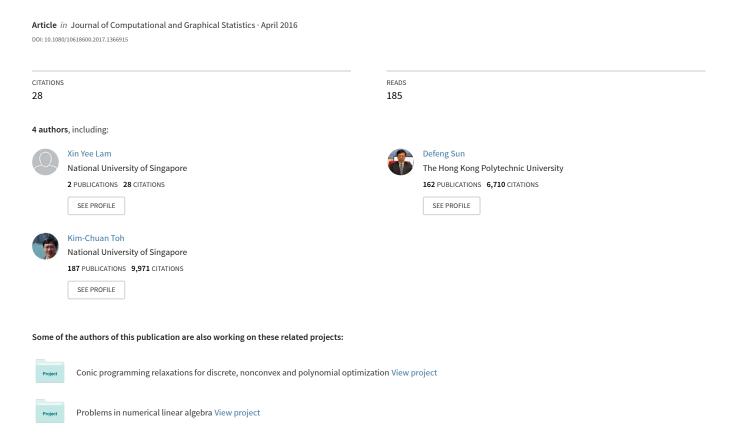
Fast Algorithms for Large-Scale Generalized Distance Weighted Discrimination



Fast algorithms for large scale generalized distance weighted discrimination

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Abstract

High dimension low sample size statistical analysis is important in a wide range of applications. In such situations, the highly appealing discrimination method, support vector machine, can be improved to alleviate data piling at the margin. This leads naturally to the development of distance weighted discrimination (DWD), which can be modeled as a second-order cone programming problem and solved by interior-point methods when the scale (in sample size and feature dimension) of the data is moderate. Here, we design a scalable and robust algorithm for solving large scale generalized DWD problems. Numerical experiments on real data sets from the UCI repository demonstrate that our algorithm is highly efficient in solving large scale problems, and sometimes even more efficient than the highly optimized LIBSVM for solving the corresponding SVM problems.

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1 Introduction

We consider the problem of finding a (kernelized) linear classifier for a training data set $\{(x_i, y_i)\}_{i=1}^n$ with $x_i \in \mathbb{R}^d$ and the class label $y_i \in \{-1, 1\}$ for all $i = 1, \ldots, n$. Here n is the sample size (the number of data vectors available) and d is the feature dimension. By far, the most popular and successful method for getting a good linear classifier from the training data is the support vector machine (SVM), originally proposed by Vapnik [24]. Indeed, it has been demonstrated in [9] that the kernel SVM is one of the best performers in the pool of 179 commonly used classifiers. Despite its success, it has been observed in Marron et al. [18] that SVM may suffer from a "data-piling" problem in the high-dimension-low-sample size (HDLSS) setting (where the sample size n is smaller than the feature dimension d). The authors proposed a new linear classifier, called "Distance Weighted Discrimination" (DWD), as a superior alternative to the SVM. DWD has become a workhorse method for a number of statistical tasks, including data

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visualization [17], hypothesis testing linked with visualization in very high dimensions [26], and adjustment for data biases and heterogeneity [1] and [16].

It is well known that there is a strong need for efficient HDLSS methods for the settings where d is large, say in the order of 10^4 – 10^5 , especially in the area of genetic molecular measurements (usually having a small sample size, where many gene level features have been measured), chemometrics (typically a small sample of high dimensional spectra) and medical image analysis (a small sample of 3-d shapes represented by high-dimensional vectors). On the other hand, given the advent of a huge volume of data collectible through various sources, especially from the internet, it is also important for us to consider the case where the sample size n is large, while the feature dimension may be moderate. Thus in this paper, we are interested in the problem of finding a linear classifier through DWD for data instances where n and/or d are large.

In [18], DWD is reformulated as a second-order cone programming (SOCP) problem, and the resulting model is solved by using a primal-dual interior-point method for SOCP problems implemented in the software SDPT3 [22]. However, the IPM based solver employed for DWD in [18] is computationally very expensive for solving problems where n or d is large, thus making it impractical for large scale problems. A recent approach to overcome such a computational bottleneck has appeared in [25], where the authors proposed a novel reformulation of the primal DWD model which consists of minimizing a highly nonlinear convex objective function subject to a ball constraint. The resulting problem is solved via its dual and an MM (minimization-majorization) algorithm is designed to compute the Lagrangian dual function for each given dual multiplier. The algorithm appears to be quite promising in theory for solving large scale DWD problems. However, the current numerical experiments and implementation of the proposed algorithm in [25] are preliminary and limited to small scale data instances, and it appears that substantial work must be done to make it efficient for large scale instances. Hence it is premature to compare our proposed algorithm with the one in [25].

The main contribution of this paper is to design a new method for solving large scale DWD problems, where we target to solve a problem with the sample size $n \approx 10^4$ – 10^6 and/or the dimension $d \approx 10^4$ – 10^5 . Our method is a convergent 3-block semi-proximal alternating direction method of multipliers (ADMM), which is designed based on the recent advances in research on convergent multi-block ADMM-type methods [21, 14, 5] for solving convex composite quadratic conic programming problems.

ADMM is an algorithm that aims to solve structured convex optimization problems by breaking them into smaller pieces and then optimizing each of them sequentially. Due to its simplicity and efficiency, it has been popular in many applications such as image restoration, constrained sparse regression, and etc. The classical ADMM was originally proposed by Glowinski and Marrocco [11] and Gabay and Mercier [10] for solving convex problems with two blocks of variables. To overcome the difficulty when the subproblems are not easy to solve exactly, Eckstein and Bertsekas [7] introduced the first inexact version of the ADMM to solve the subproblems approximately. Subsequently, Eckstein himself proposed a proximal ADMM in [6] to make the subproblems further easier to solve. In [8], the authors allowed the proximal term added to be positive semidefinite, thus introducing the semi-proximal ADMM. We note that the extension to the semi-proximal setting is non-trivial since the convergence of the algorithm is much harder to establish compared to the earlier algorithms.

The classical ADMM is initially formulated for solving a 2-block convex optimization problem with a collection of coupling linear constraints. Over the years, there are many variations of ADMM proposed and applied to a great variety of optimization problems. A natural modification is to extend the original ADMM from two-block to multi-block settings. However, in [4], it was shown that the directly extended ADMM may not be convergent, even when the subproblems are solved exactly. Thus it is necessary to make some modifications to the directly extended ADMM in order to get a convergent algorithm. In [21], the authors proposed a semiproximal ADMM for solving a convex conic programming problem with 3 blocks of variables and 4 types of constraints. The algorithm is a convergent modification of the ADMM with an additional inexpensive step in each iteration. It is usually faster than the directly extended ADMM with unit step-length. In [14], the authors proposed a Schur complement based (SCB) convergent semi-proximal ADMM for solving a multi-block linearly constrained convex programming problem whose objective function is the sum of two proper closed convex functions plus an arbitrary number of convex quadratic or linear functions. One of the key contributions in [14] is the discovery of the Schur complement based decomposition method which allows the multi-block subproblems to be solved efficiently while ensuring the convergence of the algorithm. More recently, [15] generalized the SCB decomposition method in [14] to the inexact symmetric Gauss-Seidel decomposition method, which provides an elegant framework and simpler derivation for the SCB decomposition method in [14]. Based on this previous research, in [5], the authors proposed an *inexact* symmetric Gauss-Seidel based multi-block semi-proximal ADMM for solving a class of high-dimensional convex composite conic optimization problems to moderate accuracy. It is a versatile and yet efficient algorithm, which has been demonstrated to have much better performance than the possibly non-convergent directly extended ADMM, even using the aggressive step-length of 1.618, in solving high dimensional convex quadratic semidefinite programming problems.

Inspired by the above works, we propose a convergent 3-block semi-proximal ADMM, which is a modification of the inexact sGS-ADMM algorithm designed in [5] to solve the DWD model. The first contribution we made in the paper is in reformulating the primal formulation of the generalized DWD model (using the terminology from [25]) and adapting the powerful inexact sGS-ADMM framework for solving the reformulated problem. This is in contrast to numerous SVM algorithms which are primarily designed for solving the dual formulation of the SVM model. The second contribution we made is in designing highly efficient techniques to solve the subproblems in each of the inexact sGS-ADMM iteration. If n or d is moderate, then the complexity at each iteration is $O(nd) + O(n^2)$ or $O(nd) + O(d^2)$ respectively. If both n and d are large, then we employ the conjugate gradient iterative method for solving the large linear systems of equations involved. We also devise various strategies to speed up the practical performance of the sGS-ADMM algorithm in solving large scale instances (with the largest instance having n=256,000 and $d\approx 3\times 10^6$) of DWD problems with real data sets from the UCI machine learning repository [23]. We should emphasize that the key in achieving high efficiency in our algorithm depends very much on the intricate numerical techniques and sophisticated implementation we have developed.

