

# LING/C SC/PSYC 438/538

Lecture 24

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Adminstration

538 Presentations

- Last Homework will be out Wednesday!

# 538 Presentations

Name	Presentation	Date
Alkuraydis,Ahmed		
Barner,Jacob Ryan		
Bejarano,Cielo S		
Bell,Jack T		
Converse,Amber Charlotte	8.4 HMM Part-of-Speech Tagging	7
Cox,Samantha Ann	18 Word Senses and WordNet, 18.1-18.3	7
Davis,Katherine Nicole		
Dharmala,Bayu	12.3 Some Grammar Rules for English	7
Hopper,Ashlyn Danielle		
Jain,Varshit Chirag	23 Question Answering, 23.4-23.6	5
Kankia,Kevin Pinakin		
Kleczewski,Alison	15.4 Event and State Representations, 15.4, 15.4.1, 15.4.2	7
LaScola Ochoa,Logan Michelle		
Logan,Haley Brooke	21.1 Coreference Phenomena: Linguistic Background	5
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Mangkang,Tinnawit		
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Mendoza,Freddy		
Murphy III,Michael LaMotte		
Pinto,Aayush Bernard	20 Lexicons for Sentiment, Affect, and Connotation, 20.1-20.3	5
Pipatanangkura,Leighanna D	23.2 IR-based Factoid Question Answering	5
Raju,Anish	Lexical and Vector Semantics, 6.1-6.2	7
Reeve,Keegan Austin	Coreference Resolution: Mention and Architectures 21.3-21.4	7
Ruparel,Deep Anil	6 : Vector Semantics and Embeddings, 6.3-6.5	5
Shakyam Shreya Nupur	24.2 Chatbots	7
Shu,Qiyu		
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Thompson,Brendan S		
Warrick,Baylor M		
West,Georgia Alexandra		
Willittes,Taylor		
Yuan,Alan		

# Last Time

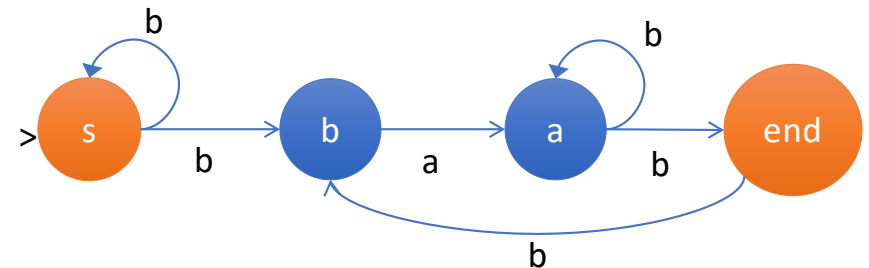
- SWI-Prolog DCG: grammar rules
  - $n \rightarrow RHS$  ( $RHS$  = sequence (,) of terminal and non-terminal symbols)
  - $n$  (non-terminal symbol)
  - $[t]$  (terminal symbol)
- Chomsky Hierarchy type 3 (equiv. FSA and regex):
  - if RHS limited to either  $[t]$ ,  $n$  or  $[t]$
- **Key Concepts so far:**
  - recursive rules
    - examples: factorial,  $\Sigma^*$  and grammar rules, e.g. multiple defs. for nonterminal  $n$
  - backtracking: *explore multiple possible paths of execution*
    - *Prolog remembers (FOR YOU) which grammar rules have been tried*

# Prolog Derivations

- Prolog's **computation rule**:
  - Try first matching rule in the database
  - Backtrack if matching rule leads to failure
  - undo and try next matching rule
- For grammars, this means:
  - Top-down left-to-right derivations
  - **left-to-right** = expand leftmost nonterminal first
  - Leftmost expansion done recursively = **depth-first**

bab2.prolog

```
1. s --> [].  
2. s --> [b], b.  
3. s --> [b], s.  
4. b --> [a], a.  
5. a --> [b].  
6. a --> [b], b.  
7. a --> [b], a.
```



# Enumeration

Using bab2.prolog

4. the set of all strings from the alphabet  $a, b$  such that each  $a$  is immediately preceded by and immediately followed by a  $b$ ;

- What happens with the query `s(List, []). ?`

```
?- [bab2].
```

```
true.
```

```
?- s(List, []).
```

```
List = [] ;
```

← ; another answer?

```
List = [b, a, b] ;
```

```
List = [b, a, b, a, b] ;
```

```
List = [b, a, b, a, b, a, b] ;
```

```
List = [b, a, b, a, b, a, b, a, b] ;
```

```
List = [b, a, b, a, b, a, b, a, b, a, b|...] [write]
```

```
List = [b, a, b, a, b, a, b, a, b, a, b] ;
```

```
List = [b, a, b, a, b, a, b, a, b, a, b, a, b] ;
```

