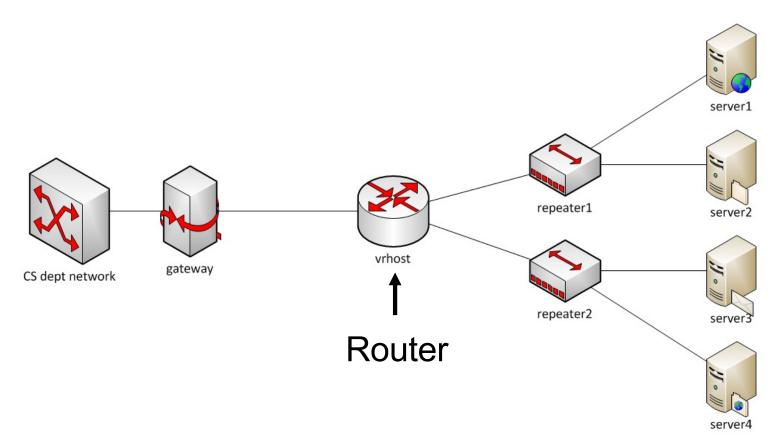
CSC 525: Computer Networks

build a simple router



- Implement basic packet forwarding functionality at the router
 - ARP, IP, and ICMP
- Will be tested with real traffic: ping and http download
- Ignore the repeaters or treat them as wires.

what you'll start with

- Virtual network topology
 - One per student. Only use your own topology number!
 - Only accessible from within the CS department.
- Stub code
 - It handles the backend setup, and send/receive pkts.
 - Its header files provide some basic data structures
 - Protocol header, interface list, etc.
 - Receive packets: sr_handlepacket()
 - Send packets: sr_send_packet()
- Your job is to implement the packet processing logic within sr_handlepacket().
 - Feel free to add new functions and new files, as long as "make" works.

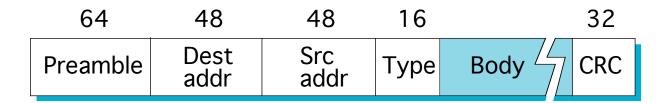
What does a packet look like?

Ethernet IP or ARP TCP or ICMP Payload

A byte pointer to the beginning of the packet

- How to parse the packet header?
 - Cast to a data structure of the protocol header, use pointer to access specific header fields.
 - E.g., (struct sr_ethernet_hdr *) pkt; where pkt is (uint8_t *) that points to the beginning of a packet.
 - Header structures are in the header files of the stub code
- How to know which protocol header to use?
 - Each protocol header contains a field that identifies the next protocol header.

Ethernet



- Ignore preamble and CRC; they're handled by hardware, not passed to router software.
- Addresses
 - 48-bit unicast address assigned to each interface
 - example: 8:0:2b:e4:b1:2
 - Broadcast address: all 1s, i.e., ff:ff:ff:ff:ff
- Type: what next protocol header is (e.g., IP or ARP)
- Body: min 46 bytes, max 1500 bytes

Let's look at a real packet

Start the router with packet logging

```
lectura 171 \stackrel{-}{>} ./sr -t 363 \stackrel{-}{-}1 log
```

ping one of your servers, then stop both ping and router.

tcpdump to display the packets in hexadecimal

Accessing header fields

```
struct sr ethernet hdr
#ifndef ETHER ADDR LEN
#define ETHER ADDR LEN 6
#endif
    uint8 t ether dhost[ETHER ADDR LEN];
    uint8_t ether shost[ETHER ADDR LEN];
    uint16 t ether type;
} attribute ((packed));
uint8 t *p;
struct sr ethernet hdr *eth;
struct ip *iphdr
void sr handlepacket(sr, p, len, iface)
  eth = (struct sr ethernet hdr *) p;
  eth->ether type is IP;
  iphdr = (struct ip*)(p + sizeof(struct sr ethernet hdr);
```

Address Resolution Protocol

- ARP: Given an IP address, finds the corresponding Ethernet address
 - broadcast the request, i.e., who knows the ethernet address of IP 1.2.3.4?
 - the target machine responds with its MAC address.
 - Cache the reply in a table to avoid asking the same question later.
 - Cache times out after some time.

