

Sorting I

- Sorting Overview
- Bubble Sort
- Selection Sort
- Insertion Sort

HW 3 → Friday } 11:59 PM
P3 → Monday }

Why do we care so much about sorting?

- Sorting is sometimes a key ingredient in other algorithms.
- Sorting is a simple problem, so it is a good problem for understanding and appreciating differences in the algorithms.



What do we care about with sorting methods?

- **Certification:** verification of correctness
- **Runtime:**
 - number of comparisons or array accesses
 - why not exchanges?
- **Extra memory:**
 - in situ (in place)
 - extra copy of array
- **Types of data:** anything that can be ordered
- **Comparable/Comparator:** Java interfaces that make it easier to write more flexible sorting algorithms

} call stack
(recursive algorithms)

Visualization

<https://www.cs.usfca.edu/~galles/visualization/ComparisonSort.html>

<https://www.youtube.com/watch?v=kPRA0W1kECg>

Bubble Sort

<u>5</u>	<u>4</u>	3	0	1	2
4	<u>5</u>	<u>3</u>			
4	3	<u>5</u>	<u>0</u>		
4	3	0	<u>5</u>	<u>1</u>	
4	3	0	1	<u>5</u>	<u>2</u>
<u>4</u>	<u>3</u>	0	1	2	5
3	<u>4</u>	<u>0</u>			
3	0	<u>4</u>	<u>1</u>		
3	0	1	<u>4</u>	<u>2</u>	<u>5</u>
3	0	1	2	4	5

Bubble Sort

Input: An array A of size N

Output: Sorted A

best: $O(N^2)$

worst: $O(N^2)$

```
for j from 1 to N-1
  i = 0
  while i+1 < N:
    if A[i] > A[i+1]
      swap A[i] and A[i+1]
    i++
  end while
end for
```

Selection Sort

Repeatedly **select** the smallest item and put it into the next unsorted spot.

		a[]										
i	min	0	1	2	3	4	5	6	7	8	9	10
		S	O	R	T	E	X	A	M	P	L	E
0	6	S	O	R	T	E	X	A	M	P	L	E
1	4	A	O	R	T	E	X	S	M	P	L	E
2	10	A	E	R	T	O	X	S	M	P	L	E
3	9	A	E	E	T	O	X	S	M	P	L	R
4	7	A	E	E	L	O	X	S	M	P	T	R
5	7	A	E	E	L	M	X	S	O	P	T	R
6	8	A	E	E	L	M	O	S	X	P	T	R
7	10	A	E	E	L	M	O	P	X	S	T	R
8	8	A	E	E	L	M	O	P	R	S	T	X
9	9	A	E	E	L	M	O	P	R	S	T	X
10	10	A	E	E	L	M	O	P	R	S	T	X
		A	E	E	L	M	O	P	R	S	T	X

