

Tutorial - 5

Q1 - What's difference between DFS & BFS. Write applications of both algorithm.

DFSBFS

- |   |  |
|---|--|
| → Uses stack to find shortest path  | → Uses Queue to find shortest path   |
| → Stands for "Depth First search"   | → Stands for "Breadth First Search"  |
| → Suitable where solution are closer to source  | → Suitable for searching vertices closer to given source   |
| → Recursive algorithm - uses backtracking   | → No concept of backtracking   |
| → Requires less memory  | → Requires more memory   |
| → Considers not all neighbours first, $\therefore$ & then explores all path through this decision<br>eg:- cyclic graph, topological order | → Considers all neighbours first. $\therefore$ not suitable for decision making<br>eg:- bipartite graph, shortest path, etc. |

Q2 - Which data structure is used to implement DFS & BFS & why?

Ans 2 - DFS algorithm traverses a graph in depthward motion & uses stack to remember to get next vertex to search when a dead end occurs in any iteration. So, basically in order to reach to the deepest node first stack - used.



→ DFS uses queue to avoid recursion. Queue keeps track of iterations remaining to be searched as opposed to needing a potentially large amount of processor stack space for recursive solution.

Q3 - What do you mean by sparse & dense graphs? Which representation of graph is better for both.

### Dense Graph

### Sparse Graph

→ A graph in which no. of edges is close to maximal no. of edges

→ A graph in which no. of edges is minimal to no. of edges

→ Graphs are densely connected

→ Graphs are sparsely connected (eg: trees)

→  $G = (V, E)$  in which  $|E| = O(|V|^2)$

→  $G = (V, E)$  in which  $|E| = O(|V|)$

→ If a graph is dense, we should store it as adjacency matrix

→ If graph - sparse - store it as a list of edge

Q4. How to detect cycle graph using BFS & DFS?  
Ans 4. Using BFS:-

i) compute Indegree (no. of incoming edges) for each vertex present in graph & take the count of visited nodes as 0

ii) Put all vertices with In-degree as 0 add to queue (in queue operation)



- iii) Remove vertex from queue  
 → increment count of visited nodes by 1
- iv) Decrease in-degree by 1 for all neighbouring nodes
- v) If in-degree of a neighbouring node is reduced to zero, add to queue

Repeat step 3 till queue's empty.

Using DFS :-

- i) Create graph using given no. of edges & vertices
- ii) Create recursive (f) that have current index or vertices, visited array & parent nodes
- iii) Mark current node as visited
- iv) Find unvisited vertices that are adjacent to current node
- v) If adjacent node not parent - already visited - return true

Q5 - What do you mean by disjoint set data structure?  
 Explain 3 operations along with examples, which can be performed on disjoint sets.

Ans - Disjoint - set data structure :- Also called Union Find Data Structure or Merge - Find Set is data structure that stores a collection of disjoint (non-overlapping) sets.

Three operations performed by Disjoint - set data structure :-



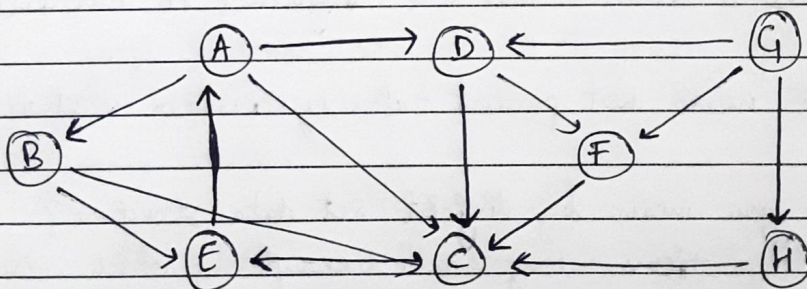
- i) Making new set containing a new element
- ii) Finding representative of set containing given element
- iii) Merging two sets

① Making New Set : Make set operation adds a new element into a new set containing only the new element & new set is added to data structure.

② Finding Set Representatives : Find operation follows chain of parent pointers from a specified query node  $u$  until it reaches a root element. Root element represents set to which  $u$  belongs & may be  $u$  itself.

