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Discrete Lab 3 Report

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Problem Statement:

• Q1:

An implementation for sieve of Eratosthenes algorithm for finding all prime numbers up to any given limit.

• Q2:

An implementation for Trial Division algorithm which finds the Prime factorization of a given number.

• Q3:

An implementation for the extended Euclidean algorithm that finds the greatest common divisor d of two positive integers a and b. In addition, it outputs Bezout's coefficients s and t such that $d = s \ a + t \ b$. where d is the GCD of a, b

• <u>Q4</u>:

An implementation for Chinese remainder theorem that takes as input m1, m2, m3,, mn that are pairwise relatively prime and (a1, a2,, an) and calculates x.

• <u>Q5:</u>

An implementation for Miller's test (a probabilistic primality test) which takes a number and determine whether it is probably prime or composite.

Used data structures:

• Q1:

• Vector <bool> primes (number + 1, true) >>>> A vector called primes with type bool with size number+1 and assign all values to true.

• Q2:

- o vector
bool> primes "which will store true in index i if i is prime"
- o vector<int> PrimeNumbers "which will contain all the prime numbers up to given limit"

• Q3:

- o vector <int> r "the remainder"
- o vector <int> q "the quotient"
- o vector <int> s "the first coefficient of Bezout's identity"
- o vector <int> t "the second coefficient of Bezout's identity"

• **O4**:

- o vector <int> a "holds a1, a2, a3 an"
- o vector <int> m "holds m1, m2, m3 mn"
- o vector $\langle int \rangle$ M "holds M1, M2, M3 Mn . But the index 0 holds M which is the total product of m1 \cdot m2 \cdot m3 $\cdot \cdot \cdot \cdot \cdot$ mn"
- o vector <int> y "holds y1, y2, y3 yn"

• Q5:

o No Data structures used just variables of type int, long long

Algorithms used:

• Q1:

- o Get the number from the user.
- o Call the sieve function and pass the number on it.
- Make a vector of type Boolean with size = number+1 and assign all values to True (except 0, 1 will be false) supposing that all the numbers are primes.
- Loop from 2:number+1 with increment +1
- o If the current number on the vector is assigned to True(meaning that it's a prime) and its square $(i^2) \le$ the number.
 - Loop from (i^2): number with increment +i
 - Assign all the multiples of i in the vector to false.
- o Print all the primes.

• Q2:

- o Take the number to be factorized from the user and store it in variable number
- Use the Sieve of Eratosthenes Algorithm to find all prime numbers until floor(sqrt(number))
- o Store the prime numbers returned from Sieve algorithm in vector PrimeNumbers
- o Iterate on PrimeNumbers elements
- o Find the first prime number which divides the number
- o If found, then this prime is in the factorization of the number. so repeat step 2 but instead of number perform it on **number/prime** recursively
- o if not found then the number is already a prime number, and it is the factorization of itself

• O3:

- o Get the 2 numbers from the user.
- o Call the exEuclidean function and pass the numbers on it.
- o Make 4 vectors r,q,s,t of type int
- o Assign the initial values on the vectors.
- o <u>If</u> one of the 2 numbers was 0 the Bezouts will be with their initials and the GCD will be the other number not equal to 0.
- o If not: Loop till the remainder reaches 0 starting from counter 0
- o Assign the quotient and the remainder obtained from dividing the two numbers.
- o When the counter reaches 2, we will start evaluating the Bezouts and assigning them to their vectors.
- A final step when the remainder reaches 0 before we exit the loop, we evaluate the last values of the Bezouts which are the final answer.
- o Print the results.

• Q4:

- o Get the number of equations from the user then a_i followed by m_i
- o Compute M and M1, M2, M3 Mn

- Calling The Chinese_Remainder_Theorem() function Which in turn calls the Inverse function to compute y_i
- Compute $X = \sum a_i M_i y_i$ and returns it

• Q5:

- o Take the number to be tested for primality from the user and store it in variable n
- Find two numbers m, k such that $n-1 = m \cdot 2^k$
- \circ Find the test variable $T = a^m \pmod{n}$ by fast modular exponentiation algorithm where a is random variable 1 < a < n-1
- o If T = 1 or T = n-1 then the number n is probably prime
- $\circ \quad Else\ compute\ T=T^2\ (mod\ n)$
- 2
- o If T = 1 then n is composite
- \circ If T = n-1 then n is probably prime
- \circ Else repeat step ② k times but in each time if T = 1 then n is composite else if T = n-1 then n is probably prime
- o If T never becomes 1 or n-1 in the previous k loops, then n is composite

