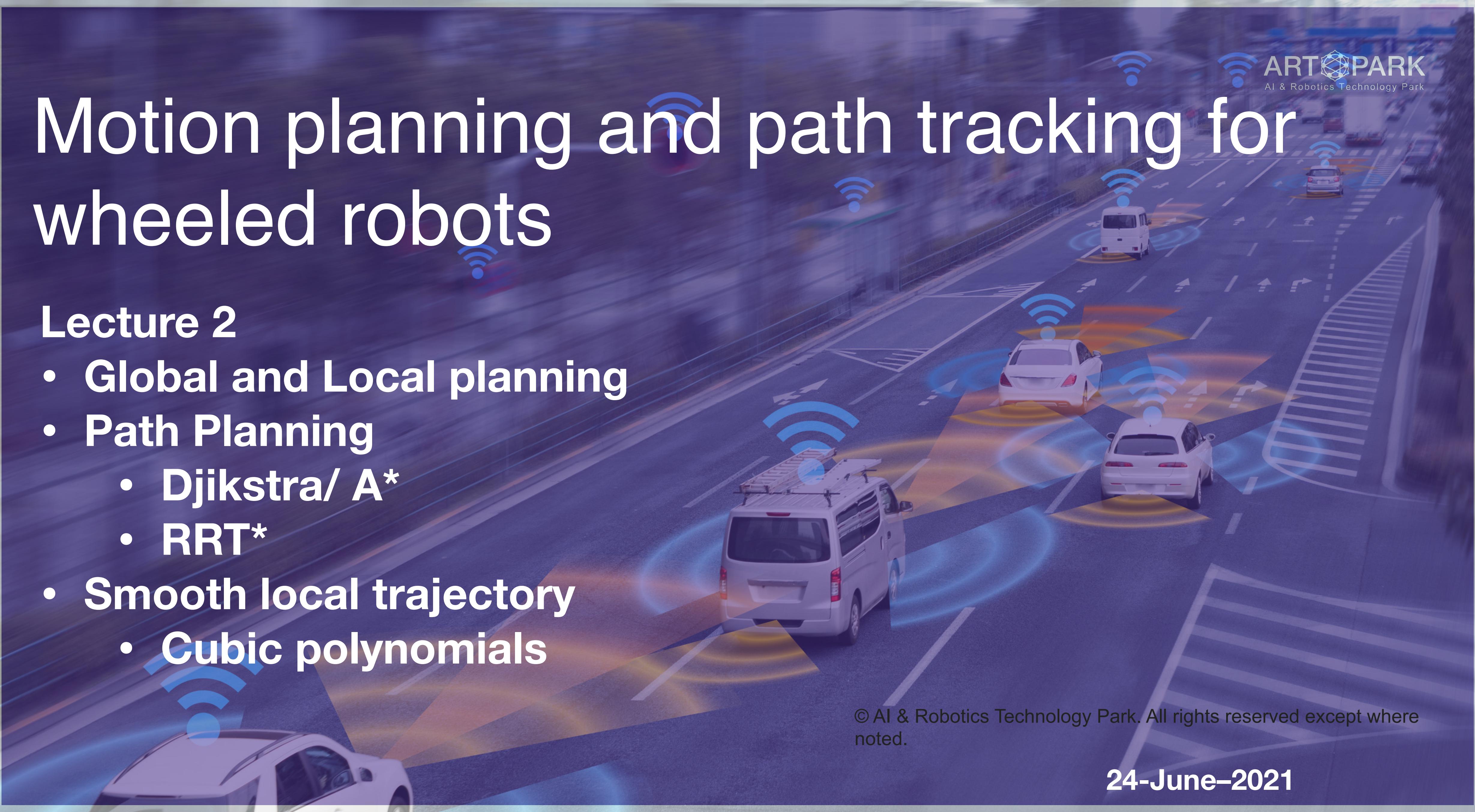


# Motion planning and path tracking for wheeled robots

## Lecture 2

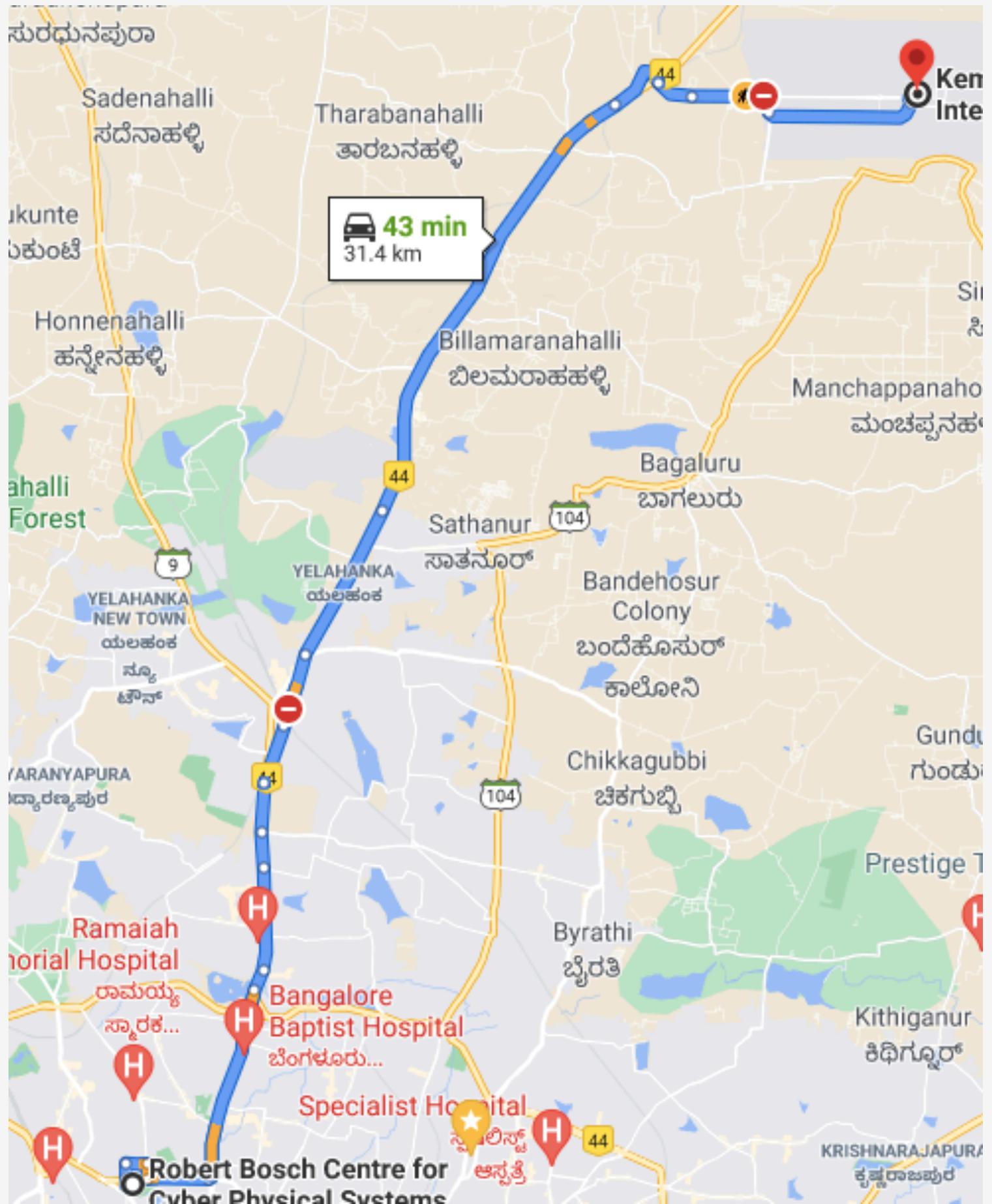
- Global and Local planning
- Path Planning
  - Djikstra/ A\*
  - RRT\*
- Smooth local trajectory
  - Cubic polynomials



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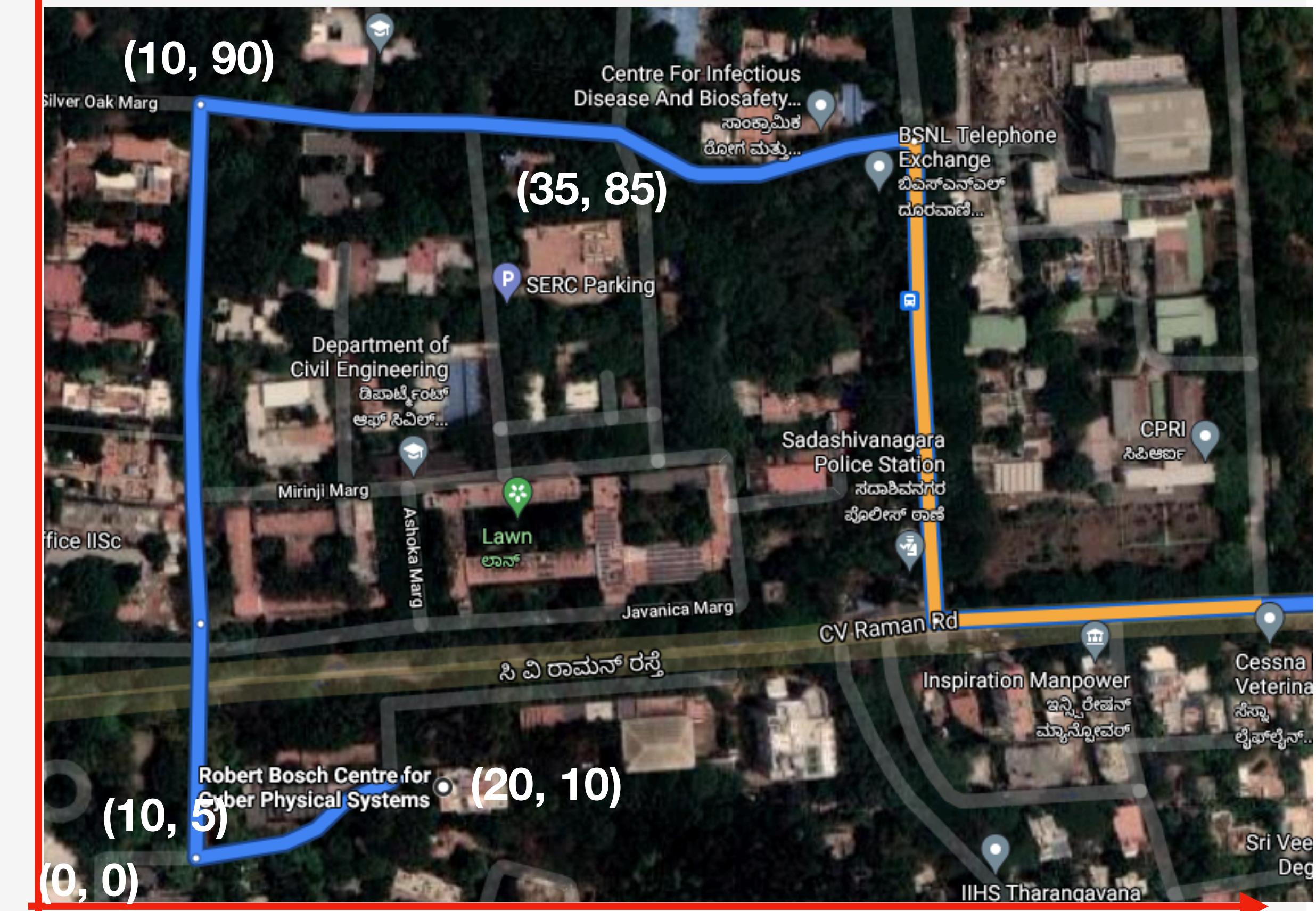
# Planning

Global planning: Go from campus to Airport



Locally safe planning

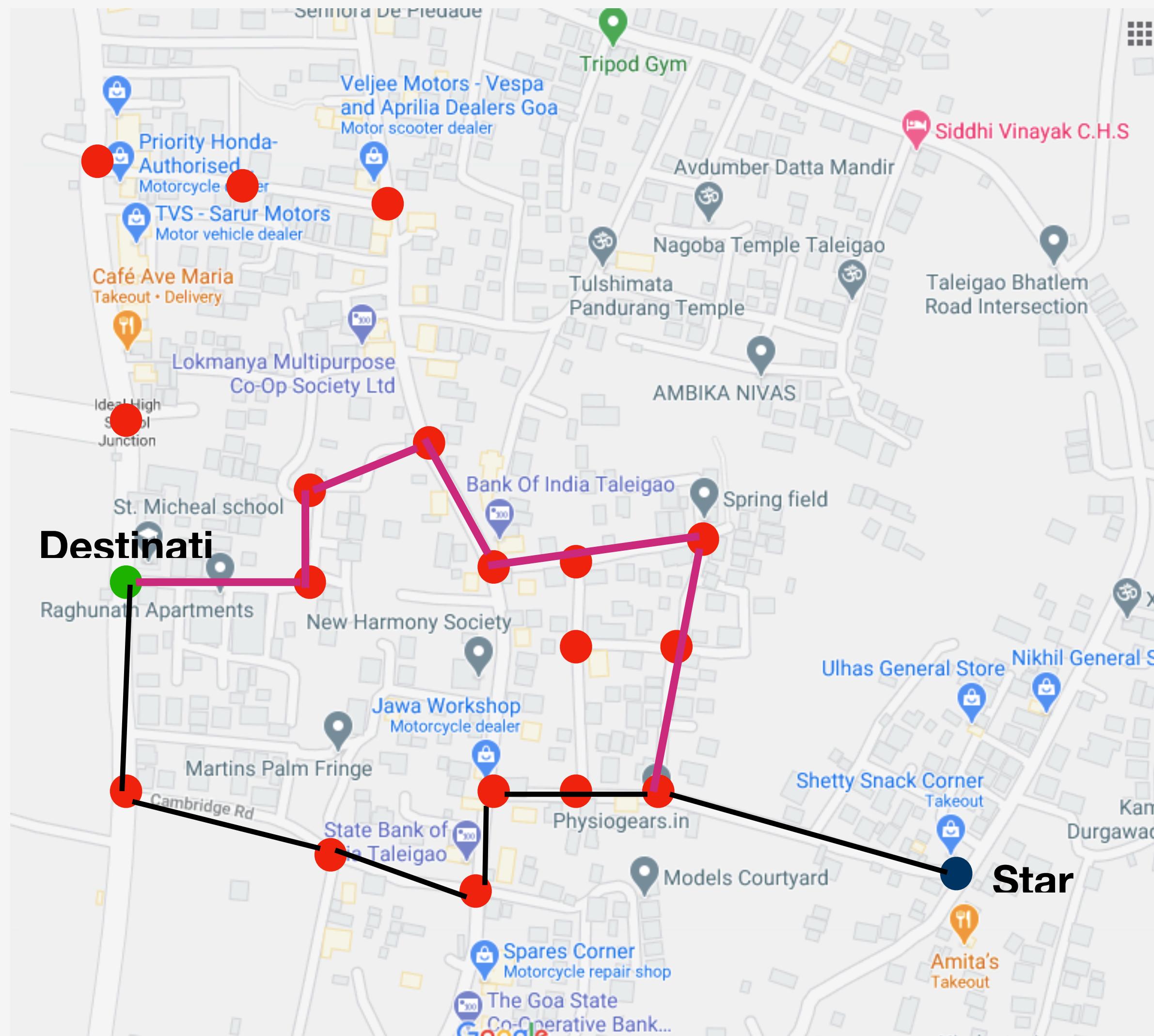
Vehicle action in immediate time horizon