③ Merging Two Sets : operation  $\text{union}(u, y)$  replaces set containing  $u$  i set containing  $y$  with their Union.

Q6 - Run DFS & BFS on graph shown on RHS.

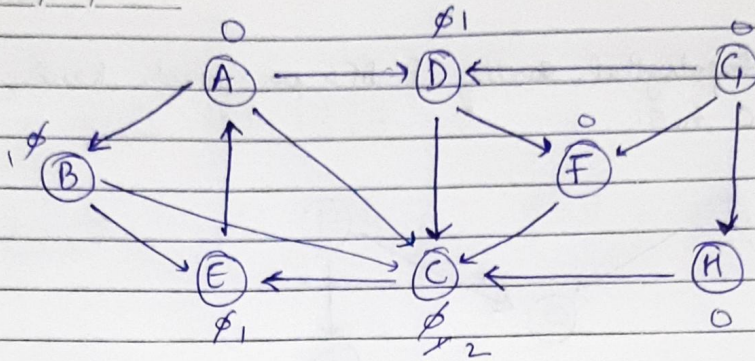


Ans 6 - BFS :-  
 A → parent node  
 B C D → child node

- taking any one child node, say 'B'

E C & furthermore (A, B)

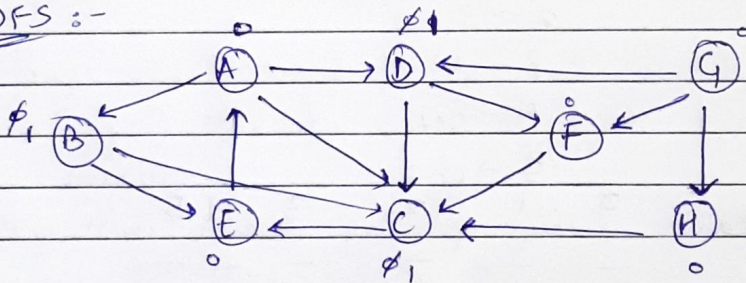




Parent	A	A	A	B	B	D	D	F
Node	B	C	D	E	C	C	F	C

path :-  $A \rightarrow B \rightarrow C$

DFS :-



Nodes stack  
Processed ←

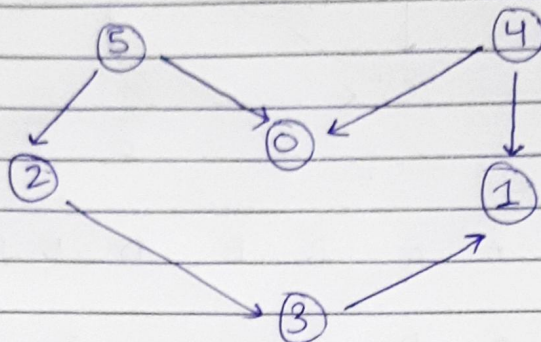
A	A
B	AB
E	BE
C	BEC
C	EC

path  $A \rightarrow B \rightarrow C$



Q7 - Apply topological sorting & DFS on graph having from 0 to 5.

Ans7 -



adjacent list (G)

0 → 1

1 →

2 → 3

3 → 1

4 → 0, 1

5 → 0, 2

visited    0    1    2    3    4    5

—    —    —    —    —    —

Stack (empty)

Step-1 : Topological sort (0), visited(0) = true



list empty (no recursion call)

Stack    0    —

Step-2 : Topological sort [2], visited[2] = true

"    "    [3]    "    "    [3] = true

- already visited, no more recursion



call

stack 

0	1	3	2
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Step-3 :-  $+1$   $[5]$  "  $[5]$  = true  
 2, 0 already visited no recursion call

stack 

0	1	2	3	4	5
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Print all elements of stack from top  $\rightarrow$  bottom

Q8 - Can heap be used to implement priority queue? Name few algorithm where you need to use priority queue?

Ans - Heaps can be used to implement priority queue. It takes  $O(\log N)$  time to insert & delete each element in priority queue.

Algorithm where priority queue can be applied are Dijkstra, Prim's Algorithm.

Q9 - Difference b/w Max & Min. Heap?

It's used to  
maximum element  
in heap

It's used to access  
minimum element in  
heap.