

# CS499 Homework 9 (First Draft)

## Interstellar

### Exercise 9.1

We define  $f_1 : \mathbb{N} \rightarrow \mathbb{N}$

$$f_1(0) = 0, f_1(1) = 1, \dots, f_1(n) = n.$$

We define  $f_2 : \mathbb{N} \rightarrow \mathbb{N}^2$  based on this graph:

	0	1	2	3	4	5	...
0	(0,0)	(0,1)	(0,2)	(0,3)	(0,4)	(0,5)	...
1	(1,0)	(1,1)	(1,2)	(1,3)	(1,4)	(1,5)	...
2	(2,0)	(2,1)	(2,2)	(2,3)	(2,4)	(2,5)	...
3	(3,0)	(3,1)	(3,2)	(3,3)	(3,4)	(3,5)	...
...	...	...	...	...	...	...	...

Figure 1:

$$f_2(0) = (0,0), f_2(1) = (0,1), f_2(2) = (1,0) \dots$$

We define  $f_3 : \mathbb{N} \rightarrow \mathbb{N}^3$  based on this graph:

	0	1	2	3	4	5	...
$f_2(0) = (0,0)$	(0,0,0)	(0,0,1)	(0,0,2)	(0,0,3)	(0,0,4)	(0,0,5)	...
$f_2(1) = (0,1)$	(0,1,0)	(0,1,1)	(0,1,2)	(0,1,3)	(0,1,4)	(0,1,5)	...
$f_2(2) = (1,0)$	(1,0,0)	(1,0,1)	(1,0,2)	(1,0,3)	(1,0,4)	(1,0,5)	...
$f_2(3) = (2,0)$	(2,0,1)	(2,0,2)	(2,0,3)	(2,0,4)	(2,0,5)	...	...
...	...	...	...	...	...	...	...

Figure 2:

$f_2(0) = (0, 0, 0), f_2(1) = (0, 0, 1), f_2(2) = (0, 1, 0) \dots$

And so on, we can define  $f_k, k \in \mathbb{N}$ . Now we can define a bijection  $\mathbb{N} \rightarrow \mathbb{N}^*$  base on this graph:

	0	1	2	3	4	5	...
$f_1$	$f_1(0)$	$f_1(1)$	$f_1(2)$	$f_1(3)$	$f_1(4)$	$f_1(5)$	...
$f_2$	$f_2(0)$	$f_2(1)$	$f_2(2)$	$f_2(3)$	$f_2(4)$	$f_2(5)$	...
$f_3$	$f_3(0)$	$f_3(1)$	$f_3(2)$	$f_3(3)$	$f_3(4)$	$f_3(5)$	...
$f_4$	$f_4(0)$	$f_4(1)$	$f_4(2)$	$f_4(3)$	$f_4(4)$	$f_4(5)$	...
$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\vdots$

Figure 3:

We have  $0 \rightarrow f_1(0), 1 \rightarrow f_1(1), 2 \rightarrow f_2(0) \dots$ . This is a bijection  $\mathbb{N} \rightarrow \mathbb{N}^*$ .

### Exercise 9.2

We can define a bijection from  $\{0, 1\}^{\mathbb{N}}$  to  $\{0, 1\}^{\mathbb{N}} \times \{0, 1\}^{\mathbb{N}}$  as follows.

Given  $A = (a_1 a_2 a_3 a_4 \dots, b_1 b_2 b_3 b_4 \dots)$ , we define  $f(A) = a_1 b_1 a_2 b_2 a_3 b_3 a_4 b_4 \dots$ . To be more precisely,

$$f(A)[i] = \begin{cases} A[1][\frac{i+1}{2}], & i \text{ is odd number} \\ A[2][\frac{i}{2}], & i \text{ is even number} \end{cases}$$

Obviously, for each  $A \in \{0, 1\}^{\mathbb{N}} \times \{0, 1\}^{\mathbb{N}}$ , there is only one  $f(A) \in \{0, 1\}^{\mathbb{N}}$ . For each  $B \in \{0, 1\}^{\mathbb{N}}$ , there is only one  $B = f^{-1}(B) \in \{0, 1\}^{\mathbb{N}} \times \{0, 1\}^{\mathbb{N}}$ . Therefore,  $f$  is a bijection and  $\{0, 1\}^{\mathbb{N}} \cong \{0, 1\}^{\mathbb{N}} \times \{0, 1\}^{\mathbb{N}}$ .

Using the fact that  $\mathbb{R} \cong \{0, 1\}^{\mathbb{N}}$ , we can get  $\mathbb{R} \cong \mathbb{R} \times \mathbb{R}$ .

### Exercises 9.3

We use the Cantor's method to prove that. For any  $A \in (\{0, 1\}^{\mathbb{N}})^{\mathbb{N}}$ , we define that  $f(A)$  is the  $\{0, 1\}^{\mathbb{N}}$  sequence we get by following the blue line as follows.

