

# CS499 Homework 4 (First Draft)

## Interstellar

### Exercise 4.1

This proof goes wrong when it calculates  $E[X_T]$  in a second way. According to this proof,  $T$  is a certain number. Thus, after walking  $T$  steps, actually we can reach  $i$  with probability  $Pr[T, i] (i = 0, 1, 2, \dots, k)$ . And obviously, none of  $Pr[T, i]$  equals 0 for a certain  $T$ . So we cannot derive the equation  $E[X_T] = k \cdot p_j + 0 \cdot (1 - p_j)$ . If we slightly change the proof, the equation can be derived correctly. We let  $T \rightarrow \infty$ , of course,

$$\lim_{T \rightarrow \infty} Pr[T, i] = 0 (i = 1, 2, \dots, k - 1)$$

$$\lim_{T \rightarrow \infty} Pr[T, k] = P_j$$

$$\lim_{T \rightarrow \infty} Pr[T, 0] = 1 - P_j$$

In this way the equation works.

### Exercise 4.2

1. Since

$$E(T) = \sum_{n=0}^{\infty} (2n+1) C_n p^{n+1} (1-p)^n$$

and

$$\begin{aligned} \sum_{n=0}^{\infty} (2n+1) C_n p^{n+1} (1-p)^n &= p \sum_{n=0}^{\infty} \frac{2n+1}{n+1} \binom{2n}{n} [p(1-p)]^n \\ &< 2p \sum_{n=0}^{\infty} \binom{2n}{n} [p(1-p)]^n \\ &< 2p \sum_{n=0}^{\infty} 2^{2n} [p(1-p)]^n \\ &= 2p \sum_{n=0}^{\infty} [4p(1-p)]^n \end{aligned}$$

we have

$$E(T) < 2p \sum_{n=0}^{\infty} [4p(1-p)]^n$$

Since  $p > \frac{1}{2}$ , then  $4p(1-p) < 1$ .

Thus,

$$E(T) < \frac{2p}{1 - 4p(1-p)}$$

$E(T)$  is finite.

2. We denote

$$g(x) = C_0 + C_1x + C_2x^2 + \dots + C_nx^n + \dots$$

where  $C_n$  is Calatan number.

We can compute that

$$[g(x)]^2 = C_0^2 + (C_0C_1 + C_1C_0)x + (C_0C_2 + C_1^2 + C_2C_0)x^2 + \dots + (C_0C_n + C_1C_{n-1} + \dots + C_nC_0)x^n + \dots$$

Since

$$C_n = C_0C_{n-1} + C_1C_{n-2} + \dots + C_{n-1}C_0$$

we have

$$[g(x)]^2 = C_0^2 + C_2x + C_3x^2 + C_4x^3 + \dots + C_{n+1}x^n + \dots$$

Since  $C_0 = C_1 = 1$ , we have  $x[g(x)]^2 = g(x) - 1$ .

We solve the equation and get

$$g(x) = \frac{1 \pm \sqrt{1 - 4x}}{2x}$$

Since  $g(0) = 1$ , thus

$$g(x) = \frac{1 - \sqrt{1 - 4x}}{2x}$$

$$g'(x) = \frac{1}{x\sqrt{1 - 4x}} - \frac{1 - \sqrt{1 - 4x}}{2x^2}$$

So,

$$\begin{aligned} E(T) &= \sum_{n=0}^{\infty} (2n+1)C_n p^{n+1} (1-p)^n \\ &= 2p \sum_{n=0}^{\infty} nC_n [p(1-p)]^n + p \sum_{n=0}^{\infty} C_n [p(1-p)]^n \\ &= 2p^2(1-p) \times (g[p(1-p)])' + p \times g[p(1-p)] \\ &= \frac{2p^2(1-p)}{p(1-p)\sqrt{1-4p(1-p)}} - \frac{1 - \sqrt{1-4p(1-p)}}{1-p} + p \times \frac{1 - \sqrt{1-4p(1-p)}}{2p(1-p)} \\ &= \frac{2p}{\sqrt{1-4p(1-p)}} - \frac{1 - \sqrt{1-4p(1-p)}}{2(1-p)} \\ &= \frac{2p}{2p-1} - \frac{2(1-p)}{2(1-p)} \\ &= \frac{1}{2p-1} \end{aligned}$$

### Exercise 4.3

This proof went wrong because when  $p < \frac{1}{2}$ ,  $T$  does not have a distribution sequence. As we have already known,  $\sum Pr[T = i] = \sum C_n p^{n+1} (1-p)^n < 1$ , and it fails to satisfy the normalization of a distribution sequence that  $\sum Pr[T = i] = 1$ , so  $T$  can never have an expectation.

