

Algorithms Project Finding median

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Introduction

This project is about finding the median using sorting algorithms, an algorithm such as QuickSort, SelectionSort, medianOfMedians and radixSort were chosen, to measure and compare between algorithms in finding the medians that are better in performance and faster in extracting the result for sorted and unsorted data. It is in a program to determine the time for the execution of the algorithm, and the graphics library was used tKinter to create the interface of the program with specifying the type of data if it was sorted or unsorted due to the natural time difference between sorted and unsorted data. The data was read through a file with a range of 10000 data.

The main goal of the project is to compare the Sorting algorithms in finding the median.

Algorithms Discerption/Definition, Time Complexity and Code

1.1 Finding Median with Median_Selection_Sort

1.1.1 Discerption/Definition

The Algorithm will use Selection Sort to sorts the data and then finding the median which is the middle of them all.

The selection sort algorithm sorts an array by repeatedly finding the minimum element (considering ascending order) from unsorted part and putting it at the beginning. The algorithm maintains two subarrays in a given array.

1.1.2 Time Complexity

The time complexity of this algorithm is $O(n^2)$, Because of calculating the median part takes O(1), What remain is the time complexity of the Selection Sort Algorithm part, which is $O(n^2)$.

1.1.3 Code

```
# This funaction will sort the data and find the median with Slection Sort Algorithm
def Median_Selection_Sort(A):
   Len = len(A) # O(1)
   # Traverse through all array elements
   for i in range(Len-1):# O(n)
       # Find the minimum element in remaining
       # unsorted array
       min_i = i# O(1)
       for j in range(i+1, Len): # O(n)
            if(A[j] < A[min_i]): # O(1)</pre>
               A[j], A[min_i] = A[min_i], A[j] # O(1)
       # Swap the found minimum element with
        # the first element
        if(min_i != i): # O(1)
           A[\min_{i}], A[i] = A[i], A[\min_{i}] # O(1)
# Calculating median O(1)
   # if the data have an even number of elements we have to get the average of the 2 elements in the middle
   if((Len \% 2) == 0): # O(1)
       m1 = math.floor((Len)/2) # O(1)
       m2 = math.floor((Len-1)/2) # O(1)
       med = (A[m1] + A[m2])/2 # O(1)
   # otherwise we get the mid of them all
   else:
       med = A[math.floor((Len)/2)] # O(1)
   return med
```

Figure 1 Median Selection Sort Algorithm code

1.2 Finding Median with medianOfMedians

1.2.1 Discerption/Definition

It uses a divide and conquer strategy to efficiently compute the i_th smallest number in an unsorted list of size n. Selection algorithms are often used as part of other algorithms; for example, they are used to help select a pivot in quicksort. The median-finding algorithm can find the i_th smallest element in a list in O(n) time.

How it works:

The median-of-medians algorithm is a deterministic linear-time selection algorithm. The algorithm works by dividing a list into sublists and then determines the approximate median in each of the sublists. Then, it takes those medians and puts them into a list and finds the median of that list. It uses that median value as a pivot and compares other elements of the list against the pivot. If an element is less than the pivot value, the element is placed to the left of the pivot, and if the element has a value greater than the pivot, it is placed to the right. The algorithm recurses on the list, until found the value it is looking for.

1.2.2 Time Complexity

This algorithm runs in O(n) time. n is divided into n/5 sublists of five elements each. If M is the list of all of the medians from these sublists, then M has n/5 median for each of the n/5 sublists. Let's call the median of this list (the median of the medians) p. Half of the n/5 elements in M are less than p.

