

Closest Pair of Points | $O(n \log n)$ Implementation

Difficulty Level : Hard • Last Updated : 15 Apr, 2020

We are given an array of n points in the plane, and the problem is to find out the closest pair of points in the array. This problem arises in a number of applications. For example, in air-traffic control, you may want to monitor planes that come too close together, since this may indicate a possible collision. Recall the following formula for distance between two points p and q .

$$\|pq\| = \sqrt{(p_x - q_x)^2 + (p_y - q_y)^2}$$

We have discussed a [divide and conquer solution](#) for this problem. The time complexity of the implementation provided in the previous post is $O(n (\log n)^2)$. In this post, we discuss implementation with time complexity as $O(n \log n)$.

Following is a recap of the algorithm discussed in the previous post.

- 1) We sort all points according to x coordinates.
- 2) Divide all points in two halves.
- 3) Recursively find the smallest distances in both subarrays.
- 4) Take the minimum of two smallest distances. Let the minimum be d .
- 5) Create an array `strip[]` that stores all points which are at most d distance away from the middle line dividing the two sets.
- 6) Find the smallest distance in `strip[]`.
- 7) Return the minimum of d and the smallest distance calculated in above step 6.

The great thing about the above approach is, if the array `strip[]` is sorted according to y coordinate, then we can find the smallest distance in `strip[]` in $O(n)$ time. In the implementation discussed in the previous post, `strip[]` was explicitly sorted in every recursive call that made the time complexity $O(n (\text{Log}n)^2)$, assuming that the sorting step takes $O(n \text{Log}n)$ time.

In this post, we discuss an implementation where the time complexity is $O(n \text{Log}n)$. The idea is to presort all points according to y coordinates. Let the sorted array be `Py[]`. When we make recursive calls, we need to divide points of `Py[]` also according to the vertical line. We can do that by simply processing every point and comparing its x coordinate with x coordinate of the middle line.

Following is C++ implementation of $O(n \text{Log}n)$ approach.