Finally, we conduct extensive numerical experiments to evaluate the performance of our proposed algorithm against a few other alternatives. Relative to the primal-dual interior-point method used in [18], our algorithm is vastly superior in terms of computational time and memory usage in solving large scale problems, where our algorithm can be a few thousands times faster. By exploiting all the highly efficient numerical techniques we have developed in the implementation of the sGS-ADMM algorithm for solving the generalized DWD problem, we can also get an efficient implementation of the possibly non-convergent directly extended ADMM for solving the same problem. On the tested problems, our algorithm generally requires fewer iterations compared to the directly extended ADMM even when the latter is convergent. On quite a few instances, the directly extended ADMM actually requires many more iterations than our proposed algorithm to solve the problems. We also compare the efficiency of our algorithm in solving the generalized DWD problem against the highly optimized LIBSVM [3] in solving the corresponding dual SVM problem. Surprisingly, our algorithm can even be more efficient than LIBSVM in solving large scale problems even though the DWD model is more complex,

and on some instances, our algorithm is 50–100 times faster.

The remaining parts of this paper is organized as follows. In section 2, we present the DWD formulation in details. In section 3, we propose our inexact sGS-based ADMM method for solving large scale DWD problems. We also discuss some essential computational techniques used in our implementation. We report our numerical results in section 4. We will also compare the performance of our algorithm to other solvers on the same data sets in this particular section. Finally, we conclude the paper in section 5.

Notation. We denote the 2-norm of a vector x by ||x||, and the Frobenius norm of a matrix M by $||M||_F$. The inner product of two vectors x, y is denoted by $\langle x, y \rangle$.

2 Generalized distance weighted discrimination

This section gives details on the optimization problems underlying the distance weighted discrimination.

Let (x_i, y_i) , i = 1, ..., n, be the training data where $x_i \in \mathbb{R}^d$ is the feature vector and $y_i \in \{+1, -1\}$ is its corresponding class label. We let $X \in \mathbb{R}^{d \times n}$ be the matrix whose columns are the x_i 's, and $y = [y_1, ..., y_n]^T$.

In linear discrimination, we attempt to separate the vectors in the two classes by a hyperplane $H = \{x \in \mathbb{R}^d \mid w^T x + \beta = 0\}$, where $w \in \mathbb{R}^d$ is the unit normal and $|\beta|$ is its distance to the origin. Given a point $z \in \mathbb{R}^d$, the signed distance between z and the hyperplane H is given by $w^T z + \beta$. For binary classification where the label $y_i \in \{-1, 1\}$, we want

$$y_i(\beta + x_i^T w) \ge 1 - \xi_i \quad \forall \ i = 1, ..., n,$$

where we have added a slack variable $\xi \geq 0$ to allow the possibility that the positive and negative data points may not be separated cleanly by the hyperplane. In matrix-vector notation, we need

$$r := Z^T w + \beta y + \xi \ge \mathbf{1}, \tag{1}$$

where $Z = X \operatorname{diag}(y)$ and $\mathbf{1} \in \mathbb{R}^n$ is the vector of ones.

In SVM, w and β are chosen by maximizing the minimum residual, i.e.,

$$\max \left\{ \delta - C\langle \mathbf{1}, \xi \rangle \mid Z^T w + \beta y + \xi \ge \delta \mathbf{1}, \ \xi \ge 0, \ w^T w \le 1 \right\}, \tag{2}$$

where C > 0 is a tuning parameter to control the level of penalization on ξ . For the DWD approach introduced in [18], w and β are chosen instead by minimizing the sum of reciprocals of the r_i 's, i.e.,

$$\min \Big\{ \sum_{i=1}^{n} \frac{1}{r_i} + C\langle \mathbf{1}, \xi \rangle \mid r = Z^T w + \beta y + \xi, \ r > 0, \ \xi \ge 0, \ w^T w \le 1, \ w \in \mathbb{R}^d \Big\}.$$
 (3)

Detailed discussions on the connections between the DWD model (3) and the SVM model (2) can be found in [18]. The DWD optimization problem (3) is shown to be equivalent to a second-order cone programming problem in [18] and hence it can be solved by interior-point methods such as those implemented in the solver SDPT3 [22].

Here we design an algorithm which is capable of solving large scale generalized DWD problems of the following form:

min
$$\Phi(r,\xi) := \sum_{i=1}^{n} \theta_q(r_i) + C\langle e, \xi \rangle$$
s.t
$$Z^T w + \beta y + \xi - r = 0$$

$$\|w\| \le 1, \ \xi \ge 0,$$

$$(4)$$

where $e \in \mathbb{R}^n$ is a given positive vector such that $||e||_{\infty} = 1$ (the last condition is for the purpose of normalization). The exponent q can be any given positive number, though the values of most interest are likely to be q = 0.5, 1, 2, 4, and $\theta_q(r_i)$ is the function defined by

$$\theta_q(t) = \begin{cases} \frac{1}{t^q} & \text{if } t > 0, \\ \infty & \text{if } t \leq 0. \end{cases}$$

Observe that in addition to allowing for a general exponent q in (4), we also allow for a nonuniform weight $e_i > 0$ in the penalty term for each ξ_i . By a simple change of variables and modification of the data vector y, (4) can also include the case where the terms in $\sum_{i=1}^{n} \frac{1}{r_i^q}$ are weighted non-uniformly. For brevity, we omit the details.

Proposition 1. The dual of problem (4) is given as follows:

$$-\min_{\alpha} \left\{ \Psi(\alpha) := \|Z\alpha\| - \kappa \sum_{i=1}^{n} \alpha_i^{\frac{q}{q+1}} \mid 0 \le \alpha \le Ce, \langle y, \alpha \rangle = 0 \right\}, \tag{5}$$

where $\kappa = \frac{q+1}{q} q^{\frac{1}{q+1}}$.

Proof. Consider the Lagrangian function associated with (4):

$$L(r, w, \beta, \xi; \alpha, \eta, \lambda) = \sum_{i=1}^{n} \theta_{q}(r_{i}) + C\langle e, \xi \rangle - \langle \alpha, Z^{T}w + \beta y + \xi - r \rangle + \frac{\lambda}{2}(\|w\|^{2} - 1) - \langle \eta, \xi \rangle$$

$$= \sum_{i=1}^{n} \theta_{q}(r_{i}) + \langle r, \alpha \rangle + \langle \xi, Ce - \alpha - \eta \rangle - \beta \langle y, \alpha \rangle - \langle w, Z\alpha \rangle + \frac{\lambda}{2}(\langle w, w \rangle - 1),$$

where $r \in \mathbb{R}^n$, $w \in \mathbb{R}^d$, $\beta \in \mathbb{R}$, $\xi \in \mathbb{R}^n$, $\alpha \in \mathbb{R}^n$, $\lambda, \eta \ge 0$. Now

$$\begin{split} &\inf_{r_i} \left\{ \theta_q(r_i) + \alpha_i r_i \right\} = \left\{ \begin{array}{ll} \kappa \, \alpha_i^{\frac{q}{q+1}} & \text{if } \alpha_i > 0 \\ 0 & \text{if } \alpha_i = 0 \\ -\infty & \text{if } \alpha_i < 0 \end{array} \right. = \left\{ \begin{array}{ll} \kappa \, \alpha_i^{\frac{q}{q+1}} & \text{if } \alpha_i \geq 0 \\ -\infty & \text{if } \alpha_i < 0 \end{array} \right. , \\ &\inf_{w} \left\{ - \langle Z\alpha, w \rangle + \frac{\lambda}{2} \|w\|^2 \right\} = \left\{ \begin{array}{ll} -\frac{1}{2\lambda} \|Z\alpha\|^2 & \text{if } \lambda > 0 \\ 0 & \text{if } \lambda = 0, \, Z\alpha = 0 \\ -\infty & \text{if } \lambda = 0, \, Z\alpha \neq 0 \end{array} \right. , \\ &\inf_{\xi} \left\{ \langle \xi, Ce - \alpha - \eta \rangle \right\} = \left\{ \begin{array}{ll} 0 & \text{if } Ce - \alpha - \eta = 0 \\ -\infty & \text{otherwise} \end{array} \right. , \\ &\inf_{\beta} \left\{ - \beta \langle y, \alpha \rangle \right\} = \left\{ \begin{array}{ll} 0 & \text{if } \langle y, \alpha \rangle = 0 \\ -\infty & \text{otherwise} \end{array} \right. . \end{split}$$

Let $F_D = \{ \alpha \in \mathbb{R}^n \mid 0 \le \alpha \le Ce, \langle y, \alpha \rangle = 0 \}$. Hence

$$\min_{r,w,\beta,\xi} L(r,w,\beta,\xi;\alpha,\eta,\lambda) = \begin{cases} \kappa \sum_{i=1}^n \alpha_i^{\frac{q}{q+1}} - \frac{1}{2\lambda} \|Z\alpha\|^2 - \frac{\lambda}{2} & \text{if } \lambda > 0, \ \alpha \in F_D, \\ \kappa \sum_{i=1}^n \alpha_i^{\frac{q}{q+1}} & \text{if } \lambda = 0, \ Z\alpha = 0, \ \alpha \in F_D, \\ -\infty & \text{if } \lambda = 0, \ Z\alpha \neq 0, \ \alpha \in F_D, \ \text{or } \alpha \not \in F_D. \end{cases}$$

Now for $\alpha \in F_D$, we have

$$\max_{\lambda \geq 0, \eta \geq 0} \left\{ \min_{r, w, \beta, \xi} L(r, w, \beta, \xi; \alpha, \eta, \lambda) \right\} = \kappa \sum_{i=1}^{n} \alpha_i^{\frac{q}{q+1}} - \|Z\alpha\|.$$

From here, we get the required dual problem.