Half of n/5 = n/10 For each of these n/10 elements, there are two elements that are smaller than it (since these elements were medians in lists of five elements two elements were smaller and two elements were larger). Therefore, there are 3n/10 < p and, in the worst case, the algorithm may have to recurse on the remaining 7n/10 elements. The time for dividing lists, finding the medians of the sublists, and partitioning takes

T(n) = T(n/5) + O(n) time, and with the recursion factored in, the overall recurrence to describe the median-of-medians algorithm is.

$$T(n) \le T\left(\frac{n}{5}\right) + T\left(\frac{7n}{10}\right) + O(n).$$

1.2.3 Code

```
### shis function do the partitioning depends on pivot value
| def partition(data, pivot); | 3 = 0 = 0(1) |
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```

Figure 2 medianOfMedians Algorithm code

1.3 Finding Median with quickSort

1.3.1 Discerption/Definition

Quick sort is based on the divide-and-conquer approach based on the idea of choosing one element as a pivot element and partitioning the array around it such that: Left side of pivot contains all the elements that are less than the pivot element Right side contains all elements greater than the pivot.

1.3.2 Time Complexity

The time complexity of this algorithm is $O(n^2)$.

1.3.3 Code

```
# -*- coding: utf-8 -*-
Created on Tue May 17 08:41:57 2022
@author: Manar
def partition(arr, low, high):
    i = (low-1)  # index of smaller element constant-time
pivot = arr[high]  # pivot constant-time
    for j in range(low, high):
         # If current element is smaller than or
        # equal to pivot
        if arr[j] \leftarrow pivot: \#O(n)
             # increment index of smaller element
            i = i+1 #constant-time
    arr[i], arr[j] = arr[j], arr[i]
arr[i+1], arr[high] = arr[high], arr[i+1] #constant-time
    return (i+1) #constant-time
#0(n)
def quickSort(arr, low, high):
    if len(arr) == 1:
        return arr
    if low < high: #constant-time</pre>
        # pi is partitioning index, arr[p] is now
        pi = partition(arr, low, high) #O(n)
        quickSort(arr, low, pi-1) # Sort left of partitioned # n-1
        quickSort(arr, pi+1, high) # Sort right of partitioned #n
# = O(n-1)+O(n)
# Worst case O(n^2)
```

Figure 3 quickSort Algorithm code

1.4 Finding Median with radixSort

1.4.1 Discerption/Definition

Radix sort is a sorting algorithm that sorts the elements by first grouping the individual digits of the same place value. Then, sort the elements according to their increasing/decreasing order. First, we will sort elements based on the value of the unit place. Then, we will sort elements based on the value of the tenth place. This process goes on until the last significant place.

1.4.2 Time Complexity

The radix sort that uses counting sort as an intermediate stable sort, the time complexity is O(d(n+k)).

1.4.3 Code

```
# Radix sort
  Using counting sort to sort the elements in the basis of significant places
    size = len(array)
     #auxiliary array for assigning sorted data
     #This array is used for storing the count of the elements in the array.
     count = [0] * 10
    for i in range(0, size):
   index = array[i] // place
   count[index % 10] += 1
     # Calculate cumulative count
    #0(size)
for i in range(1, 10):
    count[i] += count[i - 1]
    # Find the index of each element of the original array in count array # place the elements in output array
    #0(max) i = size - 1 #placing integers to it corresponding place in array in backward while i >= 0:  
          index = array[i] // place
          output[count[index % 10] - 1] = array[i] count[index % 10] -= 1 #decrease count by i -= 1 #deacrement by 1
    #filling original array with sorted data #0(size)
    for i in range(0, size):
    array[i] = output[i]
 # Main function to implement radix sort
def radixSort(array):
     # Get maximum element in data
     max_element = max(array)
    # Apply counting sort to sort elements based on place value. #0(moxDigit) place = 1 while max_element // place > 0:
    countingSort(array, place)
place *= 10 #increasse the place by 10 each loop
```

Figure 4 radixSort Algorithm code

1.5 tkinter

1.5.1 Interface Window

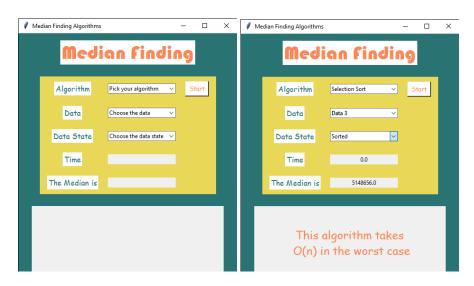


Figure 5 and 6 program interface window

1.6 Read/Write data

1.6.1 Read data

Reading the data from the file that has been chosen by the user.