### Exercise 4.4

$$E\left[\frac{1}{T+1}\right] = \sum_{n=0}^{\infty} \frac{1}{2(n+1)} C_n p^{n+1} (1-p)^n$$

We denote

$$g(x) = C_0 + C_1x + C_2x^2 + \dots + C_nx^n + \dots$$

where  $C_n$  is Calatan number.

We can compute that

$$[g(x)]^2 = C_0^2 + (C_0C_1 + C_1C_0)x + (C_0C_2 + C_1^2 + C_2C_0)x^2 + \dots + (C_0C_n + C_1C_{n-1} + \dots + C_nC_0)x^n + \dots$$

Since

$$C_n = C_0 C_{n-1} + C_1 C_{n-2} + \dots + C_{n-1} C_0$$

we have

$$[g(x)]^2 = C_0^2 + C_2 x + C_3 x^2 + C_4 x^3 + \dots + C_{n+1} x^n + \dots$$

Since  $C_0 = C_1 = 1$ , we have  $x[g(x)]^2 = g(x) - 1$ .

We solve the equation and get

$$g(x) = \frac{1 \pm \sqrt{1-4x}}{2x}$$

Since  $g(0) = 1$ , thus

$$g(x) = \frac{1 - \sqrt{1-4x}}{2x}$$

$$G(x) = \sum_{n=0}^{\infty} \frac{1}{n+1} C_n x^{n+1} = \int_0^x g(t) dt = \ln(\sqrt{1-4x} + 1) - \sqrt{1-4x} - \ln 2 + 1$$

Since  $p = \frac{1}{2}$ , then

$$\begin{aligned} E\left[\frac{1}{T+1}\right] &= \frac{1}{4} \sum_{n=0}^{\infty} \frac{1}{n+1} C_n \left(\frac{1}{4}\right)^n \\ &= \sum_{n=0}^{\infty} \frac{1}{n+1} C_n \left(\frac{1}{4}\right)^{n+1} \\ &= G\left(\frac{1}{4}\right) \\ &= 1 - \ln 2 \end{aligned}$$

#### Exercise 4.5

*Proof.* Suppose  $A_i$  whose strength is  $a_i$  and  $B_j$  whose strength is  $b_j$  fight. After fighting, for  $A_i$  we have

$$E(\text{strength}_{A_i}) = \frac{a_i}{a_i + b_j} \cdot (a_i + b_j) + \frac{b_j}{a_i + b_j} \cdot 0 = a_i$$

So after  $\forall k$  wars,

$$E\left(\sum \text{strength}_{A_i}\right) = \sum_{i=1}^m a_i$$

$$E\left(\sum \text{strength}_{B_j}\right) = \sum_{j=1}^n b_j$$

Now let's compute  $E(\sum \text{strength}_{A_i})$  in a different way:

Wars has only two results: Alice's team wins or Bob's team wins.

The former happens with probability  $p$ , the latter with probability  $1 - p$ .

Thus

$$E\left(\sum \text{strength}_{A_i}\right) = \left(\sum_{i=1}^m a_i + \sum_{j=1}^n b_j\right) \cdot p + 0 \cdot (1 - p)$$

Therefore,

$$p = \frac{\sum_{i=1}^m a_i}{\sum_{i=1}^m a_i + \sum_{j=1}^n b_j}$$

So the probability of Alice's team winning does not depend on the order in which Alice and Bob send their monsters into the arena.

#### Exercise 4.6

**proof** Define Alice's winning probability of an order  $\{a_{i_1}, a_{i_2}, \dots, a_{i_n}\}$  is  $Pr_{i_1 i_2 \dots i_n : m}$ , in which  $m$  is the number of robots in team B,  $p_{ij} = \frac{a_i}{a_i + b_j}$ ,  $q_{ij} = 1 - p_{ij}$ . We use the induction to prove.

**lemma** Consider there are only two robot in Alice's team (we call it team A) and  $m$  in Bob's team (we call it team B), the probability of team A is

$$Pr_{12:m} = \frac{q_{12}p_{21} \dots p_{2m} - q_{21}p_{11} \dots p_{1m}}{p_{21} - p_{11}} \quad (1)$$

We use induction to prove. When  $m = 2$ , we can easily get

$$Pr_{12} = Pr_{21} = \frac{a_1 a_2 + a_1 b_2 + a_2 b_2}{(a_1 + b_2)(a_2 + b_2)} = \frac{q_{12}p_{21}p_{22} - q_{21}p_{11}p_{12}}{p_{21} - p_{11}} \quad (2)$$

Assume the lemma is true when  $m = k-1$ , then

$$\begin{aligned} Pr_{12:k} &= p_{11}Pr_{12:k-1} + q_{11}p_{21} \dots p_{2k} \\ &= p_{11} \cdot \frac{q_{12}p_{22} \dots p_{2k} - q_{22}p_{12} \dots p_{1k}}{p_{22} - p_{12}} + q_{11}p_{21} \dots p_{2k} \end{aligned} \quad (3)$$

