Shared-memory Parallel Programming with Cilk Plus

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Outline for Today

- Data race detection with CilkSan
- Assessing Cilk performance with CilkScale

Concurrency Cautions

- Only limited guarantees between descendants or ancestors
 - —DAG precedence order maintained and nothing more
 - —don't assume atomicity between different procedures!

Race Conditions

Data race

- —two parallel strands access the same data
- —at least one access is a write
- —no locks held in common

General determinacy race

- —two parallel strands access the same data
- —at least one access is a write
- —a common lock protects both accesses

CilkSan

- Detects and reports <u>data races</u> when program terminates
 - —finds all data races even those by third-party or system libraries
- Does not report determinacy races
 - —e.g. two concurrent strands use a lock to access a queue
 - enqueue & dequeue operations could occur in different order potentially leads to different result

Race Detection Strategies in CilkSan

Lock covers

—two conflicting accesses to a variable don't race if some lock L is held while each of the accesses is performed by a strand

Access precedence

- —two conflicting accesses do not race if one must precede the other
 - access A is by a strand X, which precedes the cilk_spawn of strand Y which performs access B
 - access A is performed by strand X, which precedes a cilk_sync that is an ancestor of strand Y

CilkSan Limitations

- Only detects races between Cilk strands
 - —depends upon their strict fork/join paradigm
- Only detects races that occur given the input provided
 - —does not prove the absence of races for other inputs
 - —choose your testing inputs carefully!
- Runs serially, 15-30x slower
- Increases the memory footprint of an application
 - —could cause an error if memory demand is too large
- If you build your program with debug information (compile with -g), CilkSan will associate races with source line numbers

CilkSan Output

```
+ ./sum_race.cilk 1000
Running Cilksan race detector.
Race detected on location 7fff60735af0
     Write 5610cf4db29e Sum /opt/demo/race mutex reduce/sum race.cpp:31:11
         `-to variable sum (declared at /opt/demo/race mutex reduce/sum race.cpp:28)
     Spawn 5610cf4d96bf Sum /opt/demo/race mutex reduce/sum race.cpp:29:4
      Read 5610cf4db26c Sum /opt/demo/race mutex reduce/sum race.cpp:31:11
|*
         `-to variable sum (declared at /opt/demo/race mutex reduce/sum race.cpp:28)
    Spawn 5610cf4d96bf Sum /opt/demo/race mutex reduce/sum race.cpp:29:4
\| Common calling context
      Call 5610cf4da62a main /opt/demo/race mutex reduce/sum race.cpp:51:13
  Allocation context
   Stack object sum (declared at /opt/demo/race_mutex_reduce/sum_race.cpp:28)
     Alloc 5610cf4d9599 in Sum /opt/demo/race mutex reduce/sum race.cpp
      Call 5610cf4da62a main /opt/demo/race_mutex_reduce/sum_race.cpp:51:13
Race detected on location 7fff60735af0
     Write 5610cf4db29e Sum /opt/demo/race mutex reduce/sum race.cpp:31:11
         `-to variable sum (declared at /opt/demo/race mutex reduce/sum race.cpp:28)
     Spawn 5610cf4d96bf Sum /opt/demo/race_mutex_reduce/sum_race.cpp:29:4
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        `-to variable sum (declared at /opt/demo/race_mutex_reduce/sum_race.cpp:28)
     Spawn 5610cf4d96bf Sum /opt/demo/race_mutex_reduce/sum_race.cpp:29:4
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      Call 5610cf4da62a main /opt/demo/race_mutex_reduce/sum_race.cpp:51:13
```

499500

Cilksan detected 2 distinct races.
Cilksan suppressed 2995 duplicate race reports.

CilkSan Demo

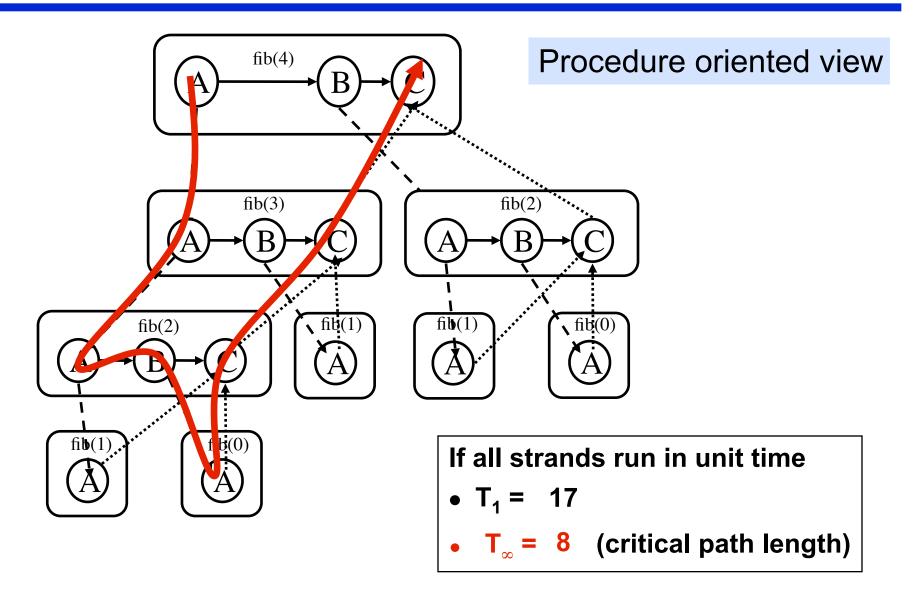
Explore CilkSan race detection

```
—cd ~/software/demo/race_mutex_reduce
—cilk_race.sh:
    - Uses -fopencilk -fsanitize=cilk
    - Run with CILK_NWORKERS=1
```

