

University of New Brunswick
Faculty of Computer Science
CS2333: Computability and Formal Languages
Homework Assignment 5, Due Time, Date 5:00 PM, March 4, 2022

Student Name: _____ Matriculation Number: _____

Instructor: Rongxing Lu

The marking scheme is shown in the left margin and [100] constitutes full marks.

- [10] 1. Verify whether the Grammar $S \rightarrow 0B|1A$, $A \rightarrow 0|0S|1AA|\varepsilon$, $B \rightarrow 1|1S|0BB$ generates the strings i) 001101011, and ii) 1100101001.
- [20] 2. Give context-free grammars that generate the following languages. In all parts the alphabet Σ is $\{0, 1\}$.
- [5] (a) $\{w | w \text{ contains at least three 1s}\}$.
- [5] (b) $\{w | w \text{ starts and ends with the same symbol}\}$.
- [5] (c) $\{w | \text{the length of } w \text{ is odd}\}$.
- [5] (d) $\{w | \text{the length of } w \text{ is odd and its middle symbol is a 0}\}$.
- [15] 3. Find a reduced grammar equivalent to the grammar G , having the production rule $P : S \rightarrow AC|B$, $A \rightarrow a$, $C \rightarrow c|BC$, $E \rightarrow aA|e$, $F \rightarrow bB|e$.
- [15] 4. Remove the unit productions from the grammar whose production rule is given by $P : S \rightarrow AB$, $A \rightarrow a$, $B \rightarrow C|b$, $C \rightarrow D$, $D \rightarrow E$, $E \rightarrow e$.
- [15] 5. Remove the Null productions from the following grammar

$$S \rightarrow XYXZ, X \rightarrow aX|\varepsilon, Y \rightarrow bY|\varepsilon, Z \rightarrow c$$

- [25] 6. Convert the following context-free grammar to the Chomsky normal form.

$$P : S \rightarrow ASA|1B, A \rightarrow B|S, B \rightarrow 0|\varepsilon$$

Solutions.

1. Verify whether the Grammar $S \rightarrow 0B|1A$, $A \rightarrow 0|0S|1AA|\varepsilon$, $B \rightarrow 1|1S|0BB$ generates the strings i) 001101011, and ii) 1100101001.

i)

S	→	0B	($S \rightarrow 0B$)
	→	00BB	($B \rightarrow 0BB$)
	→	001B	($B \rightarrow 1$)
	→	0011S	($B \rightarrow 1S$)
	→	00110B	($S \rightarrow 0B$)
	→	001101S	($B \rightarrow 1S$)
	→	0011010B	($S \rightarrow 0B$)
	→	00110101S	($B \rightarrow 1S$)
	→	001101011A	($S \rightarrow 1A$)
	→	001101011 ε	($A \rightarrow \varepsilon$)
	→	001101011	Yes!

ii)

S	→	1A	($S \rightarrow 1A$)
	→	11AA	($A \rightarrow 1AA$)
	→	110A	($A \rightarrow 0$)
	→	1100S	($A \rightarrow 0S$)
	→	11001A	($S \rightarrow 1A$)
	→	110010S	($A \rightarrow 0S$)
	→	1100101A	($S \rightarrow 1A$)
	→	11001010S	($A \rightarrow 0S$)
	→	110010100B	($S \rightarrow 0B$)
	→	1100101001	($B \rightarrow 1$) Yes!

2. Give context-free grammars that generate the following languages. In all parts the alphabet Σ is $\{0, 1\}$.

- (a) $\{w | w \text{ contains at least three 1s}\}$.

$$S \rightarrow R1R1R1R, \quad R \rightarrow 0R|1R|\varepsilon$$

- (b) $\{w | w \text{ starts and ends with the same symbol}\}$.

$$S \rightarrow 0R0|1R1|\varepsilon, \quad R \rightarrow 0R|1R|\varepsilon$$

- (c) $\{w | \text{the length of } w \text{ is odd}\}$.

$$S \rightarrow 0|1|00S|01S|10S|11S$$

- (d) $\{w | \text{the length of } w \text{ is odd and its middle symbol is a 0}\}$.

$$S \rightarrow 0|0S0|0S1|1S0|1S1$$

3. Find a reduced grammar equivalent to the grammar G , having the production rule $P : S \rightarrow AC|B$, $A \rightarrow a$, $C \rightarrow c|BC$, $E \rightarrow aA|e$, $F \rightarrow bB|e$.

Phase 1:

$$T = \{a, c, e\}$$

$$W_1 = \{A, C, E, F\}, W_2 = \{A, C, E, F, S\}, W_3 = \{A, C, E, F, S\}$$

Therefore,

$$G' = \{\{A, C, E, F, S\}, \{a, c, e\}, P, (S)\}$$

$$P : S \rightarrow AC, A \rightarrow a, C \rightarrow c, E \rightarrow aA|e, F \rightarrow bB|e$$

Phase 2:

$$Y_1 = \{S\}, Y_2 = \{S, A, C\}, Y_3 = \{S, A, C, a, c\}, Y_4 = \{S, A, C, a, c\}$$

Therefore,

$$G'' = \{\{S, A, C\}, \{a, c\}, P, (S)\}$$

$$P : S \rightarrow AC, A \rightarrow a, C \rightarrow c$$

4. Remove the unit productions from the grammar whose production rule is given by $P : S \rightarrow AB$, $A \rightarrow a$, $B \rightarrow C|b$, $C \rightarrow D$, $D \rightarrow E$, $E \rightarrow e$.

