Heap - almost complete binary tree i.e (h.1)

levels are completely filled and last level has all

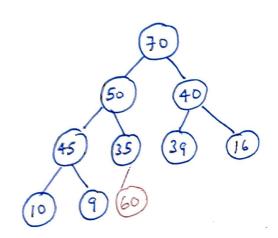
elements shifted to left

Max heap - For every node i, the value of node

is less than or equal to it is

parent value.

A [parent(i)] Z ACi]



70 50 40 45 35 39 16 10 9 60

Array representation of binary tree

Insertion in May heap

- First insert a leaf node that maintains almost complete property of binary tree.
- ex. inscrt 60
- Now, start comparing 60 with it is parent and keep swapping until parent 7 60.
- To get parent node, use away representation.
- -for ex. 60 at index 10, parent = $\begin{bmatrix} \frac{1}{2} \end{bmatrix} = 5 \rightarrow 35$ Swap(1°135), now, parent = $\begin{bmatrix} \frac{1}{2} \end{bmatrix} = 2 \rightarrow 50$, swap (5°,60), now, parent $\begin{bmatrix} \frac{2}{2} \end{bmatrix} = 1$, $\rightarrow 70$, now stop.

```
Time complexity for Presentian - O(logn),
    because maximum we can insert an element
    that is larger than all the elements already there.
                                                                                                             -> code
               insert Heap ( Arr, n, new-value) {
                                            A[n] = new-value
                              int duild = n
                                               Int parent = 0;
                                                                                                                                                                                                 while (child >1) {
                                                              parent = wild/2;
                                                                                         ef (A[parent] < A[child]) {
                                                                                  swap (A[parent], A[dild]);
child = parent;
elce s
     Now the second s
                                                         retun
```

.3.

- root node is deleted
- pick last element and place it at root
- now size of mar heap is decreased by I.
- now compare root with duildren until mar_heap
 is satisfied
- in array, keep replacing parent with larger duild. until end of array is reached or max heap property is satisfied.
- → Heap-Sort

Two steps

(ii) telete data one by one fill away is sorted.

complete heap sort implementation -

```
void
       max heapity (int arris, int n, inti) {
          int largest = 1;
           int left = 2xi+1;
           Int right = 2x1+2;
          if ( left < n la over [Left] > orr [largest]) {
                  largest = left;
          if (right <n el arr[right] > arr[largest]) {
                    largest = right;
          if (i |= largest) {
               swap (arr ["], arr [largest]);
               maxheapify ( our, n, largest);
       buildheap (int out], int n) {
        for ( int i = (1/2) +1; i>-1; i--).{
                max heapify ( aus, n, i);
```

void heapsort (Pnt ava [], Ent n) {

build heap (aru, n);

for (int i: n-1; i>-1; i--) {

Swap (arr [i], arr[o]);

Maxheapify (are, i, 0);

}.

J