It is straightforward to show that the feasible regions of (4) and (5) both have nonempty interiors. Thus optimal solutions for both problems exist and they satisfy the following KKT (Karush-Kuhn-Tucker) optimality conditions:

$$\begin{split} Z^T w + \beta y + \xi - r &= 0, & \langle y, \alpha \rangle &= 0, \\ r > 0, \ \alpha > 0, \quad \alpha \leq Ce, \quad \xi \geq 0, \quad \langle Ce - \alpha, \xi \rangle &= 0, \\ \alpha_i &= \frac{q}{r_i^{q+1}}, \ i = 1, \dots, n, & \text{either } w = \frac{Z\alpha}{\|Z\alpha\|}, \ \text{or} \ Z\alpha = 0, \ \|w\|^2 \leq 1. \end{split} \tag{6}$$

Let $(r^*, \xi^*, w^*, \beta^*)$ and α^* be an optimal solution of (4) and (5), respectively. Next we analyse some properties of the optimal solution. In particular, we show that the optimal solution α^* is bounded away from 0.

Proposition 2. There exists a positive δ such that

$$\alpha_i^* \geq \delta \ \forall i = 1, \dots, n.$$

Proof. For convenience, let $F_P = \{(r, \xi, w, \beta) \mid Z^T w + \beta y + \xi - r = 0, ||w|| \le 1, \xi \ge 0, r > 0\}$ be the feasible region of (4). Since $(\mathbf{1}, \mathbf{1}, 0, 0) \in F_P$, we have that

$$Ce_{\min}\xi_i^* \le C\langle e, \xi^* \rangle \le \Phi(r^*, \xi^*, w^*, \beta^*) \le \Phi(\mathbf{1}, \mathbf{1}, 0, 0) = n(1+C) \quad \forall i = 1, \dots, n,$$

where $e_{\min} = \min_{1 \le i \le \{e_i\}}$. Hence we have $0 \le \xi^* \le \varrho e$, where $\varrho := n(1 + C^{-1})e_{\min}^{-1}$.

Next, we establish a bound for $|\beta^*|$. Suppose $\beta^* > 0$. Consider an index i such that $y_i = -1$. Then

$$0 < \beta^* = Z_i^T w^* + \xi_i^* - r_i^* \le ||Z_i|| ||w^*|| + \xi_i^* \le K + \varrho,$$

where Z_i denotes the *i*th column of Z, $K = \max_{1 \le j \le n} \{ \|Z_j\| \}$. On the other hand, if $\beta^* < 0$. Then we consider an index k such that $y_k = 1$, and

$$0 < -\beta^* = Z_k^T w^* + \xi_k^* - r_k^* \le K + \varrho.$$

To summarize, we have that $|\beta^*| \leq K + \varrho$.

Now we can establish an upper bound for r^* . For any $i = 1, \ldots, n$, we have that

$$r_i^* = Z_i^T w^* + \beta^* y_i + \xi_i^* \le ||Z_i|| ||w^*|| + |\beta^*| + \xi_i^* \le 2(K + \varrho).$$

From here, we get

$$\alpha_i^* = \frac{q}{(r_i^*)^{q+1}} \ge \delta := \frac{q}{(2K + 2\varrho)^{q+1}} \quad \forall i = 1, \dots, n.$$

This completes the proof of the proposition.

3 An inexact SGS-based ADMM for large scale DWD problems

We can rewrite the model (4) as:

min
$$\sum_{i=1}^{n} \theta_{q}(r_{i}) + C\langle e, \xi \rangle + \delta_{B}(w) + \delta_{\mathbb{R}^{n}_{+}}(\xi)$$
s.t.
$$Z^{T}w + \beta y + \xi - r = 0, \quad w \in \mathbb{R}^{d}, \ r, \xi \in \mathbb{R}^{n},$$

where $B = \{w \in \mathbb{R}^d \mid ||w|| \le 1\}$. This is a convex minimization problem with 3 nonlinear blocks. By introducing an auxiliary variable u = w, we can reformulate the above problem as:

min
$$\sum_{i=1}^{n} \theta_{q}(r_{i}) + C\langle e, \xi \rangle + \delta_{B}(u) + \delta_{\mathbb{R}^{n}_{+}}(\xi)$$
s.t.
$$Z^{T}w + \beta y + \xi - r = 0$$

$$D(w - u) = 0, \quad w, u \in \mathbb{R}^{d}, \ \beta \in \mathbb{R}, \ r, \xi \in \mathbb{R}^{n},$$

$$(7)$$

where $D \in \mathbb{R}^{d \times d}$ is a given positive scalar multiple of the identity matrix which is introduced for the purpose of scaling the variables.

For a given parameter $\sigma > 0$, the augmented Lagrangian function associated with (7) is given by

$$L_{\sigma}(r, w, \beta, \xi, u; \alpha, \rho) = \sum_{i=1}^{n} \theta_{q}(r_{i}) + C\langle e, \xi \rangle + \delta_{B}(u) + \delta_{\mathbb{R}^{n}_{+}}(\xi) + \frac{\sigma}{2} \|Z^{T}w + \beta y + \xi - r - \sigma^{-1}\alpha\|^{2} + \frac{\sigma}{2} \|D(w - u) - \sigma^{-1}\rho\|^{2} - \frac{1}{2\sigma} \|\alpha\|^{2} - \frac{1}{2\sigma} \|\rho\|^{2}.$$

The algorithm which we will design later is based on the recent progress in algorithms for solving multi-block convex conic programming. In particular, our algorithm is designed based on the inexact ADMM algorithm in [5] and we made the essential use of the inexact symmetric Gauss-Seidel decomposition theorem in [14] to solve the subproblems arising in each iteration of the algorithm.

We can view (7) as a linearly constrained nonsmooth convex programming problem with three blocks of variables grouped as (w, β) , r, (u, ξ) . The template for our inexact sGS based ADMM is described next. Note that the subproblems need not to be solved exactly as long as they satisfy some prescribed accuracy.

Algorithm 1. An inexact sGS-ADMM for solving (7).

Let $\{\varepsilon_k\}$ be a summable sequence of nonnegative nonincreasing numbers. Given an initial iterate $(r^0, w^0, \beta^0, \xi^0, u^0)$ in the feasible region of (7), and (α^0, ρ^0) in the dual feasible region of (7). Choose a $d \times d$ symmetric positive semidefinite matrix \mathcal{T} , and perform the following steps in each iteration.