Useful for High-level estimates

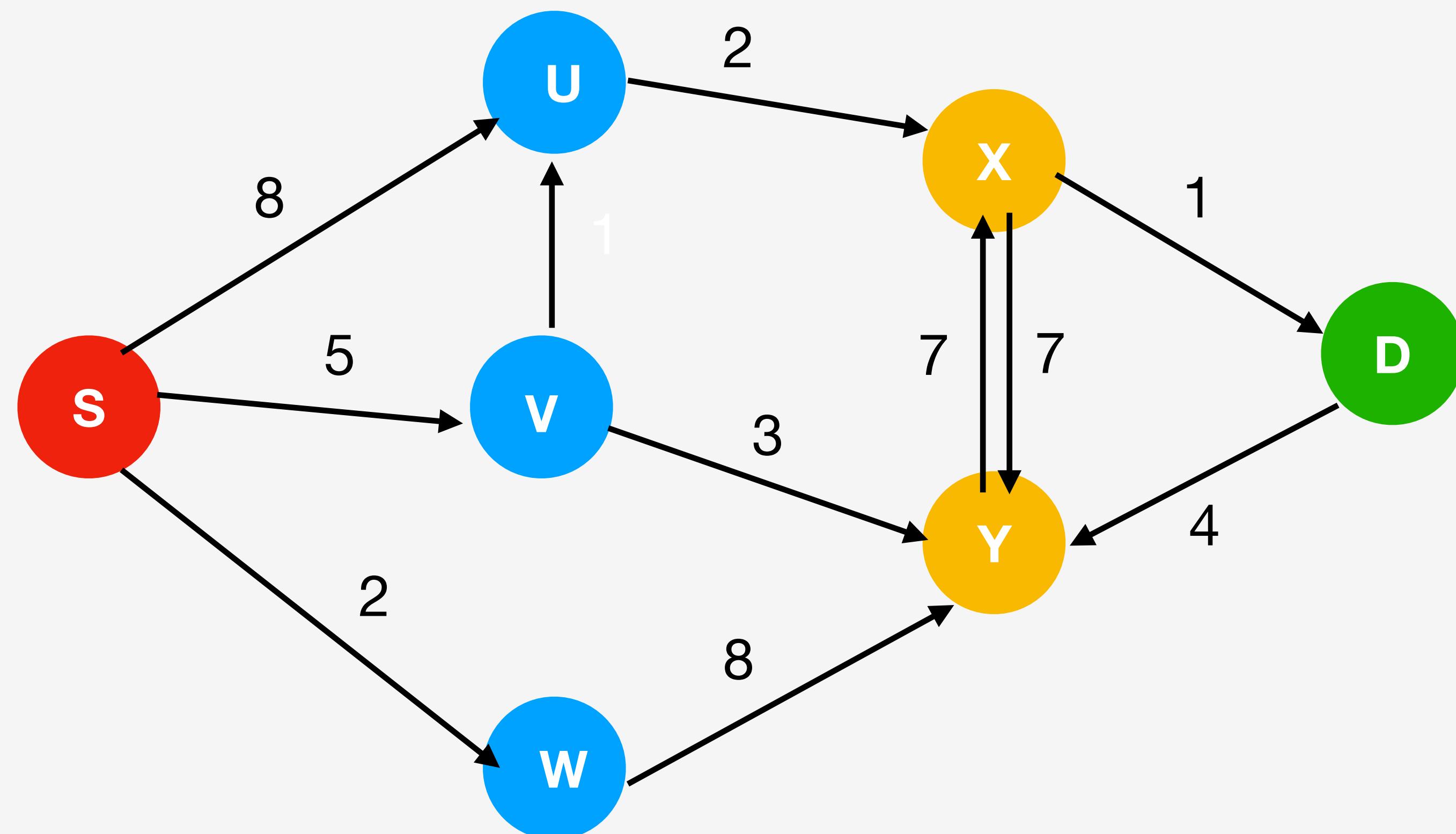
- More annotations/ Semantically richer
- Needed for Replanning

# Mission planning



- How to go from start to destination
- Focus: Speed limits, traffic flow rates, road closures
- No visibility to dynamics and obstacles
- Directed graph can represent road-network
  - One-ways may be present

# Shortest path in weighted graphs



Road network = Topological graph

Graph  $G = (V, E)$

What is the shortest path from  $S$  to  $D$ ?

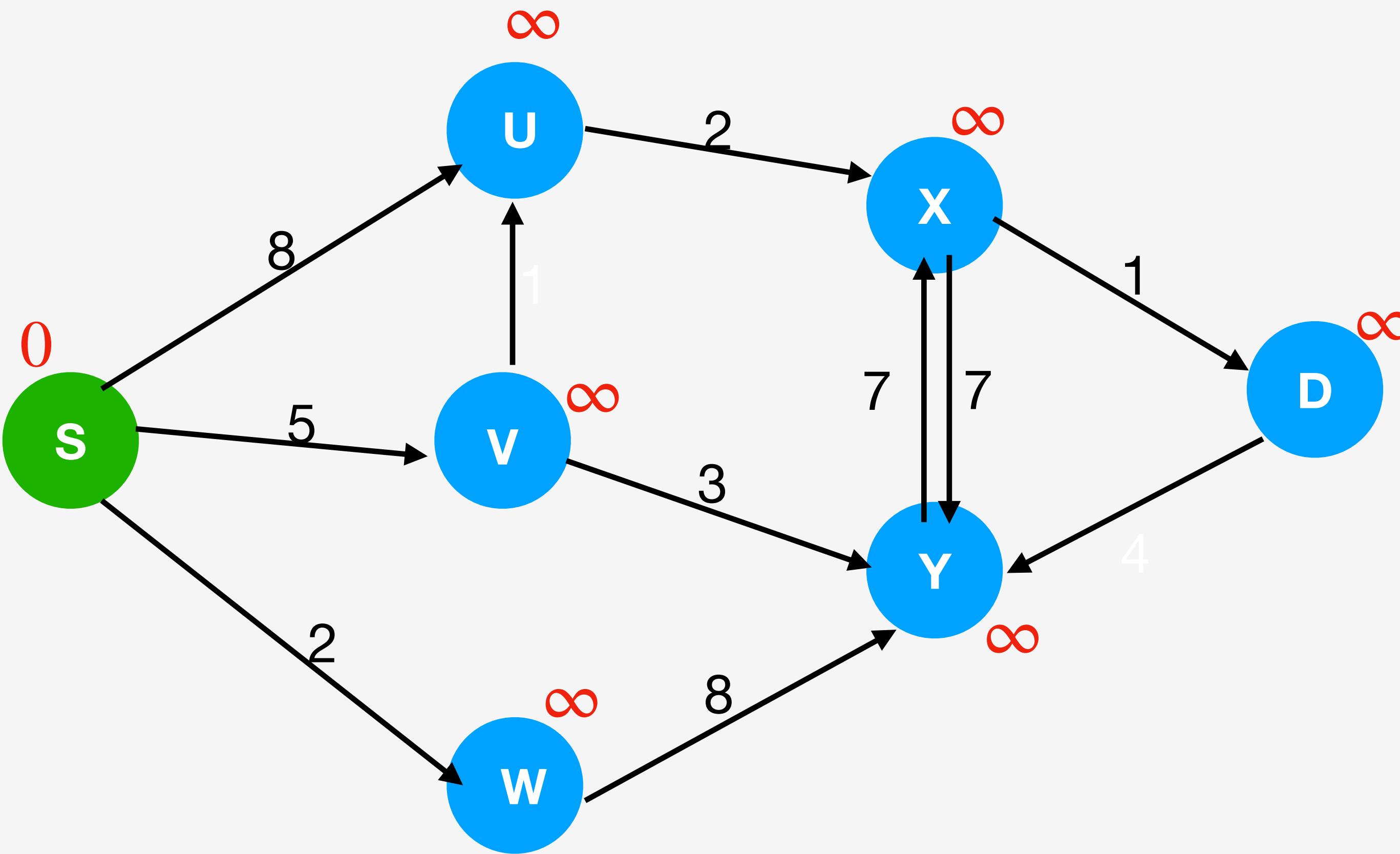
All edges may not have identical costs to traverse