entries in black are examined to find the minimum

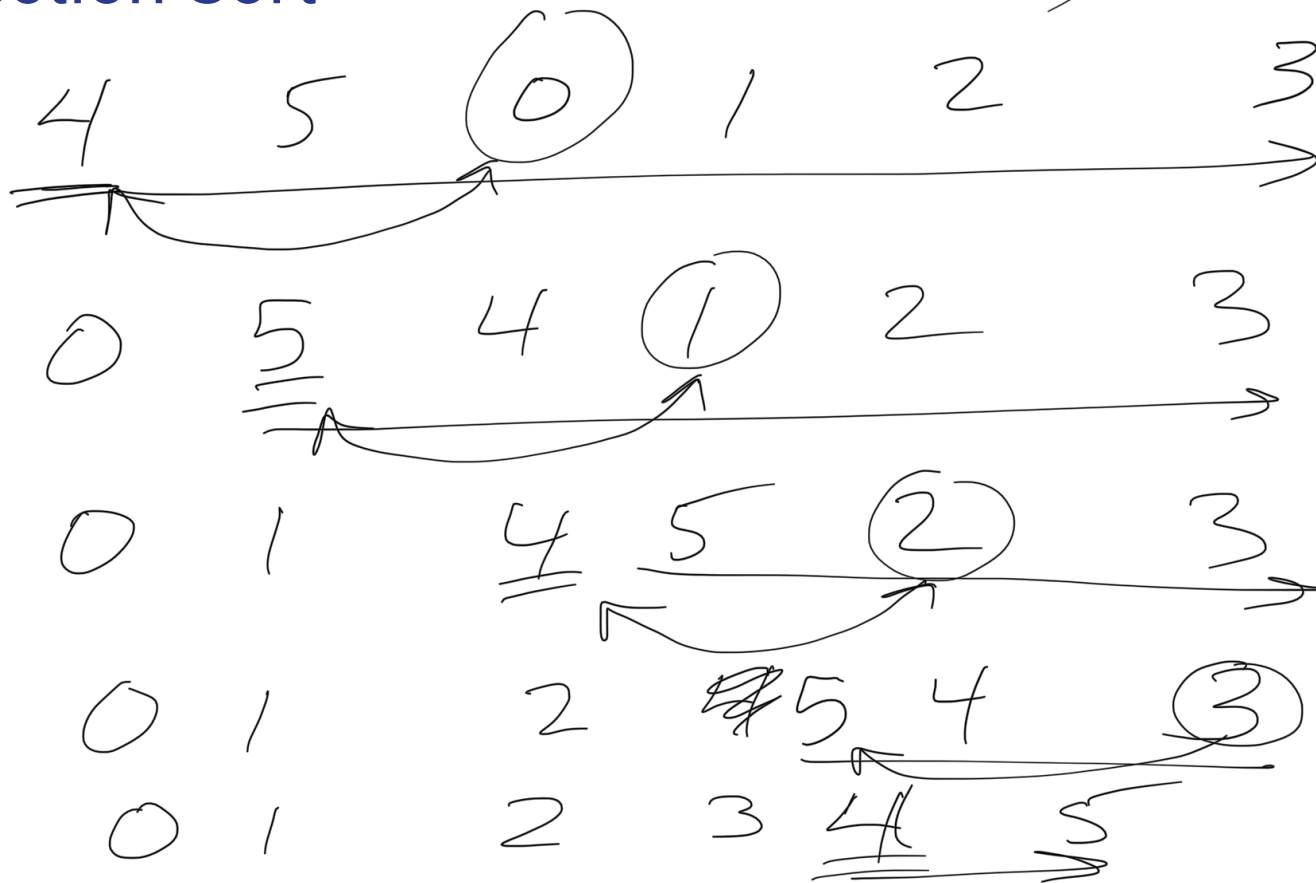
entries in red are a[min]

entries in gray are in final position

Trace of selection sort (array contents just after each exchange)

Selection Sort

$O(N^2)$



$$\begin{aligned} &N \\ &+ \\ &N-1 \\ &+ \\ &\vdots \\ &+ \\ &1 \end{aligned}$$

Selection Sort

Input: An array A of size N

Output: A , sorted

for i from 0 to $N-1$:

\rightarrow min = find the min in i to $N-1$

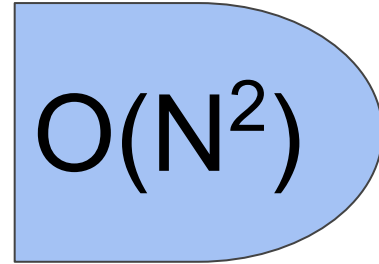
 swap the min with $A[i]$

end for

$$N + N-1 + N-2 + \dots + 1 \Rightarrow O(N^2)$$

Properties of Selection Sort

- Running time is insensitive to input
 - Input array is in random order
 - Input array is already sorted
 - Input array is sorted in opposite direction
 - Input array is all the same number
- Data movement is minimal— $O(N)$ exchanges
- Anything left of the current index is sorted and in final position


$$O(N^2)$$

Insertion Sort

Insert current item into its proper place in the sorted part of the array (on the left).

		a[]										
i	j	0	1	2	3	4	5	6	7	8	9	10
		S	O	R	T	E	X	A	M	P	L	E
1	0	O	S	R	T	E	X	A	M	P	L	E
2	1	O	R	S	T	E	X	A	M	P	L	E
3	3	O	R	S	T	E	X	A	M	P	L	E
4	0	E	O	R	S	T	X	A	M	P	L	E
5	5	E	O	R	S	T	X	A	M	P	L	E
6	0	A	E	O	R	S	T	X	M	P	L	E
7	2	A	E	M	O	R	S	T	X	P	L	E
8	4	A	E	M	O	P	R	S	T	X	L	E
9	2	A	E	L	M	O	P	R	S	T	X	E
10	2	A	E	E	L	M	O	P	R	S	T	X

entries in gray do not move

entry in red is a[j]

