
FIRST ORDER LOGIC

EXERCISES

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0.1 Warm-up exercises

Exercise 0.1. *Non Logical symbols:*

*constants a, b ; functions f, g (f is unary, g is binary); predicates P, R, Q
(unary, binary and ternary, respectively).*

Say whether the following strings of symbols are well formed FOL formulas or terms:

1. $Q(a)$
2. $P(y)$
3. $P(g(b))$
4. $\neg R(x, a)$
5. $Q(x, P(a), b)$
6. $P(g(f(a), g(x, f(x))))$
7. $Q(f(a), f(f(x)), f(g(f(z), g(a, b))))$
8. $R(a, R(a, a))$

Solution. *Well formed formulas: 2., 4., 6., and 7. All other strings are NOT well formed FOL formulas nor terms.*

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Exercise 0.2. *Non Logical symbols:*

constants a, b ; functions f^1, g^2 ; predicates P^1, R^2, Q^3 (the exponent denotes the arity).

Say whether the following strings of symbols are well formed FOL formulas or terms:

1. $R(a, g(a, a));$

0.1. WARM-UP EXERCISES

2. $g(a, g(a, a));$
3. $\forall x. \neg P(x);$
4. $\neg R(P(a), x);$
5. $\exists a. R(a, a);$
6. $\exists x. Q(x, f(x), b) \rightarrow \forall x. R(a, x);$
7. $\exists x. P(R(a, x));$
8. $\forall R(x, a);$

Solution. Well formed formulas: 1., 3., and 6. Well formed terms: 2. All other strings are NOT well formed FOL formulas nor terms.

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Exercise 0.3. Non Logical symbols:

constants a, b ; functions f^1, g^2 ; predicates P^1, R^2, Q^3 (the exponent denotes the arity).

Say whether the following strings of symbols are well formed FOL formulas or terms:

1. $a \rightarrow P(b);$
2. $R(x, b) \rightarrow \exists y. Q(y, y, y);$
3. $R(x, b) \vee \neg \exists y. g(y, b);$
4. $\neg y \vee P(y);$
5. $\neg \neg P(a);$
6. $\neg \forall x. \neg P(x);$
7. $\forall x \exists y. (R(x, y) \rightarrow R(y, x));$

$$8. \forall x \exists y. (R(x, y) \rightarrow (R(y, x) \vee (f(a) = g(a, x))));$$

Solution. Well formed formulas: 2., 5., 6., 7., and 8. All other strings are NOT well formed FOL formulas nor terms.

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Exercise 0.4. Obtain a Skolem Normal Form for the formula:

$$\exists x. \forall y. \forall z. \exists u. \forall v. \exists w. P(x, y, z, u, v, w).$$

Solution. We have the following.

1. x is not preceded by any \forall thus x becomes c .
2. u is preceded by $\forall y$. and $\forall z$. thus u becomes $f(y, z)$.
3. w is preceded by $\forall y$. and $\forall z$. and $\forall v$. thus w becomes $g(y, z, v)$.

The result is

$$\forall y. \forall z. \forall v. P(c, y, z, f(y, z), g(y, z, v)).$$

Exercise 0.5. Obtain a Skolem Normal Form for the formula:

$$\forall x. \exists y. \neg \forall z. \neg ((\neg P(x, y) \wedge Q(x, z)) \vee R(x, y, z))$$

with the matrix in CNF.

Solution. First convert the formula in prenex form with the matrix in CNF:

$$\forall x. \exists y \exists z ((\neg P(x, y) \vee R(x, y, z)) \wedge (Q(x, z) \vee R(x, y, z))).$$

Then,

1. y is preceded by $\forall x$ thus y becomes $f(x)$
2. z is preceded by $\forall x$ thus z becomes $g(x)$

We have

$$\forall x. ((\neg P(x, f(x)) \vee R(x, f(x), g(x))) \wedge (Q(x, g(x)) \vee R(x, f(x), g(x))))$$

0.2. FOL FORMALIZATION

Exercise 0.6. Obtain a Skolem Normal Form for the formula:

$$\forall x.\exists y((\neg P(x, y) \equiv \forall z.Q(x, z)) \rightarrow R(x, y, z))$$

with the matrix in CNF.