```
def readData():
     select = dataChoice.get() # get the data that have been chosen
    global data, n
     if select == "Data 1":
         with open ('data1.txt', 'r') as d:
    data0 = d.readlines() #to read the content of the data file
              data= [ x.strip() for x in data0] #to delete the new line character from the data
              for i in range (len(data)):
                  data[i] = int(data[i])
     elif select == "Data 2":
         with open ('data2.txt', 'r') as d:
              data0 = d.readlines()  #to read the content of the data file
data= [ x.strip() for x in data0]  #to delete the new line character from the data
              for i in range (len(data)):
                  data[i] = int(data[i])
    elif select == "Data 3":
    with open ('data3.txt', 'r') as d:
              data0 = d.readlines()  #to read the content of the data file
data= [ x.strip() for x in data0]  #to delete the new line character from the data
              for i in range (len(data)):
    data[i] = int(data[i])
n = len(data) # reset the value of n
    state = dataState.get() # get the state of data
     if state == "Sorted": # apply sorting if the state equal to 'Sorted'
         data.sort()
```

Figure 7 reading the data from chosen file code

1.6.2 Write data

Write down (store it) in the info.txt file.

Figure 8 write(store) the algorithm name, chosen data, state of data and total time to .txt file

167.2 .txt file

Some of data capture from info.txt file.

Figure 9 Stored data in the .txt file

Figure 10 Stored data in the .txt file

Chapter 2

Data Analysis

2.1 Analyse the data and compare

2.1.1 Analyse and Compare the data

| Data | Data State | Algorithm | Time |
|--------|------------|--------------------------|----------------------|
| Data 1 | Sorted | Selection Sort | 0.015622377395629883 |
| Data 1 | Sorted | Selection Sort | 0.015571355819702148 |
| Data 1 | Sorted | Selection Sort | 0.015573263168334961 |
| | A | VG: 0.015588998794555664 | |
| Data 1 | Unsorted | Selection Sort | 0.02383112907409668 |
| Data 1 | Unsorted | Selection Sort | 0.03119659423828125 |
| Data 1 | Unsorted | Selection Sort | 0.031198501586914062 |
| | A | VG: 0.028742074966430664 | |
| Data 2 | Sorted | Selection Sort | 0.015681743621826172 |
| Data 2 | Sorted | Selection Sort | 0.03119826316833496 |
| Data 2 | Sorted | Selection Sort | 0.015496969223022461 |
| ' | A | VG: 0.020792325337727863 | |
| Data 2 | Unsorted | Selection Sort | 0.02361917495727539 |
| Data 2 | Unsorted | Selection Sort | 0.03124213218688965 |
| Data 2 | Unsorted | Selection Sort | 0.03126931190490723 |
| ' | A | VG: 0.028710206349690754 | |
| Data 3 | Sorted | Selection Sort | 0.015677928924560547 |
| Data 3 | Sorted | Selection Sort | 0.01556849479675293 |
| Data 3 | Sorted | Selection Sort | 0.015616893768310547 |
| | A | VG: 0.015621105829874674 | |
| Data 3 | Unsorted | Selection Sort | 0.046759843826293945 |
| Data 3 | Unsorted | Selection Sort | 0.04674959182739258 |
| Data 3 | Unsorted | Selection Sort | 0.031116962432861328 |
| | A | VG: 0.04154213269551595 | 1 |

| Data 1 | Sorted | Quick Sort | 0.04717588424682617 |
|--------|----------|--------------------------|----------------------|
| Data 1 | Sorted | Quick Sort | 0.046880245208740234 |
| Data 1 | Sorted | Quick Sort | 0.04685544967651367 |
