

## 1 Assorted Joins

- Companies: (company\_id, industry, ipo\_date)
- NYSE: (company\_id, date, trade, quantity)

We have 20 pages of memory, and we want to join two tables Companies and NYSE on  $C.\text{company\_id} = N.\text{company\_id}$ . Attribute `company_id` is the primary key for Companies. For every tuple in Companies, assume there are 4 matching tuples in NYSE.

NYSE contains  $[N] = 100$  pages, NYSE holds  $pN = 100$  tuples per page.

Companies contains  $[C] = 50$  pages, C holds  $pC = 50$  tuples per page.

There are unclustered B+ tree indexes of height 1 on  $C.\text{company\_id}$  and  $N.\text{company\_id}$ .

- How many disk I/Os are needed to perform a simple nested loops join?
- How many disk I/Os are needed to perform a block nested loops join?
- How many disk I/Os are needed to perform an index nested loops join?
- For this part only, assume the index on  $NYSE.\text{company\_id}$  is clustered. What is the cost of an index nested loops join using companies as the outer relation?

- (e) How many disk I/Os are needed to perform a sort merge join without optimization? If we can perform the sort merge join optimization, how many disk I/Os are needed with optimization?
- (f) How many disk I/Os are needed to perform a hash join? Assume uniform partitioning.

## 2 Grace Hash Join

We have 2 tables – Catalog and Transactions.

Catalog has a total of 100 pages and 20 tuples per page. Transactions has a total of 50 pages and 50 tuples per page. Assume the hash functions uniformly distribute the data for both tables.

- (a) If we had 10 buffer pages, how many partitioning phases would we require for grace hash join? Consider which table we should build the hash table in the probing phase on.
- (b) What is the I/O cost for the grace hash join then?
- (c) For the above question, if we only had 8 buffer pages, how many partitioning phases would there be?
- (d) What will be the I/O cost?