

Análisis de algoritmos iterativos

ANÁLISIS Y DISEÑO DE ALGORITMOS

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Herramientas que vamos a utilizar

- Consideraciones:
 - A todas las divisiones se les aplicará la función “piso”. La función `piso(x)` devuelve el entero más pequeño o igual a x . Por ejemplo, `piso(3.3) = piso(3.9999) = piso(3.5) = 3`.
 - Todas las operaciones como asignaciones, comparaciones, decisiones, operaciones matemáticas básicas (+, -, *, /, <<, >>) se realizan en un tiempo constante.

Herramientas que vamos a utilizar

- \log_b es una función estrictamente creciente y uno a uno:
- $\log_b 1 = 0$
- $\log_b b^a = a$
- $\log_b (XY) = \log_b X + \log_b Y$
- $\log_b X^a = a \log_b X$
- $X^{\log_b Y} = Y^{\log_b X}$

Herramientas que vamos a utilizar

$$\sum_1^n c = n * c$$

$$\sum_1^n i = \frac{n * (n + 1)}{2}$$

$$\sum_1^n i^2 = \frac{2n^3 + 3n^2 + n}{6}$$

Ejemplo 1

```
procedure EXAMPLE1 (n) :  
  sum ← 0  
  for i ← 1 until n do  
    sum ← sum + 1  
  return sum
```

$$\sum_1^n c = n * c$$

$$\sum_1^n i = \frac{n * (n + 1)}{2}$$

$$\sum_1^n i^2 = \frac{2n^3 + 3n^2 + n}{6}$$

Ejemplo 2

```
procedure EXAMPLE2 (n) :  
  sum ← 0  
  for i ← 1 until n do  
    for j ← 1 until n do  
      sum ← sum + 1  
return sum
```

$$\sum_1^n c = n * c$$

$$\sum_1^n i = \frac{n * (n + 1)}{2}$$

$$\sum_1^n i^2 = \frac{2n^3 + 3n^2 + n}{6}$$

Ejemplo 3

```
procedure EXAMPLE3 (n) :  
  sum ← 0  
  for i ← 1 until n do  
    for j ← 1 until n do  
      for k ← 1 until n do  
        sum ← sum + 1  
  return sum
```

$$\sum_1^n c = n * c$$

$$\sum_1^n i = \frac{n * (n + 1)}{2}$$

$$\sum_1^n i^2 = \frac{2n^3 + 3n^2 + n}{6}$$

Ejemplo 4

```
procedure EXAMPLE4 (n) :  
  sum ← 0  
  for i ← 1 until n step +2 do  
    sum ← sum + 1  
  return sum
```

N = 4	N = 8	N = 16

$$\sum_1^n c = n * c$$

$$\sum_1^n i = \frac{n * (n + 1)}{2}$$

$$\sum_1^n i^2 = \frac{2n^3 + 3n^2 + n}{6}$$

Ejemplo 5

```
procedure EXAMPLE5 (n) :  
  sum ← 0  
  for i ← 1 until n step *2 do  
    sum ← sum + 1  
  return sum
```

N = 4	N = 8	N = 16

$$\sum_1^n c = n * c$$

$$\sum_1^n i = \frac{n * (n + 1)}{2}$$

$$\sum_1^n i^2 = \frac{2n^3 + 3n^2 + n}{6}$$

Ejemplo 6

```
procedure EXAMPLE2 (n) :  
  sum ← 0  
  for i ← 1 until n do  
    for j ← 1 until i do  
      sum ← sum + 1  
  return sum
```

$$\sum_1^n c = n * c$$

$$\sum_1^n i = \frac{n * (n + 1)}{2}$$

$$\sum_1^n i^2 = \frac{2n^3 + 3n^2 + n}{6}$$

Resumiendo

- ¿Qué sucede cuando tenemos ciclos anidados?
- ¿Qué sucede si el incremento no es ± 1 ?
 - Si es $(+ \text{ ó } -) m \Rightarrow$ ¿?
 - Si es $(* \text{ ó } /) m \Rightarrow$ ¿?

Ejercicios colaborativos

```
procedure EXERCISE1 (n) :  
  sum ← 0  
  for i ← 1 until n step +2 do  
    for j ← 1 until i do  
      sum ← sum + 1  
return sum
```

```
procedure EXERCISE2 (n) :  
  sum ← 0  
  for i ← 1 until n do  
    for j ← 1 until n step *2 do  
      sum ← sum + 1  
return sum
```

$$\sum_1^n c = n * c$$

$$\sum_1^n i = \frac{n * (n + 1)}{2}$$

$$\sum_1^n i^2 = \frac{2n^3 + 3n^2 + n}{6}$$

Algoritmo de burbuja

```
procedure BUBBLE_SORT(A) :  
for  $i \leftarrow A.length$  until 2 step -1 do  
    for  $j \leftarrow 1$  until  $i - 1$  do  
        if ( $A[j] > A[j + 1]$ ) then  
            SWAP(A,  $j$ ,  $j + 1$ )
```

$$\sum_1^n c = n * c$$

$$\sum_1^n i = \frac{n * (n + 1)}{2}$$

$$\sum_1^n i^2 = \frac{2n^3 + 3n^2 + n}{6}$$

Búsqueda binaria

```
procedure BINARY_SEARCH(A, key):  
  low ← 1  
  high ← A.length  
  while low < high do  
  begin  
    mid ← (high + low) / 2  
    if key == A[mid] then  
      return mid  
    else if key < A[mid] then  
      high = mid - 1  
    else  
      low = mid + 1  
  end  
  return low
```

$$\sum_1^n c = n * c$$

$$\sum_1^n i = \frac{n * (n + 1)}{2}$$

$$\sum_1^n i^2 = \frac{2n^3 + 3n^2 + n}{6}$$

Ejercicios colaborativos

```
procedure SELECTION_SORT(A) :  
  for i ← A.length until 2 step -1 do  
  begin  
    higher ← 1  
    for j ← 2 until i do  
      if A[higher] < A[j] then  
        higher ← j  
    if higher <> j then  
      SWAP(A, higher, i)  
  end
```

```
procedure INSERTION_SORT(A) :  
  for i ← 2 until A.length do  
    for j ← i until 2 and A[j] < A[j - 1] do  
      SWAP(A, j, j - 1)
```

$$\sum_1^n c = n * c$$

$$\sum_1^n i = \frac{n * (n + 1)}{2}$$

$$\sum_1^n i^2 = \frac{2n^3 + 3n^2 + n}{6}$$