Solution. First eliminate implications and equivalences:

$$\forall x.\exists y(((\neg P(x, y) \vee \forall z.Q(x, z)) \wedge (P(x, y) \vee \exists w.\neg Q(x, w))) \vee R(x, y, z)).$$

Notice that by eliminating the equivalence, we get two quantifiers corresponding to the initial $\forall z$. Then, write the formula in prenex form:

$$\forall x.\exists y.\exists w.\forall z.(((\neg P(x, y) \vee Q(x, z)) \wedge (P(x, y) \vee \neg Q(x, w))) \vee R(x, y, z)).$$

Finally either you Skolemize and then write the matrix in CNF or you write the matrix in CNF and then Skolemize. The result is:

$$\forall x.\forall z.((\neg P(x, f(x)) \vee Q(x, z) \vee R(x, f(x), z)) \wedge (P(x, f(x)) \vee \neg Q(x, g(x)) \vee R(x, f(x), z))).$$

Exercise 0.7. Consider a knowledge base containing just two sentences: $P(a)$ and $P(b)$. Does this knowledge base entail $\forall x. P(x)$? Explain your answer in terms of models.

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0.2 FOL Formalization

Note: Some exercises require the use of equality which can be either formalized as a congruence relation or assumed as an additional language symbol with a fixed semantics.

Exercise 0.8. What is the meaning of the following FOL formulas?

1. $bought(Frank, dvd)$
2. $\exists x.bought(Frank, x)$

3. $\forall x.(bought(Frank, x) \rightarrow bought(Susan, x))$
4. $\forall x.bought(Frank, x) \rightarrow \forall x.bought(Susan, x)$
5. $\forall x \exists y.bought(x, y)$
6. $\exists x \forall y.bought(x, y)$

Solution.

1. *"Frank bought a dvd."*
2. *"Frank bought something."*
3. *"Susan bought everything that Frank bought."*
4. *"If Frank bought everything, so did Susan."*
5. *"Everyone bought something."*
6. *"Someone bought everything."*

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Exercise 0.9. Which of the following formulas is a formalization of the sentence:

"There is a computer which is not used by any student"

- $\exists x.(Computer(x) \wedge \forall y.(\neg Student(y) \wedge \neg Uses(y, x)))$
- $\exists x.(Computer(x) \rightarrow \forall y.(Student(y) \rightarrow \neg Uses(y, x)))$
- $\exists x.(Computer(x) \wedge \forall y.(Student(y) \rightarrow \neg Uses(y, x)))$

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Exercise 0.10. Define an appropriate language and formalize the following sentences using FOL formulas.

0.2. FOL FORMALIZATION

1. *All Students are smart.*
2. *There exists a student.*
3. *There exists a smart student.*
4. *Every student loves some student.*
5. *Every student loves some other student.*
6. *There is a student who is loved by every other student.*
7. *Bill is a student.*
8. *Bill takes either Analysis or Geometry (but not both).*
9. *Bill takes Analysis and Geometry.*
10. *Bill doesn't take Analysis.*
11. *No students love Bill.*

Solution.

1. $\forall x.(Student(x) \rightarrow Smart(x))$
2. $\exists x.Student(x)$
3. $\exists x.(Student(x) \wedge Smart(x))$
4. $\forall x.(Student(x) \rightarrow \exists y.(Student(y) \wedge Loves(x, y)))$
5. $\forall x.(Student(x) \rightarrow \exists y.(Student(y) \wedge \neg(x = y) \wedge Loves(x, y)))$
6. $\exists x.(Student(x) \wedge \forall y.(Student(y) \wedge \neg(x = y) \rightarrow Loves(y, x)))$
7. $Student(Bill)$
8. $Takes(Bill, Analysis) \leftrightarrow \neg Takes(Bill, Geometry)$
9. $Takes(Bill, Analysis) \wedge Takes(Bill, Geometry)$
10. $\neg Takes(Bill, Analysis)$
11. $\neg \exists x.(Student(x) \wedge Loves(x, Bill))$

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Exercise 0.11. Define an appropriate language and formalize the following sentences using FOL formulas.

