CS 3873: Net-Centric Computing

Assignment 4: IP and Network Routing Total: 14.5

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[Mandatory] Declaration: "I warrant that	this is my own work."
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1. (2 points) IPv6 uses 128-bit addresses. If a block of 1 million addresses is allocated every picosecond, how long will the addresses last? Compare that with the age of our universe ($\approx 10^{10}$ years).

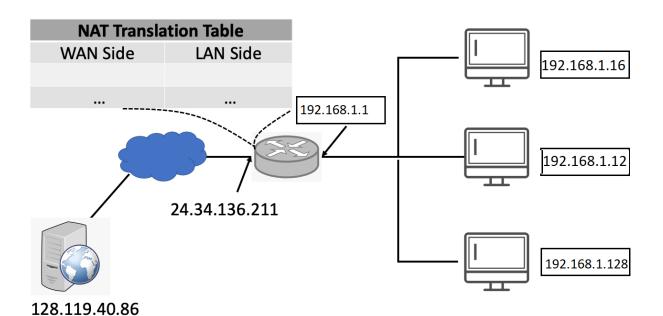
With 128 bits, there will be 3.4×10^{38} addresses. Allocating 3.4×10^{38} addresses every picosecond 10^{18} , they will last for 10^{13} years which is 1000 times the age of the universe.

2. (5 points) Consider the network setup in the following figure. Suppose that the ISP assigns the router the address 24.34.136.211 and that the network prefix of the home network is 192.168.1.0/24.

2)	192.168.1.	0 /24		Chongable
	11000000	10101000	000000/	00000000

a. Assign addresses to all 3 interfaces in the home network and the network interface of the router connected to the home network.

g)	1-11000000 10101000 00000001 00100000
	192.168.1.16
	2 - 11000000 10101000 00000001 00001100
	192.168.1.12
,	3 - 11000000 lololo00 0000000 lo000000
	192. [68.]. 128
	4-11000000 10101000 00000001 00000001
	192.[68.].]
	4-11000000 10101000 00000001 00000001



b. Suppose each host has 2 ongoing TCP connections, all to port 80 at host 128.119.40.86. Provide the 6 corresponding entries in the NAT translation table.

NA	T Translation table
WAN Side	LAN Side
24.34.136.211, 5050	192.168.1.16, 3873
24.34.136.211, 5051	192.168.1.16, 3863
24.34.136.211, 6074	192.168.1.12, 3383
24.34.136.211, 6073	192.168.1.12, 3382
24.34.136.211, 5007	192.168.1.128, 3362
24.34.136.211, 5008	192.168.1.128, 3365

3. (1 point) We made a distinction between the forwarding function and the routing function performed in the network layer. What are the key differences between routing and forwarding?

Routing: network-wide process, route datagrams among routers along and end-to-end path from source host to destination host.

Forwarding: router-local action, forward datagrams (packets) arriving on one of a router's input links to one of its output links.

4. (4 points) Consider a datagram network using 32-bit host addresses. Suppose a router has four links, numbered 0 through 3, and packets are to be forwarded to the link interfaces as follows:

Destination Address Range		Link Interface		
11100000	00000000	00000000	00000000	0
through				
11100000	00111111	11111111	11111111	
11100000	01000000	00000000	00000000	1
through				
11100000	01000000	11111111	11111111	
11100000	01110000	00000000	00000000	2
through				
11100000	01111111	11111111	11111111	
11100001	10110000	00000000	00000000	3
through				
11100001	10111111	11111111	11111111	
11100010	10000000	00000000	00000000	4
through				
11100010	11111111	11111111	11111111	

a. Complete the following forwarding table according to the above setting, assuming longest prefix matching is used to decide where to forward a packet to the correct link interface. Note that the column of network prefix should be presented in the decimal form of a.b.c.d/x. You need to decide how many entries that this table requires.

Network Prefix (Decimal)		Output Link Interface
From	То	
224.0.0.0/10	224.63.255.255/10	0
224.64.0.0/16	224.64.255.255/16	1
224.112.0.0/12	224.127.255.255/12	2
225.176.0.0/12	225.191.255.255/12	3
226.128.0.0/9	226.255.255.255/9	4

b. According to the above forwarding table, give the output link interface for each datagram with the following destination addresses:

225.180.128.1, 224.135.1.2, 224.111.0.1

225.180.128.1 \rightarrow 3 224.135.1.2 \rightarrow 2 (NS) 224.111.0.1 \rightarrow 2

- 5. (3 points) Consider a router that interconnects three subnets: Subnet 1, Subnet 2, and Subnet 3. Suppose all of the interfaces in each of these three subnets should use addresses in a large block defined by prefix 128.16.0.0/12. You need to further divide this large address block into three smaller **non-overlapping** address blocks for these three subnets. They need to further meet the following conditions:
 - a. Subnet 1 is required to support at least 300 interfaces, and the last address for this subnet is 128.31.255.255
 - b. Subnet 2 is to support at least 120 interfaces, and the last address for this subnet is 128.16.1.127
 - c. Subnet 3 is to support at least 400 interfaces, and the last address for this subnet is 128.17.7.255.

• 5)	128.16.0.0/2 = 100000000000000000000000000000000000
a)	$2^{8} < 300 < 2^{9}$ $256 < 300 < 512$
	Last address: 128.31.255.255 = 10000000 000
	Well use 29
	128.31.255.255 - 512
•	= 1000000000000000000000000000000000000
	128.31.254.0/23
b)	120 interface. Therefore, we'll use 29=128
	Last address: 128.16.1.127 = 10000000 00010000 00000001 0 111111
	128. 16.1. 127 - 128
	= start address; 128.16.1.0 = 10000000 00010000 00000001 00000000
•	128.16.1.0/24
	-0.5: 128.16.1.0/25

n

C) 4∞ interfaces, we'll use $2^9 = 512$ Last address: 128.17.7.255= 10000000 0000001 00000111 |111|11|1 128.17.7.255 = 512Start address: 128.17.6.0= 10000000 00010001 0000010 00000000