

Ans to the question no. 5

$$(1) \quad T(n) = T\left(\frac{n}{2}\right) + n - 1; \quad T(1) = 0$$

Here,

$$T(n) = T\left(\frac{n}{2}\right) + n - 1$$

$$= \left[ T\left(\frac{n}{2^2}\right) + \frac{n}{2} \right] + n - 1$$

$$= T\left(\frac{n}{2^2}\right) + \frac{n}{2} + n - 1$$

$$= \left[ T\left(\frac{n}{2^3}\right) + \frac{n}{2^2} \right] + \frac{n}{2} + n - 1$$

$$= T\left(\frac{n}{2^3}\right) + \frac{n}{2^2} + \frac{n}{2} + n - 1$$

$\vdots$

$$= T\left(\frac{n}{2^k}\right) + \frac{n}{2^{k-1}} + \frac{n}{2^{k-2}} + \dots + \frac{n}{2} + n - 1$$

Here,  $\frac{n}{2^k} = 1$

$$\Rightarrow n = 2^k$$

$$\Rightarrow k = \log n$$

Now,

$$T(n) = T(1) + n \left[ \frac{1}{2^{k-1}} + \frac{1}{2^{k-2}} + \dots + \frac{1}{2} + 1 \right] - 1$$

$$= 0 + n[1+1] - 1$$

$$= 0 + 2n - 1$$

$$= 2n - 1 \approx O(n)$$

(A-2)

$$(2) \quad T(n) = T(n-1) + n - 1; \quad T(1) = 0$$

Here,

$$T(n) = T(n-1) + n - 1$$

$$= [T(n-2) + n - 1] + n - 1$$

$$= T(n-2) + (n-1) + n - 1$$

$$= [T(n-3) + n - 2] + (n-1) + n - 1$$

$$(2) T(n) = T(n-1) + n - 1; T(1) = 0$$

Now,

$$T(n) = T(n-1) + (n-1)$$

$$= [T(n-2) + n-2] + n-1$$

$$= T(n-2) + 2n-3$$

$$= T(n-4) + 4n-10$$

⋮

$$T(n) = T(n-k) + kn - \cancel{10} c$$

[c = 10]

\*constant

Now, ~~k~~  $n-k = 1$

$$\Rightarrow k = n-1$$

$$T(n) = T(n-(n-1)) + (n-1)n - c$$

$$= T(1) + n^2 - n - c$$

$$= 0 + n^2 - n - c \approx n^2 - n \approx n^2$$

[c is negligible]

So, ~~Time~~ worst case complexity  $O(n^2)$

$$\textcircled{3} \quad T(n) = T(n/3) + 2T(n/3) + n$$

$$n + (n/3) T(3) = (n) T(1)$$

$$= 3T(n/3) + n$$

$$n + (n/3) T(3) = (n) T(1)$$

$$= 3 \left( 3T(n/3^2) + \frac{n}{3} \right) + n$$

$$n + (n/3) T(3) = (n) T(1)$$

$$= 9T(n/3^2) + 2n$$

$$= 9 \left( 3T(n/3^3) + \frac{n}{3^2} \right) + 2n$$

$$= 27T(n/3^3) + 3n$$

$$\vdots$$

$$= 3^k T(n/3^k) + kn$$

Here,  $\frac{n}{3^k} = 1$

$$\Rightarrow k = \log_3 n$$

$$T(n) = 3^k T(1)$$

$$\text{Again, } 3^k T(1) = 3^{\log_3 n} T(1) = n \cdot 1 = n$$

Therefore,  $T(n) = 3^k T(n/3^k) + kn$  ①

$$T(n) = 3^k T(n/3^k) + kn$$

$$T(1) = 1$$

$$= 3^{\log_3 n} T(1) + \log_3 n \cdot n$$

$$= n + \log_3 n \cdot n$$

$$= n + \log_3 n \cdot n$$

$$= n + \log_3 n \cdot n$$

$$\approx O(n \log n)$$

$$= n + \log_3 n \cdot n$$

$$= n + \log_3 n \cdot n$$

$$=$$

$$=$$

$$= n + \log_3 n \cdot n$$

$$L = \frac{n}{3^k}$$

$$\Rightarrow k = \log_3 n$$

$$T(1) = 1$$

$$T(n) = 3^k T(n/3^k) + kn$$



$$\textcircled{4} \quad T(n) = 2T(n/2) + n^2$$

Here, Using the master theorem, we get,

the recurrence form is,

$$T(n) = aT\left(\frac{n}{b}\right) + \Theta(n^k \log^p n)$$

where  $a \geq 1$ ,  $b > 1$ ,  $k \geq 0$  and  $p$  is real number.

~~Here, Now, we~~

we can see that,  $a=2$ ,  $b=2$ ,  $k=2$ ,  $p=0$

so, ~~mas~~ Here  $a < b^k$  and  $p=0$ .

Therefore by the master theorem

we get  $T(n) = O(n^2)$

[Proved]