| | A | VG: 0.046970526377360024 | |
| Data 1 | Unsorted | Quick Sort | 0.015588521957397461 |
| Data 1 | Unsorted | Quick Sort | 0.01573467254638672 |
| Data 1 | Unsorted | Quick Sort | 0.015735387802124023 |
| | A | VG: 0.015686194101969402 | |
| Data 2 | Sorted | Quick Sort | 0.04711294174194336 |
| Data 2 | Sorted | Quick Sort | 0.046811580657958984 |
| Data 2 | Sorted | Quick Sort | 0.046935081481933594 |

| Data 2 | Unsorted | Quick Sort | 0.015494346618652344 |
|--------|----------|---------------------------|-----------------------|
| Data 2 | Unsorted | Quick Sort | 0.01557159423828125 |
| Data 2 | Unsorted | Quick Sort | 0.0009958744049072266 |
| | | AVG: 0.01068727175394694 | |
| Data 3 | Sorted | Quick Sort | 0.04681658744812012 |
| Data 3 | Sorted | Quick Sort | 0.046808481216430664 |
| Data 3 | Sorted | Quick Sort | 0.04770660400390625 |
| | A | AVG: 0.047110557556152344 | |
| Data 3 | Unsorted | Quick Sort | 0.01557159423828125 |
| Data 3 | Unsorted | Quick Sort | 0.015708446502685547 |
| Data 3 | Unsorted | Quick Sort | 0.005028963088989258 |
| | A | AVG: 0.012103001276652018 | |
| | | | |
| Data 1 | Sorted | Radix Sort | 0.015592336654663086 |
| Data 1 | Sorted | Radix Sort | 0.015706777572631836 |
| Data 1 | Sorted | Radix Sort | 0.01560354232788086 |
| | | AVG 0.01563421885172526 | |
| Data 1 | Unsorted | Radix Sort | 0.015726566314697266 |
| Data 1 | Unsorted | Radix Sort | 0.015597105026245117 |
| Data 1 | Unsorted | Radix Sort | 0.015553951263427734 |
| I | | AVG 0.015625874201456707 | |
| Data 2 | Sorted | Radix Sort | 0.008651256561279297 |
| Data 2 | Sorted | Radix Sort | 0.002687692642211914 |
| Data 2 | Sorted | Radix Sort | 0.00099945068359375 |
| | | AVG 0.004112799962361653 | |
| Data 2 | Unsorted | Radix Sort | 0.0009958744049072266 |
| Data 2 | Unsorted | Radix Sort | 0.0019981861114501953 |
| Data 2 | Unsorted | Radix Sort | 0.0027408599853515625 |
| | A | VG 0.0019116401672363281 | <u></u> |
| Data 3 | Sorted | Radix Sort | 0.015564918518066406 |
| Data 3 | Sorted | Radix Sort | 0.015635013580322266 |
| Data 3 | Sorted | Radix Sort | 0.015699148178100586 |
| - | A | AVG 0.015633026758829754 | ' |
| Data 3 | Unsorted | Radix Sort | 0.015572071075439453 |
| Data 3 | Unsorted | Radix Sort | 0.015583276748657227 |
| Data 3 | Unsorted | Radix Sort | 0.015585899353027344 |
| | F | AVG 0.015580415725708008 | |
| D . 1 | g | N. 11 0 11 | 0.0004002500450500 |
| Data 1 | Sorted | Median of medians | 0.008489370346069336 |
| Data 1 | Sorted | Median of medians | 0.015681743621826172 |

Median of medians

Data 1

Unsorted

0.015589237213134766

| Data 1 | Unsorted | Median of medians | 0.015616893768310547 |
|--------|----------|--------------------------|----------------------|
| Data 1 | Unsorted | Median of medians | 0.015712499618530273 |
| | A | AVG 0.015639543533325195 | |
| Data 2 | Sorted | Median of medians | 0.015593290328979492 |
| Data 2 | Sorted | Median of medians | 0.016228199005126953 |
| Data 2 | Sorted | Median of medians | 0.01569533348083496 |
| | | AVG 0.0158389409383138 | |
| Data 2 | Unsorted | Median of medians | 0.015612602233886719 |
| Data 2 | Unsorted | Median of medians | 0.015568733215332031 |
| Data 2 | Unsorted | Median of medians | 0.015567541122436523 |
| | A | VG 0.015582958857218424 | |
