

From Bits and Gates to C and Beyond

Assembly Language

Chapter 7

Human-Friendly Programming

Computers need binary instruction encodings...

0001110010000110

Humans prefer symbolic languages...

a = b + c

High-level languages allow us to write programs in clear, precise language that is more like English or math. Requires a program (compiler) to translage from symbolic language to machine instructions.

Examples: C, Python, Fortran, Java, ...

We will introduce C in Chapters 11 to 19.

Assembly Language: Human-Friendly ISA Programming

Assembly Language is a low-level symbolic language, just a short step above machine instructions.

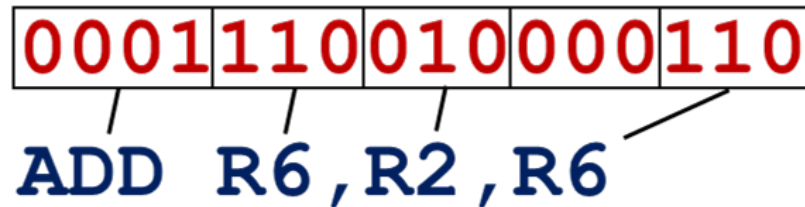
- Don't have to remember opcodes (ADD = 0001, NOT = 1001, ...).
- Give symbolic names to memory locations -- don't have to do binary arithmetic to calculate offsets.
- Like machine instructions, allows programmer explicit, instruction-level specification of program.

Disadvantage:

Not portable. Every ISA has its own assembly language.
Program written for one platform does not run on another.

Assembly Language

Very similar format to instructions -- replace bit fields with symbols.



For the most part, one line of assembly language = one instruction.

Some additional features for allocating memory, initializing memory locations, service calls.

Numerical values specified in hexadecimal (x30AB) or decimal (#10).

Example Program

```
;
; Program to multiply a number by the constant 6
;
        .ORIG    x3050
        LD       R1, SIX
        LD       R2, NUMBER
        AND      R3, R3, #0           ; Clear R3.  It will
                                       ; contain the product.
; The inner loop
;
AGAIN    ADD      R3, R3, R2
        ADD      R1, R1, #-1         ; R1 keeps track of
        BRp     AGAIN               ; the iteration.
;
        HALT
;
NUMBER   .BLKW    1
SIX      .FILL    x0006
;
        .END
```

Labels

Assembler Directives

Instructions

Comments

Assembly Language Syntax

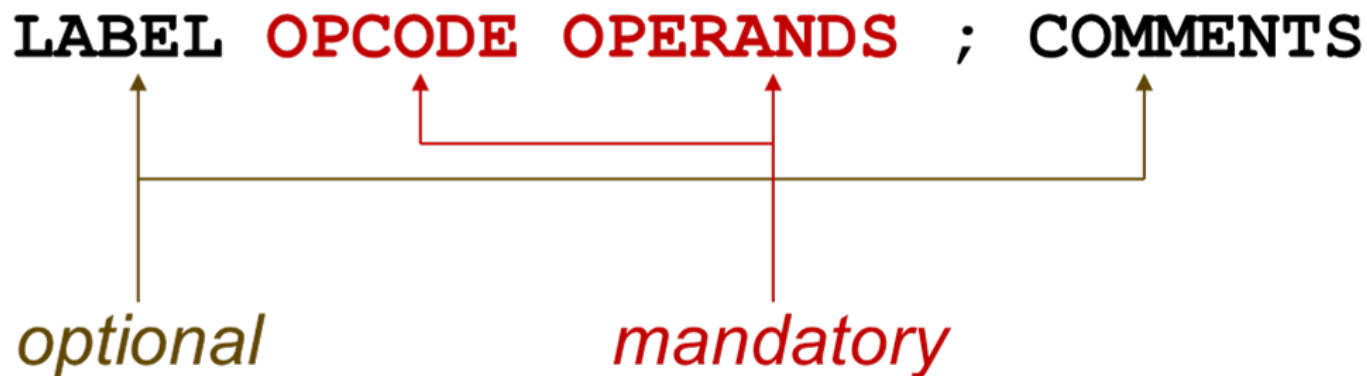
Each line of a program is one of the following:

- An instruction.
- An assembler directive (or pseudo-op).
- A comment.

Whitespace (between symbols) and case are ignored.

Comments (beginning with “;”) are also ignored.