This is not all strings in the language!  
 $\lambda | b(ab)^+$   
 $b^+$  is in the language!  
`s([b, b], []).`  
**true.**

← press w

turns off abbreviating

# Enumeration

## Why?

- **Answer:** Prolog tries first matching rule
- Order of application:

① ② ③ ④

1, 2, 4, 5,

[]

⑤ ⑥ ⑦

6, 4, 5

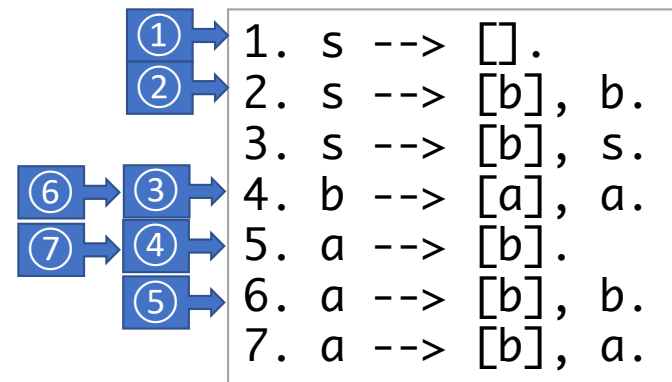
[b,a,b]

⑧ ⑨ ⑩

6, 4, 5

[b,a,b,a,b]

[b,a,b,a,b,a,b]



Rules 3 and 7 never fire!

# Enumeration

- We can swap the order of rules!

```
?- [bab2swap].
```

```
true.
```

```
?- s(List, []).
```

```
List = [] ;
```

```
List = [b, a, b] ;
```

```
List = [b, a, b, b] ;
```

```
List = [b, a, b, b, b] ;
```

```
List = [b, a, b, b, b, b]
```

Now we have  
 $\lambda | bab^+$   
cf.  $\lambda | b(ab)^+$   
before the swap

```
1. s --> [].  
2. s --> [b], b.  
3. s --> [b], s.  
4. b --> [a], a.  
5. a --> [b].  
6. a --> [b], b.  
7. a --> [b], a.
```



```
1. s --> [].  
2. s --> [b], b.  
3. s --> [b], s.  
4. b --> [a], a.  
5. a --> [b].  
6. a --> [b], a.  
7. a --> [b], b.
```



# Enumeration

What happens if we swap rules 2 and 3?

```
?- [bab2swap2].
```

```
true.
```

```
?- s(List, []).
```

```
List = [] ;
```

```
List = [b] ;
```

```
List = [b, b] ;
```

```
List = [b, b, b] ;
```

```
List = [b, b, b, b] ;
```

```
List = [b, b, b, b, b]
```

Now we have

$b^*$

cf.  $\lambda|bab^+$  and  $\lambda|b(ab)^+$

But

```
s([b,a,b,b,a,b], []).
```

```
true
```

```
1. s --> [].  
2. s --> [b], b.  
3. s --> [b], s.  
4. b --> [a], a.  
5. a --> [b].  
6. a --> [b], a.  
7. a --> [b], b.
```



```
1. s --> [].  
2. s --> [b], s.  
3. s --> [b], b.  
4. b --> [a], a.  
5. a --> [b].  
6. a --> [b], a.  
7. a --> [b], b.
```

# Breadth-first search

- Depth-first search (DFS) is **unfair** – not all rules get their turn
- Breadth-first search (BFS) is **fair**
  - computation paths are explored in order of length
  - we can simulate BFS using DFS (*at some cost to efficiency*)
  - called **Iterative Deepening** (ID) (*explanation an advanced topic*)
  - `id_meta.prolog` (*meta-level Prolog program*)