Step 1a. Compute

$$(\bar{w}^{k+1}, \bar{\beta}^{k+1}) \approx \operatorname{argmin}_{w,\beta} \left\{ L_{\sigma}(r^k, w, \beta, \xi^k, u^k; \alpha^k, \rho^k) + \frac{\sigma}{2} \|w - w^k\|_{\mathcal{T}}^2 \right\}.$$

In particular, $(\bar{w}^{k+1}, \bar{\beta}^{k+1})$ is an approximate solution to the following $(d+1) \times (d+1)$ linear system of equations:

$$\underbrace{\begin{bmatrix} ZZ^T + D^2 + \mathcal{T} & Zy \\ (Zy)^T & y^Ty \end{bmatrix}}_{(Zy)^T} \begin{bmatrix} w \\ \beta \end{bmatrix} = \bar{h}^k := \begin{bmatrix} -Z(\xi^k - r^k - \sigma^{-1}\alpha^k) + D^2u^k + D(\sigma^{-1}\rho^k) + \mathcal{T}w^k \\ -y^T(\xi^k - r^k - \sigma^{-1}\alpha^k) \end{bmatrix}.$$
(8)

We require the residual of the approximate solution $(\bar{w}^{k+1}, \bar{\beta}^{k+1})$ to satisfy

$$\|\bar{h}^k - A[\bar{w}^{k+1}; \bar{\beta}^{k+1}]\| \le \varepsilon_k. \tag{9}$$

Step 1b. Compute $r^{k+1} \approx \operatorname{argmin}_{r \in \mathbb{R}^n} L_{\sigma}(r, \bar{w}^{k+1}, \bar{\beta}^{k+1}, \xi^k, u^k; \alpha^k, \rho^k)$. Specifically, by observing that the objective function in this subproblem is actually separable in r_i for $i = 1, \ldots, n$, we can compute r_i^{k+1} as follows:

$$r_i^{k+1} \approx \arg\min_{r_i} \left\{ \theta_q(r_i) + \frac{\sigma}{2} ||r_i - c_i^k||^2 \right\}$$

$$= \arg\min_{r_i > 0} \left\{ \frac{1}{r_i^q} + \frac{\sigma}{2} ||r_i - c_i^k||^2 \right\} \quad \forall i = 1, \dots, n,$$

$$(10)$$

where $c^k = Z^T \bar{w}^{k+1} + y \bar{\beta}^{k+1} + \xi^k - \sigma^{-1} \alpha^k$. The details on how the above one-dimensional problems are solved will be given later. The solution r_i^{k+1} is deemed to be sufficiently accurate if

$$\left| -\frac{q}{(r_i^{k+1})^{q+1}} + \sigma(r_i^{k+1} - c_i^k) \right| \le \varepsilon_k / \sqrt{n} \quad \forall \ i = 1, \dots, n.$$

Step 1c. Compute

$$(w^{k+1}, \beta^{k+1}) \approx \operatorname{argmin}_{w,\beta} \left\{ L_{\sigma}(r^{k+1}, w, \beta, \xi^k, u^k; \alpha^k, \rho^k) + \frac{\sigma}{2} \|w - w^k\|_{\mathcal{T}}^2 \right\},$$

which amounts to solving the linear system of equations (8) but with r^k in the right-hand side vector \bar{h}^k replaced by r^{k+1} . Let h^k be the new right-hand side vector. We require the approximate solution to satisfy the accuracy condition that

$$||h^k - A[w^{k+1}; \beta^{k+1}]|| \le 5\varepsilon_k.$$

Observe that the accuracy requirement here is more relaxed than that stated in (9) of Step 1a. The reason for doing so is that one may hope to use the solution $(\bar{w}^{k+1}, \bar{\beta}^{k+1})$ computed in Step 1a as an approximate solution for the current subproblem. If $(\bar{w}^{k+1}, \bar{\beta}^{k+1})$ indeed satisfies the above accuracy condition, then one can simply set $(w^{k+1}, \beta^{k+1}) = (\bar{w}^{k+1}, \bar{\beta}^{k+1})$ and the cost of solving this new subproblem can be saved.

Step 2. Compute $(u^{k+1}, \xi^{k+1}) = \operatorname{argmin}_{u,\xi} L_{\sigma}(r^{k+1}, w^{k+1}, \beta^{k+1}, \xi, u; \alpha^k, \rho^k)$. By observing that the objective function is actually separable in u and ξ , we can compute u^{k+1} and ξ^{k+1} separately as follows:

$$u^{k+1} = \arg\min \left\{ \delta_B(u) + \frac{\sigma}{2} \|D(u - g^k)\|^2 \right\} = \begin{cases} g^k & \text{if } \|g^k\| \le 1\\ g^k / \|g^k\| & \text{otherwise} \end{cases},$$

$$\xi^{k+1} = \Pi_{\mathbb{R}^n_+} \left(r^{k+1} - Z^T w^{k+1} - y \beta^{k+1} + \sigma^{-1} \alpha^k - \sigma^{-1} Ce \right),$$

where $g^k = w^{k+1} - \sigma^{-1}D^{-1}\rho^k$, and $\Pi_{\mathbb{R}^n_+}(\cdot)$ denotes the projection onto \mathbb{R}^n_+ .

Step 3. Compute

$$\begin{array}{rcl} \alpha^{k+1} & = & \alpha^k - \tau \sigma(Z^T w^{k+1} + y \beta^{k+1} + \xi^{k+1} - r^{k+1}), \\ \rho^{k+1} & = & \rho^k - \tau \sigma D(w^{k+1} - u^{k+1}), \end{array}$$

where $\tau \in (0, (1+\sqrt{5})/2)$ is the steplength which is typically chosen to be 1.618.

In our implementation of Algorithm 1, we choose the summable sequence $\{\varepsilon_k\}_{k\geq 0}$ to be $\varepsilon_k = c/(k+1)^{1.5}$ where c is a constant that is inversely proportional to $\|Z\|_F$. Next we discuss the computational cost of Algorithm 1. As we shall see later, the most computationally intensive steps in each iteration of the above algorithm are in solving the linear systems of equations of the form (8) in Step 1a and 1c. The detailed analysis of their computational costs will be presented in subsection 3.3. All the other steps can be done in at most O(n) or O(d) arithmetic operations, together with the computation of $Z^T w^{k+1}$, which costs 2dn operations if we do not take advantage of any possible sparsity in Z.

3.1 Convergence results

We have the following convergence theorem for the inexact sGS-ADMM, established by Chen, Sun and Toh in [5, Theorem 1]. This theorem guarantees the convergence of our algorithm to optimality, as a merit over the possibly non-convergent directly extended semi-proximal ADMM.

Theorem 1. Suppose that the system (6) has at least one solution. Let $\{(r^k, w^k, \beta^k, \xi^k, u^k; \alpha^k, \rho^k)\}$ be the sequence generated by the inexact sGS-ADMM in Algorithm 1. Then the sequence $\{(r^k, w^k, \beta^k, \xi^k, u^k)\}$ converges to an optimal solution of problem (7) and the sequence $\{(\alpha^k, \rho^k)\}$ converges to an optimal solution to the dual of problem (7).

Proof. In order to apply the convergence result in [5], we need to express (7) in the following form:

min
$$p(r) + f(r, w, \beta) + q(\xi, u) + g(\xi, u)$$

s.t. $A_1^*r + A_2^*[w; \beta] + B^*[\xi; u] = 0,$ (11)

where

$$p(r) = \sum_{i=1}^{n} \theta_q(r_i), \quad f(r, w, \beta) \equiv 0, \quad q(\xi, u) = \delta_B(u) + C\langle e, \xi \rangle + \delta_{\mathbb{R}^n_+}(\xi), \quad g(\xi, u) \equiv 0,$$

$$A_1^* = \begin{pmatrix} -I \\ 0 \end{pmatrix}, \quad A_2^* = \begin{pmatrix} Z^T & y \\ D & 0 \end{pmatrix}, \quad B^* = \begin{pmatrix} I & 0 \\ 0 & -D \end{pmatrix}.$$

Next we need to consider the following matrices:

$$\begin{pmatrix} A_1 \\ A_2 \end{pmatrix} \begin{pmatrix} A_1^*, A_2^* \end{pmatrix} + \begin{pmatrix} 0 & 0 & 0 \\ 0 & \mathcal{T} & 0 \\ 0 & 0 & 0 \end{pmatrix} = \begin{pmatrix} I & [-Z^T, -y] \\ [-Z^T, -y]^T & M \end{pmatrix}, \quad BB^* = \begin{pmatrix} I & 0 \\ 0 & D^2 \end{pmatrix},$$

where

$$M = \begin{pmatrix} ZZ^T + D^2 + \mathcal{T} & Zy \\ (Zy)^T & y^T y \end{pmatrix} \succ 0.$$

Note that one can show that M is positive definite by using the Schur complement lemma. With the conditions that $M \succ 0$ and $BB^* \succ 0$, the conditions in Proposition 4.2 of [5] are satisfied, and hence we can establish the convergence of Algorithm 1 using Theorem 1 in [5].

3.2 Numerical computation of the subproblem (10) in Step 1b

In the presentation of Algorithm 1, we have described how the subproblem in each step can be solved except for the subproblem (10) in Step 1b. Now we discuss how it can be solved. Observe that for each i, we need to solve a one-dimensional problem of the form:

$$\min \Big\{ \varphi(s) := \frac{1}{s^q} + \frac{\sigma}{2} (s - a)^2 \mid s > 0 \Big\}, \tag{12}$$

where a is given. It is easy to see that $\varphi(\cdot)$ is a convex function and it has a unique minimizer in the domain $(0, \infty)$. The optimality condition for (12) is given by

$$s - a = \frac{q\sigma^{-1}}{s^{q+1}},$$

where the unique minimizer s^* is determined by the intersection of the line $s \mapsto s - a$ and the curve $s \mapsto \frac{q\sigma^{-1}}{s^{q+1}}$ for s > 0. We propose to use Newton's method to find the minimizer, and the template is given as follows. Given an initial iterate s^0 , perform the following iterations:

$$s_{k+1} = s_k - \varphi'(s_k)/\varphi''(s_k) = s_k \left(\frac{q(q+2)\sigma^{-1} + as_k^{q+1}}{q(q+1)\sigma^{-1} + s_k^{q+2}}\right), \quad k = 0, 1, \dots$$

Since $\varphi''(s^*) > 0$, Newton's method would have a local quadratic convergence rate, and we would expect it to converge in a small number of iterations, say less than 20, if a good initial point s^0 is given. In solving the subproblem (10) in Step 1b, we always use the previous solution r_i^k as the initial point to warm-start Newton's method. If a good initial point is not available, one can use the bisection technique to find one. In our tests, this technique is however never used.