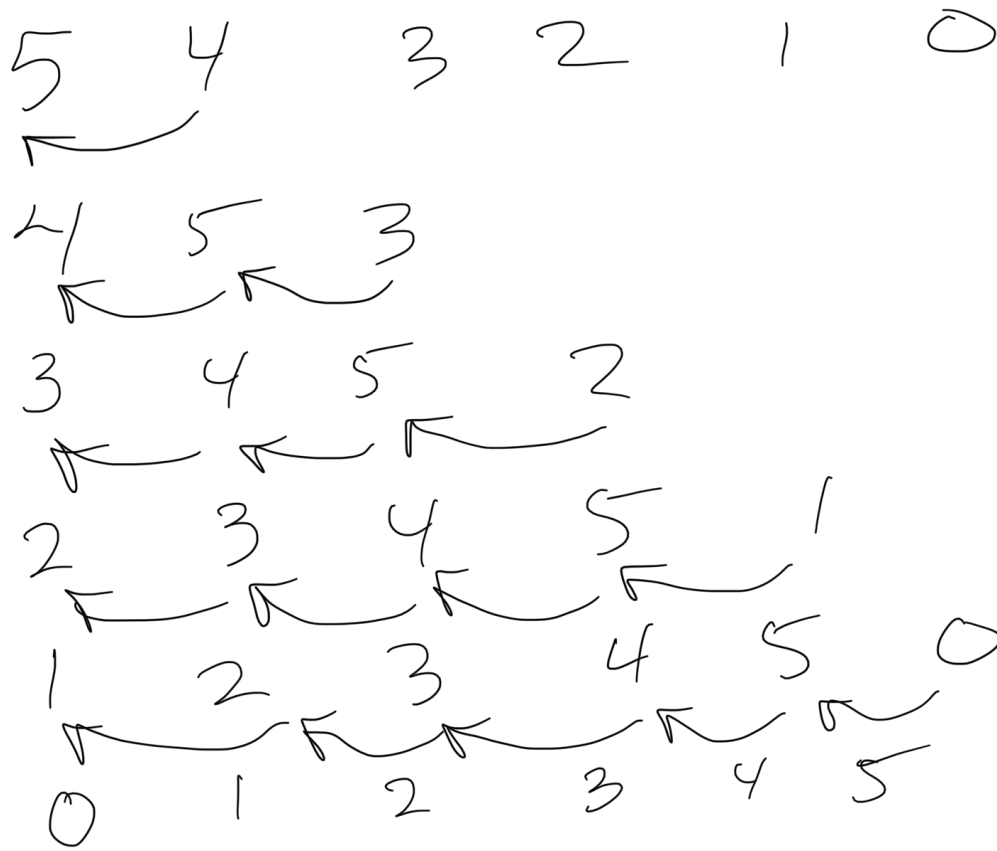
entries in black moved one position right for insertion

Trace of insertion sort (array contents just after each insertion)

Insertion Sort

best: $O(N)$

$O(N^2)$



$$\left. \begin{array}{c} 1 \\ + \\ 2 \\ + \\ 3 \\ + \\ \vdots \\ + \\ N \end{array} \right\}$$

Insertion Sort

Input: Array A of size N

Output: A , sorted

for i from 1 to $N-1$

$j = i$

while $j-1 \geq 0$:

if $A[j] < A[j-1]$:

swap $A[j]$ and $A[j-1]$

$j--$

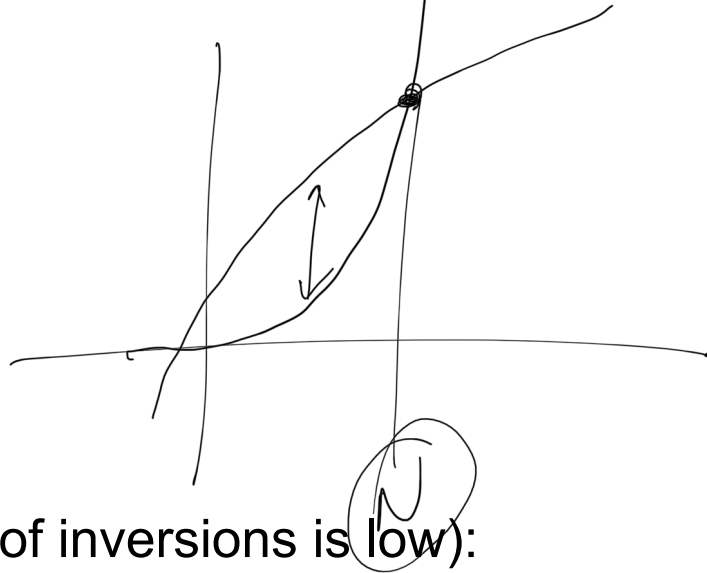
end if

else break;