Important Assumptions and details:

• Q1:

- Firstly, assume that all number less than the given number are primes then start to remove multiples of 2, 3, 5, 7 floor(sqrt(number))
- Assume that the limit number can be prime number inclusively for example if the limit number is 5 then 2, 3, 5 are the primes returned from the Sieve algorithm

• Q2:

- At first the Prime_Factorization() function is called, and the number is passed as a
 parameter then find the smallest prime number that divides the number then call the
 function Prime_Factorization() recursively but in this time number/prime is passed as
 a parameter
- o We can find the primes less than a certain number by the help of Sieve Algorithm
- o If the number is prime, then it is the Prime Factorization of itself

• Q3:

- The extended Euclidean algo. uses one pass through the steps of the Euclidean algorithm to find Bezout coefficients of a and b.
- \circ We set s0 = 1, s1 = 0, t0 = 0, and t1 = 1 (note that s is always assigned to a and t is always assigned to b)

• Q4:

- o Bezout's coefficients (s, t) are declared as global variables
- The Gcd() function finds also the bezout's coefficients by Extended Euclidean Algorithm

o The Inverse() function finds the inverse of any Linear congruence by the help of Gcd() function and Extended Euclidean Algorithm where the s Bezout's Identity is the Required Inverse

• Q5:

- o Miller's Test is probabilistic so the number can be composite, and the test says that it is prime but this happen with an absolutely very small probability
- The variable test $T = a^m \pmod{n}$ where a is random variable between 2 and n-2 inclusively but we choose a = 2