Observe that the computational cost for solving the subproblem (10) in Step 1c is O(n) if Newton's method converges within a fixed number of iterations (say 20) for all i = 1, ..., n. Indeed, in our experiments, the average number of Newton iterations required to solve (12) for each of the instances is less than 10.

3.3 Efficient techniques to solve the linear system of equations (8)

Observe that in each iteration of Algorithm 1, we need to solve a $(d+1) \times (d+1)$ linear system of equations (8) with the same coefficient matrix A. For large scale problems where n and/or d are large, this step would constitute the most expensive part of the algorithm. In order to solve such a linear system efficiently, we would have to design different techniques to solve it, depending on the dimensions n and d. We consider the following cases.

3.3.1 The case where $d \ll n$ and d is moderate

This is the most straightforward case where we set $\mathcal{T} = 0$, and we solve (8) by computing the Cholesky factorization of the coefficient matrix A. The cost of computing A is $2nd^2$ arithmetic operations. Assuming that A is stored, then we can compute its Cholesky factorization at the cost of $O(d^3)$ operations, which needs only to be performed once at the very beginning of Algorithm 1. After that, whenever we need to solve the linear system (8), we compute the right-hand-side vector at the cost of 2nd operations and solve two $(d+1) \times (d+1)$ triangular systems of linear equations at the cost of $2d^2$ operations.

3.3.2 The case where $n \ll d$ and n is moderate

In this case, we also set $\mathcal{T} = 0$. But solving the large $(d+1) \times (d+1)$ system of linear equations (8) requires more thoughts. In order to avoid inverting the high dimensional matrix A directly, we make use of the Sherman-Morrison-Woodbury formula to get A^{-1} by inverting a much smaller $(n+1) \times (n+1)$ matrix as shown in the following proposition.

Proposition 3. The coefficient matrix A can be rewritten as follows:

$$A = \widehat{D} + UEU^T, \quad U = \begin{bmatrix} Z & 0 \\ y^T & ||y|| \end{bmatrix}, \quad E = \operatorname{diag}(I_n, -1), \tag{13}$$

where $\widehat{D} = \operatorname{diag}(D, ||y||^2)$. It holds that

$$A^{-1} = \widehat{D}^{-1} - \widehat{D}^{-1}UH^{-1}U^{T}\widehat{D}^{-1}, \tag{14}$$

where

$$H = E^{-1} + U^T \widehat{D}^{-1} U = \begin{bmatrix} I_n + Z^T D^{-1} Z + y y^T / \|y\|^2 & y / \|y\| \\ y^T / \|y\| & 0 \end{bmatrix}.$$
 (15)

Proof. It is easy to verify that (13) holds and we omit the details. To get (14), we only need to apply the Sherman-Morrison-Woodbury formula in [12, p.50] to (13) and perform some simplifications.

Note that in making use of (14) to compute $A^{-1}\bar{h}^k$, we need to find H^{-1} . A rather cost effective way to do so is to express H as follows and use the Sherman-Morrison-Woodbury formula to find its inverse:

$$H = J + \bar{y}\bar{y}^T$$
, $J = diag(I_n + Z^T D^{-1}Z, -1)$, $\bar{y} = [y/||y||; 1]$.

With the above expression for H, we have that

$$H^{-1} = J^{-1} - \frac{1}{1 + \bar{y}^T J^{-1} \bar{y}} (J^{-1} \bar{y}) (J^{-1} \bar{y})^T.$$

Thus to solve (8), we first compute the $n \times n$ matrix $I_n + Z^T D^{-1}Z$ in (15) at the cost of $2dn^2$ operations. Then we compute its Cholesky factorization at the cost of $O(n^3)$ operations. (Observe that even though we are solving a $(d+1)\times (d+1)$ linear system of equations for which $d\gg n$, we only need to compute the Cholesky factorization of a much smaller $n\times n$ matrix.) Also, we need to compute $J^{-1}\bar{y}$ at the cost of $O(n^2)$ operations by using the previously computed Cholesky factorization. These computations only need to be performed once at the beginning of Algorithm 1. After that, whenever we need to solve a linear system of the form (8), we can compute \bar{h}^k at the cost of 2nd operations, and then make use of (14) to get $A^{-1}\bar{h}^k$ by solving two $n\times n$ triangular systems of linear equations at the cost $2n^2$ operations, and performing two matrix-vector multiplications involving Z and Z^T at a total cost of 4nd operations. To summarize, given the Cholesky factorization of the first diagonal block of H, the cost of solving (8) via (14) is $6nd + 2n^2$ operations.

3.3.3 The case where d and n are both large

The purpose of introducing the proximal term $\frac{1}{2}||w-w^k||_{\mathcal{T}}^2$ in Step 1a and 1c is to make the computation of the solutions of the subproblems easier. However, one should note that adding the proximal term typically will make the algorithm to converge more slowly, and the deterioration will become worse for larger $||\mathcal{T}||$. Thus in practice, one would need to strike a balance between choosing a symmetric positive semidefinite matrix \mathcal{T} to make the computation easier while not slowing down the algorithm by too much.

In our implementation, we first attempt to solve the subproblem in Step 1a (similarly for 1c) without adding a proximal term by setting $\mathcal{T}=0$. In particular, we solve the linear system (8) by using a preconditioned symmetric quasi-minimal residual (PSQMR) iterative solver [20] when both n and d are large. In each step of the PSQMR solver, the main cost is in performing the matrix-vector multiplication with the coefficient matrix A, which costs 4nd arithmetic operations. As the number of steps taken by an iterative solver to solve (8) to the required accuracy (9) is dependent on the conditioning of A, in the event that the solver requires more than 50 steps to solve (8), we would switch to adding a suitable non-zero proximal term \mathcal{T} to make the subproblem in Step 1a easier to solve.

The most common and natural choice of \mathcal{T} to make the subproblem in Step 1a easy to solve is to set $\mathcal{T} = \lambda_{\max} I - ZZ^T$, where λ_{\max} denotes the largest eigenvalue of ZZ^T . In this case the

corresponding linear system (8) is very easy to solve. More precisely, for the linear system in (8), we can first compute $\bar{\beta}^{k+1}$ via the Schur complement equation in a single variable followed by computing \bar{w}^k as follows:

$$(y^T y - (Zy)^T (\lambda_{\max} I + D)^{-1} (Zy)) \beta = \bar{h}_{d+1}^k - (Zy)^T (\lambda_{\max} I + D)^{-1} \bar{h}_{1:d}^k,$$

$$\bar{w}^{k+1} = (\lambda_{\max} I + D)^{-1} (\bar{h}_{1:d}^k - (Zy) \bar{\beta}^{k+1}),$$
(16)

where $\bar{h}_{1:d}^k$ denotes the vector extracted from the first d components of \bar{h}^k . In our implementation, we pick a \mathcal{T} which is less conservative than the above natural choice as follows. Suppose we have computed the first ℓ largest eigenvalues of ZZ^T such that $\lambda_1 \geq \ldots \geq \lambda_{\ell-1} > \lambda_{\ell}$, and their corresponding orthonormal set of eigenvectors, v_1, \ldots, v_{ℓ} . We pick \mathcal{T} to be

$$\mathcal{T} = \lambda_{\ell} I + \sum_{i=1}^{\ell-1} (\lambda_i - \lambda_{\ell}) v_i v_i^T - Z Z^T, \tag{17}$$

which can be proved to be positive semidefinite by using the spectral decomposition of ZZ^T . In practice, one would typically pick ℓ to be a small integer, say 10, and compute the first ℓ largest eigenvalues and their corresponding eigenvectors via variants of the Lanczos method. In MATLAB, such a computation can be done by using the routine eigs. To solve (8), we need the inverse of $ZZ^T + D + \mathcal{T}$. Fortunately, when $D = \mu I_d$, it can easily be inverted with

$$(ZZ^{T} + D + T)^{-1} = (\mu + \lambda_{\ell})^{-1} I_{d} + \sum_{i=1}^{\ell-1} ((\mu + \lambda_{i})^{-1} - (\mu + \lambda_{\ell})^{-1}) v_{i} v_{i}^{T}.$$

One can then compute $\bar{\beta}^k$ and \bar{w}^k as in (16) with $(\lambda_{\max}I + D)^{-1}$ replaced by the above inverse.