# BFS: enumeration is possible

- `?- [id_meta].`
- **true.**
- `?- [bab2].`
- **true.**
- `?- id(s(List, [])).`

1. List = [] ;	9. List = [b, a, b, b, b] ;	23. List = [b, a, b, b, a, b, b] ;
2. List = [b] ;	10. List = [b, b, a, b, b] ;	24. List = [b, a, b, b, b, a, b] ;
3. List = [b, a, b] ;	11. List = [b, b, b, a, b] ;	25. List = [b, a, b, b, b, b, b] ;
4. List = [b, b] ;	12. List = [b, b, b, b] ;	26. List = [b, b, a, b, a, b, b] ;
5. List = [b, a, b, b] ;	13. List = [b, a, b, a, b, b] ;	27. List = [b, b, a, b, b, a, b] ;
6. List = [b, b, a, b] ;	14. List = [b, a, b, b, a, b] ;	28. List = [b, b, a, b, b, b, b] ;
7. List = [b, b, b] ;	15. List = [b, a, b, b, b, b] ;	29. List = [b, b, b, a, b, a, b] ;
8. List = [b, a, b, a, b] ;	16. List = [b, b, a, b, a, b] ;	30. List = [b, b, b, a, b, b, b] ;
	17. List = [b, b, a, b, b, b] ;	31. List = [b, b, b, b, a, b, b] ;
	18. List = [b, b, b, a, b, b] ;	32. List = [b, b, b, b, b, a, b] ;
	19. List = [b, b, b, b, a, b] ;	33. List = [b, b, b, b, b, b]
	20. List = [b, b, b, b, b] ;	
	21. List = [b, a, b, a, b, a, b]	
	22. List = [b, a, b, a, b, b, b] ;	

# BFS: enumeration is possible

- How many a's in a string of length  $n$ ?
  - Suppose we have 3 a's in a string
  - Every  $a$  needs a  $b$  on either side, so the minimum length of the string has to be 7.
  - 1 a, minimum length 3; 2 a's, minimum length is 5.
- **General Formula:**
  - $\text{len}_{\min} = \#a * 2 + 1$ , for  $\#a \geq 1$  because
  - $\text{len}_{\min} = \#a * 3 - (\#a - 1)$ , for  $\#a \geq 1$
  - Notice that  $\text{len}_{\min}$  must always be odd
  - When length is even: must be all b's or any odd string (of length one less) in the language + one additional b.

# BFS: enumeration is possible

1. List = [] ;	16. List = [b, b, a, b, a, b] ;	31. List = [b, b, b, b, a, b, b] ;
2. List = [b] ;	17. List = [b, b, a, b, b, b] ;	32. List = [b, b, b, b, b, a, b] ;
3. List = [b, a, b] ;	18. List = [b, b, b, a, b, b] ;	33. List = [b, b, b, b, b, b] ;
4. List = [b, b] ;	19. List = [b, b, b, b, a, b] ;	34. List = [b, a, b, a, b, a, b, b]
5. List = [b, a, b, b] ;	20. List = [b, b, b, b, b] ;	
6. List = [b, b, a, b] ;	21. List = [b, a, b, a, b, a, b] ;	
7. List = [b, b, b] ;	22. List = [b, a, b, a, b, b, b] ;	
8. List = [b, a, b, a, b] ;	23. List = [b, a, b, b, a, b, b] ;	
9. List = [b, a, b, b, b] ;	24. List = [b, a, b, b, b, a, b] ;	
10. List = [b, b, a, b, b] ;	25. List = [b, a, b, b, b, b, b] ;	
11. List = [b, b, b, a, b] ;	26. List = [b, b, a, b, a, b, b] ;	
12. List = [b, b, b, b] ;	27. List = [b, b, a, b, b, a, b] ;	
13. List = [b, a, b, a, b, b] ;	28. List = [b, b, a, b, b, b, b] ;	
14. List = [b, a, b, b, a, b] ;	29. List = [b, b, b, a, b, a, b] ;	
15. List = [b, a, b, b, b, b] ;	30. List = [b, b, b, a, b, b, b] ;	

Annotations:

- 1 a: points to List 3
- 2 a's: points to List 8
- 3 a's: points to List 21

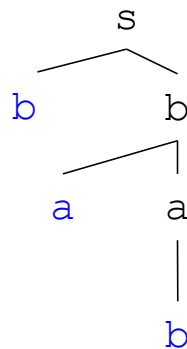
# Extra Argument: Parse Tree

- **Recovering a parse tree**

- in case of **true**, we can compute a syntax tree representation of the derivation
- simple transformation: adding an extra argument to **all** nonterminals
- applies to all grammar rules (not just regular grammars)

**Example:** bab2.prolog **again**

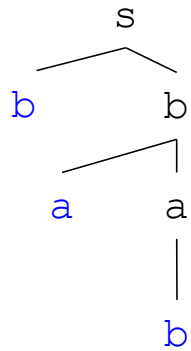
- string: **bab**



1.	s	-->	[].
2.	s	-->	[b], b.
3.	s	-->	[b], s.
4.	b	-->	[a], a.
5.	a	-->	[b].
6.	a	-->	[b], b.
7.	a	-->	[b], a.

