### [CS304] Introduction to Cryptography and Network Security

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## 1 Subgroup

A non-empty subset H in a group (G, \*) is a subgroup of G if H is itself a group with respect to the operation \* of G. If it is a proper subset and a group with respect to \* of G and  $H \neq G$ , then H is called a proper subgroup of (G, \*).

- 1.  $H \subseteq G$  or not
- 2. H is itself a group with \*

### Property:

$$(G,*)$$
 is a Group  
 $a \in G \to a*a \in G, a*a*a \in G$   
 $a*a = a^2, a*a*a = a^3 \in G$   
 $a^i = a*a*...*a \in G, (a*a)^i \to i$  operations \*

# 2 Generators and Cyclic Group

Consider a group (G,\*). Let  $\alpha \in G$ . The identity element  $\alpha^0$  belongs to G. Therefore,

$$\alpha^{0} * \alpha = \alpha^{1}$$
$$\alpha^{1} * \alpha = \alpha^{2}$$
$$\alpha^{2} * \alpha = \alpha^{3}$$

**Note:** The \* here is not multiplication, it is a binary operation not necessarily multiplication.  $\alpha^1, \alpha^2, \alpha^3$  and so on, are just notation of using the binary operation \* on same element.

Since, G is closed under \*, any two elements belonging to G, will give the result in G on performing

the binary operation \*. Since  $\alpha^0 \in G$  and  $\alpha \in G$ , therefore  $\alpha^1 \in G$ . Now, since  $\alpha^1 \in G$ , therefore  $\alpha^2 \in G$ , and so on. That means,

$$\alpha^0,\alpha^1,\alpha^2,\ldots\in G$$

The set  $\alpha^0, \alpha^1, \alpha^2, \dots$  is denoted by  $\langle \alpha \rangle$ . Also,  $\langle \alpha \rangle \subseteq G$ .  $\alpha$  is called the generator of (G, \*) iff:

for any 
$$b \in G \; \exists \; i \geq 0$$
 such that  $b = \alpha^i$  and hence  $G \subseteq \langle \alpha \rangle$ .

### We can conclude that $(G,*) = \langle \alpha \rangle$

A group is called a cyclic group if there is an element  $\alpha \in G$ , such that for every  $b \in G$ , there is an integer i with  $b = \alpha^i$ . In simple words, every element in G can be expressed as some exponent of  $\alpha$ , then  $\alpha$  is the generator of G.

#### Order of an element 2.0.1

Consider  $(G,^*)$  and |G|: finite. Let  $a \in G$ .

We already know that  $a^0$  is identity. Now, the order of an element is the least positive integer m such that  $a^m = e$ .

$$o(a) = m$$
 such that  $a^m = e$ 

Since  $a^m = e$ , so  $a^{m+1} = a$ ,  $a^{m+2} = a^2$  and so on. So we define a set H such as:

$$\mathbf{H} = \{a^0, a^1, a^2, \dots, a^{m-1}\}\$$

We understand that

- $H \subseteq G$
- H is a group under \*

## Lagrange's Theorem:

If G is a finite group and H is a subgroup of g then |H| divides |G|.

• G is a finite group

$$a \in G$$

$$O(a) \mid |G|$$

$$\Rightarrow a \in G$$

$$H = \{ e = a^0, a^1, a^2, \dots, a^{O(a)-1} \}$$

H is a subgroup of G

From Lagrange's theorem:

$$|H| \mid |G|$$
  
 $\Rightarrow O(a) \mid |G|$ 

• If the order of  $a \in G$  is t

then 
$$O(a^k) = \frac{t}{\gcd(t,k)}$$

• If 
$$gcd(t,k) = 1$$

then 
$$O(a^k) = t = O(a)$$
  
 $\Rightarrow | < a^k > | = | < a > |$ 

$$\Rightarrow | < a^n > | = | < a > |$$

 $\Rightarrow x = (a^k)^i = a^{ki} = \langle a \rangle$ 

 $x \in \langle a^k \rangle$ 

$$\langle a^k \rangle \subseteq \langle a \rangle$$
  
 $\langle a^k \rangle = \langle a \rangle$ 

$$\langle a^k \rangle \subseteq \langle a \rangle$$

$$\Rightarrow < a^k > = < a >$$

 $a^k$  is also a generator of  $\langle a \rangle$ 

$$\langle a^k \rangle = \langle a \rangle$$
 Subgroup generated by a

$$\langle a \rangle = \langle a^k \rangle$$
 Subgroup generated by  $a^k$ 

## 3 Ring:

### 3.1 Introduction:

A ring  $(R, +_R, \times_R)$  consists of one set R with two binary operations arbitrarily denoted by  $+_R$  (addition) and  $\times_R$  (multiplication) on R, satisfying the following properties:

