Homework 2: Crypto Basics

Problem 1:-

To determine the size of the key, I performed the Kasiski approach. Then I utilize all potential key combinations to deduce the plaintext.

```
substitution_cipher = 'OBRGXIMYAZZAWCATBNMUYYHAZNVGFCXPVVSIJSVLKIFAVGBIECAZSBWGRGRQWUCHMMOCYE\
FLGQQNKFSHQMGYALNKCIJQVEKVWXNFOYFYQBESGOYTXMAYTXSISNBPMSGOJBKFWRUTTMLS\
WKHTKNGNUXRNJUVGMYNYEYNLYYGPGYFSBNQQWUCHMMSLRWTFDYQRNOJUEWNLVUZIDHWXLH\
{\tt TLKNEXMALBRQGUMMGXCUFXLHTLLLRTROJYFIDHLIGADLUBVXENSCALVCDFFIQCFAHISILU \setminus {\tt TLKNEXMALBRQGUMMGXCUFXLHTLLRTROJYFIDHLIGADLUBVXENSCALVCDFFIQCFAHISILU \setminus {\tt TLKNEXMALBRQGUMMGXCUFXLHTLLRTROJYFIDHLIGADLUBVXENSCALVCDFFIQCFAHISILU \setminus {\tt TLKNEXMALBRQGUMMGXCUFXLHTLLRTROJYFIDHLIGADLUBVXENSCALVCDFFIQCFAHISILU \setminus {\tt TLKNEXMALBRQGUMMGXCUFXLHTLRTROJYFIDHLIGADLUBVXENSCALVCDFFIQCFAHISILU \setminus {\tt TLKNEXMALBRQGUMMGXCUFXLHTLRTROJYFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLIGADLUBVXENSCALVCDFIDHLUBVXENSCALVCDFIDHLUBVXENSCALVCDFIDHLUBVXENSCALVCD
XXZXNUAMZAWISRNOJYKMQYECGRSBWHAHLUFBBPDPWLJBRYOCYEAYSVYXSISPRKSNZYPHMM\
WKHXMWWMGAZNEOFMDHKORMGOKNUHTAZQRAZPWBRTQXGZFMTJAXUTRNWCAPZLUFRODLFYFL\
GUKHRODLTYRGRYWHNLRIUCNMDXOCGAKIFAQXKUQMVGZFDBVLSIJSGADLWCFGNCFMGTMWWI\
STBIMHGKXBSPVGFVWHRYHNWXSKNGHLBENHYYQPZLXUEXNHDSBGDQZIXGNQKNUXCCKUFMQI\
MMRYEYUNFHEUDIAZVUJWNGQYSFVSDNZYFNOLWGRBLJGLGTMWWISKZJAXVMXCFVEBMAAHTB\
SNGUPENMWCGBRIFFLHMYOBBBRNZIEHTAZFLTBKMUVGSYVOVMGNZYROHFKISPZLOBBVZHLB\
BKNOYBYRTHVYELSUFXGADJJISBSUTFRPZSGZPTQLQCAZHNGHGADMCCYEEODARGDLSFQHDM\
FIGKZCKYNLDWGHOEDPOHRBSBWLNKDBAMENOJDSJTETEMYHZXWXZHOYLBNGSOAWRHMWWONK\
HMVYPEZLWXUXVCDFAHSOSMGXOLWWVHTMLCZXHHOUVMHHYZBKOYAHSHOWWGRGSMFIEPHFDB\
RMTLFBVLZLESOTBEXIEYQYKBFNOJDCRLAOLWEHRMWMGADYFYZRRZJIAMHYJQVMGIMNQXKU\
ONUXUUDORHENAGRMGUL CEUDCEANEHNI ERTGYSXBYXTMI BTOTEYAMGUKWBNMNWXSHOGGI RM\
GUFYVMGYJHHFDLAWNEROHYEBNLANLHQNZYABBYKNPTKWMFNMHIFMJBSBJYTTQXLIPHLGAM\
FTQCSN'
total char = 26
# this method will give us population variance of a cipher text.
def population_variance(cipher):
       frequency = [0]*total_char
       for i in range(0,len(cipher)):
             frequency[ord(cipher[i].capitalize())-ord('A')] += 1
       sum_ = 0
       mean = 1/total_char
       for i in range(0,total_char):
              sum_ += (frequency[i]/len(cipher) - mean)**2
       return sum /total char
# this method will create chunks of cipher and calculate population variance for every chunk of cipher
def kasiski(cipher, no_of_chunks):
       chunks=[]
       for j in range(0 , no_of_chunks):
              chunks.append("")
       for j in range(0, len(cipher)):
             chunks[j%no_of_chunks] += cipher[j]
       sum_=0
       for c in chunks:
             sum_ += population_variance(c)
       return sum /no of chunks
# This method create plain text for given ciper text and key
def decrypt(cipher, key):
       plain text = '
       for i in range(0, len(cipher)):
              \verb|plain_text| = \verb|plain_text| + \verb|chr((ord(cipher[i]) - ord(key[i\%len(key)]))\%26 + ord('A'))|
       return plain_text
# this method find the letters with maximum frequencies.
{\tt def\ max\_frequency\_letter(chipher,\ select\_char):}
       frequency=[0]*total_char
       for i in range(0,len(chipher)):
              frequency[ord(chipher[i].capitalize())-ord('A')]+=1
       frequency = sorted(range(len(frequency)), \ key = lambda \ x: \ frequency[x], \ reverse = True)
       for i in range(0, select_char):
              letters.append(chr((frequency[i]-4)%26+ord('A')))
# we call kasishi method repeatedly and check for highest population variance to get the key length
max_variance = 0
for i in range(1,10):
       temp_variance = kasiski(substitution_cipher, i)
       if(max_variance < temp_variance):</pre>
          max_variance = temp_variance
```

```
key_length = i
# print(max_variance)
# print(key_length)
# This loop will create key_length different caesar ciphers
caesar_ciphers = [""]*key_length
for i in range(0, len(substitution_cipher)):
    caesar_ciphers[i % key_length] += substitution_cipher[i]
# for i in range(0, len(caesar_ciphers)):
   print(caesar_ciphers[i])
# return all the possible keys made up with letters with maximum frequencies
def possible_keys(select_char):
    letters = []
    for i in range(0, len(caesar_ciphers)):
        letters.append(max_frequency_letter(caesar_ciphers[i], select_char))
    \texttt{return} \ [ \texttt{(c0+c1+c2+c3+c4+c5)} \ \text{for c0 in letters[0] for c1 in letters[1] for c2 in letters[2] }
            for c3 in letters[3] for c4 in letters[4] for c5 in letters[5]]
#Below code trying all the possible generated keys and check for highest population variance in plain text
quessed keys = possible keys(3)
predicted_key = ''
predicted key variance = 0
predicted_plain_text =
for key in quessed_keys:
    temp_plain_text = decrypt(substitution_cipher, key)
    temp_variance = population_variance(temp_plain_text)
    if predicted_key_variance < temp_variance:</pre>
      predicted_key = key
      predicted_key_variance = temp_variance
      predicted_plain_text = temp_plain_text
print(predicted\_key + "=" + str(predicted\_key\_variance) + "\n" + predicted\_plain\_text)
```

