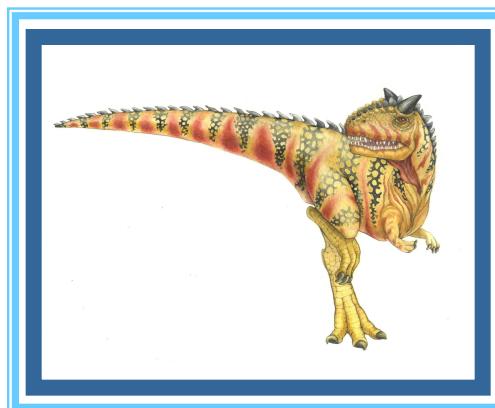


# Chapter 5: Process Synchronization

## Lecture # 14





# Previous Lecture

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- Peterson's Solution
- Synchronization Hardware
- Mutex Locks





# Semaphore

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- Synchronization tool that provides more sophisticated ways (than Mutex locks) for process to synchronize their activities.
- Semaphore **S** – integer variable
- Can only be accessed via two indivisible (atomic) operations
  - **wait()** and **signal()**
    - 4 Originally called **P()** and **V()**

- Definition of the **wait()** operation

```
wait(S) {  
    while (S <= 0)  
        ; // busy wait  
    S--;  
}
```

- Definition of the **signal()** operation

```
signal(S) {  
    S++;  
}
```





# Semaphore

```
wait(S) {  
    while (S <= 0)  
        ; // busy wait  
    S--;  
}
```

```
signal(S) {  
    S++;  
}
```

```
do {  
Entry section  wait(S)  
//critical section  
Exit section  Signal (S)  
//remainder section  
} while (true);
```

S = 1





# Classical Problems of Synchronization

- Classical problems used to test newly-proposed synchronization schemes
  - Bounded-Buffer Problem
  - Readers and Writers Problem
  - Dining-Philosophers Problem





# Classical Problems of Synchronization

- Classical problems used to test newly-proposed synchronization schemes
  - **Bounded-Buffer Problem**
  - Readers and Writers Problem
  - Dining-Philosophers Problem





# Bounded-Buffer Problem

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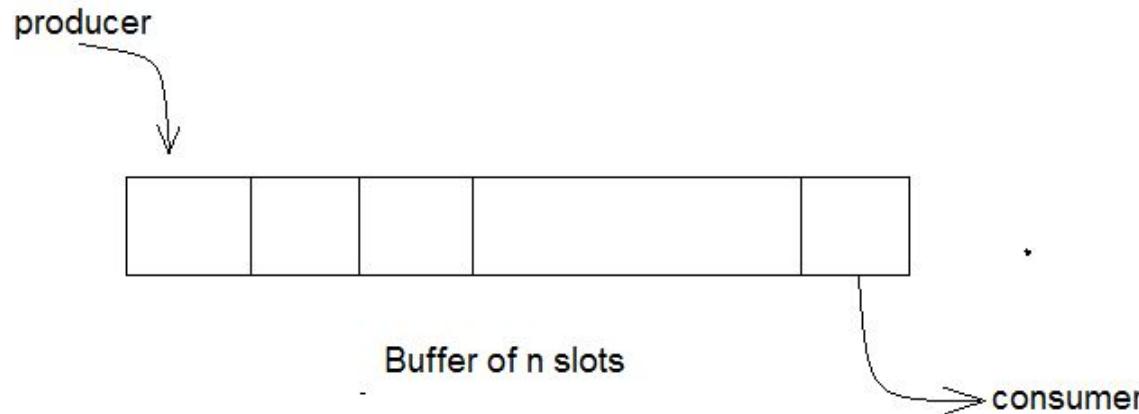
- $n$  buffers, each can hold one item
- Semaphore **S** initialized to the value 1
- Semaphore **full** initialized to the value 0
- Semaphore **empty** initialized to the value  $n$





# Bounded Buffer Problem

- ❑ There is a buffer of **n** slots and each slot can store one unit of data.
- ❑ There are two processes running, namely,
  - ❑ **producer** and **consumer**, which are operating on the buffer.



- ❑ A producer tries to insert data into an empty slot of the buffer. A consumer tries to remove data from a filled slot in the buffer. As you might have guessed by now, those two processes won't produce the expected output if they are being executed concurrently.
- ❑ There needs to be a way to make the producer and consumer work in an independent manner





# Bounded Buffer Problem (Cont.)

- The structure of the producer process

```
void producer
```

```
{
```

```
while(true)
```

```
{
```

```
    Produce(); /* produce an item in next_produced */
```

```
    wait(E);
```

```
    wait(S);
```

```
    append();
```

```
    ...
```

```
    /* add next produced to the buffer */
```

```
    ...
```

```
    signal(S);
```

```
    signal(F);
```

```
}
```

```
}
```

S	E	F
1	0	5

1	2	3	4	5
a	b	C	D	e

```
wait(S) {  
    while (S <= 0)  
        ; // busy wait  
    S--;  
}  
  
signal(S) {  
    S++;  
}
```



# Bounded Buffer Problem (Cont.)

- The structure of the consumer process

```
void consumer
{
    while(true)
    {
        wait(F);
        wait(S);
        take();
        ...
        /* consume the item in next consumed */
        ...
    }
}
```



S	E	F
1	5	0

```
wait(S) {
    while (S <= 0)
        ; // busy wait
    S--;
}

signal(S) {
    S++;
}
```

