## University of New Brunswick Faculty of Computer Science

## CS2333: Computability and Formal Languages

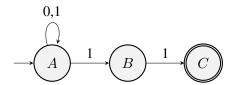
Homework Assignment 3, Due Time, Date 5:00 PM, February 11, 2022

Student Name: \_\_\_\_\_ Matriculation Number: \_\_\_\_\_

Instructor: Rongxing Lu

The marking scheme is shown in the left margin and [100] constitutes full marks.

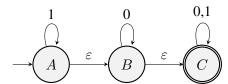
- [20] 1. Give NFAs with the specified number of states recognizing each of the following languages.
- [4] (a) The language  $\{w|w \in \{0,1\}^* \text{ end with } 00 \text{ with three states}\}.$
- [4] (b) The language  $\{0\}$  with two states.
- [4] (c) The language  $0^*1^*0^*0$  with three states.
- [4] (d) The language  $\{\varepsilon\}$  with one state.
- [4] (e) The language  $0^*$  with one state.
- [10] 2. Given below is the NFA for a language



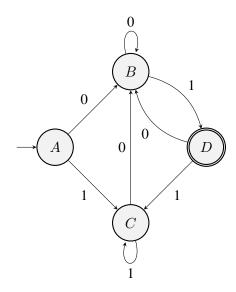
 $L = \{ \text{ set of all strings over } \{0, 1\} \text{ that end with '11'} \}$ 

Construct its equivalent DFA.

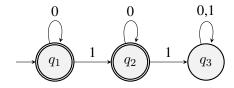
[10] 3. Convert the following  $\varepsilon$ -NFA to its equivalent NFA.



[10] 4. Minimize the following DFA with reduced states.

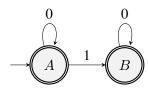


- [10] 5. Design Regular Expression for the following languages over  $\{a, b\}$ .
- [5] (a) Language accepting strings of length at least 1.
- [5] (b) Language accepting strings of length at most 3.
- [10] 6. Find the Regular Expression for the following DFA.

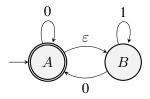


- [10] 7. Covert each of the following Regular Expression to its equivalent Finite Automata.
- [5] (a)  $0^* + 0^*10^*$
- [5] (b) 10 + (1+00)1\*0

- [20] 8. Convert each of the following NFAs to an equivalent DFA.
- [10] (a)

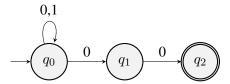


[10] (b)

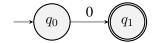


## Solutions.

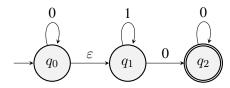
- 1. Give NFAs with the specified number of states recognizing each of the following languages.
  - (a) The language  $\{w|w\in\{0,1\}^* \text{ end with } 00 \text{ with three states}\}.$



(b) The language  $\{0\}$  with two states.



(c) The language 0\*1\*0\*0 with three states.



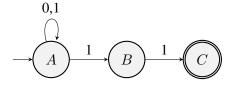
(d) The language  $\{\varepsilon\}$  with one state.



(e) The language  $0^*$  with one state.



2. Given below is the NFA for a language



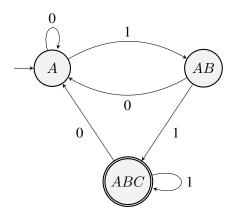
 $L=\{ \mbox{ set of all strings over } \{0,1\} \mbox{ that end with `11'} \}$ 

Construct its equivalent DFA.

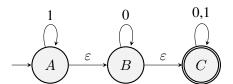
	0	1
$\rightarrow$ A	Α	A,B
В	Ø	C
<u>C</u>	Ø	Ø

Set A as the initial state,

	0	1
$\rightarrow A$	Α	AB
AB	Α	ABC
<u>ABC</u>	Α	ABC



3. Convert the following  $\varepsilon$ -NFA to its equivalent NFA.

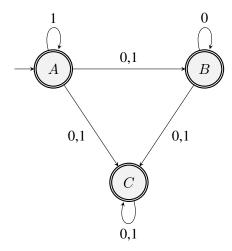


	$\varepsilon^*$	0	$\varepsilon^*$
A	A	Ø	
	В	В	B,C
	C	C	C
	$\varepsilon^*$	1	$\varepsilon^*$
A	Α	Α	A,B,C
	В	Ø	
	C	C	C
	$\varepsilon^*$	0	$\varepsilon^*$
В	В	В	B,C
	C	C	C

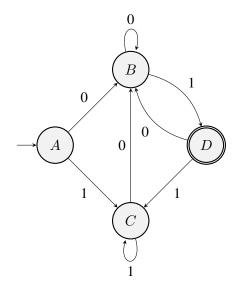
	$\varepsilon^*$	1	$\varepsilon^*$
В	В	Ø	
	C	C	C

A, B, C are all final states, since the  $\varepsilon$  can make them reach C.