ARP Packet Format

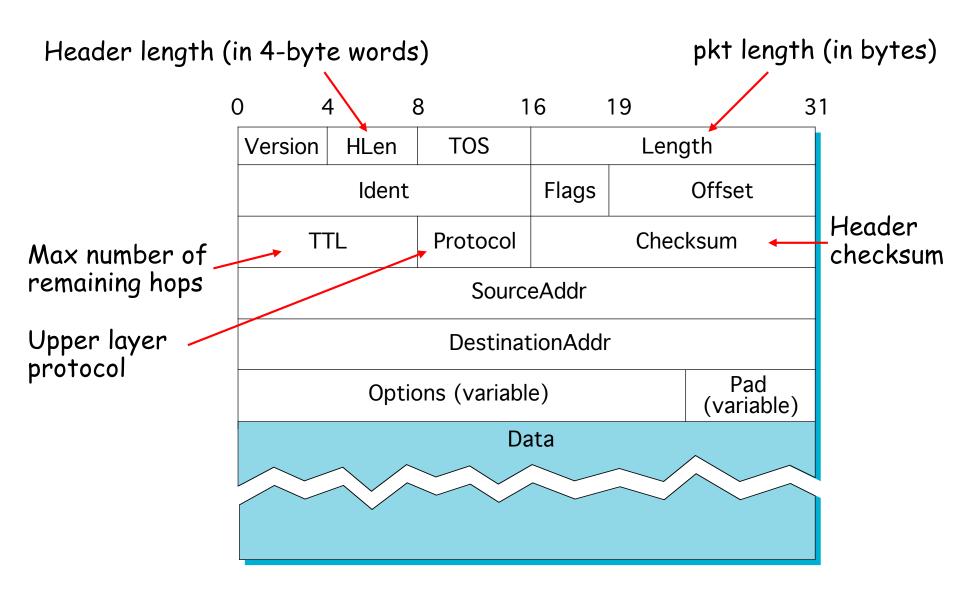
00	3 1	6 3		
Hardware type=1		ProtocolType=0x0800		
HLen= 6	PLen= 4	Operation		
SourceHardwareAddr (bytes 0–3)				
SourceHardwareAddr (bytes 4–5)		SourceProtocolAddr (bytes 0–1)		
SourceProtocolAddr (bytes 2–3)		TargetHardwareAddr (bytes 0–1)		
TargetHardwareAddr (bytes 2–5)				
TargetProtocolAddr (bytes 0–3)				

- In our context, Hardware refers to Ethernet, Protocol refers to IP.
- Hardware Type: type of physical network (e.g., Ethernet)
- ProtocolType: type of higher layer protocol (e.g., IP)
- HLEN & PLEN: length of physical and protocol addresses in bytes
- Operation: request or response
- Source/Target Physical/Protocol addresses

ARP Details

- The request is broadcasted onto the local link only.
 - Routers don't forward ARP packets.
- The reply is usually unicast to the requester, but can also be broadcasted.
- A host can learn IP/MAC mapping from incoming packets without generating separate ARP requests.
- An entry in ARP table (cache) will time out after a while.
- If you answer the request with a correct reply, you will receive an IP packet next.
 - Otherwise you'll keep getting the same ARP request.

IP Packet Header



The basic header (without options) is 20 bytes long

Basic IP packet processing

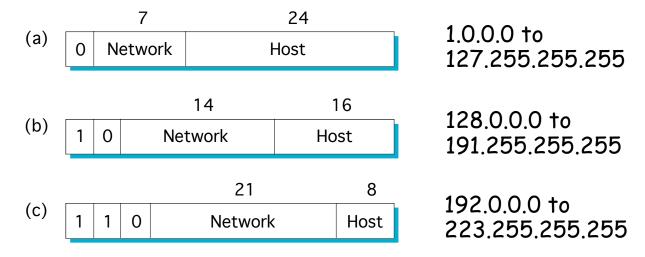
- Is the destination myself?
 - Yes. Keep processing if it's an ICMP ping, otherwise drop it.
 - No. forward the packet as follows.
- Decrement TTL by 1,
 - if the result TTL == 0, discard the packet
 - Otherwise update the checksum
 - Set checksum field to be 0, recalculate the checksum over the header, and record the result in the field.
- Use the destination address to look up the routing table to find a matching entry, which tells the next-hop for this packet.
- Send the packet to the next hop.

Network and Prefix

- In order to scale, IP routing tables record one entry per *destination network*.
 - all nodes in same network share the same IP address prefix.
- A new node gets an IP address from the network operator to ensure the address has the network's prefix.
- When a node moves from one network to another, it has to change its IP address.

How many bits in the prefix?

Original IP design: class-based address



The good: given an IP, we know its prefix.

The bad: class-B addresses are in great demand. Lots of unused addresses in larger network (i.e., shorter prefixes)

Classless Addressing

- Today we use the so-called *classless addressing*.
 - The network part or the prefix can have any number of bits, which must be explicitly specified by either a network mask or prefix length.
 - e.g., the old class B address has a prefix length /16, or equivalently, a network mask of 255.255.0.0
 - Address & mask => prefix
 - 1.2.3.4/16 or 1.2.3.4 with mask 255.255.0.0

Take a look at the interface

- Mask indicates how long the network part is
 - addr & mask = network
- Broadcast address has all ones in the host part (non prefix).

