

LET'S START WITH DBMS :).

Attribute closure

5. Now lets find candidate key, super key , prime and non-prime attributes

Super key : Set of attributes whose closure contains all the attributes given in a relation
Super set of any candidate key is super key. A key(combination of all possible attributes) which can uniquely identify two tuples.

How to find Super Key?

- Identify All Attributes
- 2. Analyze the functional dependencies and find closure.
- Generate Power Set : If A has n attributes, the power set will have 2^n subsets.
- Check for Super Key Property : For each subset in the power set, check if it can uniquely identify each tuple in the relation.

LET'S START WITH DBMS :).

Attribute closure

Q. Find all the superkeys for the relation $R(A, B, C)$ and FD : $A \rightarrow B, B \rightarrow C$.

1. Identify All Attributes:

$A = \{A, B, C\}$

2. Analyze the functional dependencies and find closure.

From $A \rightarrow B$ and $B \rightarrow C$, we can infer $A \rightarrow C$

$A^+ = \{A, B, C\}$

LET'S START WITH DBMS :).

Attribute closure

$B^+ = \{B, C\}$ because $B \rightarrow C$, but B does not determine A .

$C^+ = \{C\}$

Closure of A gives or determines all the attributes in the table so we can say it's a super key.

3. How to find all the super keys?

Find the power subset

The power set of $\{A, B, C\}$ is $2^3=8$

$\{\}, \{A\}, \{B\}, \{C\}, \{A, B\}, \{A, C\}, \{B, C\}, \{A, B, C\}$

LET'S START WITH DBMS :).

Attribute closure

A power set is the set of all subsets of a given set, including the empty set and the set itself. If you have a set X , the power set of X is denoted as 2^X

Example:

Let's take a simple set $X=\{a,b\}$

The power set of X would include the following subsets:

1. The empty set: \emptyset
2. The single-element subsets: $\{a\}$ & $\{b\}$
3. The full set itself: $\{a,b\}$

So, the power set $P(X)$ would be: $P(X)=\{\emptyset,\{a\},\{b\},\{a,b\}\}$

LET'S START WITH DBMS :).

Attribute closure

4. Verify each subset to see if it is a super key

1. $\{\}$: Not a super key.
2. $\{A\}$: Super key (as its closure determines all the attributes in relation).
3. $\{B\}$: Not a super key (does not determine A).
4. $\{C\}$: Not a super key (does not determine A or B).
5. $\{A, B\}$: Super key (A is already a super key, so adding B still keeps it a super key).
6. $\{A, C\}$: Super key (A is already a super key, so adding C still keeps it a super key).
7. $\{B, C\}$: Not a super key (B determines C, but does not determine A).
8. $\{A, B, C\}$: Super key (A is already a super key, so adding B and C keeps it a super key).

LET'S START WITH DBMS :).

Attribute closure

5. Super keys are : $\{A\}, \{A, B\}, \{A, C\}, \{A, B, C\}$

We can also say to find max number of super keys we can use the formula

where

$$2^{n-k}$$

k- candidate key with k attributes ($k < n$) in a relation

n- total no of attributes in a relation

When 1 C.K is there = 2^{n-1}

When 2 C.K is there = 2^{n-1} (for 1st ck) + 2^{n-1} (for 2nd ck) - 2^{n-2} (for combination of both)