Output:-

SUNTZU=0.0010215425227561555

WHENYOUENGAGEINACTUALFIGHTINGIFVICTORYISLONGINCOMINGTHEMENSWEAPONSWILLGROWDULLANDTHEIRARDORWILLBEDAMPENEDIFYOULAYSIEGETOATOWNYOUWILLEXHAUST
THESTATEWILLNOTBEEQUALTOTHESTRAINNEVERFORGETWHENYOURWEAPONSAREDULLEDYOURARDORDAMPENEDYOURSTRENGTHEXHAUSTEDANDYOURTREASURESPENTOTHERCHIEFTAI
HOWEVERWISEWILLBEABLETOAVERTTHECONSEQUENCESTHATMUSTENSUETHUSTHOUGHWEHAVEHEARDOFSTUPIDHASTEINWARCLEVERNESSHASNEVERBEENSEENASSOCIATEDWITHLONG
FROMPROLONGEDWARFAREONLYONEWHOKNOWSTHEDISASTROUSEFFECTSOFALONGWARCANREALIZETHESUPREMEIMPORTANCEOFRAPIDITYINBRINGINGITTOACLOSEITISONLYONEWHO
STANDTHEPROFITABLEWAYOFCARRYINGITONTHESKILLFULGENERALDOESNOTRAISEASECONDLEVYNEITHERAREHISSUPPLYWAGONSLOADEDMORETHANTWICEONCEWARISDECLAREDHEI
TURNHISARMYBACKFORFRESHSUPPLIESBUTCROSSESTHEENEMYSFRONTIERWITHOUTDELAYTHEVALUEOFTIMETHATISBEINGALITTLEAHEADOFYOUROPPONENTHASCOUNTEDFORMORET
TOCOMMISSARIAT

