(24) Round Robin (RR) Scheduling algorithm Cuteria: Time quantum' Mode: "Preemptive". Process executes upto a given time quantum êrres peulive of buist time and if the process is to ready queul. not completed. It is again sent given 78 = 2 Running in ready queue is very Sequence of processes important. (CT-AT) (TAT-BT) (Clu firstime completion TAT WT Burst Process Arrival fine tine time 0 12 8×1 12 Pi 1 10 42 11 2 4 P2 6 4 9 5 Ready gruen

-> time

context.

switching saving a process temporarily to excurte later.

In the provious example, P. is context switched again and again.

(9) How many times context switching happened.

Running process les waapas bhejo
and naye les laays

# L-2.8 Pre-emptive Priority scheduling algorithm

Criteria: "Priority"

mode: " Pre emptive"

Priority Process Arrival Burst Completion TAT WT
No. Home time time

12 7

20 P2 1 43 8 7 3

30 P3 2 2 4 2 0 40 P7 4 1 5 1 0

Hyne number = higher priority en this case.

	, .	
queue.	P1 P2 P3 P4	. *
que ue.	P1   12   P3   T1	
•	0 1 2 7	*

Running

P1 P2 P5 P3 P4 P2 P1

1 2 3 4 5 8

Mix Burst time (CAPO & 7/10 both) (26) L-2.9 Example of Sche duling

CPU, 10 , CPU Prionty AT Process 10 80 3x0 2 P, 0 Py

Mode: Pre-emphive

Criteria Priority based

lower number has

Find CT of Pr, Pz, Pz, Py

P2/270 P3110

4 5 6 7 8 9 10 RA1/6 P110 14 15 16 17

Py 210

offerent queues for different priority processes.

Every kind of process can have it is own different sheduling algorithm.

Now, the problem a Highest priority system that if there are a Process lot of system processes. Interactive Medium SJF (CPU printly 5 there would be process starvation for medium FCFS lowest Batch and low priority priority s process processes. This is solved using mulfilevel

feedback queue scheduling.

L-2.11 rulilever Feedback guen

lowest (both)

PRO, TQ=2

If completed then out

ROZ TQ=8

PROWNER

PROY FCFS

Manual

brocess does not starve because its quere keeps on changing