Because $E \rightarrow e$ and $D \rightarrow E$, we add $D \rightarrow e$ in the production rule to replace $D \rightarrow E$.

$$P : S \rightarrow AB, A \rightarrow a, B \rightarrow C|b, C \rightarrow D, D \rightarrow e, E \rightarrow e$$

Because $D \rightarrow e$ and $C \rightarrow D$, we add $C \rightarrow e$ in the production rule to replace $C \rightarrow D$.

$$P : S \rightarrow AB, A \rightarrow a, B \rightarrow C|b, C \rightarrow e, D \rightarrow e, E \rightarrow e$$

Because $C \rightarrow e$ and $B \rightarrow C$, we add $B \rightarrow e$ in the production rule to replace $B \rightarrow C$.

$$P : S \rightarrow AB, A \rightarrow a, B \rightarrow c|b, C \rightarrow e, D \rightarrow e, E \rightarrow e$$

Remove unreachable symbols

$$Y_1 = \{S\}, Y_2 = \{S, A, B\}, Y_3 = \{S, A, B, a, b\}, Y_4 = \{S, A, B, a, b\}$$

$$P : S \rightarrow AB, A \rightarrow a, B \rightarrow c|b$$

5. Remove the Null productions from the following grammar

$$S \rightarrow XYXZ, X \rightarrow aX|\varepsilon, Y \rightarrow bY|\varepsilon, Z \rightarrow c$$

Step 1: to eliminate $X \rightarrow \varepsilon$

$$\begin{aligned} S \rightarrow XYXZ &\Rightarrow S \rightarrow XYZ, S \rightarrow YXZ, S \rightarrow YZ \\ X \rightarrow aX &\Rightarrow X \rightarrow a \end{aligned}$$

New production rule becomes

$$S \rightarrow XYXZ|XYZ|YXZ|YZ, X \rightarrow aX|a, Y \rightarrow bY|\varepsilon, Z \rightarrow c$$

Step 2: to eliminate $Y \rightarrow \varepsilon$

$$\begin{aligned} S \rightarrow XYXZ|XYZ|YXZ|YZ &\Rightarrow S \rightarrow XXZ|XZ|Z \\ Y \rightarrow bY &\Rightarrow Y \rightarrow b \end{aligned}$$

New production rule becomes

$$S \rightarrow XYXZ|XYZ|YXZ|YZ|XXZ|XZ|Z, X \rightarrow aX|a, Y \rightarrow bY|b, Z \rightarrow c$$

6. Convert the following context-free grammar to the Chomsky normal form.

$$P : S \rightarrow ASA|1B, A \rightarrow B|S, B \rightarrow 0|\varepsilon$$

Step 1: Since S appears in RHS, we need to add a new state S' and $S' \rightarrow S$ is added to the production

$$P : S' \rightarrow S, S \rightarrow ASA|1B, A \rightarrow B|S, B \rightarrow 0|\varepsilon$$

Step 2: Remove the Null Production

to eliminate $B \rightarrow \varepsilon$

$$\begin{aligned} S \rightarrow 1B &\Rightarrow S \rightarrow 1 \\ A \rightarrow B &\Rightarrow A \rightarrow \varepsilon \end{aligned}$$

The production rule becomes

$$P : S' \rightarrow S, S \rightarrow ASA|1B|1, A \rightarrow B|S|\varepsilon, B \rightarrow 0$$

to eliminate $A \rightarrow \varepsilon$

$$S \rightarrow ASA \Rightarrow S \rightarrow AS|SA|S$$

The production rule becomes

$$P : S' \rightarrow S, S \rightarrow ASA|1B|1|AS|SA|S, A \rightarrow B|S, B \rightarrow 0$$

Step 3: Remove the Unit Production $S \rightarrow S, S' \rightarrow S, A \rightarrow B, A \rightarrow S$

After removing $S \rightarrow S$

$$P : S' \rightarrow S, S \rightarrow ASA|1B|1|AS|SA, A \rightarrow B|S, B \rightarrow 0$$

After removing $S' \rightarrow S$

$$P : S' \rightarrow ASA|1B|1|AS|SA, S \rightarrow ASA|1B|1|AS|SA, A \rightarrow B|S, B \rightarrow 0$$

After removing $A \rightarrow B$

$$P : S' \rightarrow ASA|1B|1|AS|SA, S \rightarrow ASA|1B|1|AS|SA, A \rightarrow 0|S, B \rightarrow 0$$

After removing $A \rightarrow S$

$$P : S' \rightarrow ASA|1B|1|AS|SA, S \rightarrow ASA|1B|1|AS|SA, A \rightarrow 0|ASA|1B|1|AS|SA, B \rightarrow 0$$

Step 4: Now find out the productions that have more than two variables in RHS

$$S' \rightarrow ASA, \quad S \rightarrow ASA, \quad A \rightarrow ASA$$

After removing these, we have

$$\begin{aligned} P : \quad & S' \rightarrow AX|1B|1|AS|SA \\ & S \rightarrow AX|1B|1|AS|SA \\ & A \rightarrow 0|AX|1B|1|AS|SA \\ & B \rightarrow 0 \\ & X \rightarrow SA \end{aligned}$$

Step 5: Now change the productions

$$S' \rightarrow 1B, \quad S \rightarrow 1B, \quad A \rightarrow 1B$$

Finally, we get

$$\begin{aligned} P : \quad & S' \rightarrow AX|YB|1|AS|SA \\ & S \rightarrow AX|YB|1|AS|SA \\ & A \rightarrow 0|AX|YB|1|AS|SA \\ & B \rightarrow 0 \\ & X \rightarrow SA \\ & Y \rightarrow 1 \end{aligned}$$