Consider that

$$\begin{aligned} \frac{p_{22}q_{12}}{p_{22} - p_{12}} &= \frac{q_{11}p_{21}}{p_{21} - p_{11}} = \frac{a_1}{a_2 - a_1} \\ \frac{p_{12}q_{22}}{p_{22} - p_{12}} &= \frac{q_{21}p_{11}}{p_{21} - p_{11}} = \frac{a_2}{a_2 - a_1} \end{aligned}$$

Plug it into the formula of  $Pr$

$$Pr_{12:k} = \frac{q_{12}p_{21} \dots p_{2k} - q_{21}p_{11} \dots p_{1k}}{p_{21} - p_{11}}$$

So the lemma proved true.

Assume that for team B with  $m - 1$  robots, the winning probability of team A is regardless of the order. Consider team B with  $m$  robots, the winning probability with the order  $\{1, 2, 3, \dots, n\}$  is

$$\begin{aligned} Pr_{123 \dots (n-1)n:m} &= p_{11}Pr_{123 \dots n:(m-1)} + q_{11}Pr_{123 \dots n-1:m} \\ &= p_{11}Pr_{123 \dots n:(m-1)} + q_{11}p_{21}Pr_{123 \dots n-1:m} + \dots \\ &\quad + q_{11}q_{12} \dots q_{(n-2)1}p_{(n-1)1}Pr_{(n-1)n:k-1} + q_{11}q_{12} \dots q_{(n-2)1}q_{(n-1)1}Pr_{n:k} \end{aligned} \quad (4)$$

exchange the order of last two robots, the winning probability is

$$\begin{aligned} Pr_{123 \dots n(n-1):m} &= p_{11}Pr_{123 \dots n:(m-1)} + q_{11}Pr_{123 \dots n-1:m} \\ &= p_{11}Pr_{123 \dots n:(m-1)} + q_{11}p_{21}Pr_{123 \dots n-1:m} + \dots \\ &\quad + q_{11}q_{12} \dots q_{(n-2)1}p_{n1}Pr_{n(n-1):k-1} + q_{11}q_{12} \dots q_{(n-2)1}q_{n1}Pr_{n-1:k} \end{aligned} \quad (5)$$

It's observed that only the last 2 terms are different, compare it by subtraction

$$Pr_{123 \dots (n-1)n:m} - Pr_{123 \dots n(n-1):m} = q_{11}q_{12} \dots q_{(n-2)1}(p_{(n-1)1}Pr_{(n-1)n:k-1} + q_{(n-1)1}Pr_{n:k} - p_{n1}Pr_{2n(n-1)k-1} - q_{n1}Pr_{n-1:k})$$

Consider that

$$Pr_{n:k} = p_{n1}p_{n2} \dots p_{nk}$$

$$Pr_{n-1:k} = p_{(n-1)1}p_{(n-1)2} \dots p_{(n-1)k}$$

and plug the lemma to rewrite  $Pr_{n(n-1):k-1}$  and  $Pr_{(n-1)n:k-1}$

$$Pr_{n(n-1):k-1} = Pr_{(n-1)n:k-1} = \frac{q_{n1}p_{(n-1)1} \dots p_{(n-1)k} - q_{(n-1)1}p_{n1} \dots p_{nk}}{p_{(n-1)1} - p_{n1}}$$

then we can get the result

$$Pr_{123 \dots (n-1)n:m} - Pr_{123 \dots n(n-1):m} = 0$$

216 iiSimilarly, when exchanging the order of any 2 robots, we can use the method showed in i to prove  
217 they have the same winning probability. Thus,the proposition proved true.  
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## 219 Questions

### 220 1.

221 When we work on Exercise 4.5 & 4.6, we find an interesting phenomenon. In Exercise 4.5, we have  
222 proved that the probability of victory only depends on the sum of robots strengths. But in 4.6, it  
223 does not work. For example, if two teams both have 1 robot with strength 1, obviously probability  
224 of victory is 50%. Then if one team separates the robot into n parts, each part with strength  $\frac{1}{n}$ , its  
225 probability of loss is  $(\frac{n}{n+1})^n$ . We can calculate that  
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$$227 \lim_{n \rightarrow \infty} (\frac{n}{n+1})^n = \frac{1}{e}$$

229 which means one can raise its winning probability by dividing its robot. So in 4.6, we cannot use  
230 the sum of robots strength exclusively as the strength of a team. Then how can we evaluate a teams  
231 strength?  
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### 233 2.

234 How can we deal with problems like Exercise 4.6, with lots of variables and circumstances?  
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