4 Experiments

In this section, we test the performance of our inexact sGS-ADMM method on several publicly available data sets. The numerical results presented in the subsequent subsections are obtained from a computer with processor specifications: Intel(R) Xeon(R) CPU E5-2670 @ 2.5GHz (2 processors) and 64GB of RAM, running on a 64-bit Windows Operating System.

4.1 Tuning the penalty parameter

In the DWD model (7), we see that it is important to make a suitable choice of the penalty parameter C. In [18], it has been noticed that a reasonable choice for the penalty parameter when the exponent q=1 is a large constant divided by the square of a typical distance between the x_i 's, where the typical distance, dist, is defined as the median of the pairwise Euclidean distances between classes. We found out that in a more general case, C should be inversely proportional to $dist^{q+1}$. On the other hand, we observed that a good choice of C also depends on the sample size n and the dimension of features d. In our numerical experiments, we empirically set the value of C to be $10^{q+1} \max \left\{1, \frac{10^{q-1} \log(n) \max\{1000, d\}^{\frac{1}{3}}}{dist^{q+1}}\right\}$, where $\log(\cdot)$ is the natural logarithm.

4.2 Scaling of data

A technique which is very important in implementing ADMM based methods in practice to achieve fast convergence is the data scaling technique. Empirically, we have observed that it is good to scale the matrix Z in (7) so that the magnitude of all the blocks in the equality

constraint would be roughly the same. Here we choose the scaling factor to be $Z_{\text{scale}} = \sqrt{\|X\|_F}$, where $\|\cdot\|_F$ is the Frobenius norm. Hence the optimization model in (7) becomes:

min
$$\sum_{i=1}^{n} \frac{1}{r_i^q} + C\langle e, \xi \rangle + \delta_{\widetilde{B}}(\tilde{u}) + \delta_{\mathbb{R}_+^n}(\xi)$$
s.t.
$$\left(\frac{Z}{Z_{\text{scale}}}\right)^T \tilde{w} + \beta y + \xi - r = 0, \ r > 0,$$

$$D(\tilde{w} - \tilde{u}) = 0, \quad \tilde{w}, \tilde{u} \in \mathbb{R}^d, \ r, \xi \in \mathbb{R}^n,$$
(18)

where $\tilde{w} = Z_{\text{scale}}w$, $\tilde{u} = Z_{\text{scale}}u$, and $\tilde{B} = \{\tilde{w} \in \mathbb{R}^d \mid ||\tilde{w}|| \leq Z_{\text{scale}}\}$. Therefore, if we have computed an optimal solution $(r^*, \tilde{w}^*, \beta^*, \xi^*, \tilde{u}^*)$ of (18), then $(r^*, Z_{\text{scale}}^{-1}\tilde{w}^*, \beta^*, \xi^*, Z_{\text{scale}}^{-1}\tilde{u}^*)$ would be an optimal solution of (7).

4.3 Stopping condition for inexact sGS-ADMM

We measure the accuracy of an approximate optimal solution $(r, w, \beta, \xi, u, \alpha, \rho)$ for (7) based on the KKT optimality conditions (6) by defining the following relative residuals:

$$\begin{split} \eta_{C_1} &= \frac{|y^T \alpha|}{1+C}, & \eta_{C_2} &= \frac{|\xi^T (Ce - \alpha)|}{1+C}, & \eta_{C_3} &= \frac{\|\alpha - s\|^2}{1+C} \text{ with } s_i = \frac{q}{r_i^{q+1}}, \\ \eta_{P_1} &= \frac{\|Z^T w + \beta y + \xi - r\|}{1+C}, & \eta_{P_2} &= \frac{\|D(w - u)\|}{1+C}, & \eta_{P_3} &= \frac{\max\{\|w\| - Z_{\text{scale}}, 0\}}{1+C}, \\ \eta_{D_1} &= \frac{\|\min\{0, \alpha\}\|}{1+C}, & \eta_{D_2} &= \frac{\|\max\{0, \alpha - Ce\}\|}{1+C}, & \end{split}$$

where Z_{scale} is a scaling factor which has been discussed in the last subsection. Additionally, we calculate the relative duality gap by:

$$\eta_{gap} := \frac{|\operatorname{obj}_{\operatorname{primal}} - \operatorname{obj}_{\operatorname{dual}}|}{1 + |\operatorname{obj}_{\operatorname{primal}}| + |\operatorname{obj}_{\operatorname{dual}}|},$$

where $\operatorname{obj}_{\operatorname{primal}} = \sum_{i=1}^n \frac{1}{r_i^q} + C\langle e, \xi \rangle$, $\operatorname{obj}_{\operatorname{dual}} = Z_{\operatorname{scale}} \| Z\alpha \| - \kappa \sum_{i=1}^n \alpha_i^{\frac{q}{q+1}}$, with $\kappa = \frac{q+1}{q} q^{\frac{1}{q+1}}$. We should emphasize that although for machine learning problems, a high accuracy solution is usually not required, it is important however to use the KKT optimality conditions as the stopping criterion to find a moderately accurate solution if one wants to design a robust solver.

We terminate the solver when $\max\{\eta_P, \eta_D\} < 10^{-5}$, $\min\{\eta_C, \eta_{gap}\} < \sqrt{10^{-5}}$, and $\max\{\eta_C, \eta_{gap}\} < 0.05$. Here, $\eta_C = \max\{\eta_{C_1}, \eta_{C_2}, \eta_{C_3}\}$, $\eta_P = \max\{\eta_{P_1}, \eta_{P_2}, \eta_{P_3}\}$, and $\eta_D = \max\{\eta_{D_1}, \eta_{D_2}\}$. Furthermore, the maximum number of iterations is set to be 5000.

4.4 Adjustment of Lagrangian parameter σ

By some preliminary experiments, we set our initial Lagrangian parameter σ to be $\sigma_0 = \min\{10C, n\}^q$, where q is the exponent in (7), and adapt the following strategy to update σ to improve the convergence speed of the algorithm in practice:

Step 1. Set $\chi = \frac{\eta_P}{\eta_D}$, where η_P and η_D are defined in subsection 4.3;

Step 2. If
$$\chi > \theta$$
, set $\sigma_{k+1} = \zeta \sigma_k$; elseif $\frac{1}{\chi} > \theta$, set $\sigma_{k+1} = \frac{1}{\zeta} \sigma_k$.

Here we empirically set θ to be 5 and ζ to be 1.1. Nevertheless, if we have either $\eta_P \ll \eta_D$ or $\eta_D \ll \eta_P$, then we would increase ζ accordingly, say 2.2 if $\max\{\chi, \frac{1}{\chi}\} > 500$ or 1.65 if

 $\max\{\chi, \frac{1}{\chi}\} > 50$. To ensure a smooth convergence of our algorithm, we will only update σ every kth iterations according the following criteria:

$$k = \begin{cases} 5 & \text{if current iteration count } \in (0, 25], \\ 10 & \text{if current iteration count } \in (25, 50], \\ 20 & \text{if current iteration count } \in (50, 100], \\ 30 & \text{if current iteration count } \in (100, 500], \\ 40 & \text{if current iteration count } \in (500, 1000], \\ 100 & \text{if current iteration count } > 1000. \end{cases}$$

Since only a finite number of updates are taken, the convergence of our algorithm is still valid.

4.5 Data Sets

In this subsection, we test our algorithm on two problem sets: the MNIST handwriting recognition and Face recognition. The details for each data set are described below:

MNIST. The MNIST database contains images of 60,000 handwritten digits and we aim to either classify the digit 4 versus digit 9 or classify the digit 4 versus other digits. For each case, we test our algorithm on 5 different training sizes while keeping the testing size to be 10,000. The dimension of features that we have used is 400. We obtain the data sets from the LASVM website [13].