1. *Bill has at least one sister.*
2. *Bill has no sister.*
3. *Bill has at most one sister.*
4. *Bill has (exactly) one sister.*
5. *Bill has at least two sisters.*
6. *Every student takes at least one course.*
7. *Only one student failed Geometry.*
8. *No student failed Geometry but at least one student failed Analysis.*
9. *Every student who takes Analysis also takes Geometry.*

Solution.

1. $\exists x. \text{SisterOf}(x, \text{Bill})$
2. $\neg \exists x. \text{SisterOf}(x, \text{Bill})$
3. $\forall x \forall y. (\text{SisterOf}(x, \text{Bill}) \wedge \text{SisterOf}(y, \text{Bill}) \rightarrow x = y)$
4. $\exists x. (\text{SisterOf}(x, \text{Bill}) \wedge \forall y. (\text{SisterOf}(y, \text{Bill}) \rightarrow x = y))$
5. $\exists x \exists y. (\text{SisterOf}(x, \text{Bill}) \wedge \text{SisterOf}(y, \text{Bill}) \wedge \neg(x = y))$
6. $\forall x. (\text{Student}(x) \rightarrow \exists y. (\text{Course}(y) \wedge \text{Takes}(x, y)))$
7. $\exists x. (\text{Student}(x) \wedge \text{Failed}(x, \text{Geometry}) \wedge \forall y. (\text{Student}(y) \wedge \text{Failed}(y, \text{Geometry}) \rightarrow x = y))$

0.2. FOL FORMALIZATION

8. $\neg \exists x.(Student(x) \wedge Failed(x, Geometry)) \wedge \exists x.(Student(x) \wedge Failed(x, Analysis))$

9. $\forall x.(Student(x) \wedge Takes(x, Analysis) \rightarrow Takes(x, Geometry))$

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Exercise 0.12. Define an appropriate language and formalize the following sentences in FOL:

- *someone likes Mary.*
- *nobody likes Mary.*
- *nobody loves Bob but Bob loves Mary.*
- *if David loves someone, then he loves Mary.*
- *if someone loves David, then he (someone) loves also Mary.*
- *everybody loves David or Mary.*

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Exercise 0.13. Define an appropriate language and formalize the following sentences in FOL:

1. *there is at least one person who loves Mary.*
2. *there is at most one person who loves Mary.*
3. *there is exactly one person who loves Mary.*
4. *there are exactly two persons who love Mary.*
5. *if Bob loves everyone that Mary loves, and Bob loves David, then Mary doesn't love David.*
6. *Only Mary loves Bob.*

Solution. We assume to have a predicate “=” defined as a congruence relation on the elements of the domain.

1. $\exists x. \text{Person}(x) \wedge \text{Loves}(x, \text{Mary})$
2. $\exists x. (\text{Person}(x) \wedge \text{Loves}(x, \text{Mary}) \wedge \forall y. ((\text{Person}(y) \wedge \neg(y = x)) \rightarrow \neg \text{Loves}(y, \text{Mary}))) \vee \forall x. (\text{Person}(x) \rightarrow \neg \text{Loves}(x, \text{Mary}))$
3. $\exists x. (\text{Person}(x) \wedge \text{Loves}(x, \text{Mary}) \wedge \forall y. ((\text{Person}(y) \wedge \neg(y = x)) \rightarrow \neg \text{Loves}(y, \text{Mary})))$
4. $\exists x. \exists y. (\text{Person}(x) \wedge \text{Person}(y) \wedge \neg(x = y) \wedge \text{Loves}(x, \text{Mary}) \wedge \text{Loves}(y, \text{Mary}) \wedge \forall z. ((\text{Person}(z) \wedge \neg(z = x) \wedge \neg(z = y)) \rightarrow \neg \text{Loves}(z, \text{Mary})))$
5. $(\forall x. (\text{Person}(x) \wedge \text{Loves}(\text{Mary}, x) \rightarrow \text{Loves}(\text{Bob}, x)) \wedge \text{Loves}(\text{Bob}, \text{David})) \rightarrow \neg \text{Loves}(\text{Mary}, \text{David})$
6. $\text{Loves}(\text{Mary}, \text{Bob}) \wedge \forall x. (\text{Person}(x) \wedge \neg(x = \text{Mary}) \rightarrow \neg \text{Loves}(x, \text{Bob}))$