end while

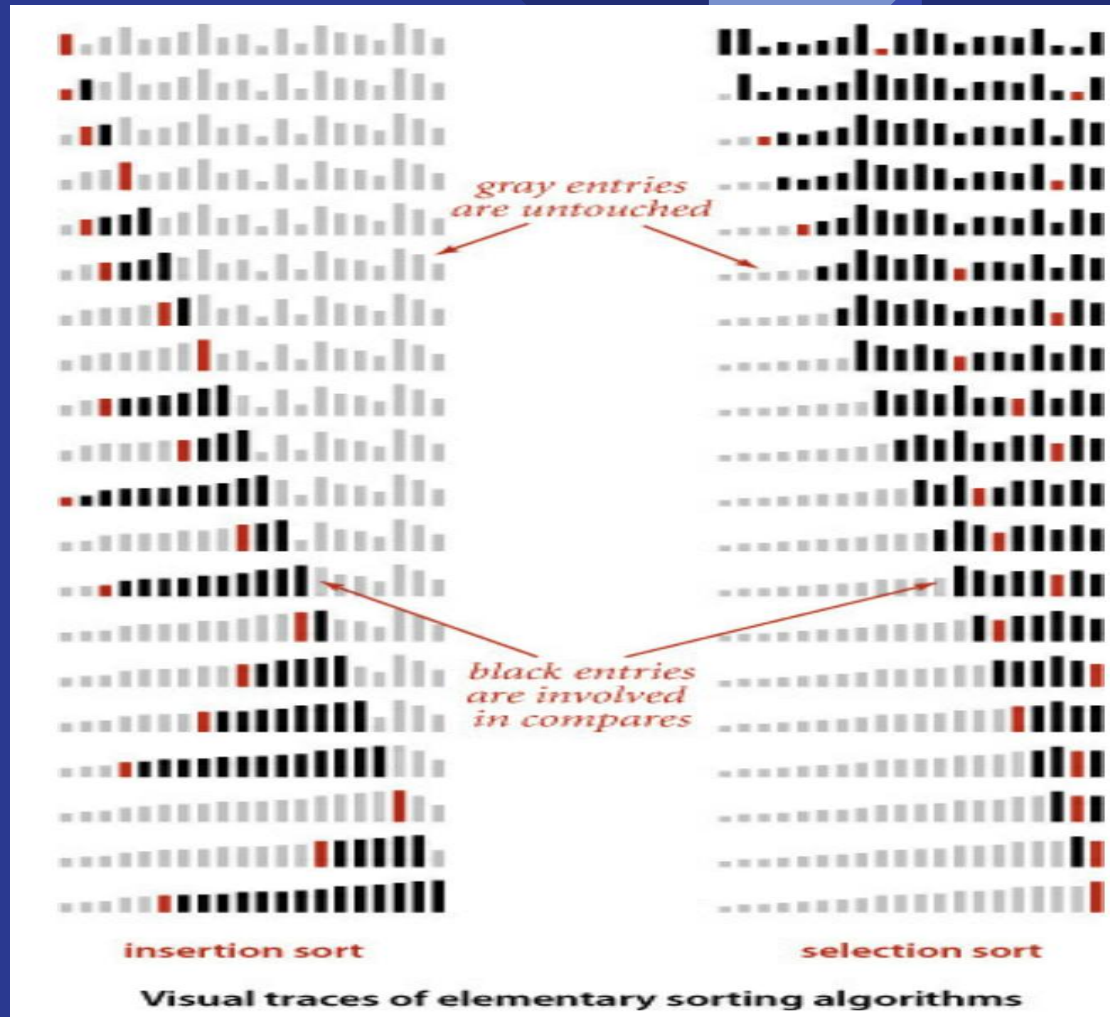
Properties of Insertion Sort

- Running time is sensitive to input
- Works well for:
 - Tiny arrays
 - Partially sorted arrays (number of inversions is low):
 - each entry is not far from its final position
 - small array appended to a large sorted array
 - array only has a few elements out of place
- Anything left of the current index is sorted but not necessarily in final position



Visual Representation of Sorting Methods:

- Shows sorting of bars by length
- Note:
 - For insertion sort, nothing right of the current index is touched
 - For selection sort, nothing left of the current index is touched.
 - Insertion sort tends to require fewer compares.



Analysis

Algorithm	Best	Average	Worst	Extra Space	Comments
Bubble Sort	$O(N)$	$O(N^2)$	$O(N^2)$	$O(1)$	<ul style="list-style-type: none">• similar to Insertion Sort with worse average case• easy to implement
Selection Sort	$O(N^2)$	$O(N^2)$	$O(N^2)$	$O(1)$	<ul style="list-style-type: none">• minimal data movement ($O(N)$ in the worst case)• does not take advantage of the input
Insertion Sort	$O(N)$	$O(N^2)$	$O(N^2)$	$O(1)$	<ul style="list-style-type: none">• takes advantage of input• works well on small arrays• works well on partially sorted arrays

References

[1] *Algorithms, Fourth Edition*; Robert Sedgewick and Kevin Wayne (and associated slides)