrow \text{CCW} \end{aligned}$$

$$\Delta(a, b, c) = \begin{vmatrix} x_a & x_b & x_c \\ y_a & y_b & y_c \\ 1 & 1 & 1 \end{vmatrix}$$

* Convex Hull (CH)

- CH of a pt. set P , $\text{CH}(P)$:
 → smallest convex set containing P
 → intersectⁿ of all convex sets containing P
- Give an algo. that computes convex hull of any given set of n pts. in the plane efficiently.
 i/p : locatⁿ of n pts.
 o/p : a convex polygon \Rightarrow a sorted sequence of pts. CW or CCW along the boundary.
- i/p = $P_1, P_2, P_3, \dots, P_n$



$$\text{o/p} = P_2, P_9, P_4, P_5, P_8$$

* Graham's Scan Algo.

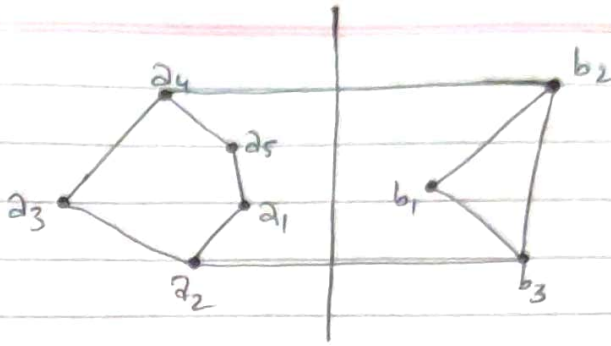
- Step 1: Find lowest pt p (guaranteed to be in)
- Step 2: Sort pts. arnd. p (in polar angle) in increasing order of angle.
- Step 3: Walk arnd. to remove concave angle.
(keep pts. with left turns & drop those with right turns)
- Step 4: Working arnd all pts. by sorting order.

* Divide & Conquer Approach

- Sort pts by x-coordinates
- For i/p set S ,
 - divide into left half A & right half B using x-coords.
 - Compute $CH(A)$
 - " $CH(B)$
 - Combine
- Solve smaller problem using brute force.

How to combine? Finding tangents

Brute force: generate pairwise (from A to B) segments & check



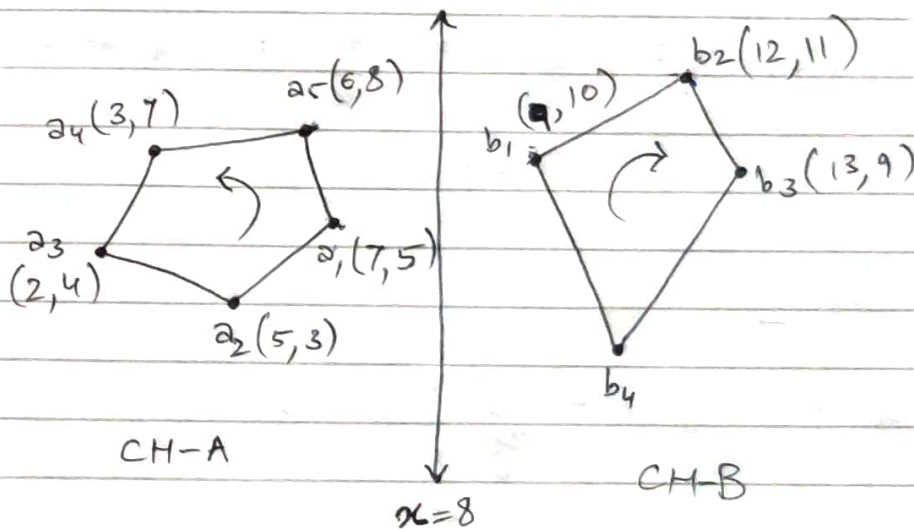
a_4, b_2 is called the upper tangent; $\text{Max } y(i, j)$
 a_2, b_3 is called the lower tangent; $\text{Min } y(i, j)$
 Complexity: $O(n^2)$

Picking max y for both isn't enough.

Need to maximise $y(i, j)$ for upper tangent & minimise in case of lower tangent.

