

# M183 Applikationssicherheit Implementieren # 13

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# Recap # 12

?

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## **Data Access**

- Databases
  - Types
  - How to store, read & modify Data (CRUD)
  - Attacks
- Data-Ressources (HTTP)
  - How to store, read & modify Data
  - Attacks

Data Access & Manipulation –  
what's next?

# Data Integrity ...

*“**Data integrity** is the maintenance of, and the assurance of the accuracy and consistency of, [data](#) over its entire [life-cycle](#), and is a critical aspect to the design, implementation and usage of any system which stores, processes, or retrieves data”*

# Integrity Domains

- **Physical** (Storing & Accessing Physical Data)
  - Altered Traffic (Physical Layer, Network Layer)
  - Material fatigue (Hardware)
  - Corrosion (Hardware)
  - Electromechanical faults (Heisenbugs :-)
  - Environmental Hazards
  - ...
- **Logical** (Concerns: Correctness & Rationality)
  - Referential Integrity (Application Layer)
  - Entity Integrity (Application Layer)
  - Programm Assertions (Application Layer)
  - ...

# How to achieve Integrity?

**Physical Layer ?**

**Logical Layer ?**

# How to achieve Integrity?

## **Physical Layer**

- Checksums
- Hash-Functions (Message verification)
- Error Correcting (Memory & Data Transport)
- **Encryption** (Prevent unwanted Data Alteration and Information gathering by third parties)
- ...

## **Logical Layer**

- Entity Integrity (Primary Keys)
- Referential Integrity (Foreign-Keys)
- Check Constraints
- Application Rules / Business Logic
- ...



# Encryption Domains

**Data «in transit»** - data being transferred via networks

Issues: Eavesdropping, Data Alteration (Planned & Unplanned) ...

**Data «at rest»** - Information stored on computers and storage devices

Issues: cryptography, brute-force, stolen cipher texts, attacks on encryption keys,  
data corruption, data destruction ...

# Data Encryption - Overview

- Block Ciphers vs Stream Ciphers
- Substitution Ciphers vs Transposition Ciphers
- Monoalphabetic Substitution
- Polyalphabetic Substitution
- Symmetric Key Systems
- Public Key Systems

# Block Ciphers vs Stream Ciphers

An important distinction in (symmetric) cryptographic algorithms is between **stream** and **block ciphers**.

**Stream ciphers** convert one symbol of plaintext directly into a symbol of ciphertext.

**Block ciphers** encrypt a group of plaintext symbols as one **block**. Simple substitution is an example of a **stream cipher**

# Block Ciphers vs Stream Ciphers 2

## **Stream Ciphers**

- + Speed of transformation: algorithms are linear in time and constant in space
- + Low error propagation: an error in encrypting one symbol likely will not affect subsequent symbols.
- Low diffusion: all information of a plaintext symbol is contained in a single ciphertext symbol
- No Immunity to tampering: easy to insert symbols without detection

## **Block Ciphers**

- + High diffusion: information from one plaintext symbol is diffused into several ciphertext symbols
- + Immunity to tampering: difficult to insert symbols without detection
- Slowness of encryption: an entire block must be accumulated before encryption / decryption can begin
- Error Propagation: An error in one symbol may corrupt the entire block

# Transposition vs Substitution Ciphers

In a **transposition cipher**, the units of the plaintext are rearranged in a different and usually quite complex order, but the units themselves are left unchanged.

By contrast, in a **substitution cipher**, the units of the plaintext are retained in the same sequence in the **ciphertext**, but the units themselves are altered.

# Transposition Cipher – Rail Fence

**Idea:** In the rail fence cipher, the plaintext is written downwards on successive "rails" of an imaginary fence, then moving up when we get to the bottom

**Example:**

Plaintext -> Ciphertext

```
W . . . E . . . C . . . R . . . L . . . T . . . E
. E . R . D . S . O . E . E . F . E . A . O . C .
. . A . . . I . . . V . . . D . . . E . . . N . .
```

```
WECRL TEERD SOEEF EAOCA IVDEN
```

# Transposition Cipher – Route Cipher

**Idea:** the plaintext is first written out in a grid of given dimensions, then read off in a pattern given in the key

**Example:**

Plaintext -> Ciphertext.

