

Week 2 Report: Algorithmic Complexity and Dynamic Programming

Introduction

This report summarizes the work completed in Week 2 of Algorithm Design (DTE-3611). The main focus was on exploring computational complexity—particularly NP-hard and NP-complete problems—and applying dynamic programming (DP) as a practical technique to address such challenges. Implementations of the Subset Sum, 0/1 Knapsack, and Traveling Salesman Problem (TSP) were developed in C++, and their performance was benchmarked to compare theoretical complexity with real execution results.

Algorithms Implemented

Three main algorithms were implemented and tested: Subset Sum, 0/1 Knapsack, and TSP using the Held–Karp method. Each demonstrates the power of dynamic programming in reducing exponential complexity to manageable computational time for small to medium input sizes.

Subset Sum Problem

Problem: Given a set of integers and a target sum, determine whether any subset of the numbers sums exactly to that target.

Approach: A bottom-up DP table $dp[i][j]$ was constructed representing whether a subset of the first i elements can form a sum of j . Tested using $\{3, 34, 4, 12, 5, 2\}$ with target = 9.

Knapsack Problem

Problem: Given a set of items with weights and values, determine the maximum value achievable within a knapsack of capacity W .

Approach: DP table storing the optimal value for each sub-capacity and sub-item combination.

Benchmark: Tested with $values = \{60, 100, 120\}$, $weights = \{10, 20, 30\}$, and $W = 50$.

Traveling Salesman Problem (TSP)

Problem: Find the shortest route visiting all cities exactly once and returning to the start.

Approach: Implemented the Held–Karp dynamic programming algorithm using bitmasking. Tested on a 5-city distance matrix.

NP-Hardness and NP-Completeness

Subset Sum is NP-complete (via reduction from 3-SAT). Knapsack is NP-hard in optimization form and NP-complete in decision form. TSP is NP-hard, with its decision variant NP-complete. All algorithms exemplify classical NP-hard problem-solving using DP.

Algorithm	Input Description	Approach	Time Complexity	Observed Time (μs)	Output
Subset Sum	Nums = {3, 34, 4, 12, 5, 2}, target = 9	DP	$O(n \times S)$	13μs	YES
0/1 Knapsack	Values = {60, 100, 120}, weights = {10, 20, 30}, capacity = 50	DP	$O(n \times W)$	13μs	Max = 220
TSP (Held–Karp)	5-city matrix	DP (bitmasking)	$O(n^2 \times 2^n)$	23μs	Cost = 95

Conclusion

Dynamic programming proved highly effective for Subset Sum and Knapsack, drastically reducing runtime compared to brute-force. The Held–Karp TSP algorithm, while optimal, revealed DP's exponential scaling limitations. Overall, the work bridged theoretical and practical understanding of algorithmic complexity and NP-hardness.

Repository

All source code and benchmark scripts are available at:

<https://github.com/Mahid-D/DTE-3611-Week2.git>