# Extra Argument: Parse Tree

- **Tree:**



`s(b, b(a, a(b)))`

- **Prolog term data structure:**

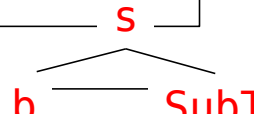
- hierarchical
- allows sequencing of arguments
- `functor(arg1, .., argn)`
- each `argi` could be another term or simple atom

## Extra Argument: Parse Tree

- Transform each rule:

1.  $s \rightarrow []$ .
2.  $s \rightarrow [b], b$ .
3.  $s \rightarrow [b], s$ .
4.  $b \rightarrow [a], a$ .
5.  $a \rightarrow [b]$ .
6.  $a \rightarrow [b], b$ .
7.  $a \rightarrow [b], a$ .

1.  $s(Tree) \rightarrow []$ .
2.  $s(Tree) \rightarrow [b], b(SubTree)$ .
- ...
5.  $a(Tree) \rightarrow [b]$ .

- 
1.  $s(s([], b)) \rightarrow []$ .
  2.  $s(s(b, SubT)) \rightarrow [b], b(SubT)$ .
  - ...
  5.  $a(a(b)) \rightarrow [b]$ .



# Extra Arguments: Parse Tree

- **Prolog grammar**

`bab2.prolog`

1. `s --> [].`
2. `s --> [b], b.`
3. `s --> [b], s.`
4. `b --> [a], a.`
5. `a --> [b].`
6. `a --> [b], b.`
7. `a --> [b], a.`

- **Prolog grammar computing a parse**

`bab2tree.prolog`

## Extra Arguments: Parse Tree

?- [bab2tree].

**true.**

?- s(Tree, [b,a,b], []).

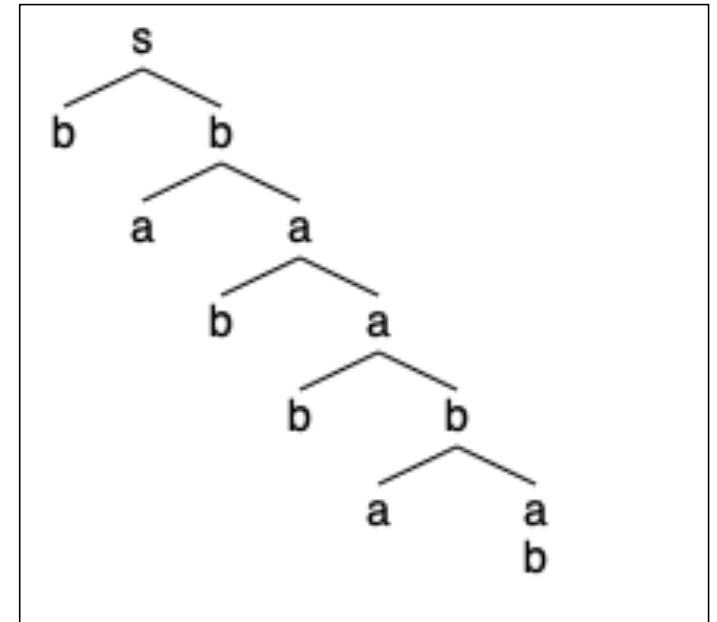
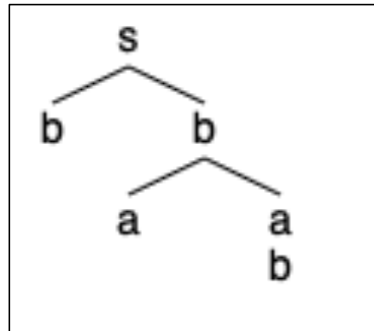
Tree = s(b, b(a, a(b))) ;

**false.**

?- s(Tree, [b,a,b,b,a,b], []).

Tree = s(b, b(a, a(b, a(b, b(a, a(b)))))) ;

**false.**



The extra argument *Tree* precedes the string lists in the call to nonterminal S

# Extra Arguments: Parse Tree

- Useful in natural language grammars
- ?- [psg2].

