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Equivalence of Functional Dependencies

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Functional Dependency: A functional dependency $X \rightarrow Y$ holds on a relation R if, for any two tuples t1 and t2 in R, whenever $t1[X]=t2[X]$, then $t1[Y]=t2[Y]$

Two sets of functional dependencies, F and G, are equivalent if the following conditions hold:

1. F implies G : Every functional dependency in G can be derived from F.

2. G implies F: Every functional dependency in F can be derived from G.

If both conditions are true, then we say that F and G are equivalent.

$$\begin{aligned} F &\supseteq G \\ G &\supseteq F \end{aligned}$$

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Equivalence of Functional Dependencies

How to find if two FD are equivalent

Step 1: Compute the Closure of both the sets

Step 2: Ensure that every functional dependency in set1 is in set2 closure

Step 3: Ensure that every functional dependency in set2 is in set1 closure

Step 4: If both subset checks pass, then set1 and set2 are equivalent.

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Step 1: Compute the Closure of both the sets

Step 2: Ensure that every functional dependency in set1 is in set2 closure

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Step 4: If both subset checks pass, then set1 and set2 are equivalent.

Equivalence of Functional Dependencies

Consider two sets of functional dependencies: $F = \{A \rightarrow B, B \rightarrow C\}$, $G = \{A \rightarrow C, A \rightarrow B\}$
Check if F and G are equivalent

$F = \{A \rightarrow B, B \rightarrow C\}$

Closure of F, attributes $\rightarrow A, B, C$

$A^+ = \{A, B, C\}$

$B^+ = \{B, C\}$

$C^+ = \{C\}$

$G = \{A \rightarrow C, A \rightarrow B\}$

Closure of G, attributes $\rightarrow A, B, C$

$A^+ = \{A, C, B\}$

$B^+ = \{B\}$

$C^+ = \{C\}$

F and G are not equivalent because $B \rightarrow C$ is not implied by G