Breadth-First Search

 $\mathbf{Input} \colon \mathbf{An}$ unweighted graph and a start vertex \mathbf{u}

Idea:

- Maintain a set R of vertices that have been reached but not searched and
- a set S of vertices that have been searched.
- The set R is maintained as a First-In First-Out list (queue)

Initialization: $R = u, S = \emptyset, d(u, u) = 0$

Iteration: As long as $R \neq \emptyset$, we search from the first vertex v of R. The neighbors of v not in $S \cup R$ are added to the back of R and then v is removed from the front of R and placed in S.

Example

Let G be the adjacency graph of the following Go shape:

```
1 2 3 4 5 6
1 . o o o o .
2 . . . . . <--- black component
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v	R	S
(2,1)	[(2,1)]	{}
(2,1)	[(2,2)]	$\{(2,1)\}$
(2,2)	[(2,3),(3,2)]	$\{(2,1),(2,2)\}$
(2,3)	[(3,2)]	$\{(2,1),(2,2),(2,3)\}$
(3,2)	[(4,2)]	$\{(2,1),(2,2),(2,3),(3,2)\}$
(4,2)	[(5,2)]	$\{(2,1),(2,2),(2,3),(3,2),(4,2)\}$
(5,2)	[(6,2)]	$\{(2,1),(2,2),(2,3),(3,2),(4,2),(5,2)\}$
(6,2)	[(6,1)]	$\{(2,1),(2,2),(2,3),(3,2),(4,2),(5,2),(6,2)\}$
(6,1)		$\{(2,1),(2,2),(2,3),(3,2),(4,2),(5,2),(6,2),(6,1)\}$

References

 \bullet West. D. Introduction to Graph Theory (second edition) 2.3.8