Key: "spiral inwards, clockwise, starting from the top right"

W	R	I	O	R	F	E	O	E
E	E	S	V	E	L	A	N	J
A	D	C	E	D	E	T	C	X

EJXCTEDECDAEWRIORFEONALEVSE

# Encryption - Monoalphabetic Substitution

**Idea:** Every Letter in an Alphabet gets exactly one substitution letter!

## **Variants**

- Cesar-Cipher: Shift (= Key) the Alphabet by n Elements
- Scramble the Alphabet: A -> B, B -> F, C -> Y, etc. (but 1:1)
- Use artificial Characters ...

## **Example**

Cesar Cipher (Shift = Key: 13, «Key»: NOP...)

Plaintext:

«Hello World»

Ciphertext:

«Uryyb Jbeyq»



# Decryption - Monoalphabetic Substitution

## Example

Cesar Cipher (Shift = Key = 13 **backwards** using the **same** «Key»: NOP...)

Ciphertext:

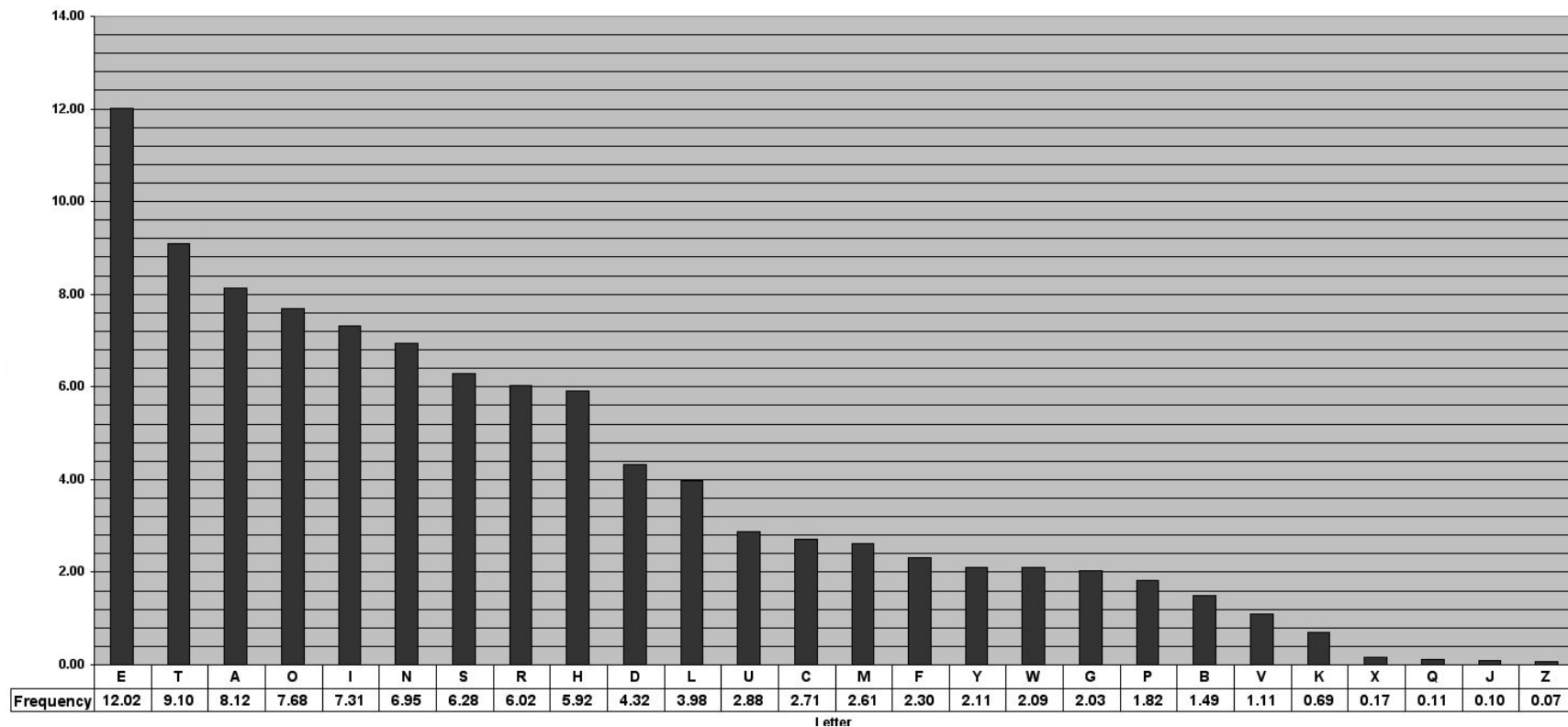
«Uryyb Jbeyq»