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0.3 Formalization and reasoning

Exercise 0.14. Consider the following sentences:

1. All actors and journalists invited to the party are late.
2. There is at least a person who is on time.
3. There is at least an invited person who is neither a journalist nor an actor.

Formalize the sentences and prove that 3. is not a logical consequence of 1. and 2.

Solution. Let A stand for “Actor”, J stand for “Journalist”, I for “Invited” and L for “Late”. Then we have the following:

1. $\varphi_1 := \forall x. ((A(x) \vee J(x)) \wedge I(x) \rightarrow L(x))$

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$$2. \quad \varphi_2 := \exists x. \neg L(x)$$

$$3. \quad \varphi_3 := \exists x. (I(x) \wedge \neg A(x) \wedge \neg J(x))$$

It is sufficient to find an interpretation $\mathcal{I} = \langle D, g \rangle$ for which the logical consequence does not hold. Let Bob, Tom and Mary be three constants. The following table:

	$L(x)$	$A(x)$	$J(x)$	$I(x)$
Bob	F	T	F	F
Tom	T	T	F	T
Mary	T	F	T	T

describes the function g such that, e.g., $g(L) = \{\text{Tom}, \text{Mary}\}$ and so on. Clearly $\models_{\mathcal{I}} \varphi_1 \wedge \varphi_2$ but $\not\models_{\mathcal{I}} \varphi_3$.

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Exercise 0.15. Consider the following sentences:

1. All actors and journalists invited to the party are late.
2. There is at least a person who is on time.
3. There is at least an invited person who is neither a journalist nor an actor.

Formalize the sentences and prove that 3. is not a logical consequence of 1. and 2. using Herbrand theorem specifying an Herbrand domain.

Solution. Let A stand for “Actor”, J stand for “Journalist”, I for “Invited” and L for “Late”. Then we have the following:

$$1. \quad \forall x. ((A(x) \vee J(x)) \wedge I(x) \rightarrow L(x))$$

$$2. \quad \exists x. \neg L(x)$$

$$3. \quad \exists x. (I(x) \wedge \neg A(x) \wedge \neg J(x))$$

We show that the third is entailed by the first two by proving that the negation of the third and the first two

1. $\forall x.((A(x) \vee J(x)) \wedge I(x) \rightarrow L(x))$
2. $\exists x.\neg L(x)$
3. $\forall x.(\neg I(x) \vee A(x) \vee J(x))$

are inconsistent, which is true if and only if their skolemize form

1. $\forall x.((A(x) \vee J(x)) \wedge I(x) \rightarrow L(x))$
2. $\neg l(C)$
3. $\forall x.(\neg I(x) \vee A(x) \vee J(x))$

are inconsistent, which is true, by Hebrand theorem with domain $D_H = \{c\}$, if and only if

1. $((A(c) \vee J(c)) \wedge I(c) \rightarrow L(c))$
2. $\neg L(c)$
3. $(\neg I(c) \vee A(c) \vee J(c))$

is unsatisfiable. The three formulas are satisfiable (a model is obtained by assigning $l(C)$ and $i(C)$ to \perp) and thus the third formula is not a logical consequence of the first 2.

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