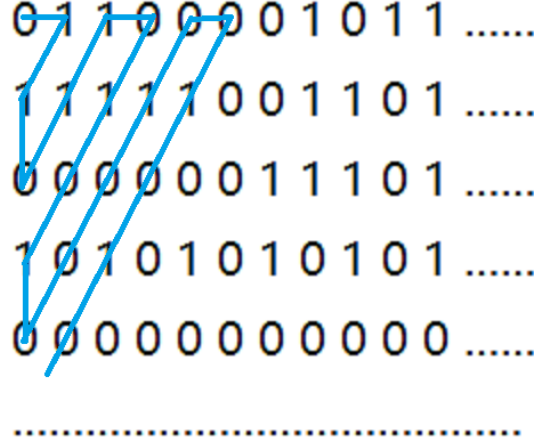


Figure 4:

It is obvious that for any  $B \in \{0, 1\}^{\mathbb{N}}$ , we can get to  $f^{-1} \in (\{0, 1\}^{\mathbb{N}})^{\mathbb{N}}$  by writing it down following the blue line. Therefore,  $f$  is a bijection and  $\{0, 1\} \cong (\{0, 1\}^{\mathbb{N}})^{\mathbb{N}}$ . Using the fact that  $\mathbb{R} \cong \{0, 1\}^{\mathbb{N}}$ , we can get  $\mathbb{R} \cong \mathbb{R}^{\mathbb{N}}$ .

#### Exercise 9.4

We know that one continuous function can be expressed by an infinite sequence of real numbers. Thus we have  $\mathcal{F} \cong \mathbb{R}^{\mathbb{N}}$ . According to **Exercise 9.3**,  $\mathbb{R} \cong \mathbb{R}^{\mathbb{N}}$ . So we have  $\mathcal{F} \cong \mathbb{R}^{\mathbb{N}} \cong \mathbb{R}$ .

#### Exercise 9.5

$000 \dots, 100 \dots, 1100 \dots, 11100 \dots$  According to this rule, the first  $n$  bits of the  $n_{th}$  sequence are 1, and the remaining bits are 0. Obviously, these sequences constitute a countably infinite chain.

#### Exercise 9.6

$100 \dots, 0100 \dots, 00100 \dots, 000100 \dots$  According to this rule, the  $n_{th}$  bit of the  $n_{th}$  sequence is 1, and the remaining bits are 0. Obviously, these sequences constitute a countably infinite antichain.

#### Exercise 9.7

We can define  $A$  using a bijection  $f$  from  $\{0, 1\}^{\mathbb{N}}$  to  $A$ , which is a subset of  $\{0, 1\}^{\mathbb{N}}$  as follows.

Assuming a string  $s$  is an element of  $\{0, 1\}^{\mathbb{N}}$  and  $f(s) = t$ , let  $s_k$  determines  $t_{2k-1}$  and  $t_{2k}$  by the following rule.

If  $s_k = 0$ , then  $t_{2k-1} := 0, t_{2k} := 1$ . If  $s_k = 1$ , then  $t_{2k-1} := 1, t_{2k} := 0$ .

For example:

$$s = 1011010 \dots$$

$$f(s) = 10, 01, 10, 10, 01, 10, 01, \dots$$

Obviously,  $f$  is a bijection and  $A$  is uncountable. Also, any two elements  $t_a, t_b$  of  $A$  is not comparable since "01" and "10" is not comparable. Thus,  $A$  is an required antichain.

#### Exercise 9.8

We can define  $A$  using a bijection  $f$  from  $\{0, 1\}^{\mathbb{N}}$  to  $A$ , which is a subset of  $\{0, 1\}^{\mathbb{N}}$  as follows.

Assuming  $s \in \{0, 1\}^{\mathbb{N}}$  and  $f(s) = t$ , the first  $k$  digit of  $s$  (we call it  $s_{1:k}$ ) determines  $t_{(2^k+1):(2^{k+1})}$  by the following rule.

Consider the  $s_{1:k}$  as a binary number  $a_k$ , then  $t_{(2^k+1):(2^k+a_k)} := 1$  and the  $t_{(2^k+a_k+1):(2^{k+1})} := 0$ . Specially, we define that the first 2 digits of  $t$  are always 0.

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For example:

$$s = 1011.....$$

$$f(s) = 00, 10, 1100, 11111000, 1111111111100000, .....$$

Obviously,  $f$  is a bijection and  $A$  is uncountable. Also, any two elements  $t_a, t_b \in A$  is comparable. Assuming that two elements  $s_a, s_b$  is different, their first different digit is the  $k_{th}$  digit and the  $k_{th}$  digit of  $s_a$  is 1, then for any  $m$  such that  $m \geq k$ , the binary number of the first  $m$  digit of  $s_a$  is greater than that of  $s_b$ , which leads to the conclusion that string  $t_a$  is "greater" than  $t_b$ . Thus,  $A$  is the required chain.