	0	1
$\rightarrow \underline{A}$	B,C	A,B,C
$\underline{\mathbf{B}}$	B,C	C
<u>C</u>	C	C



4. Minimize the following DFA with reduced states.

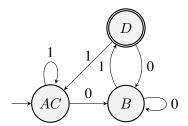


	0	1
$\rightarrow$ A	В	С
В	В	D
C	В	C
<u>D</u>	В	C

- 0-Equivalence  $\{ABC\}$  ,  $\{D\}$
- 1-Equivalence  $\{AC\},\,\{B\}$  ,  $\{D\}$
- 2-Equivalence  $\{AC\},\,\{B\}$  ,  $\{D\}$

As there is no difference between 1-Equivalence and 2-Equivalence, we have

	0	1
$\rightarrow$ AC	В	AC
В	В	D
$\underline{\mathbf{D}}$	В	AC



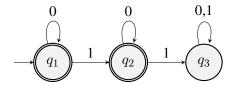
- 5. Design Regular Expression for the following languages over  $\{a,b\}$ .
  - (a) Language accepting strings of length at least 1.

$$L_1 = \{a, b, aa, ab, ba, bb, aaa, \dots\}$$
  
 $R_1 = (a+b)(a+b)^*$ 

(b) Language accepting strings of length at most 3.

$$L_2 = \{\varepsilon, a, b, aa, ab, ba, bb, aaa, aab, aba, abb, baa, bab, bba, bbb \cdots \}$$
$$R_2 = (\varepsilon + a + b)(\varepsilon + a + b)(\varepsilon + a + b)$$

6. Find the Regular Expression for the following DFA.



$$q_1 = \varepsilon + q_1 \cdot 0 \tag{1}$$

$$q_2 = q_1 \cdot 1 + q_2 \cdot 0 \tag{2}$$

$$q_3 = q_2 \cdot 1 + q_3 \cdot 0 + q_3 \cdot 1 \tag{3}$$

Based on the Arden Theorem, we know if R = Q + RP, we have  $R = QP^*$ .

Final state  $q_1$ 

From (1), we have

$$\underbrace{q_1}_{R} = \underbrace{\varepsilon}_{Q} + \underbrace{q_1}_{R} \cdot \underbrace{0}_{P}$$

and then  $q_1 = \varepsilon 0^*$ , that is  $q_1 = 0^*$ .

Final state  $q_2$ 

From (2), we have

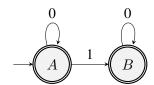
$$q_2 = q_1 \cdot 1 + q_2 \cdot 0$$
  $q_2 = 0^* \cdot 1 + q_2 \cdot 0$   $\underbrace{q_2}_R = \underbrace{0^* \cdot 1}_Q + \underbrace{q_2}_R \cdot \underbrace{0}_P$ 

and then  $q_2 = 0^* 10^*$ .

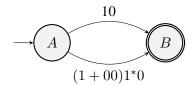
Then, R is the union of both final states

$$R = 0^* + 0^*10^*$$
 or  $R = 0^* + 0^*10^* = 0^*(\varepsilon + 10^*)$ 

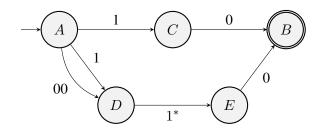
- 7. Covert each of the following Regular Expression to its equivalent Finite Automata.
  - (a)  $0^* + 0^*10^*$



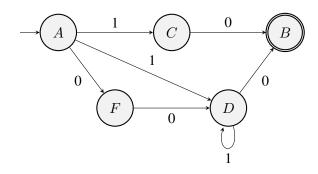
(b) 10 + (1+00)1\*0



 $\Rightarrow$ 

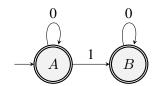


 $\Rightarrow$ 

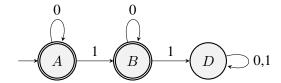


8. Convert each of the following NFAs to an equivalent DFA.

(a)



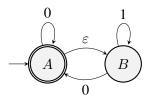
	0	1
$\rightarrow \underline{A}$	A	В
<u>B</u>	В	Ø



	0	1
$\rightarrow \underline{A}$	A	В
<u>B</u>	В	D
D	D	D

D is a dead state.

(b)



Step 1: convert  $\varepsilon$ -NFA to NFA.

	$\varepsilon^*$	0	$\varepsilon^*$
A	A	A	A,B
	В	A	A,B
	$\varepsilon^*$	1	$\varepsilon^*$
A	Α	Ø	
	В	В	В
	$\varepsilon^*$	0	$\varepsilon^*$
В	В	A	A,B
	$\varepsilon^*$	1	$\varepsilon^*$
В	В	В	В
Step	2:		

	0	1
$\rightarrow \underline{A}$	A,B	В
В	A.B	В

Step 3: set AB as the initial state (Note that, if we set A as the initial state, we can use the minimization of DFA to get the same result.)

	0	1
$\rightarrow \underline{AB}$	AB	В
В	AB	В