- road terrain conditions
- road length/ time to travel
- pay tolls

Dijkstra's algorithm

Factors in cost of edge traversal

# Dijkstra's algorithm: Iteration 0

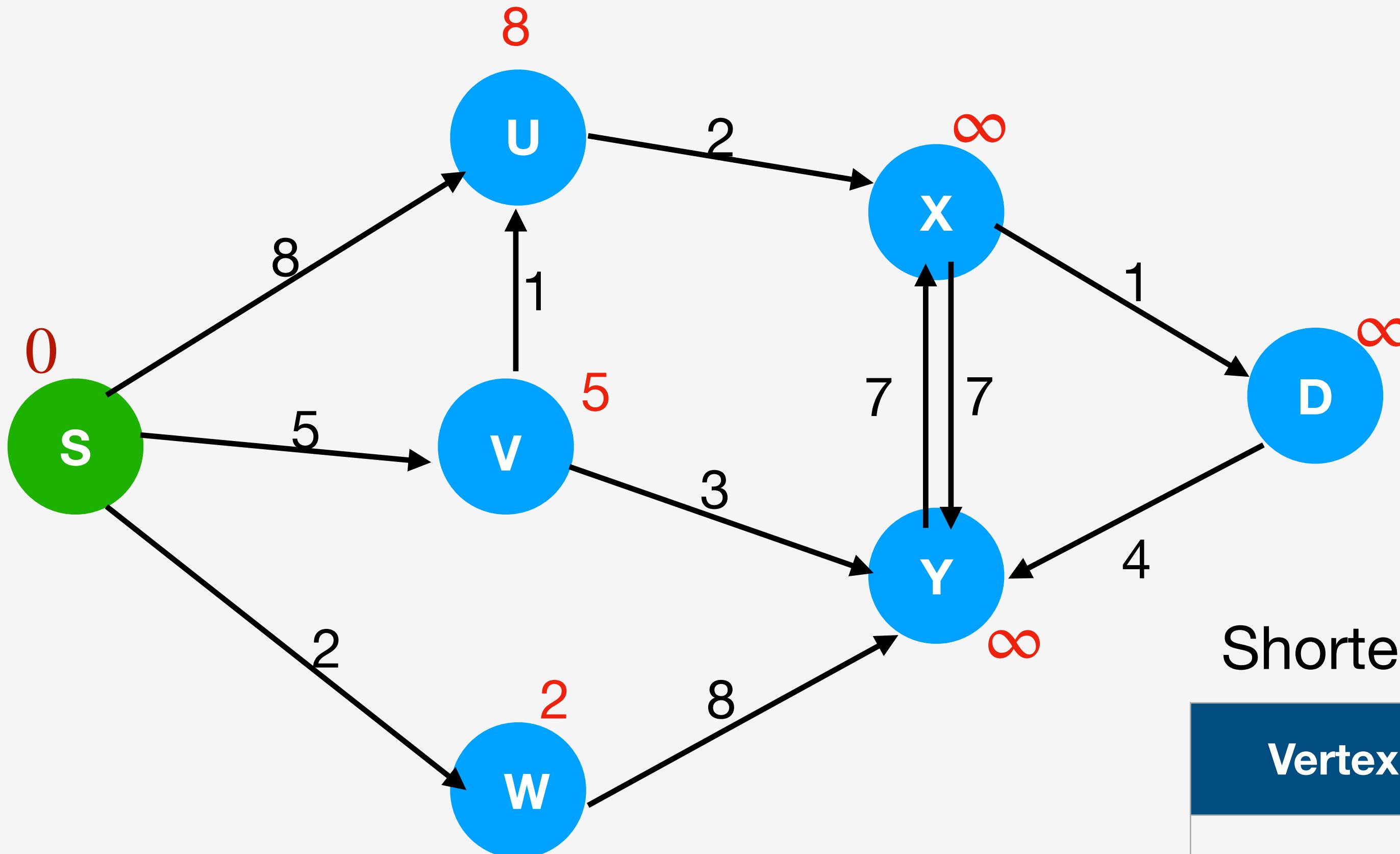


Shortest path set: empty

# Dijkstra's algorithm: Iteration 1

Min cost table

Vertex	Min-cost
U	8
V	5
W	2



Shortest path table

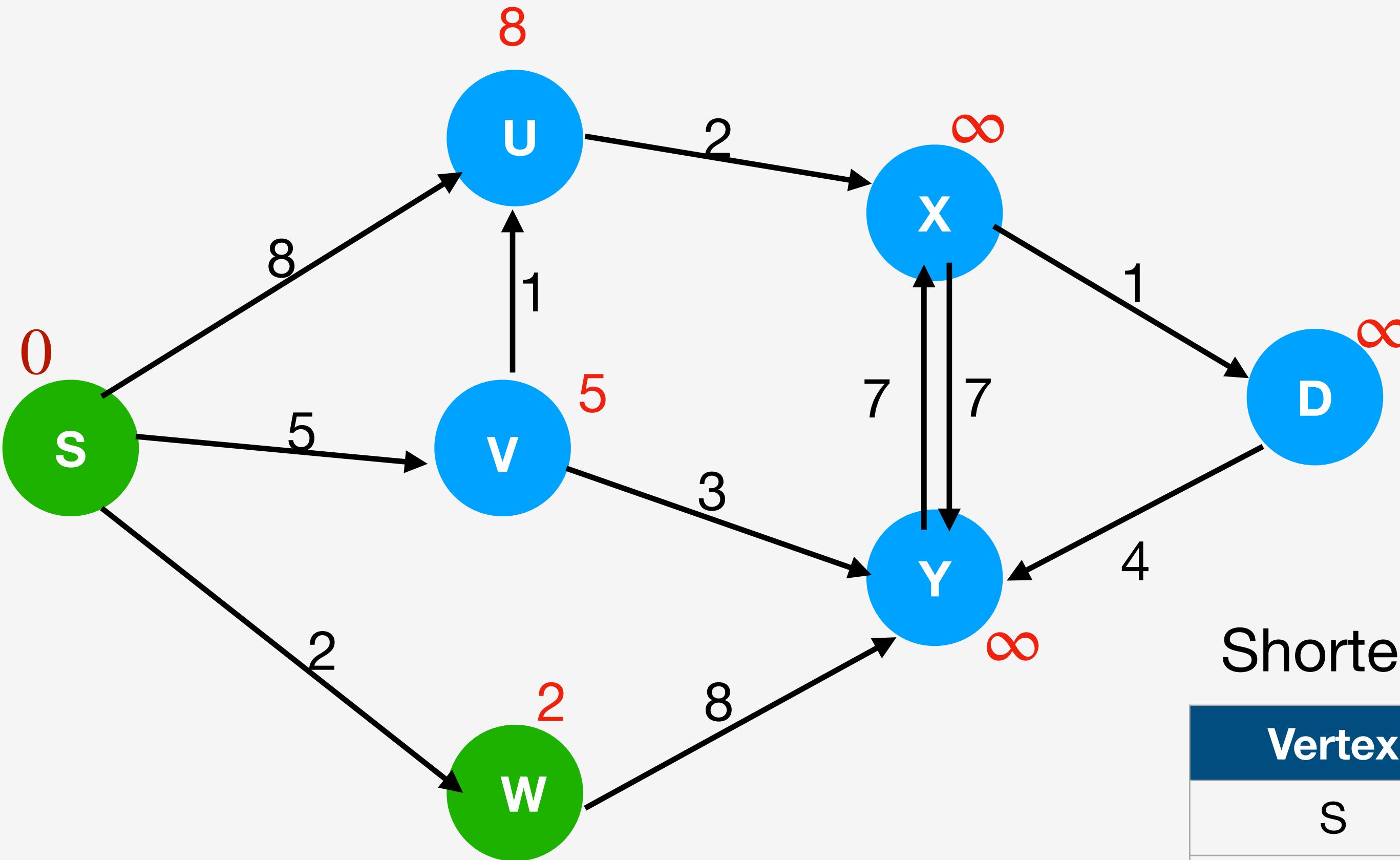
Vertex	Predecessor
S	-

Min cost of U, V and W got updated after S was promoted to Shortest path set

# Dijkstra's algorithm: Iteration 2

Min cost table

Vertex	Min-cost
U	8
V	5
Y	10



Shortest path table

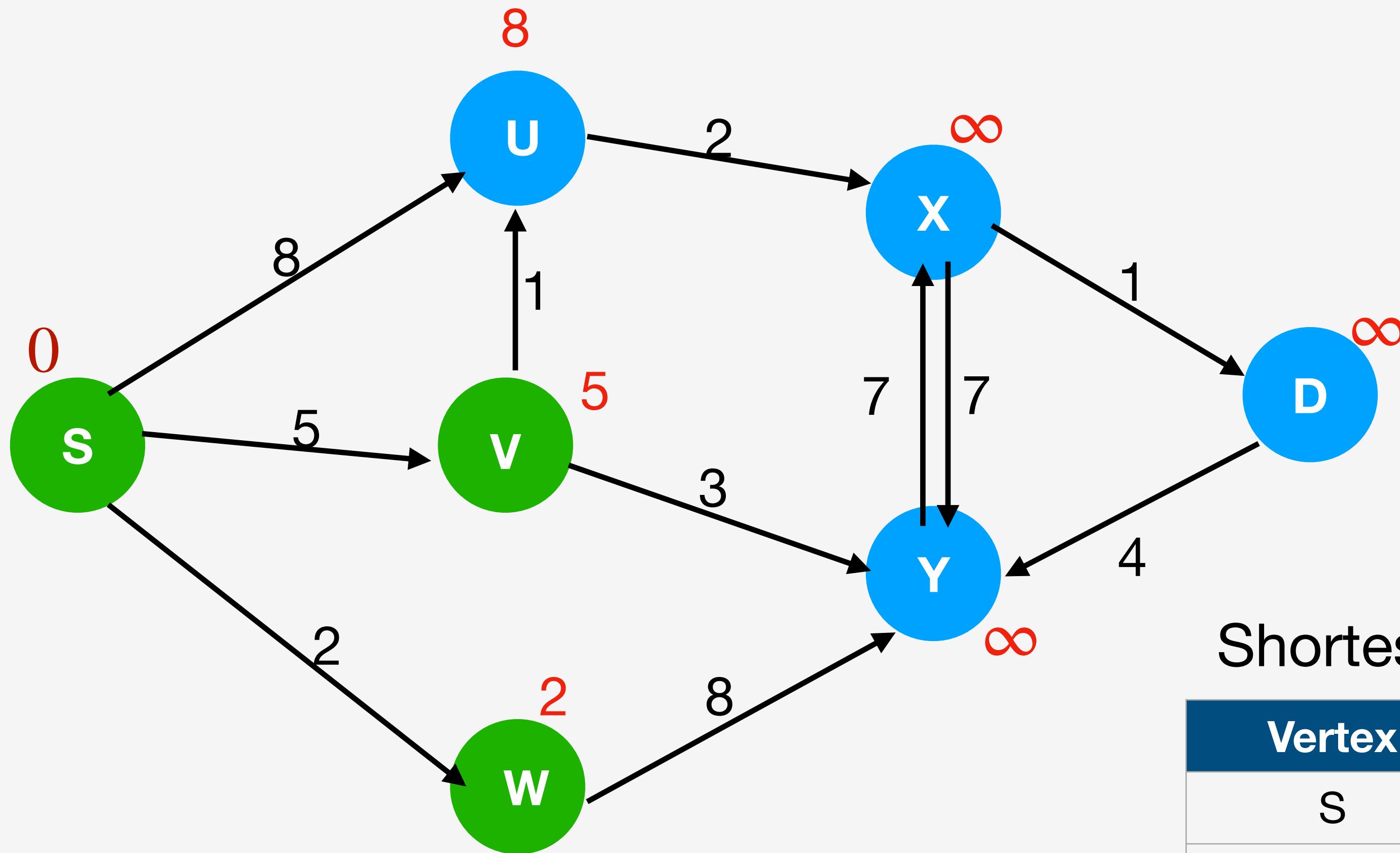
Vertex	Predecessor
S	-
W	S

Min cost of Y got updated after W was promoted to Shortest path set

# Dijkstra's algorithm: Iteration 3

Min cost table

Vertex	Min-cost
U	6
Y	8



Min cost of U, Y got updated after V was promoted to Shortest path set

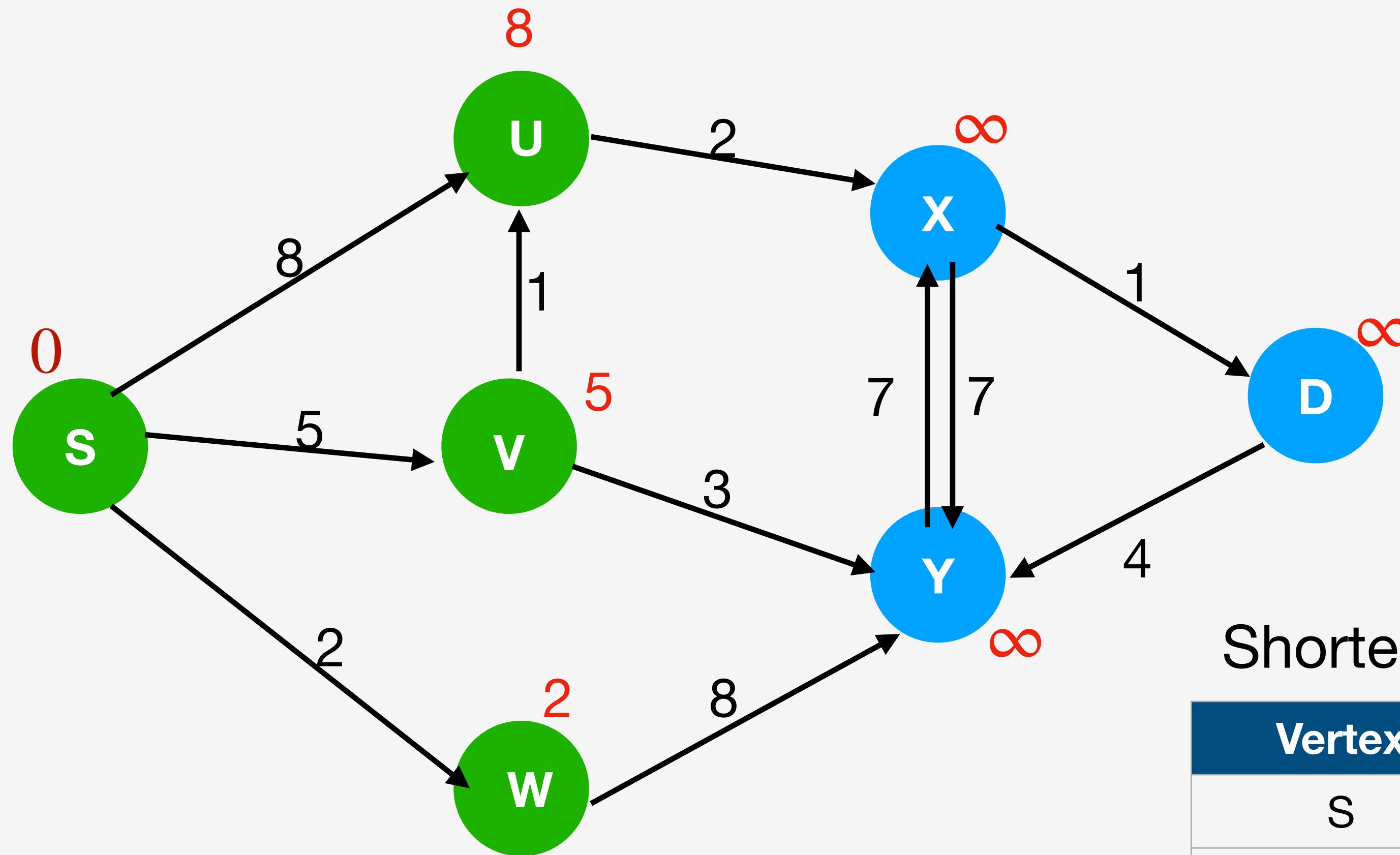
Shortest path table

Vertex	Predecessor
S	-
W	S
V	S

# Dijkstra's algorithm: Iteration 4

Min cost table

Vertex	Min-cost
X	8
Y	8



