

78. Subsets

Given an array of unique elements, we need to find all the subsets. This is a backtracking problem.

Explanation is best seen in code itself.

CODE

```
vector <vector <int>> powerset;
```

```
vector <int> subset;
```

```
void backtrack ( vector <int> & nums, int start) {
```

```
    powerset.push_back(subset);
```

```
    for(int i = start, i < nums.size(); i++) {
```

```
        subset.push_back(nums[i]);
```

```
        backtrack (nums, i+1);
```

```
        subset.pop_back();
```

```
    }
```

```
vector <vector <int>> subsets (vector <int> & nums) {
```

```
    backtrack (nums, 0);
```

```
    return powerset;
```

```
};
```

In this solution, we ~~just~~ use a depth first ⁽²⁾ approach. We dive into let's say [1] and then add all the subsets that contain 1 before moving on ~~to~~ to the subsets that do not contain 1.

Time complexity $\rightarrow O(2^n)$

space complexity $\rightarrow O(n)$

\rightsquigarrow excluding the space taken by powerset.

\rightarrow Another method ^{is} to use a bitstring

For example for $n=3$, we generate all the bitstrings using backtracking

$(0,0,0)$, ~~$(1,0,0)$~~ , $(1,0,1)$

0 0 1

0 1 1

1 1 1

1 1 0

0 1 0

1 0 0

1 0 1

now, when we have all the bitstrings, we just loop ϕ on each bitstring and add the element to subset if $\text{bit} == 1$.