This is the problem of calculating connected components. We maintain a counter, a contra an array that necords whether the node has been visited and a queue that records neighbours of a node. We are geven adjacency matrix and not List. This affect time complexity.

- (i) Traverse for i from 1 to n. If i has not been villed, invielle the counter. Let villed (i) - true Push i to quem.
- (ii) NOW, while queue is not empty (BFS), keep visiting the neighbours and add them to queme using adjacency matrix.
- An empty queue means one connected component & maversed.
- (in) Return the final counter value.

worst case occurs when Time complexity -> all the components are disconnected. complexity for that case becomes $O(n^2)$.

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space complexity -> O(n), visited owney and stack/ (2)
we can me DFS. BFS both for this problem, below
 implementation uses BFS
# CODE
   int num Provinces (vector < vector < int>> & isconnected){
           int n = is (onnected. size ();
           vector <bool> visited (n, false);
            queue <int> 2; int counter = 0;
           for ( int i=0; i/n; i++) {
                  4 (1 visited[i]){
                        q.push(i);
                        visited[i] =true;
                        counter ++',
                 while (19.empty()){
                        int node = q. front();
                        q.pop()
                        for (int j=0; j<n; j++) {
                            if ( nadi )= 1 22 is (onnected [node][]]
                                  &L visited [j] == false) {
                                  q. push (j);
                                  visited [j] = true:
          return counter; ,
```