Min cost of X got updated after U was promoted to Shortest path set

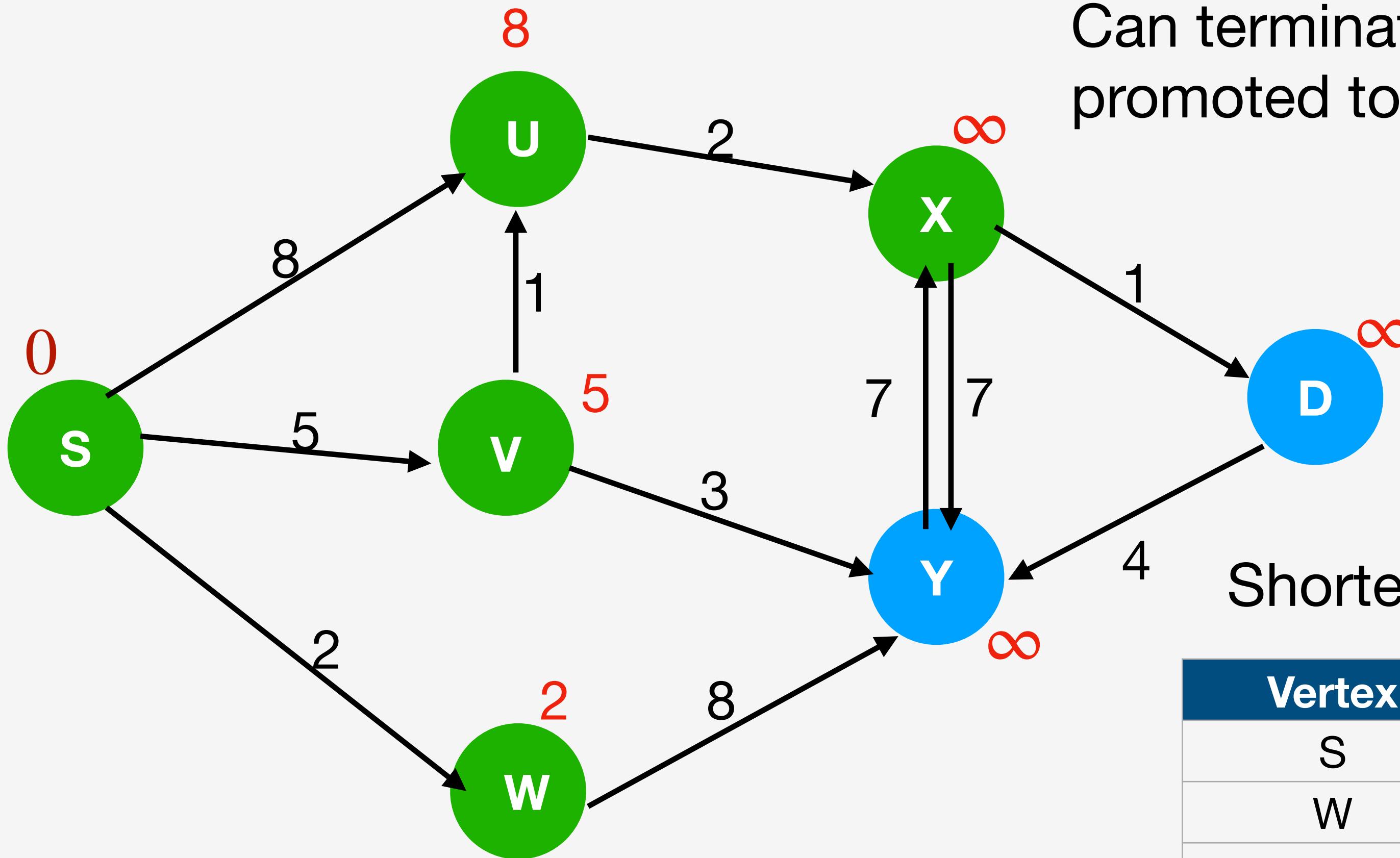
Shortest path table

Vertex	Predecessor
S	-
W	S
V	S
U	V

# Dijkstra's algorithm: Iteration 5

Min cost table

Vertex	Min-cost
D	9
Y	8



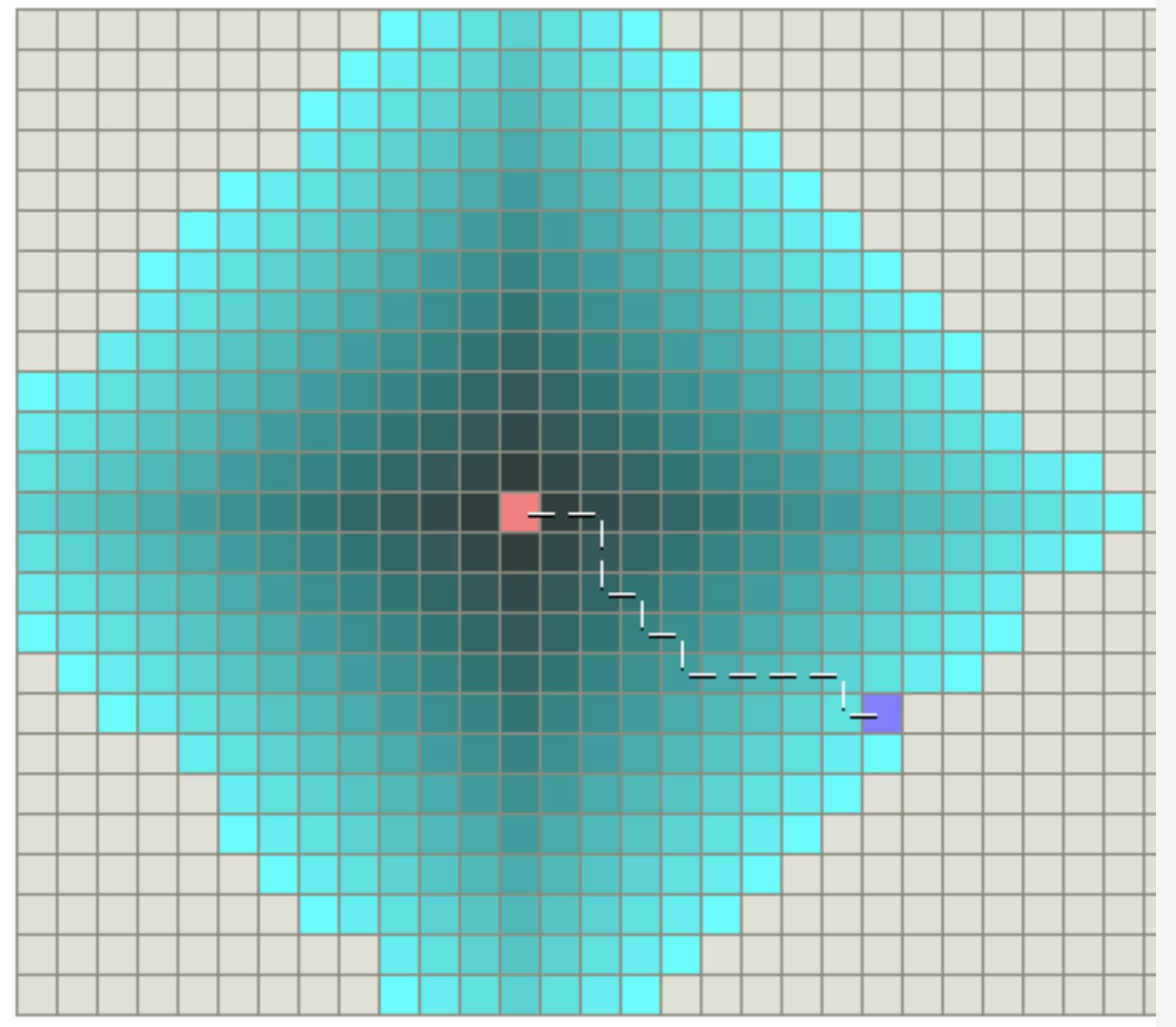
Can terminate once D is promoted to Shortest path

Vertex	Predecessor
S	-
W	S
V	S
U	V
X	U

Note that min cost of D got updated after X was promoted to Shortest path set

# Dijkstra

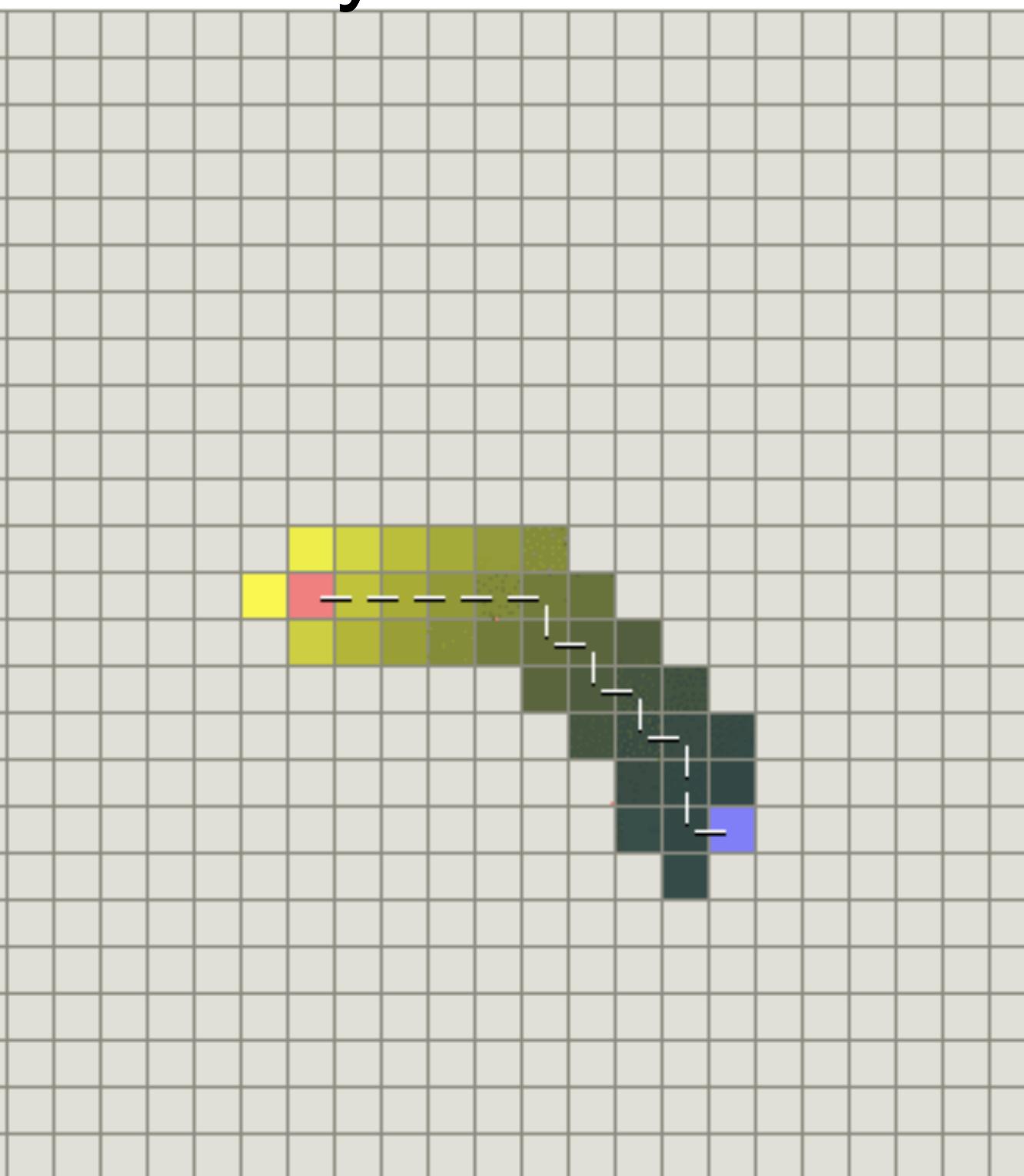
- Outward expansion from source node till destination is reached
- Slow but guaranteed shortest-path



Remember - Convergence only after  
D is promoted to Shortest path

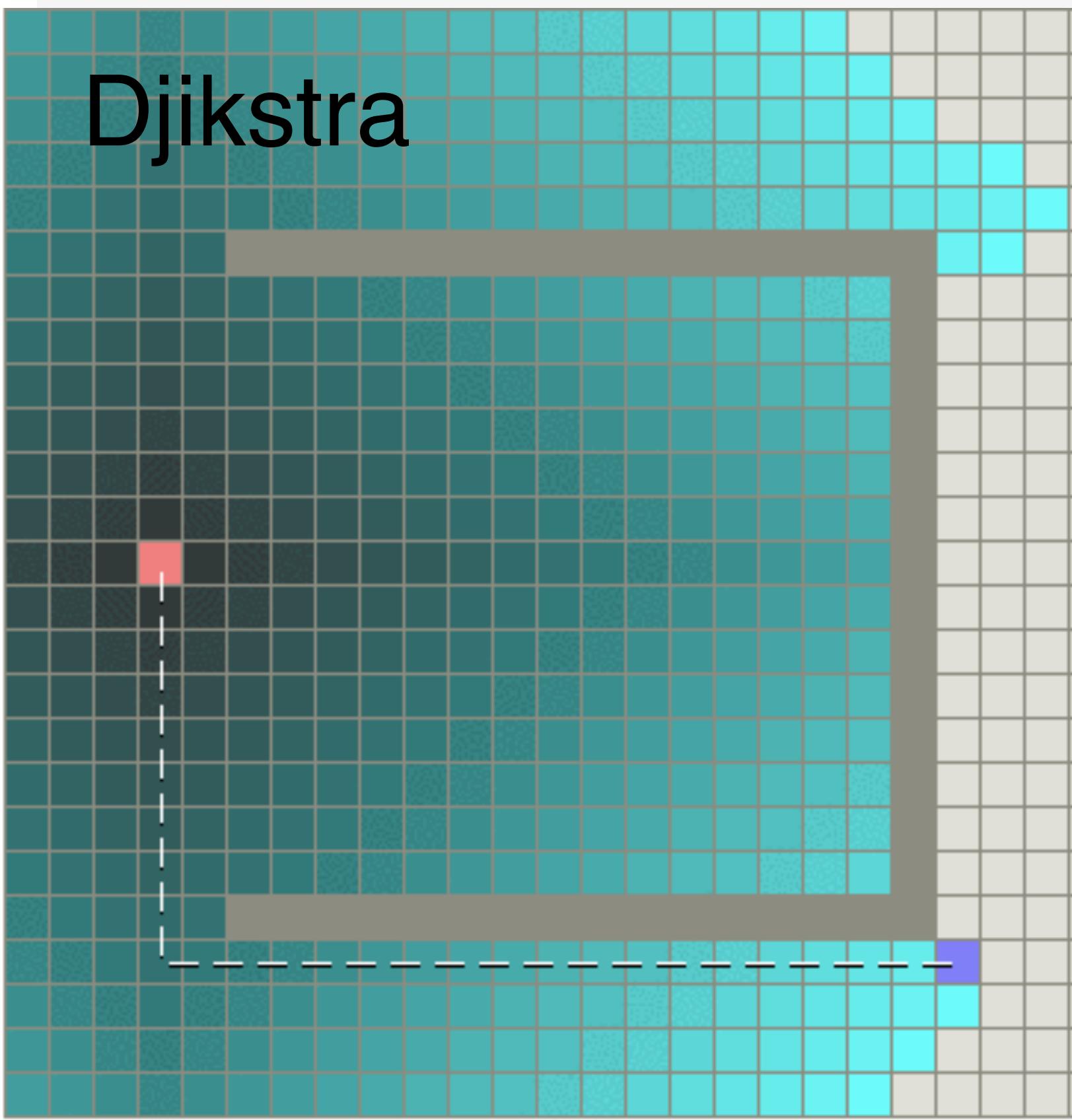
# Greedy best-first

- Select the neighbor which has lowest distance to goal
- Distance to goal = some heuristic function
- Can be very fast

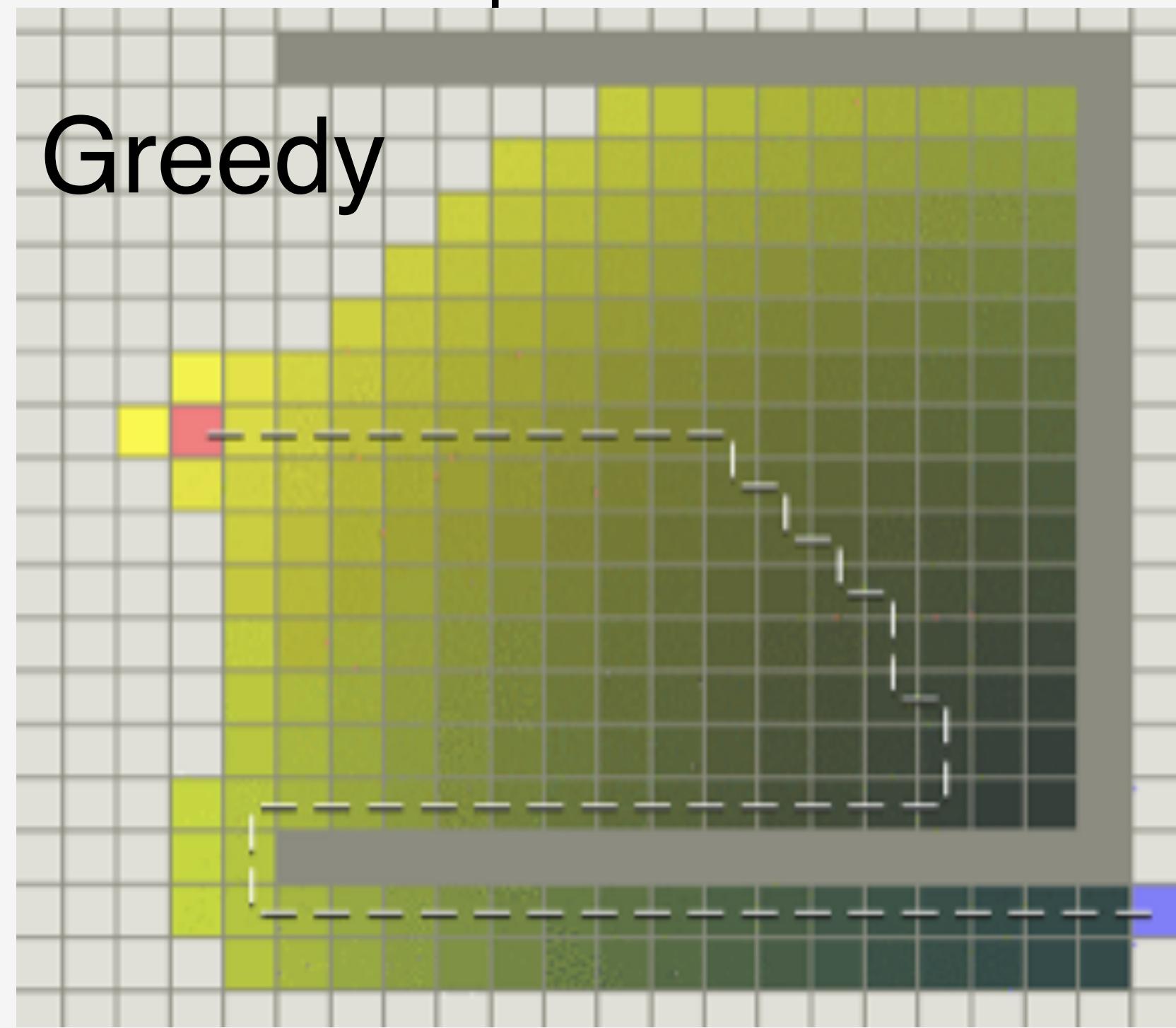


# A\*: Djikstra + Greedy best-first

- Optimal paths
- Slow-ish for straight-forward courses
- Fast for straight-forward cases
- Not so desirable paths in face of some convex obstacles