Face. The faces dataset is a collection of 19×19 images, some of which represent faces and some do not. There are a training set of 6,977 images and a testing set of 24,045 images. The number of features is 361 and we obtain the data from the MIT CBCL web site [2].

exponent $q = 1$			Т	esting					
Data	n	d	C	Iter	Time (s)	double	Error (%)	n	Error (%)
MNIST-49	3959	400	1.02e+02	81	0.22	56	2.25	1991	3.87
MNIST-49	5929	400	1.10e+02	81	0.36	55	2.56	1991	3.37
MNIST-49	7925	400	1.13e+02	81	0.52	57	2.57	1991	3.42
MNIST-49	9847	400	1.15e+02	83	0.63	55	2.77	1991	3.37
MNIST-49	11791	400	1.23e+02	92	0.95	60	2.77	1991	3.16
MNIST-4others	20000	400	1.01e+02	187	3.09	158	1.58	10000	2.20
MNIST-4others	30000	400	1.05e+02	292	6.89	239	1.65	10000	2.09
MNIST-4others	40000	400	1.05e+02	284	9.44	238	1.71	10000	1.95
MNIST-4others	50000	400	1.08e + 02	265	11.15	265	1.80	10000	1.98
MNIST-4others	60000	400	1.12e+02	269	13.37	269	1.77	10000	1.91
Face	6977	361	1.00e+03	121	0.50	89	2.45	24045	3.23

Table 1: The performance of our inexact sGS-ADMM method on 11 data sets. In the table, "double" is the iteration count needed for the extra step 1c in algorithm 1.

Table 1 shows the number of iterations and the runtime required for our DWD method in solving 11 data sets. Here, we report the result for the exponent q = 1 and weight e = 1, the vector of ones. The training error as well as the prediction error are also presented in the table.

Denote the index set $S = \{i = 1, \dots, n \mid y_i[\operatorname{sgn}(\beta + x_i^T w)] \leq 0\}$ for which the data instances are categorized wrongly, where $\operatorname{sgn}(x)$ is the sign function given by

$$sgn(x) = \begin{cases} 1 & \text{if } x > 0, \\ 0 & \text{if } x = 0, \\ -1 & \text{if } x < 0. \end{cases}$$

The training error and testing error are both defined by $\frac{|S|}{n} \times 100\%$, where |S| is the cardinality of the set S.

We observe that the testing error is reasonable and the training time is generally short, even when the training size is large. Under the column "double" in Table 1, we also record the number of iterations for which the extra Step 1c is executed to ensure the convergence of Algorithm 1.

4.6 UCI Data

In this subsection, we test our algorithm on the UCI data set. It is a compilation of data consists of a wide variety of domains collected in [23]. The datasets we have chosen here are all classification problems with two classes. However, the size for each class may not be balanced. To tackle the case of uneven class proportions, we would use the weighted DWD model discussed in [19]. Specifically, we consider the model (4) using e = 1 and the term $\sum_{i=1}^{n} 1/r_i^q$ is replaced by $\sum_{i=1}^{n} \tau_i^q/r_i^q$, with the weights τ_i given as follows:

$$\tau_i = \begin{cases} \frac{\tau_-}{\max\{\tau_+, \tau_-\}} & \text{if } y_i = +1\\ \frac{\tau_+}{\max\{\tau_+, \tau_-\}} & \text{if } y_i = -1 \end{cases}$$

where $\tau_{\pm} = (|n_{\pm}|K^{-1})^{\frac{1}{1+q}}$. Here n_{\pm} is the number of data points with class label ± 1 respectively and $K := n/\log(n)$ is a normalizing factor.

Data	n	d	C	Iter	Time (s)	psqmr double	Train-error (%)
a8a	22696	123	6.27e + 02	212	1.48	0 212	15.10
a9a	32561	123	6.49e + 02	238	2.25	0 238	14.93
covtype	581012	54	3.13e+03	1327	207.74	0 191	23.74
gisette	6000	5000	1.00e+02	112	27.63	0 112	0.00
ijcnn1	35000	22	4.23e+03	515	4.04	0 515	7.77
mushrooms	8124	112	3.75e + 02	81	0.33	0 81	0.00
real-sim	72309	20958	1.55e + 04	237	35.59	983 237	1.44
w7a	24692	300	5.95e + 02	846	6.19	0 846	1.17
w8a	49749	300	6.36e + 02	1230	15.14	0 1230	1.20
rcv1	20242	47236	1.55e + 03	65	5.84	385 49	0.72
leu	38	7129	1.00e+02	354	1.23	0 354	0.00
colon	62	2000	1.68e + 02	41	0.76	0 41	0.00
prostate	102	6033	1.69e + 02	41	0.11	0 41	0.00
srbct	63	2308	1.37e + 02	113	2.39	0 113	0.00
farm-ads	4143	54877	2.13e+03	81	3.60	434 49	0.34
dorothea	800	100000	1.80e + 02	384	10.79	0 384	0.00
url-svm	256000	3231961	3.96e + 05	201	498.17	571 169	0.01

Table 2: The performance of our inexact sGS-ADMM method on the UCI data sets.

Table 2 presents the number of iterations and runtime required, as well as training error produced when we perform our inexact sGS-ADMM algorithm to solve 17 data sets. The results are generated using the exponent q=1. In the table, "psqmr" is the iteration count for the preconditioned symmetric quasi-minimal residual method for solving the linear system (8). A '0' for "psqmr" means that we are using a direct solver as mentioned in subsection 3.3.1 and 3.3.2.

Again, our algorithm is capable of solving all the data sets, even when the size of the data matrix is huge. In addition, for data with unbalanced class size, such as w7a and w8a, our algorithm is able to produce a classifier with small training error.

4.7 Comparison with other solvers

In this subsection, we compare our inexact sGS-ADMM method for solving (4) via (7) with the primal-dual interior-point method implemented in [22] and used in [18]. We also compare our method with the directly extended (semi-proximal) ADMM (using the aggressive step-length 1.618) even though the latter's convergence is not guaranteed. Note that the directly extended ADMM we have implemented here follows exactly the same design used for sGS-ADMM, except that we switch off the additional Step 1c in the Algorithm 1. We should emphasize that our directly extended ADMM is not a simple adaption of the classical ADMM, but instead incorporates all the sophisticated techniques we have developed for sGS-ADMM.

We will report our computational results for two different values of the exponent, q = 1 and q = 2, in Tables 3 and 4, respectively.

exponent $q = 1$			sGS-ADMM			dire	MM	IPM				
Data	n	d	C	Time (s)	Iter	Error (%)	Time (s)	Iter	Error (%)	Time (s)	Iter	Error (%)
a8a	22696	123	6.27e + 02	1.48	212	15.10	0.81	170	15.10	1367.56	47	15.09
a9a	32561	123	6.49e + 02	2.25	238	14.93	1.79	281	14.93	3075.26	43	14.93
covtype	581012	54	3.13e+03	207.74	1327	23.74	183.65	1284	23.74	-	-	-
gisette	6000	5000	1.00e+02	27.63	112	0.00	102.10	1657	0.00	2764.09	62	20.50*
ijcnn1	35000	22	4.23e+03	4.04	515	7.77	3.45	503	7.77	1704.26	34	7.77
mushrooms	8124	112	3.75e + 02	0.33	81	0.00	0.89	303	0.00	139.43	50	0.00
real-sim	72309	20958	1.55e + 04	35.59	237	1.44	68.86	702	1.42	-	-	-
w7a	24692	300	5.95e + 02	6.19	846	1.17	25.90	5000	1.17	4544.73	53	1.17
w8a	49749	300	6.36e + 02	15.14	1230	1.20	42.89	5000	1.18	-	1	-
rcv1	20242	47236	1.55e + 03	5.84	65	0.72	4.66	81	0.72	9310.52	43	0.72
leu	38	7129	1.00e+02	1.23	354	0.00	5.43	2661	0.00	1.40	12	0.00
colon	62	2000	1.68e + 02	0.76	41	0.00	2.39	301	0.00	1.12	16	0.00
prostate	102	6033	1.69e + 02	0.11	41	0.00	0.17	121	0.00	3.99	12	0.00
srbct	63	2308	1.37e + 02	2.39	113	0.00	14.11	1581	0.00	1.13	12	0.00
farm-ads	4143	54877	2.13e+03	3.60	81	0.34	2.86	81	0.34	514.04	36	0.34
dorothea	800	100000	1.80e + 02	10.79	384	0.00	77.70	5000	0.00	11.85	18	0.00
url-svm	256000	3231961	3.96e + 05	498.17	201	0.01	924.31	541	0.00	-	1	-

Table 3: Comparison between the performance of our inexact sGS-ADMM, directly Extended ADMM "directADMM", and the interior point method "IPM" on the UCI data sets. A '*' next to the error in the table means that the problem set cannot be solved properly by the respective solver. '-' means the algorithm cannot solve the dataset due to the constraint of computer memory.

Table 3 reports the runtime, number of iterations required as well as the training error of 3 different solvers for solving the UCI data sets. We can observe that the interior point method is almost the slowest to achieve optimality compared to the other two solvers, despite requiring the least number of iterations, especially when the sample size n is large. The inefficiency of the interior-point method is caused by its need to solve an $n \times n$ linear system of equations in each iteration, which could be very expensive if n is large. In addition, it cannot solve the DWD problem where n is huge due to the excessive computer memory needed to store the large $n \times n$ matrix.