Plaintext:

«Hello World»

# Cryptoanalysis - Monoalphabetic Substitution - 1

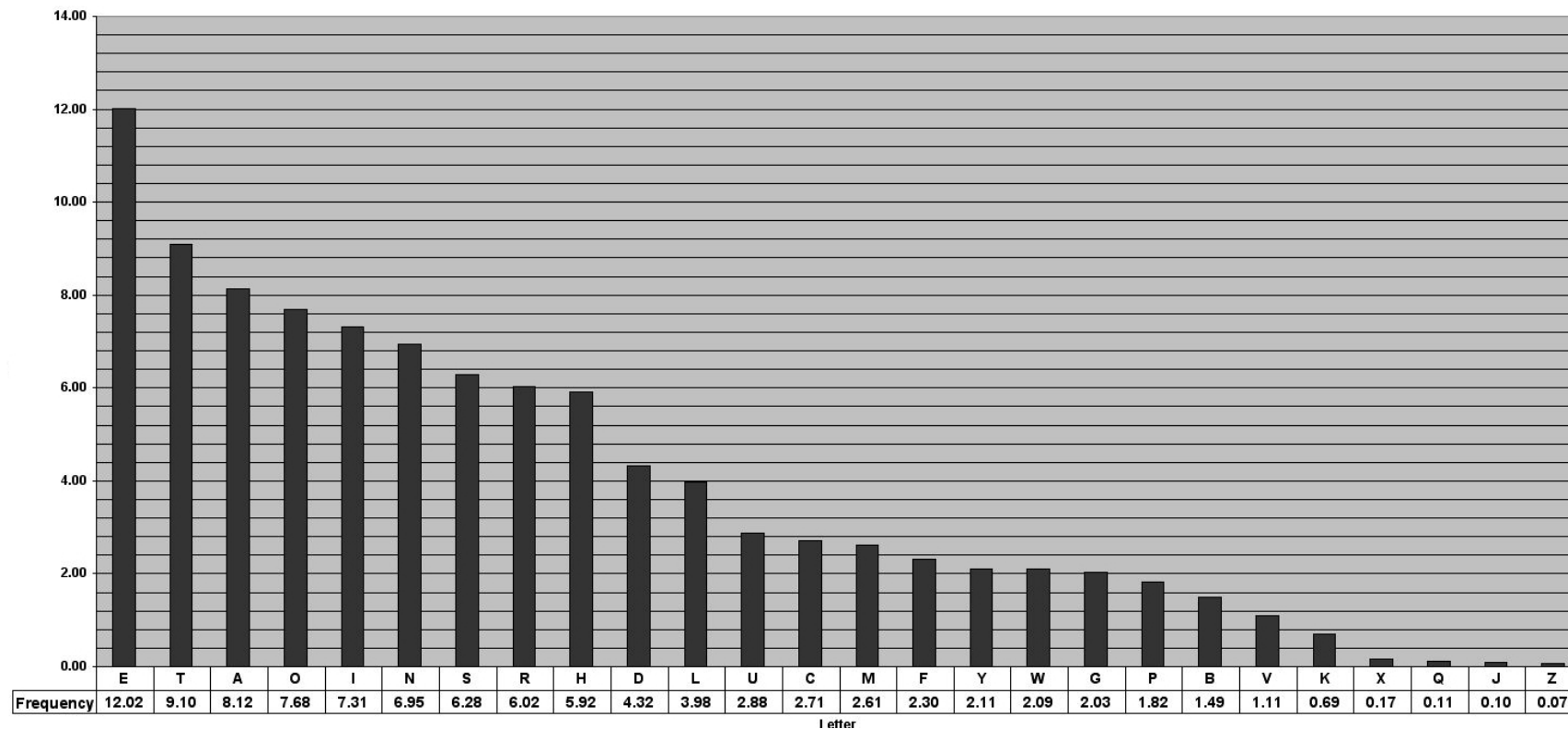
Consider Character Frequencies in English / German Texts