Performance Measures

- T_s = serial execution time
- T₁ = execution time on 1 processor (total work), T₁ ≥ T_s
- T_p = execution time on P processors
- T_m = execution time on infinite number of processors
 - longest path in DAG
 - length reflects the cost of computation at nodes along the path
 - known as "critical path length"

Work and Critical Path Example



CilkScale - Scalability Analysis Tool

- Collects stats of parallel performance
- Measures work, span, parallelism
- Automatically benchmarks program in different number of processors
- Produce tables and graphical plots with the above performance and scalability measurements
- Predicts speedup on various numbers of processors based on work and span

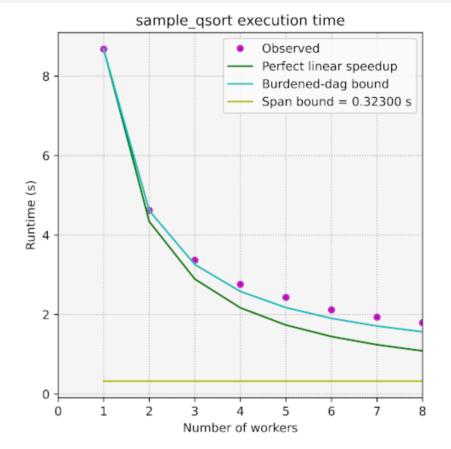
CilkScale - Scalability Analysis Tool

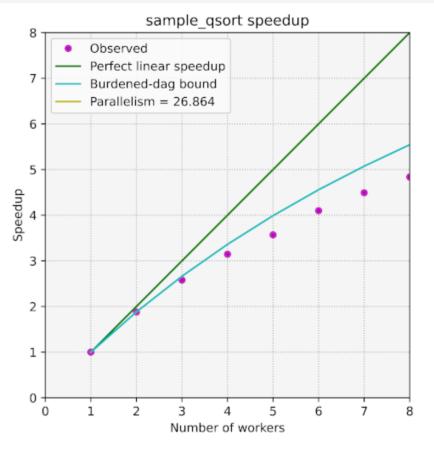
- Instruments Cilk program to collect performance measurements during its execution
- Cilkscale instrumentation operates in one of two modes
 - Benchmarking mode: Cilkscale measures the wall-clock execution time of your program
 - compile with -fopencilk -fcilktool=cilkscale-benchmark
 - Work/span analysis mode: Cilkscale measures the work, span, and parallelism of your program
 - compile with -fopencilk -fcilktool=cilkscale
- By default, Cilkscale only reports performance results for whole-program execution
 - fine-grained analyses of specific sub-computations is possible

Automatic benchmarks and visualization

ref: https://www.opencilk.org/doc/users-guide/cilkscale/

```
$ python3 /opt/opencilk/share/Cilkscale_vis/cilkscale.py \
    -c ./qsort_cs -b ./qsort_cs_bench \
    -ocsv cstable_qsort.csv -oplot csplots_qsort.pdf \
    -args 100000000
```





Rule of Thumb

As a rule of thumb, the parallelism of a computation is deemed sufficient if it is about 10x larger (or more) than the number of available processors. On the other hand, if the parallelism is too high — say, several orders of magnitude higher than the number of processors — then the overhead for spawning tasks that are too fine-grained may become substantial

Burdened Dag Bound:

Slowed execution if every continuation incurs a steal (incurs overhead of parallelism)

CilkScale Demo

Explore CilkScale for performance analysis using fib example

```
—cilkscale_run.sh fib 47
—cilkscale_run.sh fib_cutoff 47 2
—cilkscale_run.sh fib_cutoff 47 5
—cilkscale_run.sh fib_cutoff 47 10
—cilkscale_run.sh fib_cutoff 47 20
```

References - I

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References - II

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- Charles Leiserson, Bradley Kuzmaul, Michael Bender, and Hua-wen Jing. MIT 6.895 lecture notes - Theory of Parallel Systems. http://bit.ly/mit-6895-fall03
- Intel Cilk++ Programmer's Guide. Document # 322581-001US.