On the other hand, our inexact sGS-ADMM method outperforms the directly extended (semi-proximal) ADMM for 11 out of 17 cases in terms of runtime. For the other cases, we are only slower by a relatively small margin. Furthermore, when our algorithm outperforms the directly extended ADMM, it often shortens the runtime in a large margin. In terms of number of iterations, for 14 out of 17 cases, the directly extended ADMM requires at least the same number of iterations as our inexact sGS-ADMM method. We can say that our algorithm is remarkably efficient and it further possesses convergence guarantee. In contrast, the directly extended ADMM is not guaranteed to converge although it is also very efficient when it does converge. We can observe that the directly extended ADMM sometimes would take many more iterations to solve a problem compared to our inexact sGS-ADMM, especially for the instances in Table 4, possibly because the lack of convergence guarantee makes it difficult for the method to find a sufficiently accurate approximate optimal solution.

To summarize, our inexact sGS-ADMM method is an efficient yet convergent algorithm for solving the primal form of the DWD model. It is also able to solve large scale problems which cannot be handled by the interior point method.

	exponent $q=2$				sGS-ADMM			directADMM				IPM		
Data	n	d	C	Time (s)	Iter	Error (%)	Time (s)	Iter	Error (%)	Time (s)	Iter	Error (%)		
a8a	22696	123	1.57e + 04	2.11	248	15.10	3.29	531	15.11	7034.94	45	15.17		
a9a	32561	123	1.62e+04	2.79	279	14.94	8.86	1183	14.94	-	-	-		
covtype	581012	54	1.52e + 05	74.69	368	23.74	336.85	2343	23.74	-	-	-		
gisette	6000	5000	1.00e+03	48.04	312	0.00	273.49	5000	1.07	2508.89	55	0.00		
ijcnn1	35000	22	2.69e + 05	2.95	233	7.94	8.28	1004	7.88	2002.77	38	7.98		
mushrooms	8124	112	7.66e + 03	1.40	301	0.00	16.33	5000	0.10	650.16	52	0.00		
real-sim	72309	20958	1.10e+06	29.25	180	1.51	19.26	181	1.51	-	-	-		
w7a	24692	300	1.44e+04	4.21	473	1.15	32.87	5000	2.47	5372.32	48	1.18		
w8a	49749	300	1.54e + 04	8.41	543	1.13	52.77	5000	2.66	-	-	-		
rcv1	20242	47236	3.21e+04	20.61	121	0.83	25.44	267	0.78	7547.28	43	0.79		
leu	38	7129	1.00e+03	1.64	461	0.00	10.55	5000	0.00	1.48	15	0.00		
colon	62	2000	3.02e+03	0.67	41	0.00	7.57	986	0.00	1.82	26	0.00		
prostate	102	6033	2.40e+03		41	0.00	0.41	301	0.00	8.78	28	0.00*		
srbct	63	2308	2.17e + 03		136	0.00	51.28	5000	0.00	1.69	21	0.00		
farm-ads	4143	54877	5.51e+04	5.60	80	0.41	5.36	113	0.41	383.43	41	0.41*		
dorothea	800	100000	1.37e + 03	6.41	253	0.00	80.14	5000	0.00	25.89	35	0.00		
url-svm	256000	3231961	1.84e + 07	1198.81	384	0.50	1406.38	770	0.50	-	-	-		

Table 4: Comparison between the performance of our inexact sGS-ADMM, directly Extended ADMM "directADMM", and the interior point method "IPM" on the UCI data sets for q=2. A '*' next to the error in the table means that the problem set cannot be solved properly by the respective solver.

Table 4 reports the runtime, number of iterations required as well as the training error of 3 different solvers for solving the UCI data sets for the case when q = 2. Again, we can see that the interior point method is almost always the slowest to converge to optimality.

Our sGS-ADMM algorithm outperforms the directly extended ADMM algorithm in 15 out of 17 data sets in terms of runtime. For the other 2 cases, our algorithm is only slightly slower. In terms of the number of iterations, it has the best performance among all the data sets. On the other hand, for 7 data sets, the number of iterations required by the directly extended ADMM hits the maximum iterations allowed, probably implying nonconvergence of the method. For the interior point method, it takes an even longer time to solve the problems compared to the case when q = 1. This is due to an increase in the number of constraints generated in the second-order cone programming formulation of the DWD model with q = 2.

The numerical result we obtained in this case is consistent with the one we obtained for the case q = 1. This further shows the merit of our algorithm in a more general setting. We could also expect the similar result when the exponent is 4 or 8.

4.8 Comparison with LIBSVM

The support vector machine (SVM) is a state-of-the-art method for linear classification problem. In this subsection, we will compare the performance of our algorithm (applying to the DWD model) with the performance of the well-known LIBSVM solver in [3] for solving the corresponding SVM model. Our main aim is to compare the runtime taken by both algorithms to get a good classifier as well as the training classification error produced.

			DWD v	ia sGS-A	DMM	(q=1)	LIBSVM					
Data	n	d	C	Time (s)	Iter	Error (%)	C	Time (s)	Iter	Error (%)		
a8a	22696	123	6.27e + 02	1.48	212	15.10	0.041	50.43	34811	15.44		
a9a	32561	123	6.49e + 02	2.25	238	14.93	0.025	92.65	25721	15.24		
covtype	581012	54	3.13e+03	207.74	1327	23.74	0.032	18463.71	224517	23.70		
gisette	6000	5000	1.00e+02	27.63	112	0.00	0.006	79.56	14818	0.40		
ijcnn1	35000	22	4.23e+03	4.04	515	7.77	0.126	26.34	10137	9.21		
mushrooms	8124	112	3.75e + 02	0.33	81	0.00	0.179	0.47	1003	0.00		
real-sim	72309	20958	1.55e + 04	35.59	237	1.44	0.608	3615.18	52591	1.37		
w7a	24692	300	5.95e + 02	6.19	846	1.17	0.356	14.45	62830	1.37		
w8a	49749	300	6.36e + 02	15.14	1230	1.20	0.160	62.49	124373	1.36		
rcv1	20242	47236	1.55e + 03	5.84	65	0.72	0.148	692.39	33517	0.62		
leu	38	7129	1.00e+02	1.23	354	0.00	0.140	0.03	146	0.00		
colon	62	2000	1.68e + 02	0.76	41	0.00	0.160	0.02	295	4.84		
prostate	102	6033	1.69e + 02	0.11	41	0.00	0.364	0.09	708	0.00		
srbct	63	2308	1.37e + 02	2.39	113	0.00	0.235	0.02	250	0.00		
farm-ads	4143	54877	2.13e+03	3.60	81	0.34	9.496	14.61	9574	0.17		
dorothea	800	100000	1.80e + 02	10.79	384	0.00	0.031	5.06	2646	0.00		
url-svm	256000	3231961	3.96e + 05	498.17	201	0.01	-	-	-	-		

Table 5: Comparison between the performance of our inexact sGS-ADMM on DWD model (q = 1) with LIBSVM. '-' means the result cannot be obtained within 24 hours.

Table 5 shows the runtime and number of iterations needed for solving the binary classification problem via DWD model and SVM model respectively on the UCI datasets. It also gives the training classification error produced by the two models. We use our proposed inexact sGS-ADMM algorithm to solve the DWD model and the LIBSVM to solve the corresponding SVM model. The tolerance for terminating LIBSVM is set to be 10⁻⁵.

We may observe from the table that the DWD and SVM models produced comparable training classification errors, although there are some discrepancies due to the differences in the models and penalty parameters used. Our algorithm almost always solves the problem faster than LIBSVM. Although LIBSVM is faster than sGS-ADMM for small datasets (5 cases), our algorithm can be 50-100 times faster than LIBSVM when solving large data instances. Furthermore, LIBSVM may have the risk of not being able to handle extremely large-scaled datasets. For example, it cannot solve the biggest dataset (url-svm) within 24 hours. This also demonstrated the efficiency of our algorithm.

5 Conclusion

In this paper, by making use of the recent advances in ADMM from the work in [5], [14] and [21], we proposed a convergent 3-block inexact symmetric Gauss-Seidel-based semi-proximal ADMM algorithm for solving large scale DWD problems. We applied the algorithm successfully to the primal formulation of the DWD model and designed highly efficient routines to solve the subproblems arising in each of the inexact sGS-ADMM iterations. Numerical experiments for the cases when the exponent equals to 1 and 2 demonstrated that our algorithm is capable of solving large scale problems, even when the sample size and/or the feature dimension is huge. In addition, it is also highly efficient while guaranteeing the convergence to optimality. As a conclusion, we have designed an efficient method for solving the binary classification model through DWD.

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