Two Finger Algo.:

- Start with line seg. using rightmost of a & leftmost of b
- Move cw for b → update if it improves $y(i, j)$.
- Move ccw for a → update if it improves $y(i, j)$.
- Repeat until $y(i, j)$ converges



$a_1 b_1$:

$$m = \frac{10 - 5}{9 - 7} = \frac{5}{2}$$

$$c = y_1 - mx_1 = 5 - 2.5(7) \\ = -12.5 \\ j = 7.5$$

$$a_1, b_2 : (y - 5) = \frac{6}{5} (x - 7)$$

$$y - 5 = \frac{6}{5} \Rightarrow y = 6.2$$

$$a_5(6, 8) \quad b_1(9, 10) : y - 8 = \frac{2}{3} (x - 6)$$

$$y = 8 + \frac{4}{3} = 9.33$$

$$y_{\max}(i, j) = y(5, 1) = 9.33 \Rightarrow \text{Set } a_5, b_1$$

$$a_5(6, 8) \quad b_2(12, 11) : y - 8 = 0.5(x - 6)$$

$$y = 9$$

Set a_5, b_1

$$a_4, b_1 : y - 7 = \frac{2.5}{8} (x - 3) \\ = 2.5$$

$$y = 9.5$$

Set a_4, b_1

$$a_4(3,7) \quad b_2(12,11) : y-7 = \frac{4}{9}(x-3)$$

$$y = \frac{20}{9} + 7$$

$$= \frac{83}{9} = 9.2...$$

set $a_4 b_1$

$$a_3(2,4) \quad b_1(9,10) : y-4 = \frac{6}{7}(x-2)$$

$$y = \frac{36+28}{7}$$

$$= \frac{64}{7} = 9.1...$$

set $a_4 b_1$

$a_4 b_2 :$

\Rightarrow conversion

\Rightarrow tangent : $a_4 b_1$

Exp. 2B

Convex Hull

Set of pts. : $\{(1,1), (2,4), (3,2), (5,3), (4,1)\}$

$$\text{Orientat}^n(\text{pt } a, \text{pt } b, \text{pt } c) = \begin{vmatrix} a.x & b.x & c.x \\ a.y & b.y & c.y \\ 1 & 1 & 1 \end{vmatrix}$$

1) Brute

For line $(1,1) \rightarrow (2,4)$:
or $((1,1), (2,4), (3,2)) = -1 \Rightarrow \text{CCW}$
or $((1,1), (2,4), (5,3)) = -1 \Rightarrow \text{CCW}$
or $((1,1), (2,4), (4,1)) = 1 \Rightarrow \text{CW}$

\therefore All pts. don't lie on the same side of the line, the line seg. can't be part of the hull.

2) Graham's Scan

Pt. with lowest y-coord : $(1,1)$ $\{ \because x\text{-coord of pt. } (1,1) \text{ is less than that of pt. } (4,1) \}$

Now, we need to sort all the pts. by polar angle w.r.t. anchor pt.
 \Rightarrow Sorted pts.: $[(1,1), (4,1), (3,2), (5,3), (2,4)]$

Build Hull: start with $\{(1,1)\}$

Add $(4,1) \Rightarrow \text{hull} = \{(1,1), (4,1)\}$

" $(3,2) \Rightarrow \text{check orientation } ((1,1), (4,1), (3,2)) = -1 \Rightarrow \text{CCW}$

\Rightarrow Part of the hull $\Rightarrow \text{hull: } \{(1,1), (4,1), (3,2)\}$

Add $(5,3) \Rightarrow \text{check or } ((4,1), (3,2), (5,3)) = -1 \Rightarrow \text{add to hull}$

Add $(2,4) \Rightarrow \text{check or } ((3,2), (5,3), (2,4)) = -1 \Rightarrow \text{CCW}$

But check or $((4,1), (3,2), (2,4)) = 1 \Rightarrow \text{CW}$

\therefore pop $(3,2) \Rightarrow \text{check or } ((1,1), (4,1), (2,4)) = -1$

so, add to hull \Rightarrow final hull: $\{(1,1), (4,1), (5,3), (2,4)\}$

3) Divide & Conquer

Divide { Left = $[(1,1), (2,4)]$
 Right = $[(3,2), (4,1), (5,3)]$

Conquer { Compute left hull \Rightarrow \because only two pts, hull = $\{(1,1), (2,4)\}$
 " right " \Rightarrow for 3 pts. we check their orientations.
 or $((3,2), (4,1), (5,3)) = -1$
 or $((5,3), (3,2), (4,1)) = -1$
 or $((4,1), (5,3), (3,2)) = -1$ } All CCW \Rightarrow hull = $\{(3,2), (4,1), (5,3)\}$

Merge: Find rightmost pt. in left hull : $(2,4)$
 Find leftmost " " right " : $(3,2)$
 " upper tangent : $(2,4) \rightarrow (5,3)$
 " lower " : $(1,1) \rightarrow (4,1)$
 Merged hull = $\{(1,1), (4,1), (5,3), (2,4)\}$