

# The Cedilleum Language Specification

## Syntax, Typing, Reduction, and Elaboration

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## 1 Syntax

$id$	identifiers for definitions
$u$	term variables
$X$	type variables
$k$	kind variables
$x ::= id \mid u \mid X$	non-kind variables

Figure 1: Identifiers

$uterm s ::=$	$u$
	$\lambda u. uterm$
	$uterm uterm$

Figure 2: Untyped terms

<i>mod</i>	<code>::= <b>module</b> <i>id</i> . <i>imprt</i>* <i>cmd</i>*</code>	module declarations
<i>imprt</i>	<code>::= <b>import</b> <i>id</i> .</code>	module imports
<i>cmd</i>	<code>::= <i>defTermOrType</i>       <i>defDataType</i>       <i>defKind</i></code>	definitions
<i>defTermOrType</i>	<code>::= <i>id</i> <i>checkType</i>? = <i>term</i> .</code>	term definition
	<code>      <i>id</i> : <i>kind</i> = <i>type</i> .</code>	type definition
<i>defDataType</i>	<code>::= <b>data</b> <i>id</i> <i>param</i>* : <i>kind</i> = <i>constr</i>* .</code>	datatype definitions
<i>defKind</i>	<code>::= <i>k</i> = <i>kind</i></code>	kind definition
<i>checkType</i>	<code>::= : <i>type</i></code>	annotation for term definition
<i>param</i>	<code>::= (<i>x</i> : <i>typeOrKind</i>)</code>	
<i>typeOrKind</i>	<code>::= <i>type</i>       <i>kind</i></code>	
<i>constr</i>	<code>::=   <i>id</i> : <i>type</i></code>	

Figure 3: Modules and definitions

<i>kind</i>	<code>::= <math>\Pi x : \text{typeOrKind} . \text{kind}</math></code>	explicit product
	<code>      <math>\text{typeOrKind} \rightarrow \text{kind}</math></code>	kind arrow
	<code>      <math>\star</math>       (<i>kind</i>)</code>	
<i>type</i>	<code>::= <math>\Pi x : \text{type} . \text{type}</math></code>	explicit product
	<code>      <math>\forall x : \text{typeOrKind} . \text{type}</math></code>	implicit product
	<code>      <math>\lambda x : \text{typeOrKind} . \text{type}</math></code>	type-level function
	<code>      <math>\text{type} \Rightarrow \text{type}</math></code>	arrow with erased domain
	<code>      <math>\text{type} \rightarrow \text{type}</math></code>	normal arrow type
	<code>      <math>\text{type} \cdot \text{type}</math></code>	application to another type
	<code>      <math>\text{type } \text{term}</math></code>	application to a term
	<code>      <math>\{ \text{uterm} \simeq \text{uterm} \}</math></code>	untyped equality
	<code>      (<i>type</i>)</code>	
	<code>      <i>X</i></code>	type variable
	<code>      •</code>	hole for incomplete types

Figure 4: Kinds and types

$term$	$::=$	$\lambda x \text{ class}^? . term$	normal abstraction
		$\Lambda x \text{ class}^? . term$	erased abstraction
		$[ \text{defTermOrType} ] - term$	let
		$\rho \text{ term} - term$	equality elimination by rewriting
		$\phi \text{ term} - term \{term\}$	type cast
		$\chi \text{ type}^? - term$	check a term against a type
		$\delta - term$	ex falso quodlibet
		$\theta \text{ term term}$	elimination with a motive
		$term \text{ term}$	applications
		$term - term$	application to an erased term
		$term \cdot type$	application to a type
		$\beta \{term\}$	reflexivity of equality
		$\varsigma \text{ term}$	symmetry of equality
		$\mu \text{ term motive}^? \{ \text{case}^* \}$	pattern match and fixpoint
		$u$	term variable
		$(term)$	
		$\bullet$	hole for incomplete term
$vararg$	$::=$	$u$	normal constructor argument
		$- u$	erased constructor argument
		$\cdot X$	type constructor argument
$class$	$::=$	$: typeOrKind$	
$motive$	$::=$	$@ \text{ type}$	motive for induction
$case$	$::=$	$  id \text{ arg}^* \mapsto term$	pattern-matching cases

Figure 5: Annotated Terms