| Data 3 | Sorted | Median of medians | 0.015494585037231445 |
| Data 3 | Sorted | Median of medians | 0.01562809944152832 |
| Data 3 | Sorted | Median of medians | 0.0157470703125 |
| | A | VG 0.015623251597086588 | |
| Data 3 | Unsorted | Median of medians | 0.01550436019897461 |
| Data 3 | Unsorted | Median of medians | 0.015582561492919922 |
| Data 3 | Unsorted | Median of medians | 0.015591144561767578 |
| | I | AVG: 0.01555935541788737 | |

3 set of different data of size n = 1000

Selection:

AVG sorted data = 0.0173341433207194.

AVG unsorted data=0.03299813800387912.

Quicksort:

AVG sorted data = 0.04701142840915256.

AVG unsorted data=0.012825489044189453.

Radix:

AVG sorted data = 0.011793348524305558.

AVG unsorted data=0.011039310031467015

Median of medians:

AVG sorted data = 0. 01490473747253418.

AVG unsorted data=0.015593952602810329

2.1.2 Conclusion

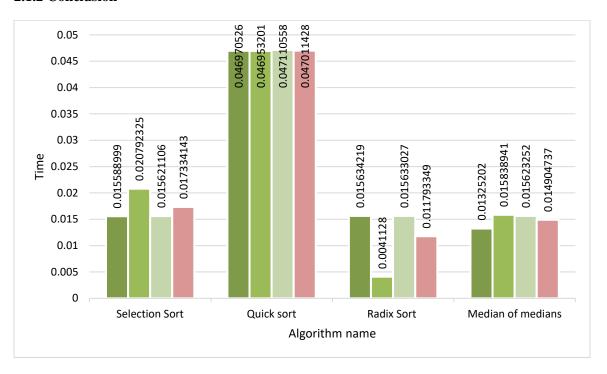


Figure 11 Chart of Sorted data

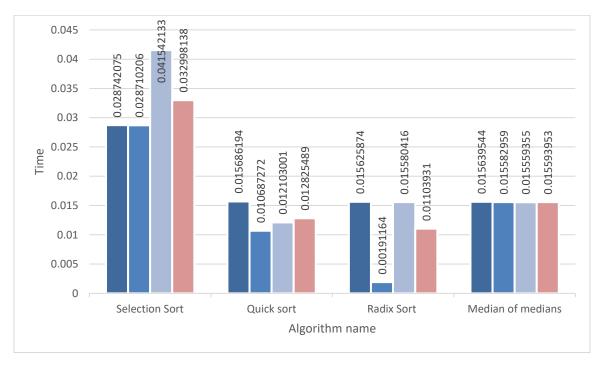


Figure 12 Chart of Unsorted data

A comparison of algorithms has been studied by a data size of 1000 divided into 3 data files was used for measurement and analysis to find the median. The sample was tested once for each datafile and once for its state Stored and Unsorted, then it's identified the average of each datafile of each algorithm, After that the average for each sorted and unsorted data of each algorithm. We noticed through the graph that Quicksort has different time complexity depend on unsorted or sorted data, it shows that the Quicksort algorithm If the data was sorted in the process of finding the median in the it takes 0.047011428 seconds, and if it is unsorted it takes the time 0.012825489044189453 seconds, this is a large time difference compared to other algorithms. On the other hand, Radix sort has a small difference between the time in the case of sorted 0.011793348524305558 seconds, and unsorted 0.011039310031467015 seconds. In a same condition, the Selection Sort have a big difference in the case of the data is unsorted 0.03299813800387912 again, and if it is sorted 0.0173341433207194 second. As for the median of medians, they have the same time, but varying in the case that they are unsorted 0.015593952602810329 seconds, and sorted 0.01490473747253418 seconds, it can be because the algorithm is special to find the median, so it has a convergent and varying time, unlike Quicksort.

Closure

depending on what we studied and learned and what we saw when we test these algorithms and analyzed them, we can confirm that some of them take linear time such as Radix sort and Median of medians algorithms and they could be better choice than others to find the median in large data sets, on the other hand there are some of the algorithms take more than O(n) such as Quick sort and Selection sort and they're not the best choice for large data sets.

| Student name | Responsibilities |
|--------------|--|
| Razan | medianOfMedians Algorithm All GUI file Analysis |
| Heba | radixSort Algorithm Presentation Analysis |
| Manar | <i>quicksort</i> Algorithm Report Analysis |
| Shaymaa | Median_Selection_Sort Algorithm Report Only read/write data, data file and state selection in the GUI file |

Reference:

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