## MTL103: Optimization Methods and Applications

*Major*(2023)

Time: 2 hrs, Total marks: 40.

**Read the Honour Code:** "As a student of IIT Delhi, I will not give or receive aid in examinations. I will do my share and take an active part in seeing to it that others as well as myself uphold the spirit and letter of the Honour Code."

## Instructions

- Write your answers neatly and to the point.
- Remember that you will be graded on what you write and not what you intend to write.

## **Questions:**

1. Consider the following graph G where the capacity of each edge is given.

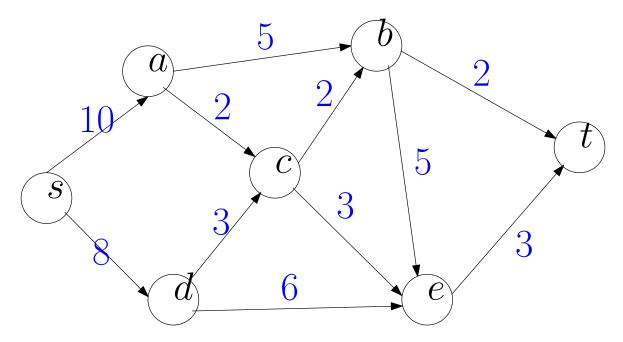


Figure 1: Directed Graph G

- a. Write down the incidence matrix of this graph.
- b. Write down the LP formulation for finding the maximum flow from s to t.

(3+5=8)

2. Suppose we have a polynomial time subroutine **EM** that takes an ellipsoid E = E(z, D) in  $\mathbb{R}^n$  and a halfspace  $H = \{x \in \mathbb{R}^n | a^T x \ge a^T z\}$  (where a is a non-zero vector) as input; and returns an ellipsoid  $E' = E(\overline{z}, \overline{D})$  satisfying the following:

- $\bullet \ \overline{z} = z + \frac{1}{n+1} \frac{Da}{\sqrt{a^T Da}},$
- $\overline{D} = \frac{n^2}{n^2 1} (D \frac{2}{n+1} \frac{Daa^T D}{a^T Da})$  is symmetric positive definite,
- $E \cap H \subset E'$ ,
- $Vol(E') < e^{-1/(2(n+1))} Vol(E)$ .

Consider the following LP:

Maximize  $c^T x$ Subject to  $Ax \leq b$ ,

where A, b have integer entries with magnitude bounded by some U. Using **EM** as a subroutine, propose an algorithm to find out the optimum solution of the given LP. What is the time complexity of your algorithm?

(10)

- 3. For a given set of points  $x_0, x_1, x_m \in \mathbb{R}^n$ , the Voronoi cell  $V(x_i)$  associated with the point  $x_i$  is defined as follows:  $V(x_i) = \{x \in \mathbb{R}^n | d(x_i, x) \leq d(x_j, x) \forall j \neq i\}$ . Here, d(x, y) is the Euclidean distance between the points x and y. Prove that the set  $V(x_i)$  is a convex set.
- 4. Let  $A = [a_{ij}]$  be a pay-off matrix in a two-person zero-sum game. Let us assume that two elements  $a_{ij}$  and  $a_{hk}$  are saddle points. Prove or disprove that  $a_{ik}$  and  $a_{hj}$  are also saddle points. (6)
- 5. Consider the following LP for some fixed  $0 < \epsilon < 1/2$ .

Minimize 
$$x_3$$
  
Subject to  $x_1 - r_1 = \epsilon$ ,  
 $x_1 + s_1 = 1$ ,  
 $x_j - \epsilon x_{j-1} - r_j = 0$ ,  $j = 2, 3$   
 $x_j + \epsilon x_{j-1} + s_j = 1$ ,  $j = 2, 3$ ,  
 $x_j, r_j, s_j \ge 0$ ,  $j = 1, 2, 3$ .

- (a). How many BFS can it have? Are all of them non-degenerate? Justify your answer.
- (b) Let  $N \in \mathbb{Z}$  be the answer to the part (a). Enumerate these BFS in an order  $S_1, S_2, \ldots S_i, \ldots, S_N$  such that the cost of the objective function is strictly decreasing and for each  $i \in [N-1]$ , the basic feasible solutions  $S_i$  and  $S_{i+1}$  are adjacent BFS. Here each BFS  $S_i$  is a vector of the form  $(x_1, x_2, x_3, r_1, r_2, r_3, s_1, s_2, s_3)$ . You need to specify their values to distinguish one from the other. (4+8=12)