

# LET'S START WITH DBMS :).

## Lossy and Lossless decomposition

### Lossy Decomposition

In normalisation we generally break/decompose the table into 2 or more tables.

Consider there is a relation R,

**Lossy decomposition** occurs when a relation R is decomposed or broken into two or more relations, but data is lost, and the original relation R can't be reconstructed by joining these decomposed relations.

There are 2 rules for a decomposition to be lossy

1. Some data from the original relation R is lost after decomposition
2. Join of the decomposed relations( $R_1, R_2..R_n$ ) is not equal to the original relation R

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## Lossy and Lossless decomposition

### Lossy decomposition

There is a relation  $R(A,B,C)$

A	B	C
1	2	3
4	2	6

R

Step 1: Lets decompose the relation based on any attribute and keep that attribute as common, for now lets use B as common attribute

Decomposed relations:  $R1(A,B)$  and  $R2(B,C)$

A	B
1	2
4	2

R1

B	C
2	3
2	6

R2

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## Lossy and Lossless decomposition

### Lossy decomposition

Step 2: Lets perform a natural join between R1 and R2

When we do R1 natural join R2 we won't get the original relation back(lossy)

A	B	C
1	2	3
4	5	6

R

A	B	C
1	2	3
1	2	6
4	2	3
4	2	6

We can see some additional tuples that were not in the original relation R (lossy decomposition)

R1 natural join R2

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## Lossy and Lossless decomposition

### Lossy decomposition

How to ensure a decomposition is lossless

1. Divide or decompose the table on basis of CK or SK present in the relation so that there is no duplicacy
2. For a decomposition to be lossless
  - a.  $R1 \cup R2 = R$
  - b.  $R1 \cap R2 = \text{common attribute}$
3. To ensure that a decomposition is lossless, a common approach is to use the dependency preservation property

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## Lossy and Lossless decomposition

### Lossless Decomposition

In normalisation we generally break/decompose the table into 2 or more tables.

Consider there is a relation R,

**Lossless decomposition** ensures that when a relation R is decomposed/breaked into two or more relations, no data is lost, and the original relation R can be again reconstructed by joining these decomposed relations.

There are 2 rules for a decomposition to be lossless

1. All data in the original relation R should be preserved after decomposition
2. Join of the decomposed relations  $(R_1, R_2..R_n) = \text{original relation } R$

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## Lossy and Lossless decomposition

### Lossless Decomposition

So if the table is decomposed and we want to query the attributes present in both the tables we will use the join operation.

#### Natural Join :

- The natural join operation combines tuples(rows) from two relations based on common attributes.
- It only includes those combinations of tuples that have the same values for the common attributes.

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## Lossy and Lossless decomposition

### Lossless decomposition

There is a relation  $R(A,B,C)$  with CK as A

A	B	C
1	2	3
4	5	6

R

Step 1: Lets decompose the relation based on the CK or SK of the given R

Decomposed relations:  $R1(A,B)$  and  $R2(A,C)$

A	B
1	2
4	5

R1

A	C
1	3
4	6

R2

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## Lossy and Lossless decomposition

### Lossless decomposition

There is a relation  $R(A,B,C)$  with CK as A

A	B	C
1	2	3
4	5	6

R

Step 2: Lets perform a natural join between R1 and R2

When we do R1 natural join R2 we get the original relation back(lossless)

A	B	C
1	2	3
4	5	6

R1 natural join R2