Routing Table Lookup

```
dest nexthop (gw) mask interface 172.24.74.64 0.0.0.0 255.255.255.248 eth1 172.24.74.80 0.0.0.0 255.255.255.248 eth2 0.0.0.0 172.24.74.17 0.0.0.0 eth0
```

- Lookup match if (packet's dest & mask) == (table's dest & mask)
- Longest prefix match: if there're multiple matches in the table, pick the one that has the longest prefix.
- 0.0.0.0 as the dest means that this is the default route. It matches every address, but has the shortest prefix. Used only if no other match.
- Nexthop and interface is for where to forward the packet.
- 0.0.0.0 as the nexthop means that the nexthop IP is the same as the destination IP of the packet.

How to send to next-hop?

- After looking up the routing table, we know the nexthop's IP address and the interface.
- Now need to send the packet to the nexthop via the interface.
- But how to target the packet to the nexthop?
 - Cannot modify dest address in IP header since it has to remain the same for final destination.
- Use next-hop's MAC address as the destination in Ethernet header!
 - IP destination remain the same the entire journey, but Ethernet destination address changes at every hop.
- But, how to get next-hop's MAC address?
 - We've already covered it.

Put it together

- IP packet processing
- IP packet forwarding
 - Routing table lookup to find out nexthop
 - ARP to find out MAC address of nexthop
 - Maintain ARP cache
 - Send packet to nexthop using its MAC address.

Internet Control Message Protocol

- ICMP: used by hosts & routers for network-level error reporting and diagnosis
 - unreachable host, network, port, protocol, echo request/reply, and many more.
- ICMP msgs are carried in IP packets
- ICMP message format

IP header				
type	code	checksum		
unused (or used by certain ICMP types)				
IP header and first 64bits of data				
Or				
data (according to ICMP types)				

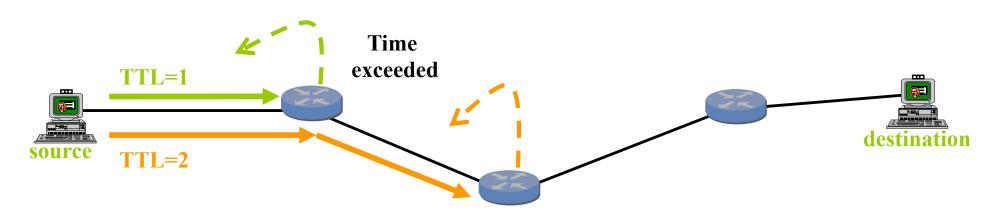
<u>Type</u>	<u>Code</u>	<u>description</u>
0	0	echo reply (ping)
3	0	dest. network unreachable
3	1	dest host unreachable
3	2	dest protocol unreachable
3	3	dest port unreachable
3	6	dest network unknown
3	7	dest host unknown
4	0	source quench (congestion
		control - not used)
8	0	echo request (ping)
9	0	route advertisement
10	0	router discovery
11	0	TTL expired
12	0	bad IP header

ping

- If received an ICMP echo request whose IP destination is the router itself, reply with an ICMP echo reply.
- Don't need to do anything to echo requests destined to other hosts, just normal IP packet forwarding.

Traceroute

- Not required for this project
- Measuring the forwarding path
- Time-To-Live field in IP packet header
 - Source sends a packet with a TTL of n
 - Each router along the path decrements the TTL
 - "TTL exceeded" sent when TTL reaches 0
- Traceroute exploits this behavior by sending a sequence of packets with increasing initial TTL.



Send packets with TTL=1, 2, 3, ... and record source of "TTL exceeded" mesg

Traceroute Example

```
bobcat2 117 > traceroute www.google.com
 1
    * * *
    tcsn-dsl-qw04-196.tcsn.qwest.net (168.103.240.196) 52.629 ms
 3
    * tcsn-agw1.inet.gwest.net (168.103.240.93) 52.172 ms
    tus-core-01.inet.qwest.net (205.171.149.33) 51.762 ms
 4
   svl-core-01.inet.gwest.net (67.14.12.6) 74.450 ms
   pax-edge-01.inet.gwest.net (205.171.214.30) 76.415 ms
   * * *
   209.85.130.6 (209.85.130.6) 76.325 ms
   66.249.94.226 (66.249.94.226) 77.372 ms
10
   216.239.49.66 (216.239.49.66) 76.611 ms
11
   mc-in-f99.google.com (66.102.7.99) 76.144 ms
```

The Implementation

ARP implementation

- Be able to respond to ARP request, send ARP request when needed, and maintain ARP cache.
- IP forwarding: given an incoming packet
 - Decrement TTL and update checksum.
 - Lookup the routing table to find the nexthop's IP.
 - Invoke ARP to find out the nexthop's MAC address.
 - Send the packet to the nexthop.

ICMP

- Send echo reply in response to echo request.

Four steps

- Test the stub code
- Implement ARP
- Implement IP forwarding
- Implement ICMP ping

• Test the code along the way; don't test a big chunk of code all at once.