An instruction has the following format:



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Mandatory: Opcode and Operands

Opcodes

Reserved symbols that correspond to LC-3 instructions.
Listed in Appendix A and Figure 5.3.

- For example: ADD, AND, LD, LDR, ...

reserved means that it
cannot be used as a label

Operands

- Registers -- specified by Rn, where n is the register number.
- Numbers -- indicated by # (decimal) or x (hex).
- Label -- symbolic name of memory location.
- Separated by comma (whitespace ignored).
- Number, order, and type correspond to instruction format.

```
ADD R1, R1, R3      ; DR, SR1, SR2
ADD R1, R1, #3       ; DR, SR1, Imm5
LD R6, NUMBER       ; DR, address (converted to PCoffset)
BRz LOOP            ; nzp becomes part of opcode, address
```

Optional: Label and Comment

Label

- Placed at the beginning of the line.
- Assigns a symbolic name to the address corresponding to that line.

```
LOOP    ADD    R1,R1,#-1    ; LOOP is address of ADD
        BRp    LOOP
```

Comment

A semicolon, and anything after it on the same line, is a comment.

Ignored by assembler.

Used by humans to document/understand programs.

Tips for useful comments:

- Avoid restating the obvious, as “decrement R1.”
- Provide additional insight, as in “accumulate product in R6.”
- Use comments and empty lines to separate pieces of program.