Code Snippets:

```
1
       #include <iostream>
 2
       #include <vector>
 3
       #include <cmath>
 4
       using namespace std;
 5
 6
       void sieve (long long number)
 7
    ⊟ {
 8
                                                                     // suppose all numbers are primes
           vector<bool> primes(number+1, true);
                                                                      // 0 , 1 are not primes
 9
           primes[0] = primes[1] = false;
           for (long long i = 2; i \le number; i++)
10
11
12
                if (primes[i] && (i * i) <= number)</pre>
13
                    for (long long j = i * i ; j <= number ; j += i) // removing all multiples of a prime
14
15
                        primes[j] = false;
16
17
18
19
           for(long long i = 2 ; |i <= number ; i++)</pre>
20
                                                                        // printing the primes
    白
21
22
               if(primes[i])
                   cout << i << endl;
23
24
25
26
27
       int main()
   ⊟ {
28
           long long input;
29
           cout << "Enter the limit number:";</pre>
30
31
           cin >> input;
32
           cout << "The primes are:" << endl;</pre>
33
           sieve(input);
34
    • <u>Q2:</u>
```

```
1 #include <bits/stdc++.h>
       using namespace std;
5
      //Perform Sieve Algorithm on the given number and return a vector of integers //which contains all the Primes less than the passed number
6
       vector<int> Sieve (long long number)
8
            vector<bool> primes(number+1, true);
                                                                           // suppose all numbers are primes
10
           for (long long i = 2; i \le sqrt(number); i++)
11
                if (primes[i])
                    for (long long j = i * i ; j <= number ; j += i) primes[j] = false; // removing all multiples of a prime
12
13
                                                //the vector which will contain all the prime numbers up to a given limit
           for(long long i = 2; i <= number; i++)  //inserting the primes in the vector
  if(primes[i]) PrimeNumbers.push_back(i);</pre>
15
16
            return PrimeNumbers;
```

```
22
23
      void Prime_Factorization (long long number)
           vector<int> PrimeNumbers = Sieve(floor(sqrt(number))); //finding all the primes less than squareRoot the number
25
          for(int i = 0; i < PrimeNumbers.size(); i++) {
   if(number % PrimeNumbers[i] == 0) {
      cout << PrimeNumbers[i] << " ";</pre>
                                                                 //if the number is divisible by a prime then print it //and find the \underbrace{\text{Prime}}_{\text{Factorization}} for number/prime
26
27
28
                  Prime_Factorization(number/PrimeNumbers[i]);
29
                  return:
 30
 31
                                     //if the number is not divisible by any prime less than sqrt(number)
32
33
          cout << number:
                                     //then the number is already prime so print it
 35
      int main()
 36
 37
          long long number;
 38
          cout << "Enter the desired number to be factorized:";</pre>
 39
   -1
          cin >> number;
 40
          cout << "The factorization is:" << endl;</pre>
   1
 41
          Prime_Factorization(number);
 42
          return 0;
43
        <u>Q3:</u>
  1
           .ude <iostream>
  2
          .ude <vector>
  3
   4
           namespace std;
  5
   6
            exEuclidean (int numl , int num2)
  8
           '/ 4 vectors for the remainder , quotient , first bezont coeff. , second bezont coeff
  9
           rector <int> r;
 10
           rector <int> q;
          rector <int> s:
 11
 12
          rector <int> t:
 13
            / assigning the initials
 14
          .f(num1 >= num2)
 15
 16
               r.push_back(numl);
 17
               r.push back(num2);
               s.push back(1);
 18
               s.push_back(0);
t.push_back(0);
 19
 20
               t.push_back(1);
 21
 22
               q.push_back(0);
 23
 24
          :lse if (numl < num2)
 25
 26
               r.push back(num2);
 27
               r.push back(numl);
               t.push_back(1);
 28
 29
               t.push back(0);
               s.push_back(0);
 30
 31
               s.push_back(1);
 32
               q.push_back(0);
 33
 34
 35
              our stopping condition is when the reminder = 0
 36
          thile(r.back() != 0)
 37
               q.push_back(r.at(j) / r.at(j+1));
 38
               r.push_back(r.at(j) % r.at(j+1));
 39
 40
```

```
// we start to assign new values to the coeff. vectors when we reach iteration number 2 if (j\,>\,1)
 43
                                          s.push_back(s.at(j-2) - q.at(j-1)*s.at(j-1));
t.push_back(t.at(j-2) - q.at(j-1)*t.at(j-1));
46
 47
48
                                  \texttt{cout} \ll \texttt{j} \ll \texttt{"} \quad \texttt{"} \ll \texttt{r.at(\texttt{j})} \ll \texttt{"} \quad \texttt{"} \ll \texttt{r.at(\texttt{j}+1)} \ll \texttt{"} \quad \texttt{"} \ll \texttt{q.at(\texttt{j}+1)} \ll \texttt{"} \quad \texttt{"} \ll \texttt{r.at(\texttt{j}+2)} \ll \texttt{"} \quad \texttt{"} \ll \texttt{"} \quad \texttt{"} \ll \texttt{s.at(\texttt{j})} \ll \texttt{"} \quad \texttt{"} \ll \texttt{t.at(\texttt{j})} \ll \texttt{end1};
                                 j++;
// assigning the last values "Berout's coeff.s"
 49
 51