Problem 2:-

```
frequency_in_english ={ "A": .08167, "B": .01492, "C": .02782, "D": .04253, "E": .12702, "F": .02228,
                                                                                                                                                          "G": .02015, "H": .06094, "I": .06996, "J": .00153, "K": .00772, "L": .04025,
                                                                                                                                                          "M": .02406, "N": .06749, "O": .07507, "P": .01929, "Q": .00095, "R": .05987, "S": .06327, "T": .09056, "U": .02758, "V": .00978, "W": .02360, "X": .00150,
                                                                                                                                                        "Y": .01974, "Z": .00074 }
plain text='ethicslawanduniversitypolicieswarningtodefendasystemyouneedtobeabletot\
hinklikeanattackerandthatincludesunderstandingtechniquesthatcanbeusedt\
 ocompromisesecurityhoweverusingthosetechniquesintherealworldmayviolate\
 thelawortheuniversitysrulesanditmaybeunethicalundersomecircumstancesev\
 enprobingforweaknessesmayresultinseverepenaltiesuptoandincludingexpuls\
ioncivilfinesandjailtimeourpolicyineecsisthatyoumustrespecttheprivacya\
nd proper tyrights of other satall times or else you will fail the course acting lawf \land \\
ully and ethically is your responsibility carefully read the computer fraud and ab \verb|\|
use act c faa a federal statute that broadly criminalizes computer in trusion this i \verb|\| like the computer of the computer 
 so ne of several laws that govern hacking understand what the law prohibits if indou \verb|\| \\
btwe can refer you to an attorney please review its spolicies on responsible use of the contraction of the
 echnology resources and caen spolicy documents for guide {\tt lines} concerning proper{\tt 'lines} and {\tt caen} and {\tt caen}
plain_text = plain_text.upper()
 total char = 26
```

```
# Below code is simple implementation of population variance using given english letter frequencies
sum_{-} = 0
for char in frequency_in_english:
 sum_ += frequency_in_english[char]
mean = sum_/total_char
variance = 0
for char in frequency_in_english:
 xi = frequency_in_english[char]
 variance += (xi-mean)**2
part_a = variance/total_char
print("a. Population variance of the relative letter frequencies in English text = "+str(part_a))
\# this method will give us population variance of a given text.
def population_variance(cipher):
    frequency = [0]*total_char
    for i in range(0, len(cipher)):
       frequency[ord(cipher[i])-ord('A')] += 1
    sum_{-} = 0
    mean = 1/total_char
    for i in range(0,total_char):
       sum_ += (frequency[i]/len(cipher) - mean)**2
    return sum /total char
print("\nb. Population variance of the relative letter frequencies in given Plain text = "+str(population_variance(plain_text)))
# below code generate cipher text with given keys and calculate population variance for each cipher text
\label{lem:print('\nc. Population variance of the cipher text : ')}
for key in ("YZ", "XYZ", "WXYZ", "VWXYZ", "UVWXYZ"):
    cipher_text = ""
    for i in range(0,len(plain_text)):
        \label{eq:cipher_text} cipher\_text + chr((ord(plain\_text[i]) + ord(key[i\%len(key)]) - 2*ord('A')) \% 26 + ord('A'))
    print("(using "+ key + ") = " +str(population_variance(cipher_text)))
# this will calculate mean population variance of cipher text using all the keys
print('\nd. Mean population variance of the cipher text :')
for key in ("YZ", "XYZ", "WXYZ", "VWXYZ", "UVWXYZ"):
    cipher_text = ""
    # make cipher text
    for i in range(0,len(plain_text)):
        cipher_text = cipher_text + chr((ord(plain_text[i])+ord(key[i%len(key)])-2*ord('A'))%26+ord('A'))
    key_length = len(key)
    ciphers = []
    # break the cipher text into key length chunks
    for _ in range(0, key_length):
     ciphers.append("")
    for c in range(0,len(cipher_text)):
     ciphers[c%key_length] += cipher_text[c]
    # calculate variance for every chunk and mean the variances
    variance_sum = 0
    for i in range(0, len(ciphers)):
    variance_sum += population_variance(ciphers[i])
print("(using "+ key + ") = " +str(variance_sum/len(ciphers)))
# Below code calculate mean population variance of the cipher text using UVWXYZ with different key length(2-5)
print("\ne. Mean population variance of the cipher text using UVWXYZ with different key length: ")
cipher_text = ""
key = "UVWXYZ"
for i in range(0,len(plain_text)):
    \label{eq:cipher_text} {\tt cipher_text + chr((ord(plain_text[i]) + ord(key[i\%len(key)]) - 2*ord('A'))\%26 + ord('A'))} \\
for key_length in range(2, 6):
  ciphers = []
  # break the cipher text into key length chunks
  for \_ in range(0, key_length):
    ciphers.append("")
  for c in range(0, len(cipher_text)):
   ciphers[c%key_length] += cipher_text[c]
  variance_sum = 0
   # calculate variance for every chunk and mean the variances
  for i in range(0, len(ciphers)):
    variance_sum += population_variance(ciphers[i])
  print("(using key length " + str(key_length) + " ) = " + str(variance_sum/key_length))
```