Djikstra



Greedy

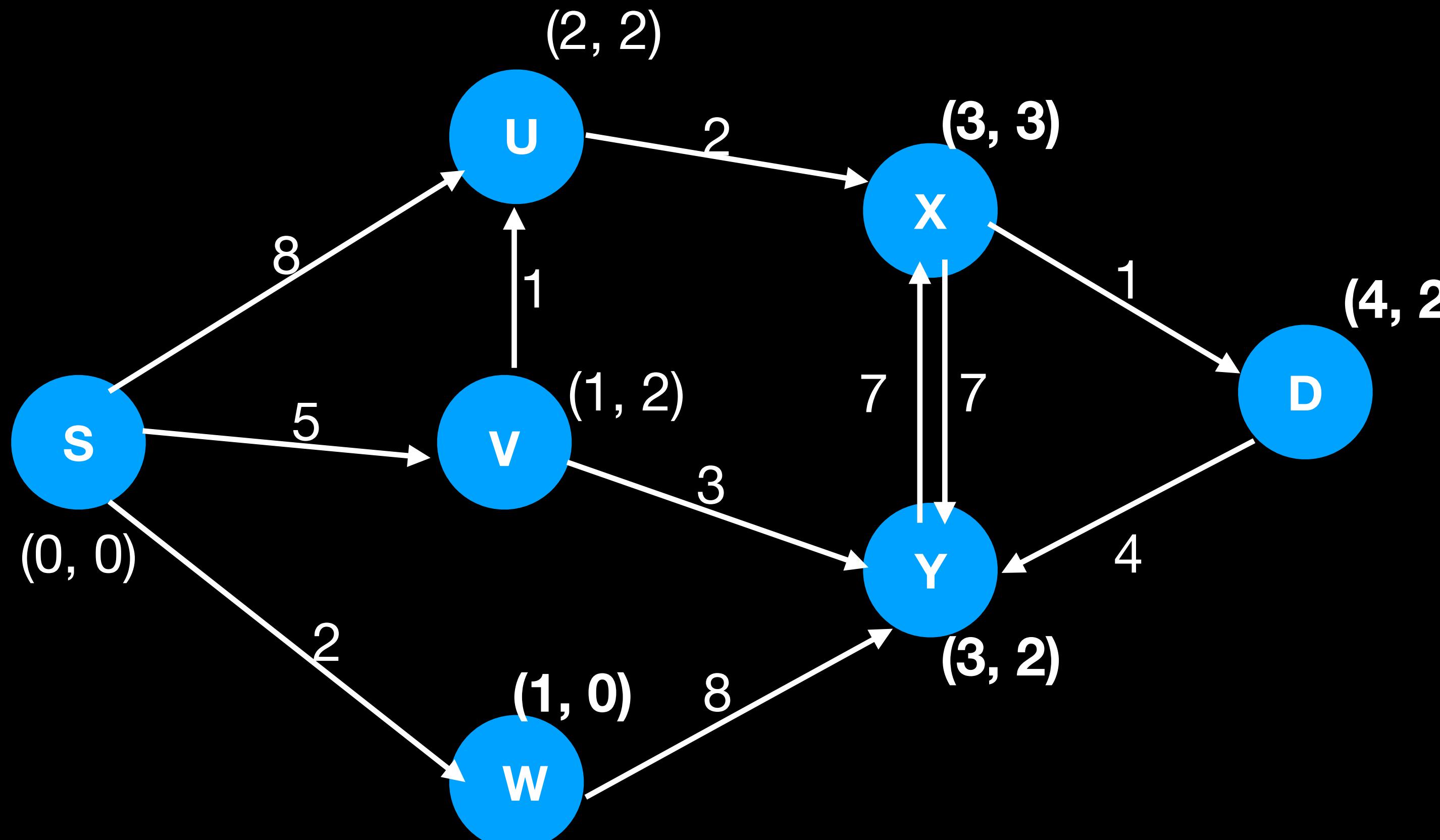
**A\* = combines the best features**

**Node Cost = Cost to Source +  
Heuristic cost to Destination**

$$g(n) = c(n) + w^*h(n)$$

**$h(n) = 0 \Rightarrow$  Djikstra**  
 **$h(n) \gg c(n) \Rightarrow$  Greedy**

# A\* search: Iteration 0

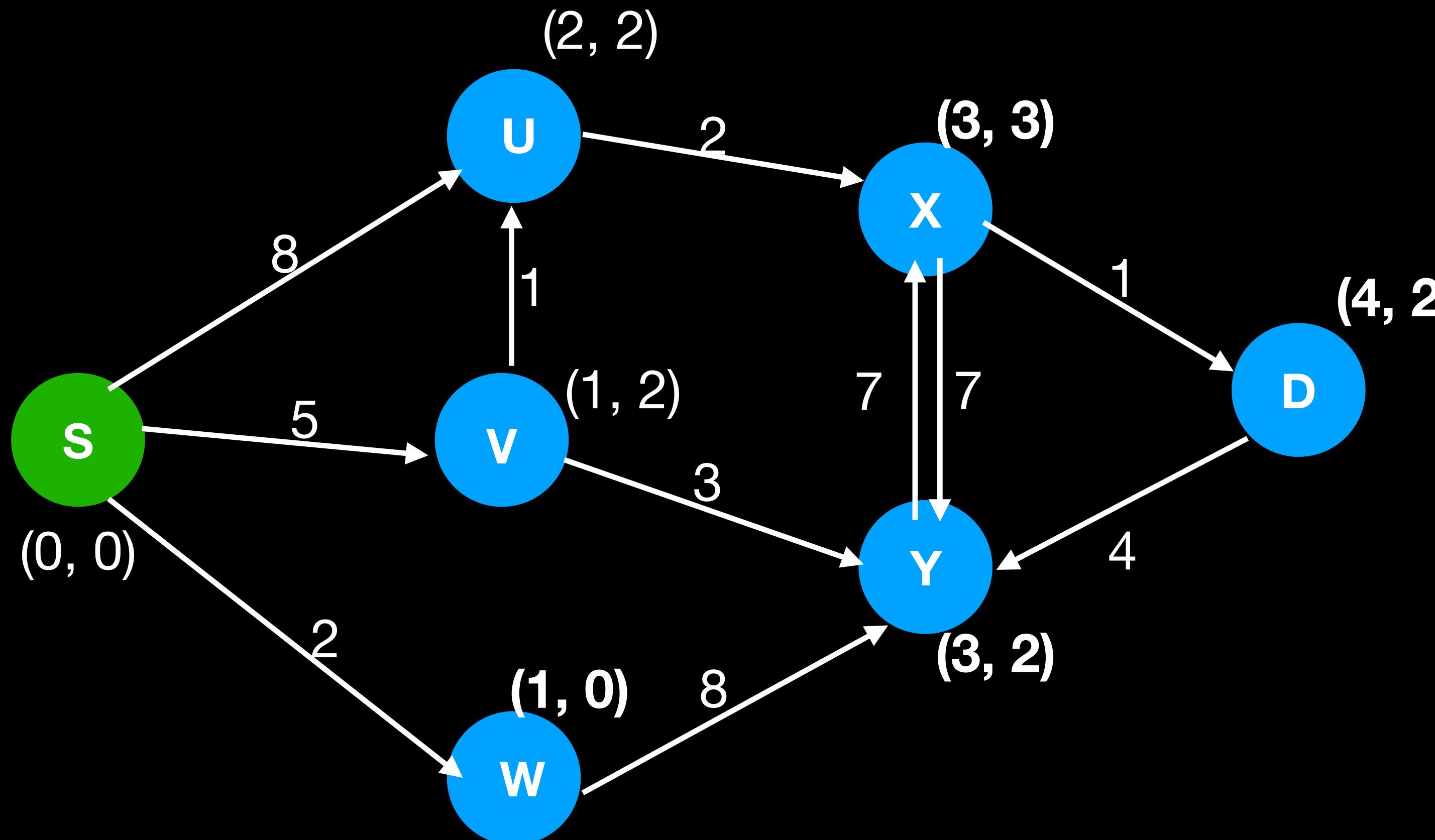


Shortest path set: empty

Vertex	Cost	Heuristic	Total
S	0	4.89	4.89

Simplest heuristic to destination = Euclidean distance  
Can also use time estimate as heuristic