Assembler Directive

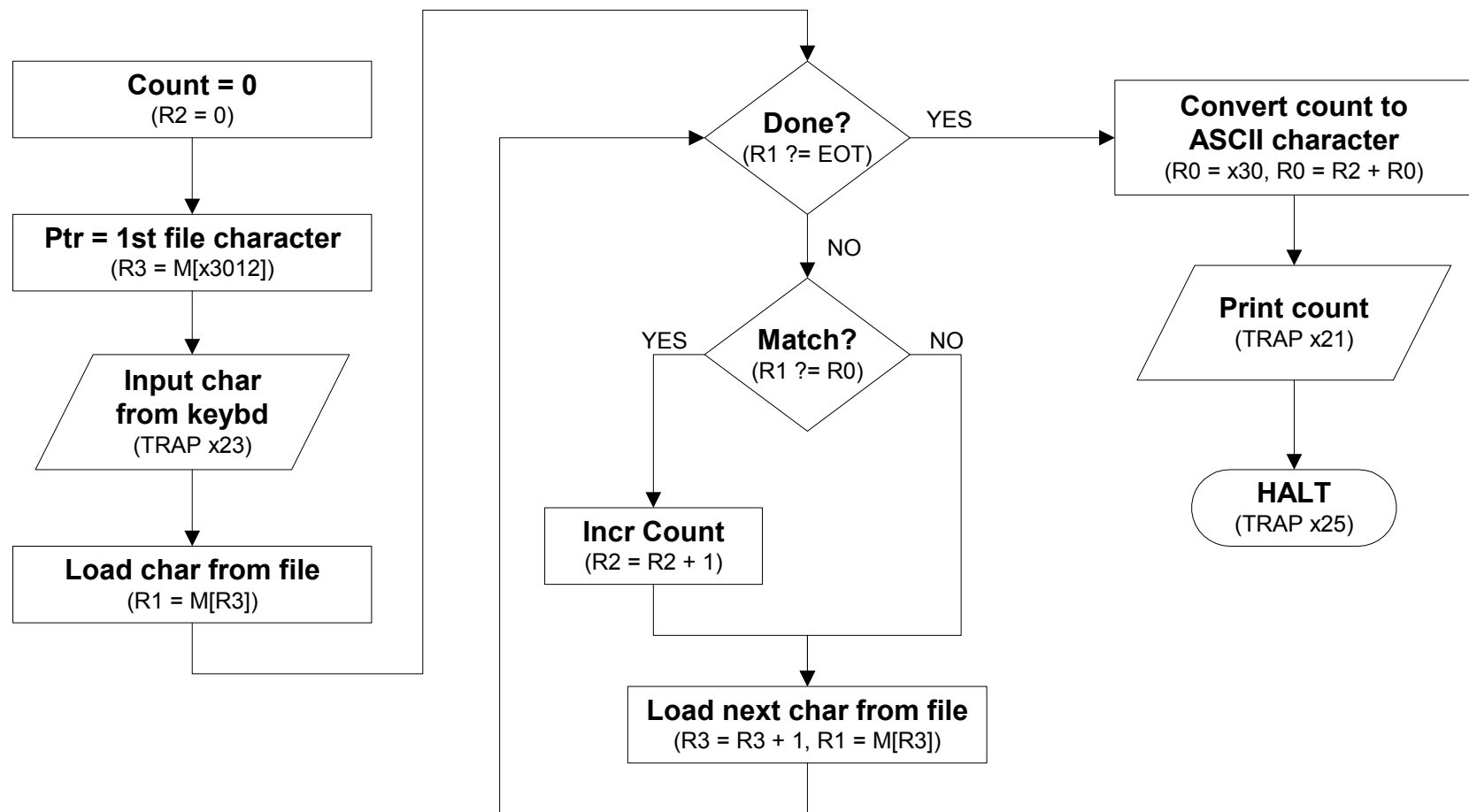
Pseudo-operation

- Does not refer to an actual instruction to be executed.
- Tells the assembler to do something.
- Looks like an instruction, except "opcode" starts with a dot.

<i>Opcode</i>	<i>Operand</i>	<i>Meaning</i>
.ORIG	address	starting address of program
.END		end of program
.BLKW	n	allocate n words of storage
.FILL	n	allocate one word, initialize with value n
.STRINGZ	n-character string	allocate n+1 locations, initialize w/ characters and null terminator

Sample Program: Counting Occurrences in a File

Once again, we show the program that counts the number of times (up to nine) a user-specified character appears in a file.



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Assembly Language Program ¹

```
;
; Program to count occurrences of a character in a file.
; Character to be input from the keyboard.
; Result to be displayed on the monitor.
; Program only works if no more than 9 occurrences are found.
;
;
; Initialization
;
        .ORIG      x3000
        AND        R2, R2, #0           ; R2 is counter, initially 0
        LD         R3, PTR              ; R3 is pointer to characters
        TRAP       x23                  ; R0 gets character input
        LDR        R1, R3, #0           ; R1 gets first character
;
; Test character for end of file
;
TEST     ADD        R4, R1, #-4          ; Test for EOT (ASCII x04)
        BRz        OUTPUT              ; If done, prepare the output
;
; Test character for match.  If a match, increment count.
;
        NOT        R1, R1
        ADD        R1, R1, #1
        ADD        R1, R1, R0           ; Compute R0-R1 to compare
        BRnp       GETCHAR             ; If no match, do not increment count
        ADD        R2, R2, #1
```

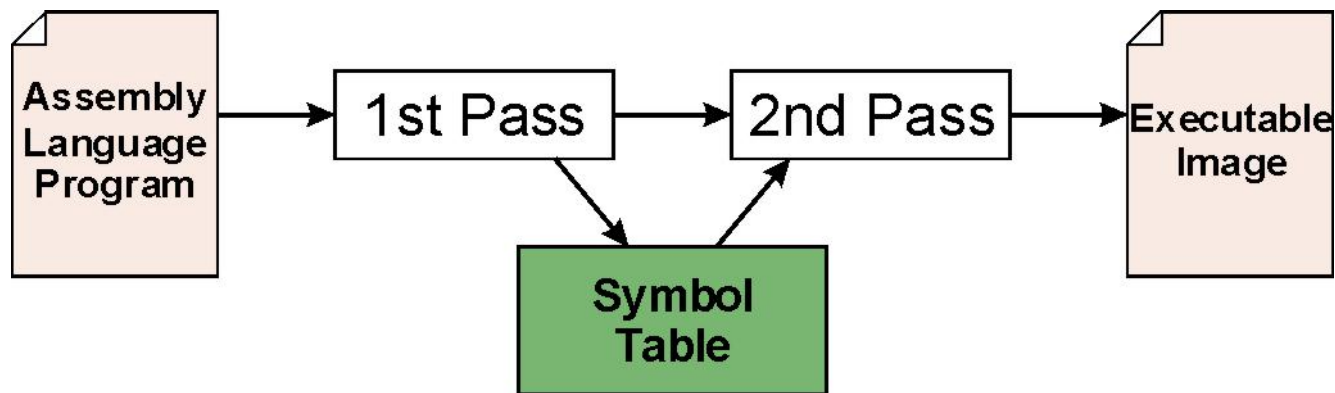
Assembly Language Program ₂

```
;
; Get next character from file.
;
GETCHAR      ADD      R3, R3, #1      ; Point to next character.
              LDR      R1, R3, #0      ; R1 gets next char to test
              BRnzp    TEST
;
; Output the count.
;
OUTPUT       LD        R0, ASCII      ; Load the ASCII template
              ADD      R0, R0, R2      ; Covert binary count to ASCII
              TRAP     x21             ; ASCII code in R0 is displayed.
              TRAP     x25             ; Halt machine

;
; Storage for pointer and ASCII template
;
ASCII        .FILL     x0030
PTR          .FILL     x4000
              .END
```

