

Assignment - 2

1. bool linearSearch(int* arr, int key, int size)

```
{  
    for (int i = 0; i < size; i++ && arr[i] <= key && arr[i] <= key; i++)  
    {  
        if (arr[i] == key)  
            return true;  
    }  
    return false;  
}
```

2. void insertionSort(int* arr, int n)

```
{  
    for (int i = 1; i < n; i++)  
    {  
        k = arr[i];  
        j = i - 1;  
        while (arr[j] > k && j >= 0)  
        {  
            arr[j+1] = arr[j];  
            j--;  
        }  
        arr[j+1] = k;  
    }  
}
```

```
arr[j+1] = k;
```

```
}
```

```
}
```

Recursive

```
void RInsertion (vector<int> a, int n)
```

```
{
```

```
    if (n <= 1)
```

```
        return;
```

```
    RInsertion(a, n-1);
```

```
    int k = a[n-1];
```

```
    int j = n-2;
```

```
    while (j >= 0 && a[j] > k)
```

```
    {
```

```
        a[j+1] = a[j];
```

```
    }
```

```
        j--;
```

```
    arr[j+1] = k;
```

```
}
```

Insertion sort is called online sorting because we don't require the ~~whole~~ whole array instead ~~we~~ ~~the~~ ~~number~~ it can work even if the no. of elements keep increasing. But other algorithms like ~~is~~ selection and bubble sort require the whole array.

3. Bubble

$B \rightarrow O(n)$

$A \rightarrow O(n^2)$

$W \rightarrow O(n^2)$

selection

$B \rightarrow O(n^2)$

$A \rightarrow O(n^2)$

$W \rightarrow O(n^2)$

Insertion

$B \rightarrow O(n)$

$A \rightarrow O(n^2)$

$W \rightarrow O(n^2)$

Merge

$B \rightarrow O(n \log n)$
 $A \rightarrow O(n \log n)$
 $W \rightarrow O(n \log n)$

Quick

$B \rightarrow O(n \log n)$
 $A \rightarrow O(n \log n)$
 $W \rightarrow O(n \log n)$

Count

$B \rightarrow O(n+k)$
 $A \rightarrow O(n+k)$
 $W \rightarrow O(n+k)$

4. Inplace Sort : Bubble, Selection, Insertion, Quick, heap.

Stable \rightarrow Insertion, Selection, Merge

Online : Insertion

5. Iterative

```
bool binarySearch (int* arr, int n, int key)
```

```
{
```

```
    int low = 0, up = n-1;
```

```
    while (low <= up)
```

```
    {
```

```
        int mid = low + (up-low)/2;
```

```
        if (arr[mid] == key)
```

```
            return true;
```

```
        else if (arr[mid] < key)
```

```
            low = mid + 1;
```

```
        else
```

```
            up = mid - 1;
```

```
    }
```

```
    return false;
```

```
}
```

Recursive

~~best case~~ ~~binary~~ binarySearch (int* arr, int low, int up)

{

if (low > up)
return false;

int mid = low + (up - low) / 2;

if (key > arr[mid])

~~low = mid + 1;~~ binarySearch (arr, mid + 1, up);

else if (key < arr[mid])

~~up = mid - 1;~~ binarySearch (arr, low, mid - 1);

else

return true;

}

~~arr~~

T.C

~~Time~~ $O(\log n)$

S.C

$O(1)$

6. $T(n) = T(n/2) + O(1)$

quicksort is best for practical use. It is considered as one of the fastest sorting algorithms for avg. case. Its average case time complexity is $O(n \log n)$ and worst case is $O(n^2)$ but occurs rarely as the pivot element is chosen randomly every time. Also, it is an in-place sort.

9. In an array, an inversion occurs when two elements are not in sorted order.


```
int mergeSort (int arr[], int temp[], int l, int r)
```

```
{
```

```
    int mid, inv = 0;
```

```
    if (r > l) {
```

```
        mid = (r + l) / 2;
```

```
        inv = inv + mergeSort(arr, temp, l, mid);
```

```
        inv = inv + mergeSort(arr, temp, mid + 1, r);
```

```
        inv = inv + merge(arr, temp, l, mid + 1, r);
```

```
    }
```

```
    return inv;
```

```
}
```

```
int merge (int arr[], int temp[], int l, int mid, int r)
```

```
{
```

```
    int i, j, k;
```

```
    int inv = 0;
```

```
    i = l;
```

```
    j = mid;
```

```
    k = l;
```

```
    while ((i <= mid - 1) && (j <= r))
```

```
    {
```

```
        if (arr[i] <= arr[j])
```

```
            temp[k++] = arr[i++];
```

```
        else
```

```
        {
```

```
            temp[k++] = arr[j++];
```

```
            inv += mid - i + 1;
```

```
        }
```

```
    }
```

```
    while (i <= mid - 1)
```

```
        temp[k++] = arr[i++];
```

```
    while (j <= mid - r)
```

```
        temp[k++] = arr[j++];
```



```

    return arr;
}
int mergeSort(int arr[], int size) {
    int temp[size];
    return mergeSort(arr, temp, 0, size-1);
}

```

10. The best case is when the pivot element divides the array into two equal halves ($T.C = O(n \log n)$)
 • The worst case is when the pivot element divides the array into two unbalanced halves ($T.C = O(n^2)$).

19. Merge Sort

Best Case: $T(n) = 2T(n/2) + O(n)$

Worst Case: $T(n) = 2T(n/2) + O(n)$

Quick Sort

Best: $T(n) = 2T(n/2) + O(n)$

Worst: $T(n) = T(n-1) + O(n)$

Similarities

- Both algo. work on divide and conquer strategy.
- Both algo. have best and worst case time complexity of $O(n \log n)$.

Differences

- Merge sort is more efficient on larger arrays.
- Merge sort requires extra space proportional to the size of input array, whereas quicksort is an in-place sort.

12. void selection (int arr[])

{

int n = arr.size();

int i = 0;

for (i = 0; i < n - 1; i++)

{

int min = i;

for (int j = i + 1; j < n; ++j)

{

if (arr[j] < arr[min])

min = j;

}

int m = arr[min];

while (min > i)

{

arr[min] = arr[min - 1];

min--;

}

arr[i] = m;

}

}

13. void bubble (int arr[])

{

int n = arr.size(), i, j = 0;

bool swap;

for (i = 0; i < n - 1; i++)

{

swap = false;

for (j = 0; j < n - i - 1; j++)

{

if (arr[j] > arr[j + 1])

{

swap (arr[j], arr[j + 1]);

}

swap = true;

}


```
if (!swap)
```

```
{
```

```
break;
```

```
}
```

```
}
```

```
}
```