# A\* search: Iteration 1

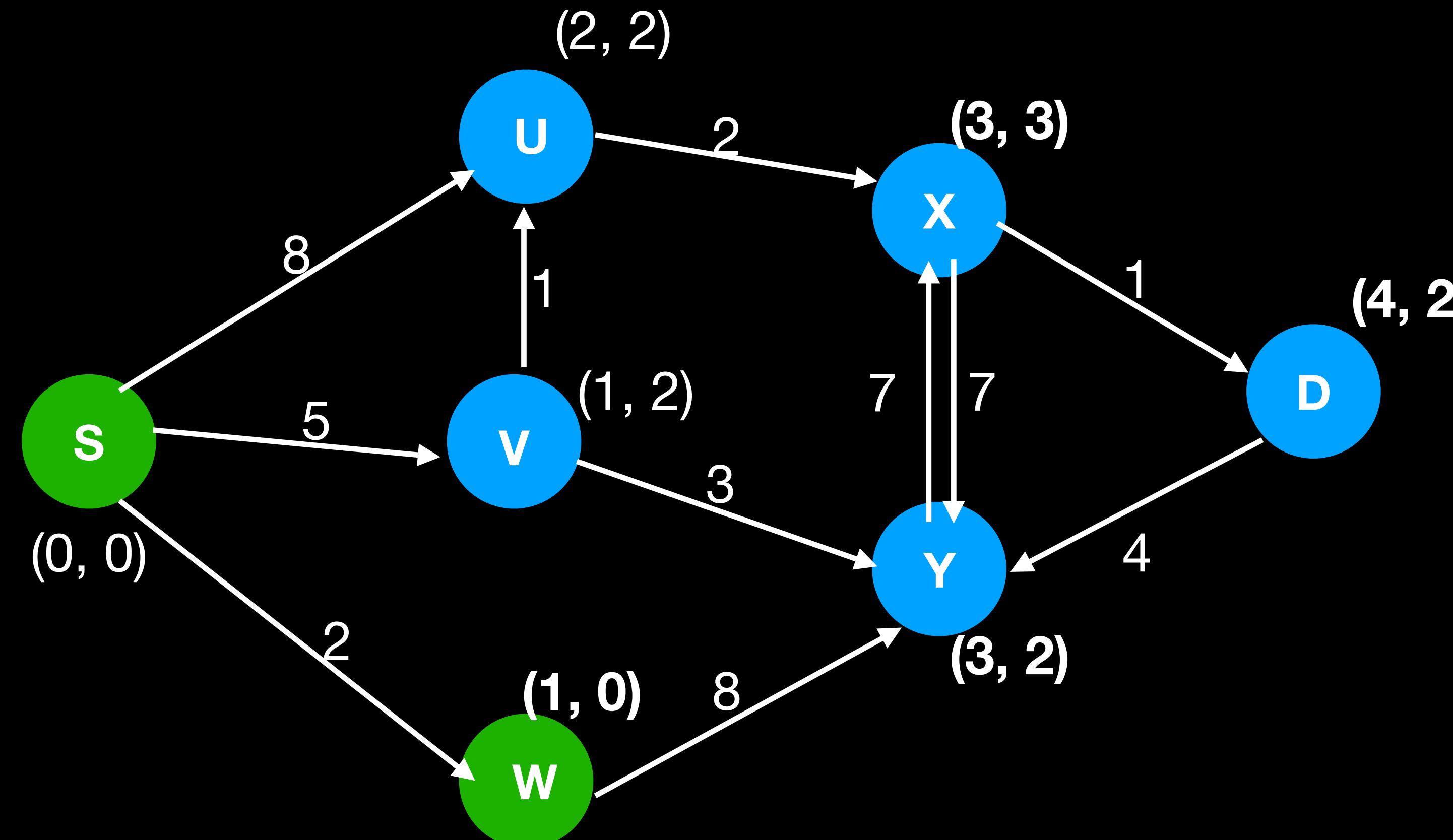


Vertex	Cost	Heuristic	Total
U	8	2	10
V	5	3	8
W	2	3.6	5.6

Shortest path table

Vertex	Predecessor
S	-

# A\* search: Iteration 2

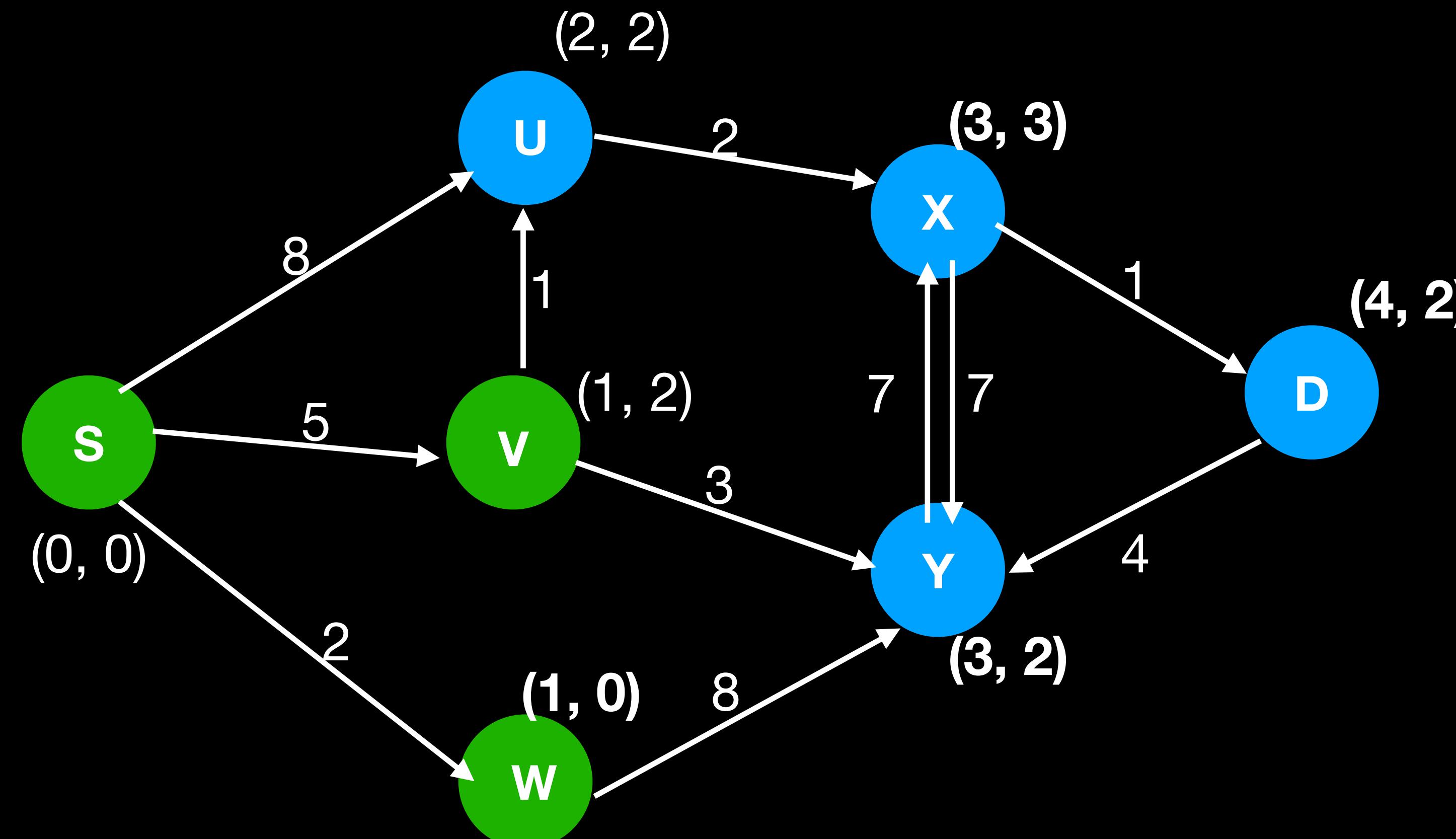


Vertex	Cost	Heuristic	Total
U	8	2	10
V	5	3	8
Y	10	1.414	11.414

Shortest path table

Vertex	Predecessor
S	-
W	S

# A\* search: Iteration 3

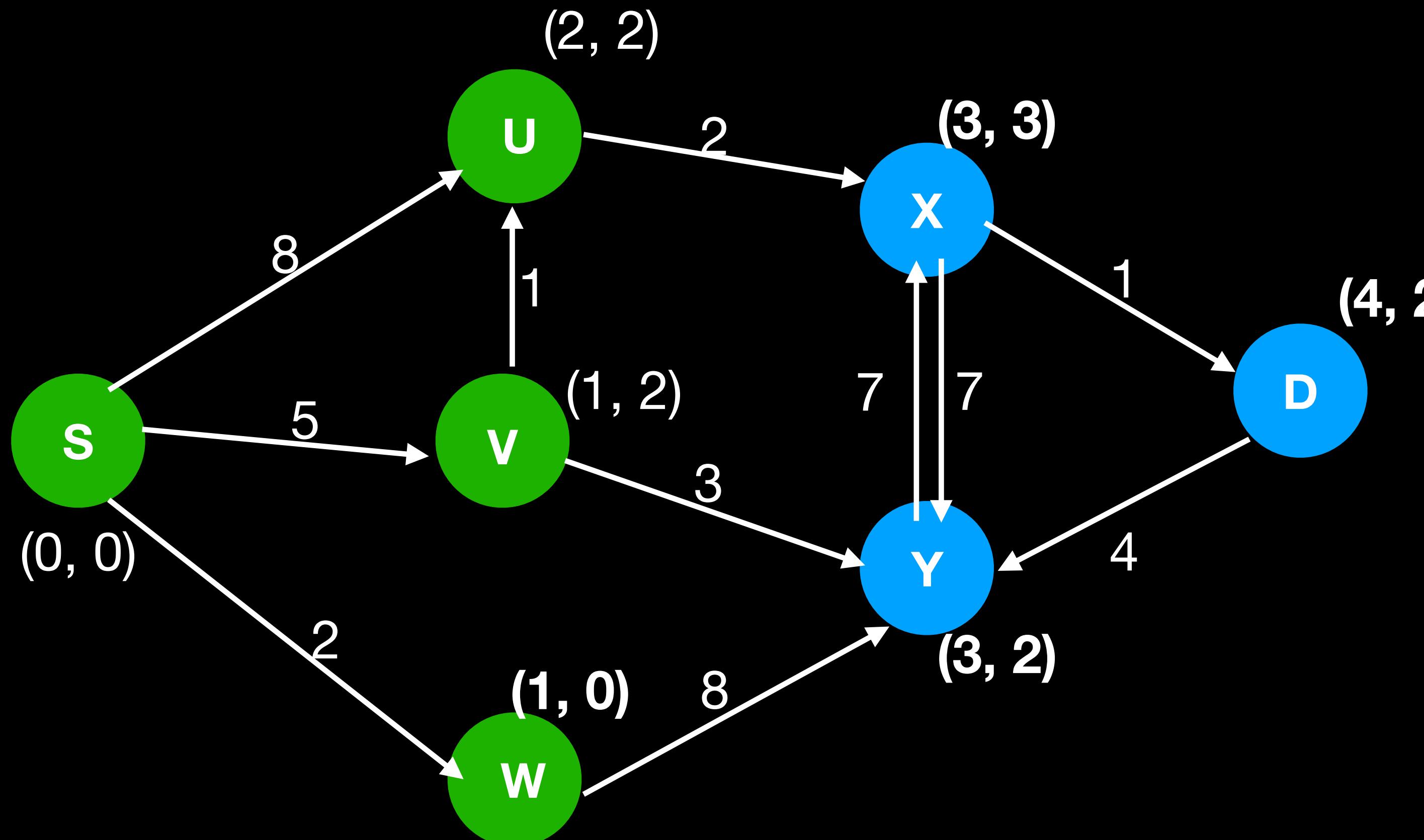


Vertex	Cost	Heuristic	Total
U	6	2	8
Y	8	1	9

Shortest path table

Vertex	Predecessor
S	-
W	S
V	S

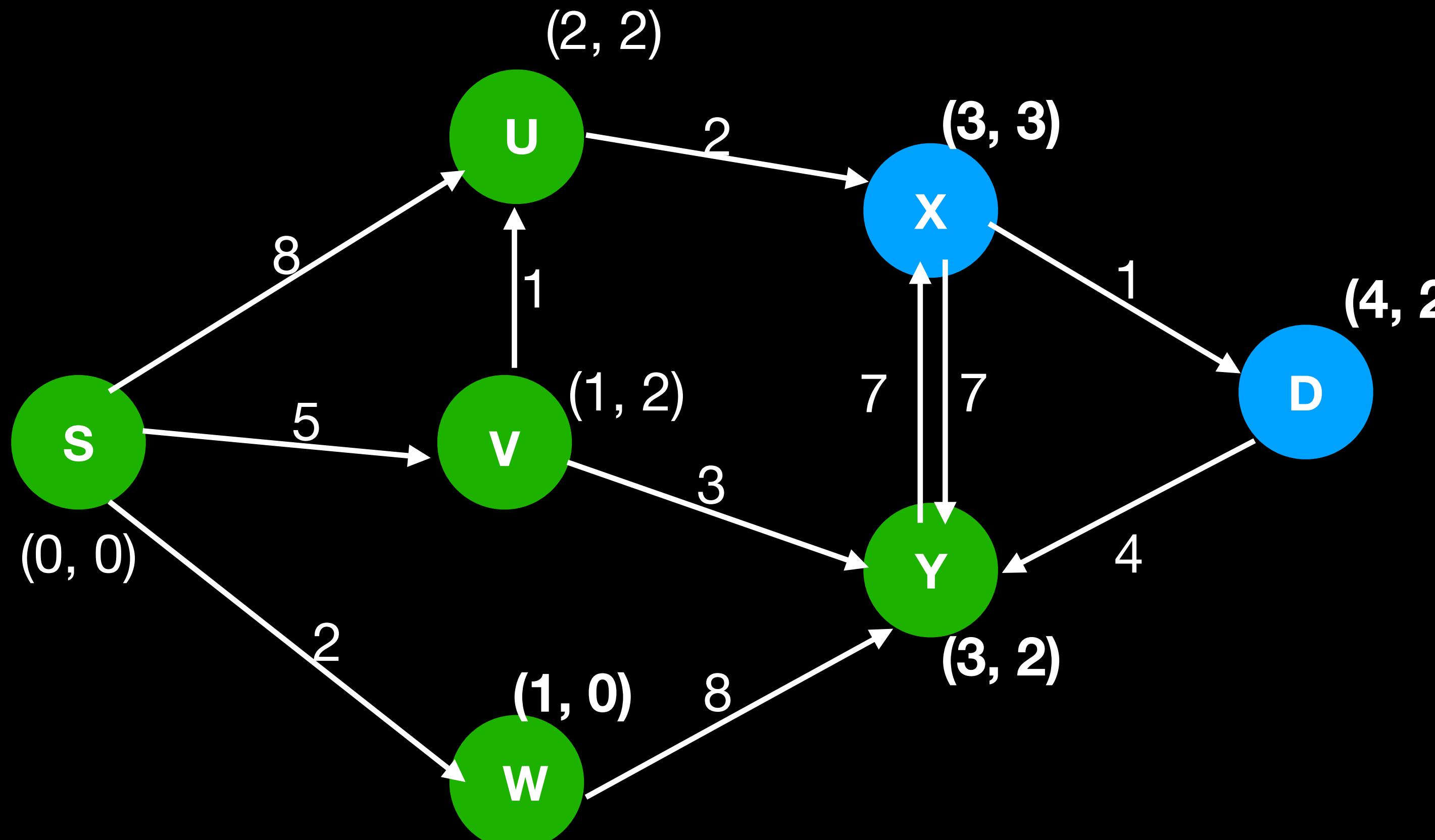
# A\* search: Iteration 4



Shortest path table

Vertex	Predecessor
S	-
W	S
V	S
U	V

# A\* search: Iteration 4

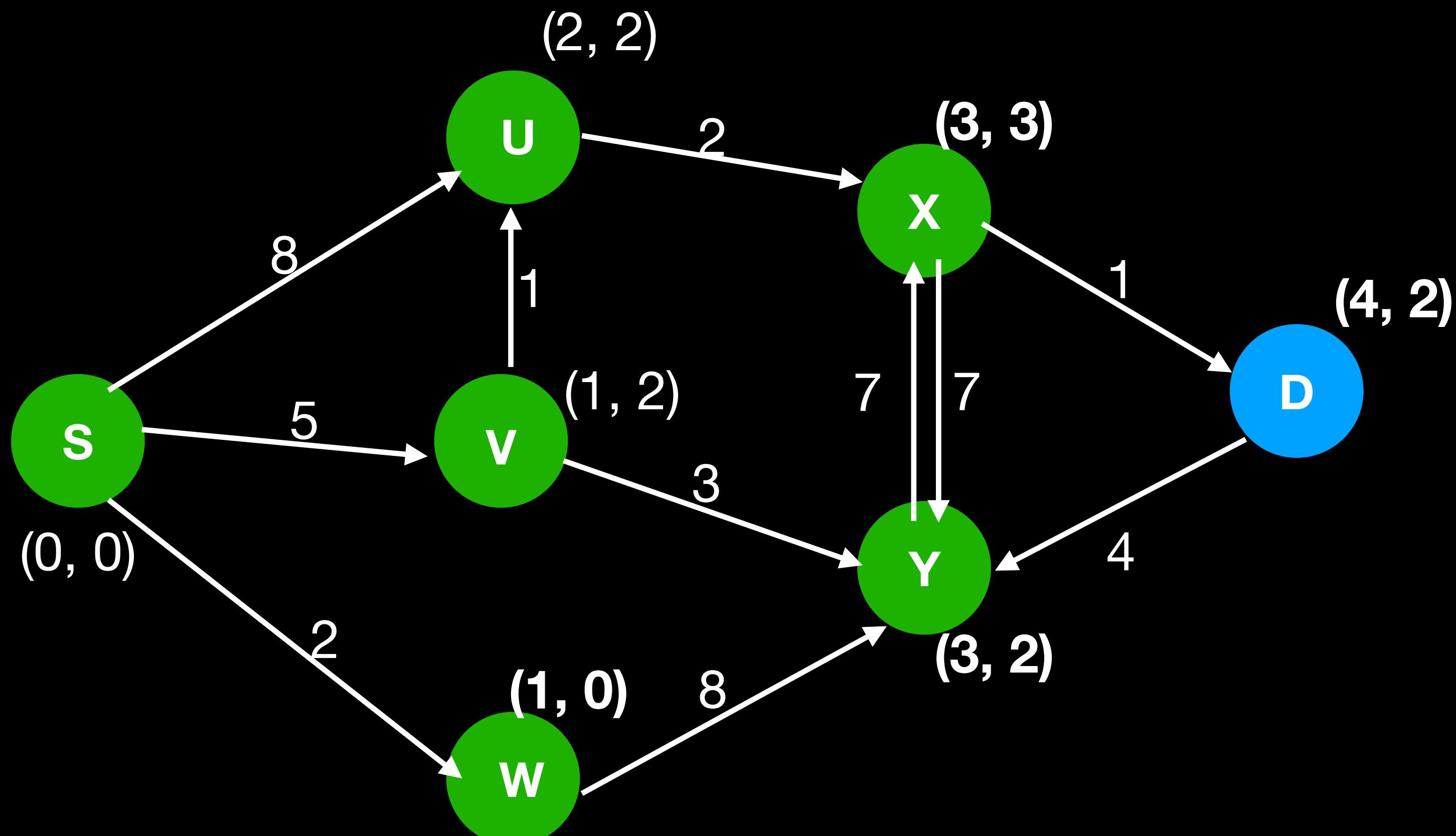


Vertex	Cost	Heuristic	Total
X	8	1.414	9.414

Shortest path table

Vertex	Predecessor
S	-
W	S
V	S
U	V
Y	V

# A\* search: Iteration 4



Vertex	Cost	Heuristic	Total
D	8	1	9

Shortest path table

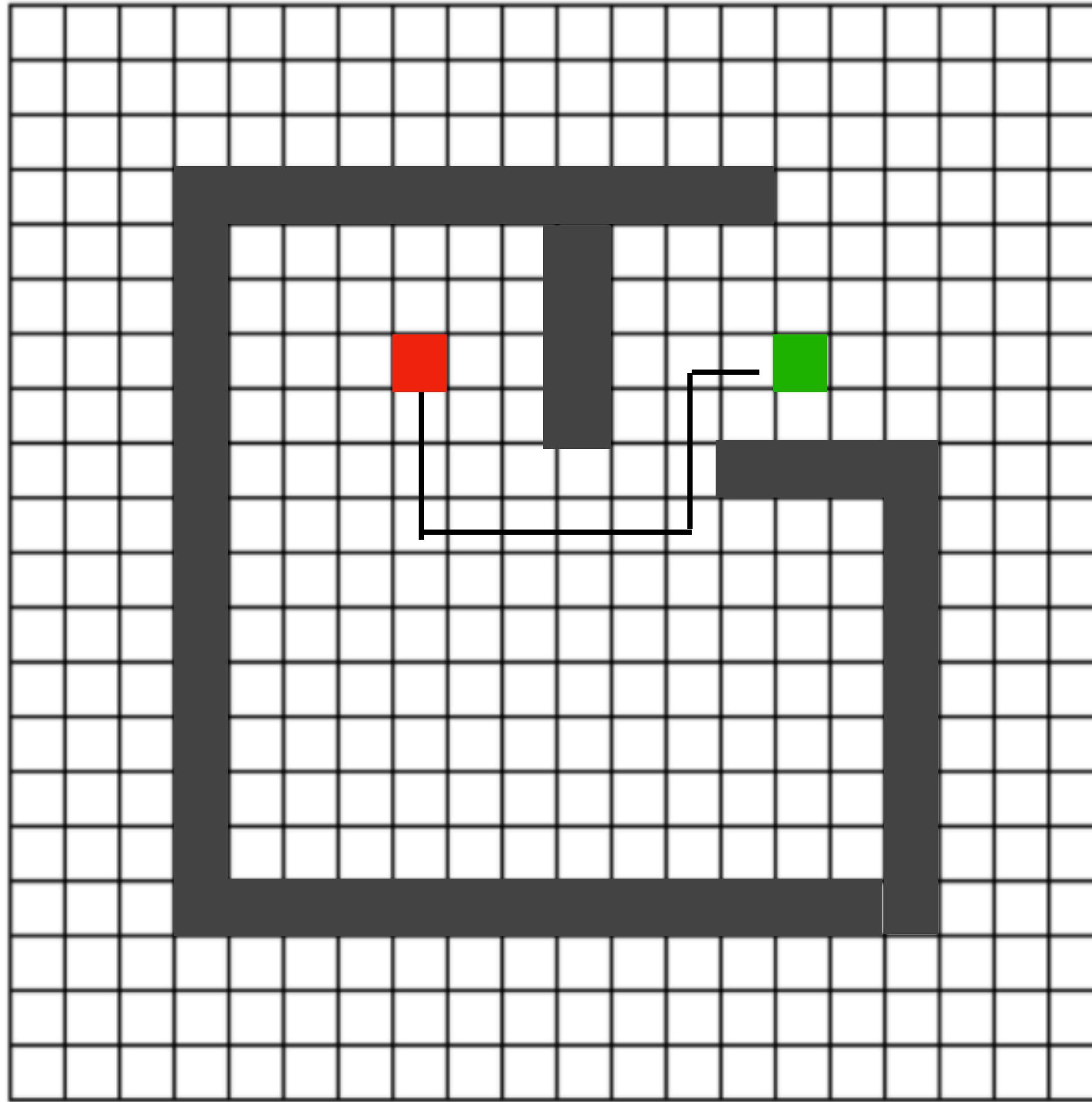
Vertex	Predecessor
S	-
W	S
V	S
U	V
Y	V
X	U

# Lab exercise: Mission-planning in city maps



- Open-street maps
- OSNX package to load map into a graph
- Small cities ~2000 edges, 5000 vertices
- Lat-Long based node query

# Occupancy grids to graphs



- Occupancy grid is a common representation of the environment for autonomous robots
- Every free cell is a vertex
- Edges to only adjacent free cells
- Occupied cells are not added
- Easy to convert from occupancy grid to graph