Assembly Process

The assembler is a program that translate an assembly language (.asm) file to a binary object (.obj) file that can be loaded into memory.



First Pass:

- Scan program file, check for syntax errors.
- Find all labels and calculate the corresponding addresses: the symbol table.

Second Pass:

- Convert instructions to machine language, using information from symbol table.

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First Pass: Construct the Symbol Table

1. Find the .ORIG statement,
which tells us the address of the first instruction.
 - Initialize location counter (LC), which keeps track of the current instruction.
2. For each non-empty line in the program:
 - If line contains a label, add label and LC to symbol table.
 - Increment LC.
 - NOTE: If statement is .BLKW or .STRINGZ, increment LC by the number of words allocated.
3. Stop when .END statement is reached.

NOTE: A line that contains only a comment is considered an empty line.

First Pass on Sample Program (Comments Removed)

```
--
x3000      .ORIG    x3000
x3001      AND      R2, R2, #0
x3002      LD        R3, PTR
x3003      TRAP      x23
x3004      TEST     ADD      R1, R3, #0
x3005      BRz       OUTPUT
x3006      NOT       R1, R1
x3007      ADD       R1, R1, #1
x3008      ADD       R1, R1, R0
x3009      BRnp      GETCHAR
x300A      ADD       R2, R2, #1
x300B      GETCHAR  ADD      R3, R3, #1
x300C      LDR       R1, R3, #0
x300D      BRnzp     TEST
x300E      OUTPUT   LD        R0, ASCII
x300F      ADD       R0, R0, R2
x3010      TRAP      x21
x3011      TRAP      x25
x3012      ASCII    .FILL    x0030
x3013      PTR      .FILL    x4000
--
--      .END
```

Label	Address
TEST	x3004
GETCHAR	x300B
OUTPUT	x300E
ASCII	x3012
PTR	x3013

Second Pass: Convert to Machine Instructions

1. Find the .ORIG statement,
which tells us the address of the first instruction.
 - Initialize location counter (LC), which keeps track of the current instruction.
2. For each non-empty line in the program:
 - If line contains an instruction, translate opcode and operands to binary machine instruction. For label, lookup address in symbol table and subtract (LC+1). Increment LC.
 - If line contains .FILL, convert value/label to binary. Increment LC.
 - If line contains .BLKW, create n copies of x0000 (or any arbitrary value). Increment LC by n.
 - If line contains .STRINGZ, convert each ASCII character to 16-bit binary value. Add null (x0000). Increment LC by n+1.
3. Stop when .END statement is reached.

Errors during Code Translation

While assembly language is being translated to machine instructions, several types of errors may be discovered.

- Immediate value too large -- can't fit in Imm5 field.
- Address out of range -- greater than $LC+1+255$ or less than $LC+1-256$.
- Symbol not defined, not found in symbol table.

If error is detected, assembly process is stopped and an error message is printed for the user.

Beyond a Single Object File

Larger programs may be written by multiple programmers, or may use modules written by a third party. Each module is assembled independently, each creating its own **object file** and **symbol table**.

To execute, a program must have all of its modules combined into a single **executable** image.

Linking is the process to combine all of the necessary object files into a single executable.

External Symbols

In the assembly code we're writing, we may want to symbolically refer to information defined in a different module.

For example, suppose we don't know the starting address of the file in our counting program. The starting address and the file data could be defined in a different module.

We want to do this:

```
PTR    .FILL    STARTofFILE
```

To tell the assembler that `STARTofFILE` will be defined in a different module, we could do something like this:

```
.EXTERNAL    STARTofFILE
```

This tells the assembler that it's not an error that `STARTofFILE` is not defined. It will be up to the linker to find the symbol in a different module and fill in the information when creating the executable.



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