Output:-

```
a. Population variance of the relative letter frequencies in English text = 0.0010405667735207099
b. Population variance of the relative letter frequencies in given Plain text = 0.001006973795435334
c. Population variance of the ciphertext :
(using YZ) = 0.0005425572595902266
(using XYZ) = 0.0003474151001623527
(using WXYZ) = 0.0002459193736666264
(using VWXYZ) = 0.00018148975566557982
(using UVWXYZ) = 0.00016797152786163775
Explanation:- More non-uniform shifts in the plaintext are introduced with a longer key. This lowers variance by increasing the randomness
So, the key with a small length has high variance but low randomness.
d. Mean population variance of the cipher text :
(using YZ) = 0.0011098867554911514
(using XYZ) = 0.0011034546954876623
(using WXYZ) = 0.0012505199318386133
(using VWXYZ) = 0.0011889247809577482
(using UVWXYZ) = 0.0013157126756027854
Explanation: When I compare the variance to the other three parts, I find that part a, which was the population variance of relative lette
This is due to the fact that the variance was computed across the set of characters in the key that were all moved by the same character.
e. Mean population variance of the ciphertext using UVWXYZ with different key lengths:
(using key length 2 ) = 0.00046995129412711825
(using key length 3 ) = 0.0005082166002495672
(using key length 4 ) = 0.0005558574514618472
(using key length 5 ) = 0.0003314984771028727
Explanation: This is a very effective Kasiski assault variation. We must break ciphertexts into blocks of a specific length and search for
population variance. We are confident that if we predict the key length incorrectly, the average of population variations will be significa
but if we guess the key length correctly, the average of population variances will almost certainly increase to 0.001 level.
```

Explanation:-

- **c.** More non-uniform shifts in the plaintext are introduced with a longer key. This lowers variance by increasing the randomness of relative letter frequency. So, the key with a small length has high variance but low randomness.
- **d.** When I compare the variance to the other three parts, I find that part a, which was the population variance of relative letter frequencies in English text, is comparable. This is due to the fact that the variance was computed across the set of characters in the key that were all moved by the same character. As a result, each set performs as a Caesar cipher.
- e. This is a very effective Kasiski assault variation. We must break ciphertexts into blocks of a specific length and search for commonly occurring patterns in the Kasiski attack using population variance. We are confident that if we predict the key length incorrectly, the average of population variations will be significantly lower than the level of regular English texts, but if we guess the key length correctly, the average of population variances will almost certainly increase to 0.001 level.

Problem 3:-

No, all the 56 bits of the DES key is used an unequal number of times in K_i . Like zeroth location, the bit is missing only once, in key 3 but other keys are missing more than once.

Here are the missing bits of the K_i representing the key of round i

I'm assuming a bit position starting from 0 in 56 bits (0-55 locations).

Round Key Missing bits position from initial 56 bits in the K_i

 K_1 9 18 22 25 35 38 43 54 K_2 10 19 23 26 36 39 44 55

K_3	0 12 21 25 29 38 41 46
K_4	2 14 23 27 31 40 43 48
K_5	1 4 16 25 33 42 45 50
K_6	3 6 18 27 35 44 47 52
K_7	1 5 8 20 37 46 49 54
K_8	3 7 10 22 28 39 48 51
K_9	4 8 11 23 29 40 49 52
K10	6 10 13 25 31 42 51 54
K11	8 12 15 27 28 33 44 53
K12	1 10 14 17 30 35 46 55
K13	3 12 16 19 29 32 37 48
K14	5 14 18 21 31 34 39 50
K15	7 16 20 23 33 36 41 52
K16	8 17 21 24 34 37 42 53