# Rapidly-exploring random trees (RRT)

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## Djikstra/ A\*

- Accepts pre-built graph with vertices/ edges
- Computes the shortest path between 2 nodes

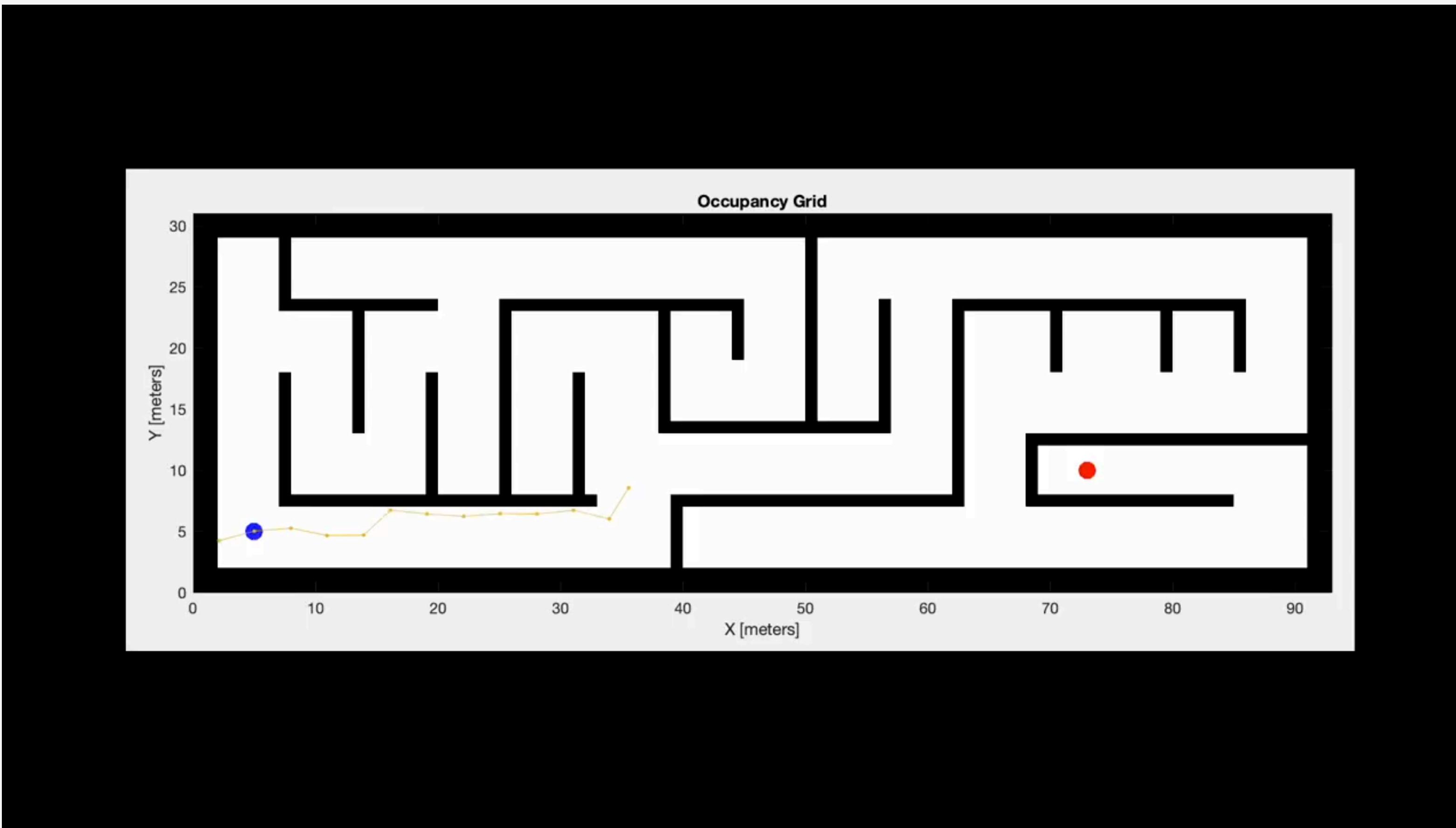
How to construct the graph?

What happens in case of high-dimensional grids?

## RRT

- Randomly generate a node
- Find closest node in current graph
  - Method to madness: insert within some known distance of the nearest node
- Add edge to closest node
- Node/ Edge should not collide with obstacle
- Rinse and repeat

# RRT\* visualization



Source: <https://www.youtube.com/watch?v=QR3U1dgc5RE>

# RRT\* pseudo-code

$G(V,E) = \{V = [S], E = [ ]\}$

Iterate over n steps

$X_{new} = \text{FindRandomPosition}()$

If  $X_{new}$  close to obstacle, try again

$X_{nearest} = \text{Nearest}(G, X_{new})$

$\text{Cost}(X_{new}) = \text{Distance}(X_{new}, X_{nearest})$

$X_{best}, X_{neighbors} = \text{findNeighbors}(G, X_{new}, \text{radius})$

$\text{Link} = \text{Chain}(X_{new}, X_{best})$

For  $x'$  in  $X_{neighbors}$

If  $\text{Cost}(X_{new}) + \text{Distance}(X_{new}, x') < \text{Cost}(x')$

$\text{Cost}(x') = \text{Cost}(X_{new}) + \text{Distance}(X_{new}, x')$

$\text{Parent}(x') = X_{new}$

$G[\text{Edges}] += (X_{new}, x')$

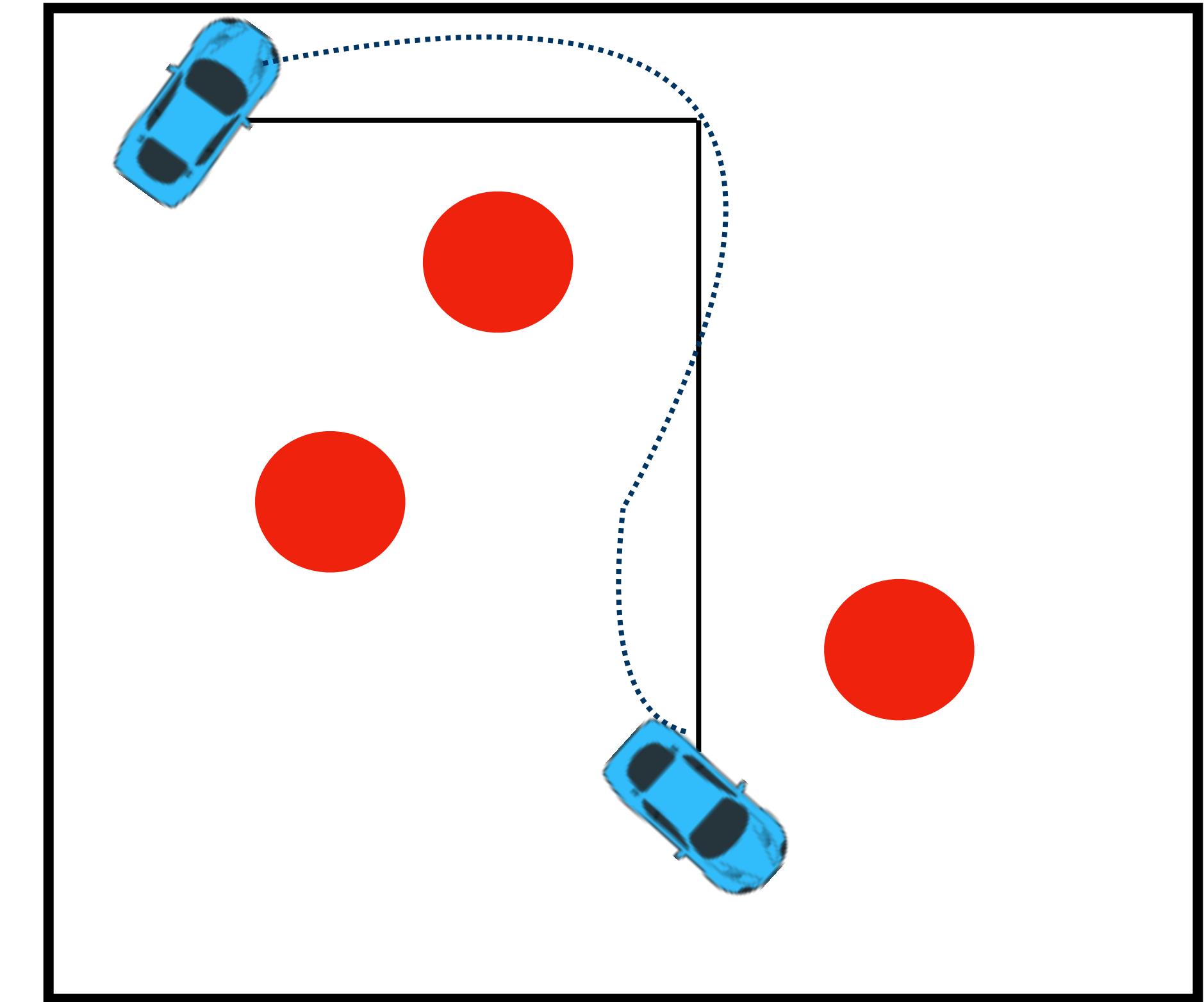
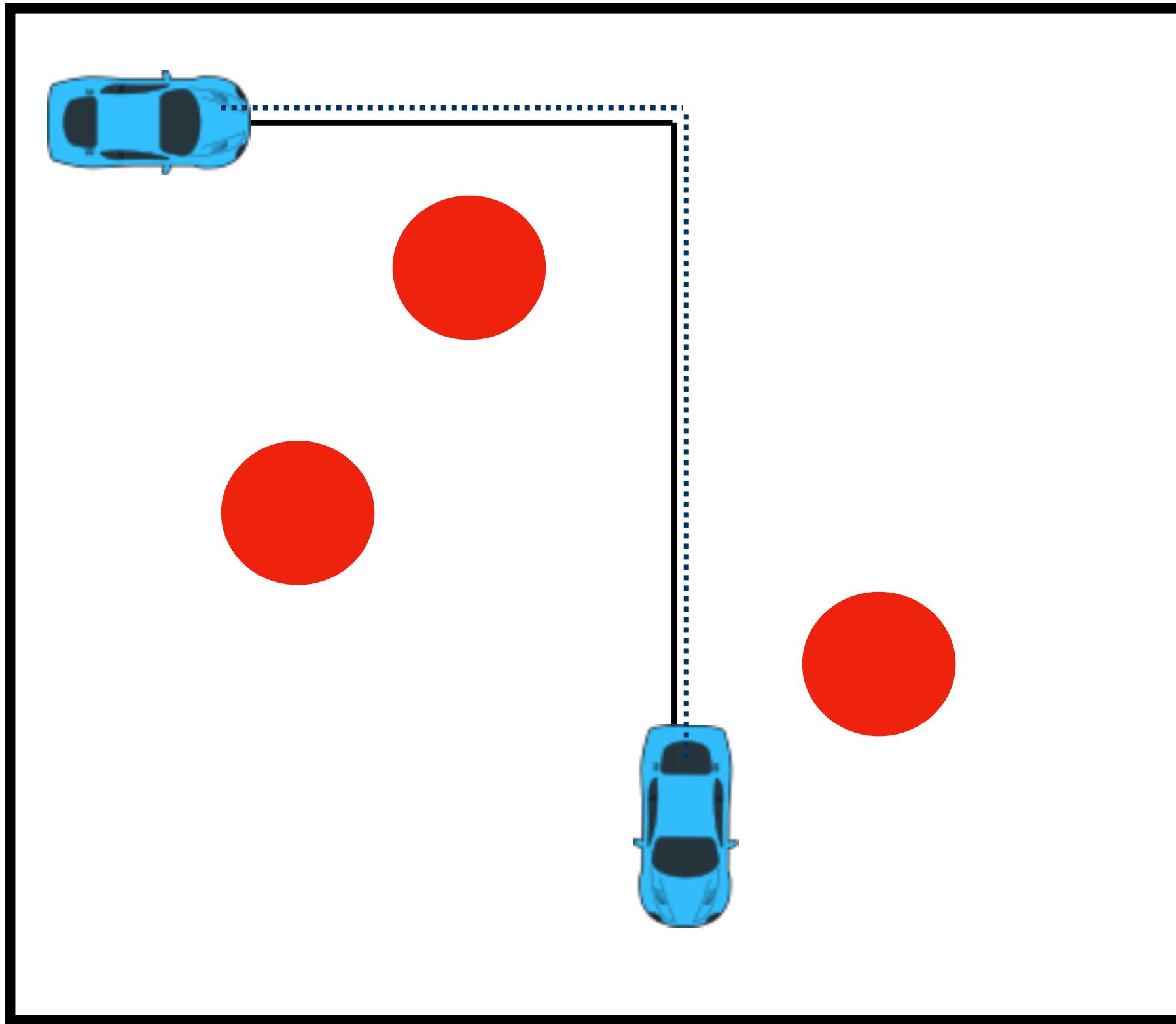
$G[\text{Edges}] += \text{Link}$

