

# The Cedilleum Language Specification

## Syntax, Typing, Reduction, and Elaboration

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## 1 Syntax

$id$	identifiers for definitions
$u$	term variables
$X$	type variables
$k$	kind variables
$x ::= id \mid u \mid X \mid k$	any variable

Figure 1: Identifiers

$uterm s ::= u$
$\lambda u. uterm$
$uterm uterm$

Figure 2: Untyped terms

<i>mod</i>	::=	<b>module</b> <i>id</i> . <i>imprt</i> * <i>cmd</i> *	module declarations
<i>imprt</i>	::=	<b>import</b> <i>id</i> .	module imports
<i>cmd</i>	::=	<i>defTermOrType</i> <i>defDataType</i> <i>defKind</i>	definitions
<i>defTermOrType</i>	::=	<i>id</i> <i>checkType</i> ? = <i>term</i> .	term definition
		<i>id</i> $\triangleleft$ <i>kind</i> = <i>type</i> .	type definition
<i>defDataType</i>	::=	<b>data</b> <i>id</i> <i>param</i> * : <i>kind</i> = <i>constr</i> * .	datatype definitions
<i>defKind</i>	::=	<i>k</i> = <i>kind</i>	
<i>checkType</i>	::=	: <i>type</i>	annotation for term definition
<i>param</i>	::=	( <i>x</i> : <i>typeOrKind</i> )	
<i>typeOrKind</i>	::=	<i>type</i> <i>kind</i>	
<i>constr</i>	::=	<i>id</i> : <i>type</i>	

Figure 3: Modules and definitions

<i>kind</i>	::=	$\star$	
		$\prod x : \text{typeOrKind} . \text{kind}$	explicit product
		$\text{typeOrKind} \rightarrow \text{kind}$	kind arrow
		<i>k term</i>	
		<i>k</i> · <i>type</i>	
<i>type</i>	::=	<i>X</i>	
		$\prod x : \text{type} . \text{type}$	explicit product
		$\forall x : \text{typeOrKind} . \text{type}$	implicit product
		$\lambda x : \text{typeOrKind} . \text{type}$	type-level function
		$\text{type} \rightarrow \text{type}$	normal arrow type
		$\text{type} \Rightarrow \text{type}$	arrow with erased domain
		<i>type</i> · <i>type</i>	
		<i>type term</i>	
		{ <i>uterm</i> $\simeq$ <i>uterm</i> }	untyped equality

Figure 4: Kinds and types

<i>term</i>	<i>::=</i>	<i>x</i>	
		$\lambda x \text{ class}^? . \text{term}$	normal abstraction
		$\Lambda x \text{ class}^? . \text{term}$	erased abstraction
		$[ \text{defTermOrType} ] - \text{term}$	let
		$\text{term } \text{arg}^*$	applications
		$\beta \{ \text{term} \}$	reflexivity of equality
		$\varsigma \text{ term}$	symmetry of equality
		$\rho \text{ term} - \text{term}$	equality elimination by rewriting
		$\delta - \text{term}$	ex falso quodlibet
		$\phi \text{ term} - \text{term} \{ \text{term} \}$	type cast
		$\chi \text{ type}^? - \text{term}$	check a term against a type
		$\mu \text{ term } \text{motive}^? \{ \text{case}^* \}$	pattern match and fixpoint
<i>arg</i>	<i>::=</i>	<i>term</i>	normal application
		$- \text{term}$	application to erased term
		$\cdot \text{term}$	application to type
<i>class</i>	<i>::=</i>	$: \text{typeOrKind}$	
<i>motive</i>	<i>::=</i>	$@ \text{type}$	motive for induction
<i>case</i>	<i>::=</i>	$  \text{id } \text{arg}^* . \text{term}$	

Figure 5: Annotated Terms