- 1.  $(R, +_R)$  is an abelian group with the identity element  $0_R$
- 2. The operation  $\times_R$  is associative,i.e,  $a \times_R (b \times_R c)$
- 3. There is a multiplication identity denoted by  $1_R$  with  $1_R \neq 0_R$  such that
  - $1_R \times_R a = a \times_R 1_R = a$   $\forall a \in R$
- 4. The operation  $\times_R$  is distributive over  $+_n$ , i.e,  $(b +_R c) \times_R a = (b \times_R a) +_R (c \times_R a)$   $a \times_R (b +_R c) = (a \times_R b) +_R (a \times_R c)$

### Field

A field is a non-empty set F together with two binary operation +(addition) and \*(multiplication) fow which the following properties are satisfied

- (F, +) is an abelian group
- If  $0_F$  denotes the additive identity element of (F,+) then  $(F \setminus \{0_F\},*)$  is a commutative/abelian group.
- $\forall$  a,b,c  $\in$  F, we have,

$$a^*(b+c) = (a^*b) + (a^*c)$$

#### Note:

- $(Z, +, \cdot)$  is not a field because inverse does not exist
- $\bullet$   $(Q,+,\cdot)$

(Q, +): abelian group

0: additive identity

1: multiplicative identity

 $(Q \setminus \{0\}, \cdot \text{ forms an abelian group.})$ 

Hence, it is a field.

**Example:** Is  $(\mathbb{F}_p, +_p, *_p)$  a field, where p is a prime number?

**Solution:** We know that  $(\mathbb{F}_p, +_p)$  an abelian group with identity element 0. Now, the set  $\mathbb{F}_p - \{0\}$  has existing multiplicative inverse iff gcd(x, p) = 1 for each  $x \in \mathbb{F}_p - \{0\}$ . Since, p is prime, gcd(x, p) = 1 for all possible integers that x can take. Hence,  $(\mathbb{F}_p, +_p, *_p)$  is a field.

### 4 Field Extension

Suppose  $K_2$  is a field with addition(+) and multiplication(\*).

Suppose  $K_1K_2$  is closed under both these operations such that  $K_1$  itself is a field with the restriction of + and \* to the set  $K_1$ . Then  $K_1$  is called a subfield of  $K_2$  and  $K_2$  is called a field extension of  $K_1$ .

As  $K_1$  is a subset of  $K_2$ . Let F be a field (F, +, \*). Consider the polynomial ring F[x], which consists of all polynomials with coefficients in the field F:

$$F[x] = \{a_0 + a_1x + a_2x^2 + \dots | a_i \in F\}$$

The addition operation of two polynomials in F[x]:

$$(a_0 + a_1x + a_2x^2 + \dots + a_{n-1}x^{n-1}) + (b_0 + b_1x + b_2x^2 + \dots + b_{n-1}x^{n-1})$$

results in:

$$(a_0 + b_0) + (a_1 + b_1)x + (a_2 + b_2)x^2 + \dots + (a_{n-1} + b_{n-1})x^{n-1}$$

where  $a_i+b_i$  is the additive operation in the field F. The multiplication operation of two polynomials in F[x]:

$$(a_0 + a_1x + a_2x^2 + \ldots + a_{n-1}x^{n-1}) * (b_0 + b_1x + b_2x^2 + \ldots + b_{n-1}x^{n-1})$$

results in:

$$(a_0b_0) + (a_0b_1 + a_1b_0)x + \ldots + (a_{n-1}b_{n-1})x^{2n-2}$$

### 5 Irreducible Polynomial

A polynomial  $P(x) \in F[x]$  of degree  $n \ge 1$  is called irreducible if it cannot be written in the form of  $P_1(x) * P_2(x)$  with  $P_1(x), P_2(x) \in F[x]$  and degree of  $P_1(x), P_2(x)$  must be greater than or equal to 1. It means that P(x) is irreducible if it can not be factorised.

**Example:**  $x^2 + 1 \in \mathbb{F}_2[x]$ .

**Solution:**  $(x+1)*(x+1) = x^2 + (1+1) \cdot x + 1 = x^2 + 1$ . Therefore,  $(x^2+1) = (x+1)*(x+1)$  in  $\mathbb{F}_2[x]$ . Hence,  $(x^2+1)$  is reducible in  $\mathbb{F}_2[x]$ . Note that it is not possible to factor  $x^2+1$  in  $\mathbb{R}[x]$ , where  $\mathbb{R}$  is set of real numbers.

# 6 Advanced Encryption Standard:

- It is Standardized by NIST.
- Rijndael winner of Advanced Encryption Standard Competition.
- Winner of the Competition was named AES.

AES is based on -

- 1. Iterative block cipher.
- 2. It is based on SPN.

# 6.1 Types of AES:

- 1. AES 128
  - (a) Block size = 128 bit
  - (b) Number of Rounds = 10
  - (c) Secret key size = 128 bit
- 2. AES 192
  - (a) Block size = 128 bit
  - (b) Number of Rounds = 12
  - (c) Secret key size = 192 bit
- 3. AES 256
  - (a) Block size = 128 bit
  - (b) Number of Rounds = 14
  - (c) Secret key size = 256 bit