Return G

# Lab exercise

<https://theclassytim.medium.com/robotic-path-planning-rrt-and-rrt-212319121378>

# Path planning vs Trajectory planning

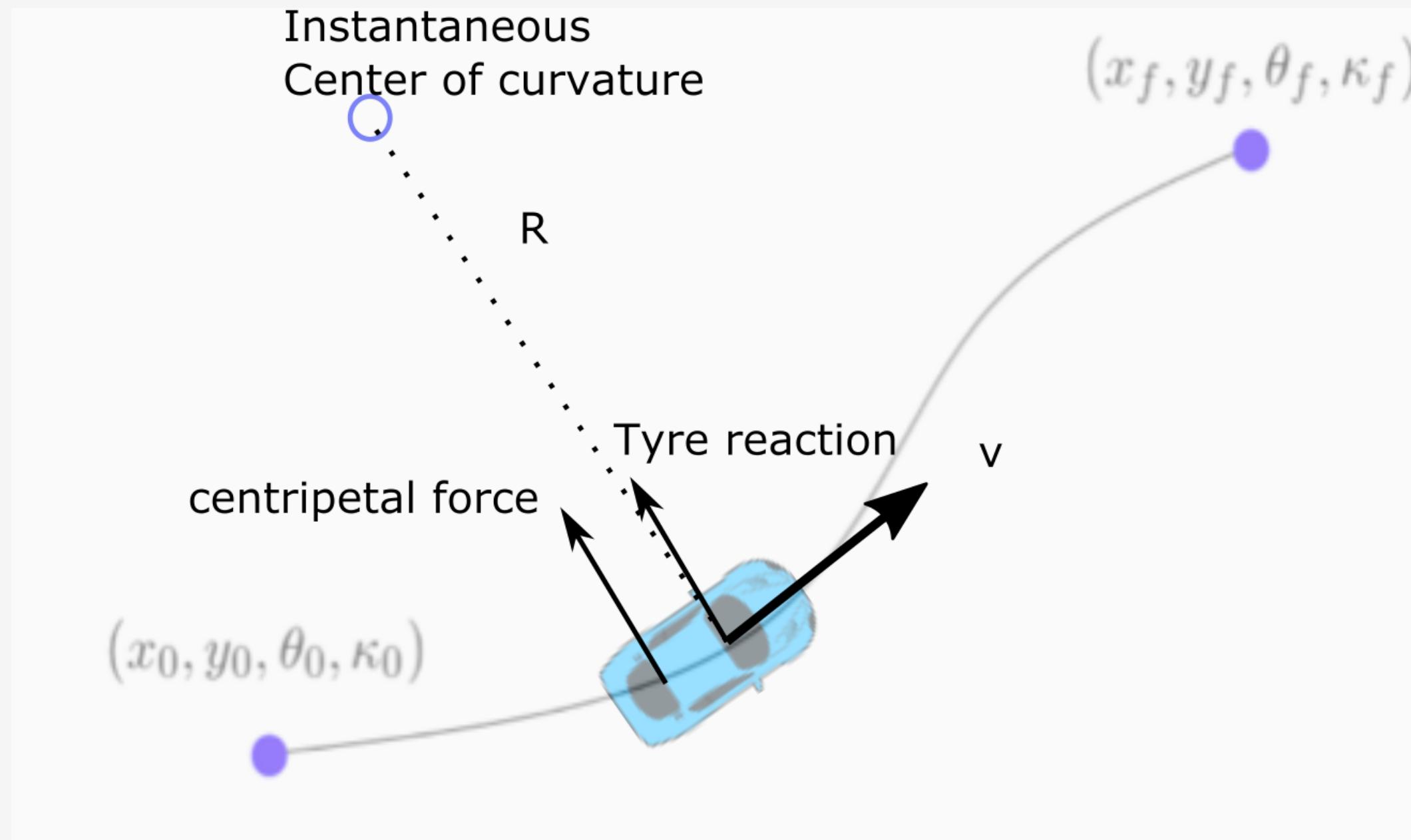


A\*/ Djikstra/ RRT\* are path planners

No notion of robot orientation/ velocities/ acceleration

Does not solve the parking problem

# Path curvature



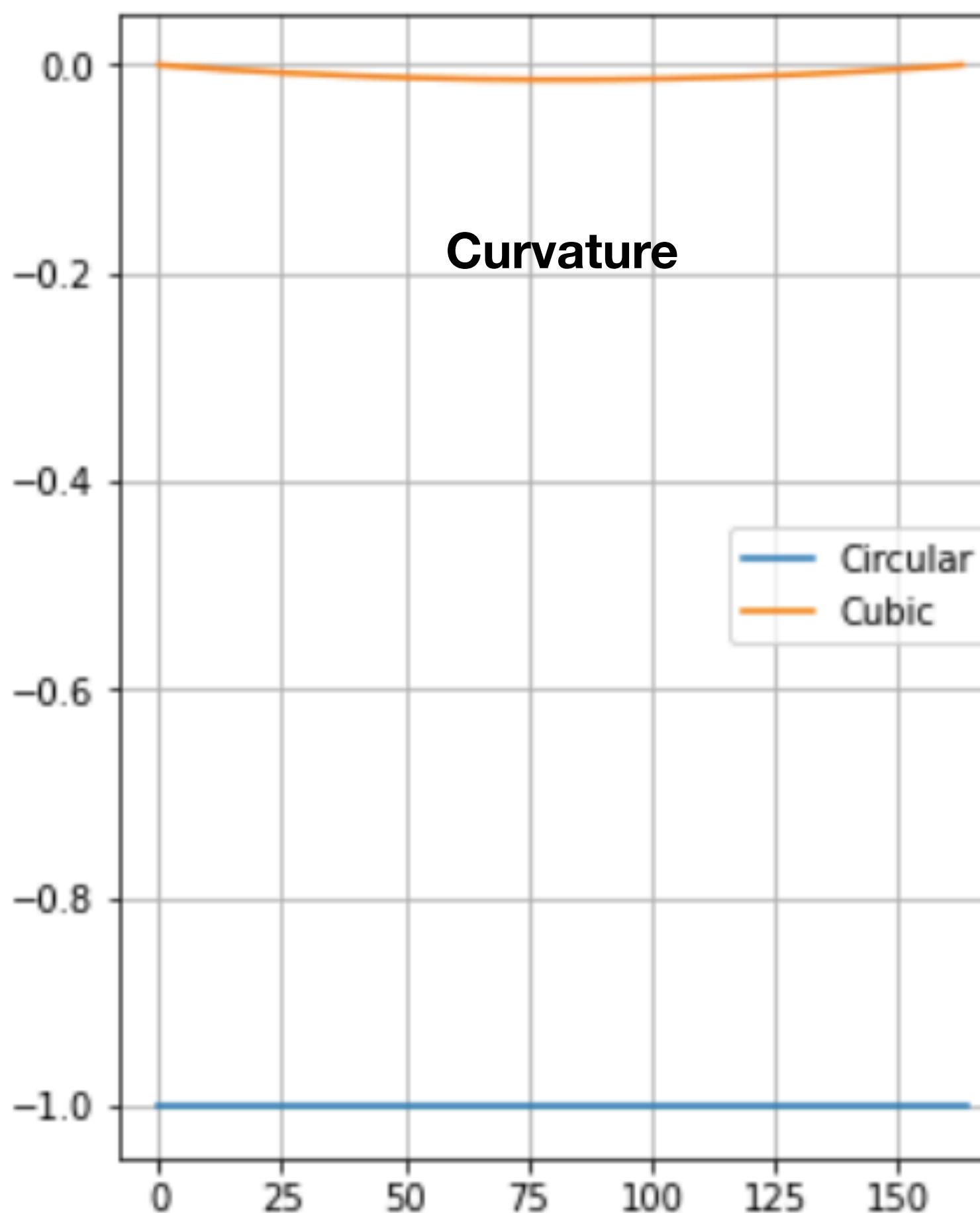
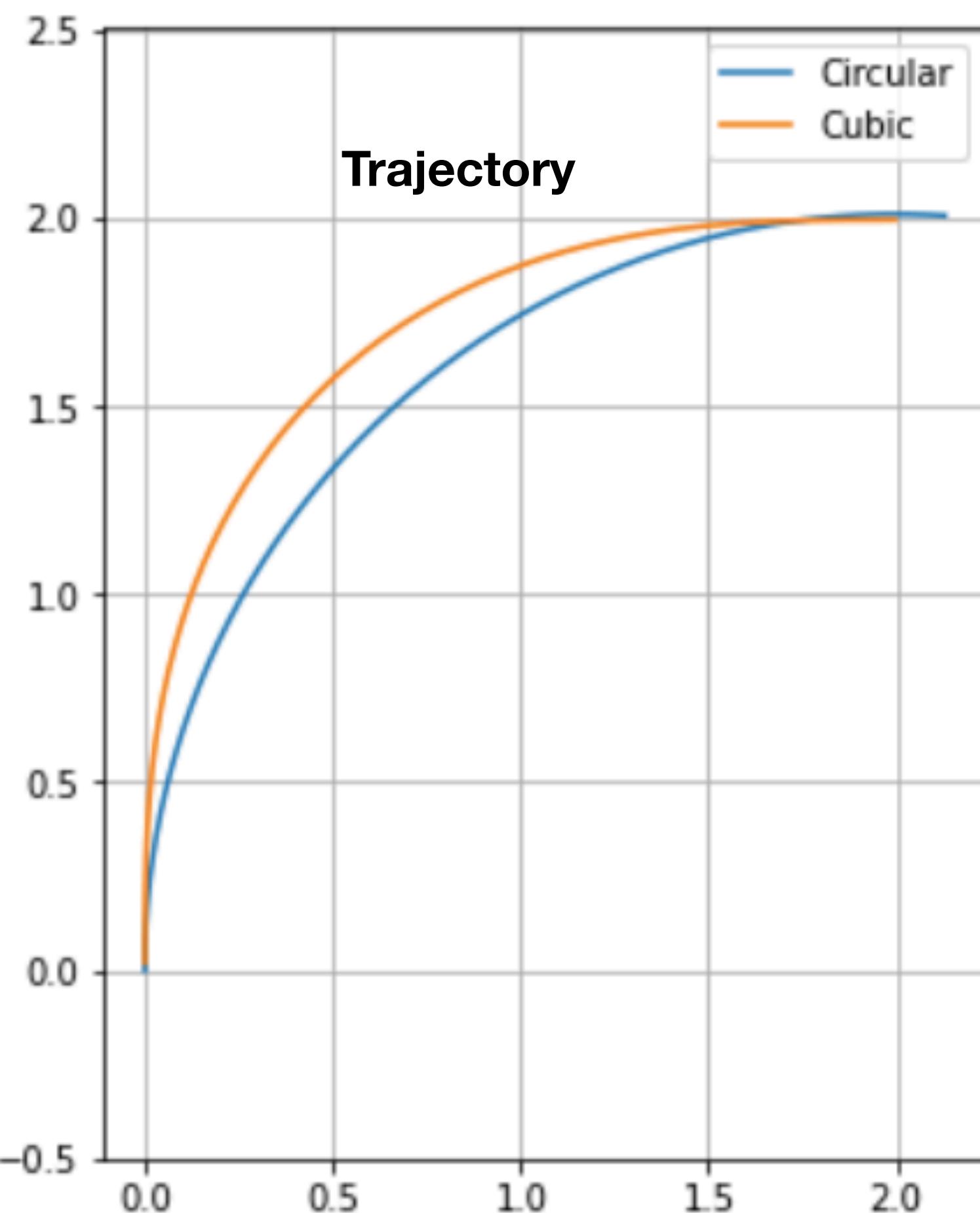
- Curved Path
  - Forward plus rotational motion
  - Weight shifts to the inner wheels
  - Lateral tyre reaction forces produces centripetal acceleration
  - $\alpha = \frac{v^2}{R} = \kappa v^2$

- Curvature: time-derivative of orientation
  - Higher-curvature means sharper turns, higher accelerations
- Higher-order derivatives of vehicle position => velocity, acceleration, jerk, snap
- Passenger comfort may impose boundary constraints on higher-order terms
- Path smoothness / continuity are much desired properties

# Path Curvature

$$\bullet \quad R = \frac{V}{\omega} \quad \kappa = \frac{1}{R}$$

- Straight line:  $\kappa = 0$
- Circular path:  $\kappa = \text{constant}$



- 2 m radius turn compared
- Circular trajectories have a fixed curvature of 1
- Cubic polynomials
  - Similar turn as circular (not as economic)
  - Curvature can start/ end at 0.0 - seamlessly merge with straight segments

# Simple Curve-fitting approach

Fit a polynomial for smooth transition from  $\theta_i$  to  $\theta_f$

For simplicity,  $\theta_i = 1$  and  $\theta_f = 0$

Our function can be parameterized as

$$f'(\theta) = K(\theta(1 - \theta))^{2n}$$

$$f'(0) = f'(1) = 0$$

- $n = 1 \rightarrow$  cubic polynomial
  - Example in week 1:  $K = 6$
- $n = 2 \rightarrow$  quintic polynomial

$\{\theta_i\}$  is fixed

$\{x_i\}$  and  $\{y_i\}$  can be computed for a fixed  $\nu$