                                         \begin{array}{lll} s.push\_back\,(s.at\,(j-2) &- q.at\,(j-1)\,{}^*s.at\,(j-1)\,);\\ t.push\_back\,(t.at\,(j-2) &- q.at\,(j-1)\,{}^*t.at\,(j-1)\,);\\ cout &<< "\tttt" &<< s.at\,(j) &<< " " << t.at\,(j) << endl; \end{array}
 54
 55
56
 57
58
59
                        cout << "GCD is:" << r.at(j) <<endl;
cout << "Resout's coefficients are:" << endl;
// special case if the user insert a 0
if(num2 == 0 || num1 == 0)
 60
                                 cout << "s:" << s.at(0) << endl;
cout << "t:" << t.at(0) << endl;</pre>
 63
 64
65
 66
67
                          else
                                 cout << "s:" << s.back() << endl;
cout << "t:" << t.back() << endl;</pre>
 68
 69
70
71
                int main()
                        int num1 , num2;
cout << "Enter the 2 numbers:" << end1;
cin >> num1 >> num2;
 75
76
77
                          exEuclidean(numl,num2);
                         return 0;
80
                    Q4:
       •
```

```
1
       #include <bits/stdc++.h>
 2
 3
       using namespace std;
 5
      int s,t; //Bezout's Coefficients
 6
 8
      //returns the greatest common divisor of two numbers and sets their Bezout's identity recursively
     ☐int Gcd(int a,int b) {
10
           if(b == 0){
               s = 1;
11
                t = 0;
12
13
                return a;
14
           int gcd = Gcd(b,a % b);
15
16
           int temp = s;
17
           s = t;
18
           t = temp - t * (a/b);
19
           return gcd;
21
22
       // finds the inverse of a (mod m) where the inverse is the s bezout's identity
     ☐ int Inverse(int a, int m) {
23
24
           Gcd(a.m):
2.5
           return s;
26
      // it takes vector a as parameter which contains al,a2,a3, .... an
       // it takes vector M as parameter which contains M in index 0 and M1, M2, M3, · · · · · · Mn
29
   // vector y contains v1.v2.v3. ····· vn. where each individual y is the inverse of the corresponding M (mod m) // vector m contains m1.m2.m3, ···· mn // n is the number of equations of linear congruences
30
31
32
33
     int Chinese_Remainder_Theorem(vector<int> a , vector<int> M , vector<int> y , vector<int> m , int n ){
34
           int X = 0;
           for(int i=1;i<=n;i++){</pre>
35
              y[i] = Inverse(M[i], m[i]);

X += a[i] * M[i] * y[i];
                                                 //Calling the Inverse Function To compute vi
36
                                                 //x is the Summation of al·Ml·vl + a2·M2·v2 ..... an·Mn·vn
38
           return X < 0 ? X%M[0] + M[0] : X%M[0];
39
40
```

```
int main()
  45
46
           cout << "Enter number of equations" << endl;</pre>
  47
48
           int n;
cin >> n;
  49
           vector<int> a(n+1);
  50
51
           vector<int> m(n+1);
vector<int> M(n+1);
  52
           vector<int> y(n+1);
  53
54
                           4[0] will hold the the product of m1 · m2 · m3 ····· mn
          55
56
57
                                                                  //taking the equations from the user
  58
59
  60
          for(int i=1;i<=n;i++)
    M[i] = M[0]/m[i];</pre>
                                         //Computing M1, M2, M3 .... Mn
  61
  62
  63
64
           int X = Chinese_Remainder_Theorem(a ,M ,y ,m ,n );
cout << "X = " << X << " (mod " << M[0] << ")" << endl;</pre>
                                                                   //Calling The Chinese_Remainder_Theorem function to compute X
         Q5:
    •
 1
          #include <bits/stdc++.h>
  2
  3
         using namespace std;
  4
  5
  6
         // computes a power m (mod n) by modular exponentiation
 7

☐int Modular Exponentiation(int a , int m , int n) {
 8
                int result = 1;
 9
                int power = a % n;
10
                while (m) {
11
                                                                                //if ai = 1
12
                             result = (result * power) % n;
13
                      power = (power * power) % n;
14
                      m >>= 1;
15
16
                return result;
17
    // perform the Miller's Rabin Primality test on the given number n and return true if it is prime , false otherwise
     bool Miller Test(int n) {
          if(n == 1) return false;
if(n == 2 || n == 3) return true;
 21
2.2
 23
           long long k,m;
 24
    中
           for(k = 1 ; k < log2(n) ; k++){
                                                     //finds m, k such that n-1 = m * 2^k
 25
              if((n-1) % (long long) pow(2,k) != 0){
 26
27
                  m = (n-1)/pow(2, --k);
                  break;
 28
 29
 30
          int T = Modular_Exponentiation(2, m, n);
                                                     //find the Test variable which is T = a^m \pmod{n} where we choose a = 2
 31
          if(T == 1 || T == n-1) return true;
                                                     //if T = 1 \text{ or } T = n-1 \text{ then } n \text{ is probably prime}
 32
 33
           for(int i = 0; i < k; i++) {</pre>
              T = Modular Exponentiation(T, 2, n); //T = T^2 (mod n)
if(T == 1) return false;
if(T == n-1) return true;
 34
 35
 36
 37
 38
           return false;
39
  44
            int main()
  45
   46
                  cout << "Enter number to be tested for primality"<<end1;</pre>
   47
                  int n;
  48
                  cin >> n;
                  if(Miller Test(n)) cout << n << " is probably a Prime number";</pre>
   49
                  else cout << n << " is a composite number";</pre>
   50
   51
            }
```