What happens to these Frequencies when doing monoalphabetic substitutions?

# Cryptoanalysis - Monoalphabetic Substitution - 2

Only the Letter – Label changes – not the Frequencies!



Instead of E, with a Shift of 13, R will be the most frequent character!

# Lab – Monoalphabetic Substitution

1. Monoalphabetic Encryption Engine (Parameter: Shift)
2. Monoalphabetic Cryptoanalysis Engine (Frequency Analysis with Tables / BarChart)
3. Monoalphabetic Decryption Engine (Parameter: Back-Shift)

# Encryption - Polyalphabetic Substitution

## Idea

One single Character of an alphabet can be mapped to different Characters of the same alphabet.

## Example

Vigenère Cipher: using a Key K for encryption of a plaintext

Plaintext: A T T A C K A T D A W N

Key: L e m o n L e m o n L e

Ciphertext: L X F O P V E F R N H R

Letter A is Mapped to L, O, E and N

# Decryption - Polyalphabetic Substitution

## Example

Vigenère Cipher: using a Key K for decryption of a ciphertext

Ciphertext: L X F O P V E F R N H R

Key: L e m o n L e m o n L e

Plaintext: A T T A C K A T D A W N

# Cryptoanalysis - Polyalphabetic Substitution

## Idea

The Fact that there are certain mappings of letters or words, that stay the same (a.k.a Repetitions)

## Example 1

Vigenère Cipher: using a Key K for encryption of a plaintext

Plaintext: A T **T** A C K A **T** D A W N

Key: L e m o n L e m o n L e

Ciphertext: L X **F** O P V E **F** R N H R

# Cryptoanalysis - Polyalphabetic Substitution

## Example 2

Vigenère Cipher: using a Key K for encryption of a plaintext

Key:	ABCDABCDABCDABCDABCDABCDABCD
Plaintext:	<b>CRYPTO</b> ISSHORTFOR <b>CRYPTO</b> GRAPHY
Ciphertext:	<b>CSASTP</b> KVSIQUTGQU <b>CSASTP</b> IUAQJB

Key-Length may be: 16 or every factor of 16: 8, 4, 2, 1



# Cryptoanalysis - Polyalphabetic Substitution

When the Key Length  $L$  is known  $\rightarrow$  display Ciphertext in Columns of  $L$

L1	L2	L3	L4	L5	L6
C	V	J	T	N	A
F	E	N	M	C	D
M	K	B	X	F	S
T	K	L3	H	G	S
O	J	W	H	O	F

Perform Frequency analysis within each Column and calculate the shift to the regular frequency of the language.

L1 corresponds to the letter of the shifted position.

# Lab – Polyalphabetic Substitution

1. Polyalphabetic Encryption Engine (Parameter: Key as Character String)
2. Polyalphabetic Cryptoanalysis Engine (Frequency Analysis with Tables / BarChart)
3. Polyalphabetic Decryption Engine (Parameter: Key as Character String)
4. Perform Frequency Analysis per Key Column