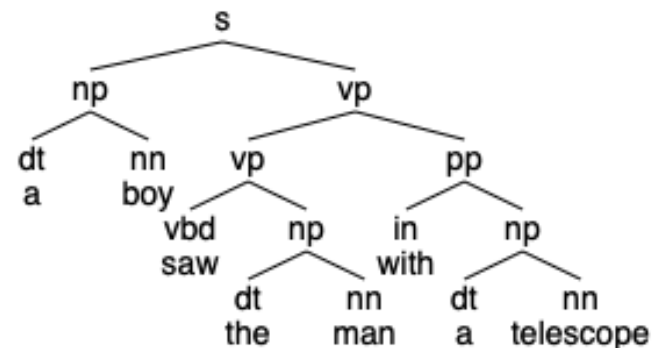
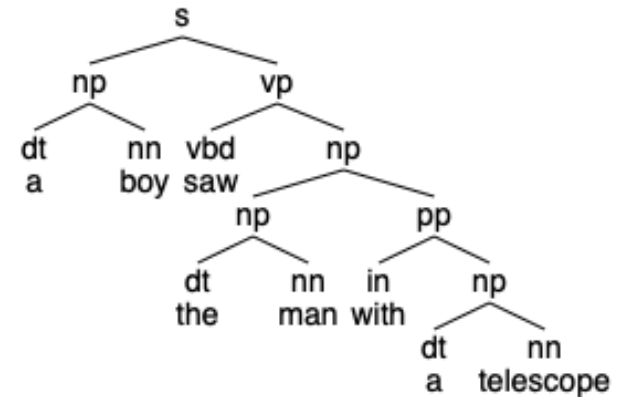
**true.**

?- s(Tree, [a, boy, saw, the, man, with, a, telescope], []).

```
Tree = s(np(dt(a), nn(boy)),  
vp(vbd(saw), np(np(dt(the), nn(man)),  
pp(in(with), np(dt(a),  
nn(telescope)))))) ;
```

```
Tree = s(np(dt(a), nn(boy)),  
vp(vp(vbd(saw), np(dt(the), nn(man))),  
pp(in(with), np(dt(a), nn(telescope)))))  
;
```

**false.**



# Extra Arguments

- Extra arguments are powerful
  - they allow us to impose (grammatical) constraints and change the expressive power of the system
    - if used as read-able memory (cf. *Turing Machine discussion*)
- **Example:**
  - $a^n b^n c^n$   $n > 0$  is not a context-free language (type-2)
  - *i.e. you cannot write rules of the form  $n \rightarrow \text{RHS}$  to generate this language*
  - in fact, it's context-sensitive (type-1)

# Beyond Regular Languages

- Language
  - $a^n b^n = \{ab, aabb, aaabbb, aaaabbbb, \dots\} \ n \geq 1$
- A regular grammar extended to allow both left and right recursive rules can accept/generate it:
- `anbn.prolog`
  1. `a --> [a], b.`
  2. `b --> [b].`
  3. `b --> a, [b].`

Set  
enumeration

```
?- a(L, []).  
L = [a, b] ;  
L = [a, a, b, b] ;  
L = [a, a, a, b, b, b] ;  
L = [a, a, a, a, b, b, b, b] ;  
L = [a, a, a, a, a, b, b, b, b, b]
```

- Example:

```
?- a([b,b,a,a], []).  
false.  
  
?- a([a,b], []).  
true ;  
false.  
  
?- a([a,a,b], []).  
false.  
  
?- a([a,a,b,b], []).  
true ;  
false.  
  
?- a([a,a,b,b,b], []).  
false.  
  
?- a([a,b,a,b], []).  
false.
```

Set  
membership

# Beyond Regular Languages

- Language

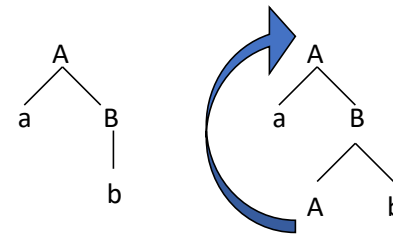
- $a^n b^n = \{ab, aabb, aaabbb, \dots\} \ n \geq 1$

- A regular grammar extended to allow both left and right recursive rules can accept/generate it:

1.  $a \rightarrow [a], b.$
2.  $b \rightarrow [b].$
3.  $b \rightarrow a, [b].$

- Intuition:

- grammar implements the stacking of partial trees balanced for a's and b's:



# Beyond Regular Languages

- Language
  - $a^n b^n = \{ab, aabb, aaabbb, aaaabbbb, \dots\} \ n \geq 1$
- A regular grammar extended to allow both left and right recursive rules can accept/generate it:
  1.  $a \rightarrow [a], b.$
  2.  $b \rightarrow [b].$
  3.  $b \rightarrow a, [b].$
- A type-2 or context-free grammar (CFG) has no restrictions on what can go on the RHS of a grammar rule
- **Note:**
  - *CFGs still have a single nonterminal limit for the LHS of a rule*
- **Example:**
  1.  $s \rightarrow [a], [b].$
  2.  $s \rightarrow [a], s, [b].$