Sample Runs:

• <u>Q1:</u>

```
Enter the limit number:200
The primes are:
                                                                                                                       Enter the limit number:15
                                                      The primes are:
Enter the limit number:100
The primes are:
2
5
7
11
13
                                                      11
13
                                                      Process returned 0 (0x0)
                                                                                        execution time : 4.616 s
                                                      Press any key to continue.
19
23
                                                      Enter the limit number:50
The primes are:
                                                                                                                         103
107
109
113
127
131
137
139
149
151
157
31
37
41
                                                     3
7
11
13
17
19
23
29
31
41
43
47
43
47
53
59
61
67
                                                                                                                         163
167
173
179
181
191
71
73
83
                                                                                                                         193
197
89
                                                      Process returned 0 (0x0) execution time : 3.987 s
97
                                                      Press any key to continue.
```

• O2:

Enter the desired number to be factorized: 889 The factorization is: 7 127 Enter the desired number to be factorized: 1024 The factorization is: 2 2 2 2 2 2 2 2 2

Enter the desired number to be factorized: 12911 The factorization is: 12911 Enter the desired number to be factorized: 729 The factorization is: 3 3 3 3 3

Enter the desired number to be factorized: 45 The factorization is: 3 3 5 Enter the desired number to be factorized: 18 The factorization is: 2 3 3

• Q3:

```
Enter the 2 numbers:
Enter the 2 numbers:
                                    66
24
5
                                    0
                                        66
                                                 9
                                                      3
                                                              1
0
    24
         5
              4
                  4
                              0
                          1
                                                     1
                                                             0
                                            3
                                                 2
                                                                 1
    5
             1
                 1
                         0
                             1
         4
                                        3
                                                 3
                                                     0
                                                             1
                                                                 -9
             4
                 0
                         1
                             -4
                                                              -2
                                                                   19
                          -1
                               5
                                    GCD is:1
GCD is:1
                                    Bezout's coefficients are:
Bezout's coefficients are:
                                    5:-2
s:-1
                                    t:19
t:5
                                    Enter the 2 numbers:
Enter the 2 numbers:
                                    7007
804
                                    140
14
                                    0
                                        7007
                                                140
                                                      50
0
    804
          14
                57
                     6
                                 0
                             1
                                        140
                                               7
                                                   20
                                                        0
                                                                0
         6
             2
                  2
                         0
                             1
                                                                   1
    14
                                                                  -50
        2
            3
                 0
                        1
                             -57
                               115
                                    GCD is:7
                          -2
                                    Bezout's coefficients are:
GCD is:2
                                    s:1
Bezout's coefficients are:
5:-2
                                    t:-50
t:115
                                    Enter the 2 numbers:
                                    84562
Enter the 2 numbers:
                                    62
100
                                                              56
                                                                           0
                                    ø
                                        84562
                                                62
                                                      1363
                                            56
6
                                        62
                                                              0
                                        56
                                                 9
                                                                  -1363
    100
          5
               20
                                0
                    0
                            1
                                            2
                                                                  1364
                                        6
                                                             -1
GCD is:5
                                                              10
                                                                   -13639
Bezout's coefficients are:
                                    GCD is:2
                                    Bezout's coefficients are:
5:0
                                    s:10
t:1
                                    t:-13639
```

• Q4:

Enter number of equations Enter number of equations Enter a1 followed by m1 Enter a1 followed by m1 5 7 Enter a2 followed by m2 Enter a2 followed by m2 2 19 3 5 Enter a3 followed by m3 Enter a3 followed by m3 5 11 2 7 Enter a4 followed by m4 X = 23 (mod 105) 9 23 $= 2700 \pmod{33649}$

```
Enter number of equations

5
Enter a1 followed by m1
1 3
Enter a2 followed by m2
2 5
Enter a3 followed by m3
3 7
Enter a4 followed by m4
4 11
Enter a5 followed by m5
5 13
X = 14227 (mod 15015)
```

```
Enter number of equations
3
Enter a1 followed by m1
1 2
Enter a2 followed by m2
2 3
Enter a3 followed by m3
3 5
X = 23 (mod 30)
```

• <u>Q5:</u>

Enter number to be tested for primality

4
4 is a composite number

Enter number to be tested for primality 541 541 is probably a Prime number

Enter number to be tested for primality 97 97 is probably a Prime number

Enter number to be tested for primality 19211 19211 is probably a Prime number

Enter number to be tested for primality 633 633 is a composite number