

# Searching VI (Skiplists)

# Motivation

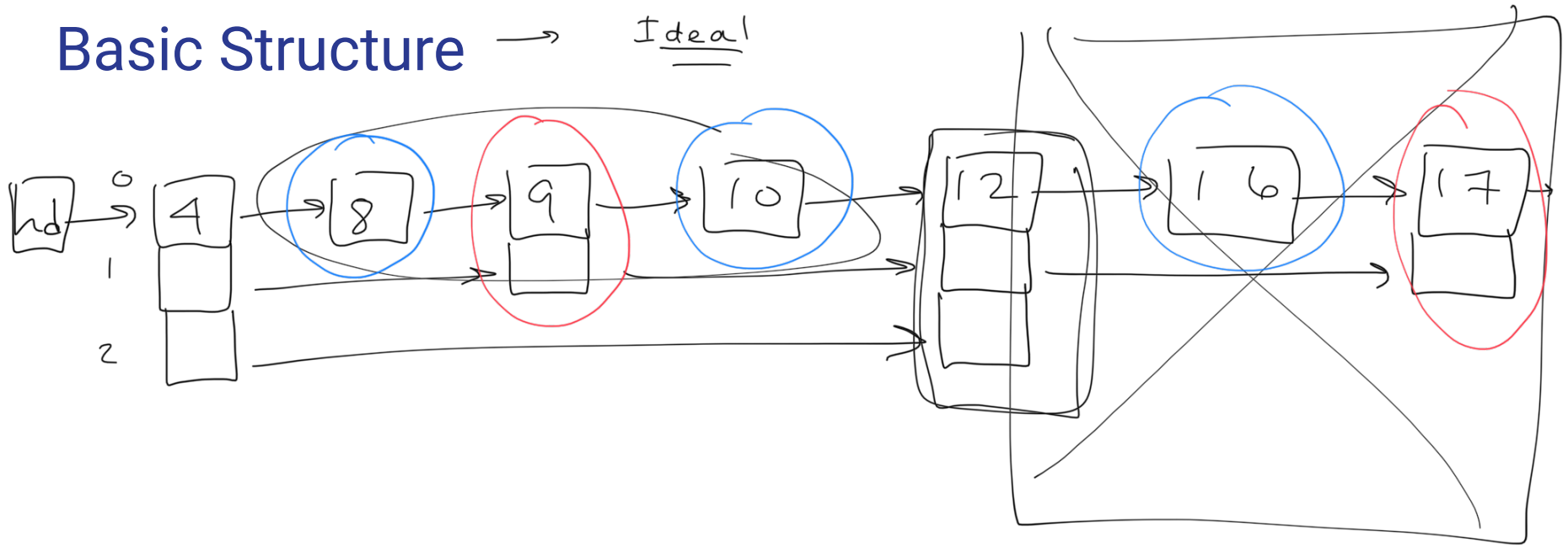
BST pros : ordered, time is good for random data,  
simple

BST cons : worst-case is  $O(N)$

build randomization into structure



# Basic Structure → Ideal



# Operations

insert  $\rightarrow$  at what level?

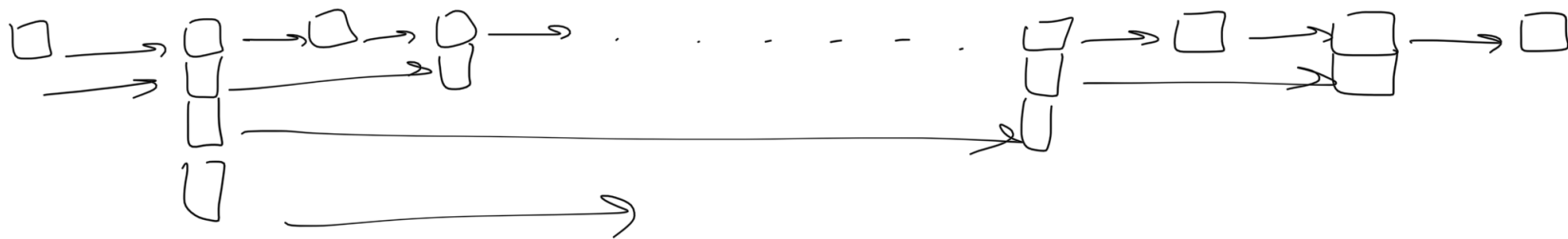
search  $\rightarrow$  start at bottom level

delete



# The Skiplist...

- ...does not *guarantee*  $O(\log N)$  runtime for operations,...
- ...does provide  $O(\log N)$  performance *with a high probability*.
- ...can provide a balance between good performance and ease of implementation.



N 10 2 4

Ordered [Hashtables]

- BST
- 2-3
- RB
- Skip lists

# Ideally...

- You would have enough levels so that the ideal search would skip half the values, then a quarter of the remaining values, then an eighth of the remaining values, etc...
- Which would mean that a *search* (or an *insert*) would be  $O(\log N)$  with high probability.
- Space: If the head node isn't counted,
  - $\frac{1}{2}$  the nodes should have 1 pointer...
  - $\frac{1}{4}$  should have 2 pointers...
  - $\frac{1}{8}$  should have 3 pointers...
  - and so on...
- However, we don't enforce the *ideal* situation.
- Instead, we rely on probability and randomization.

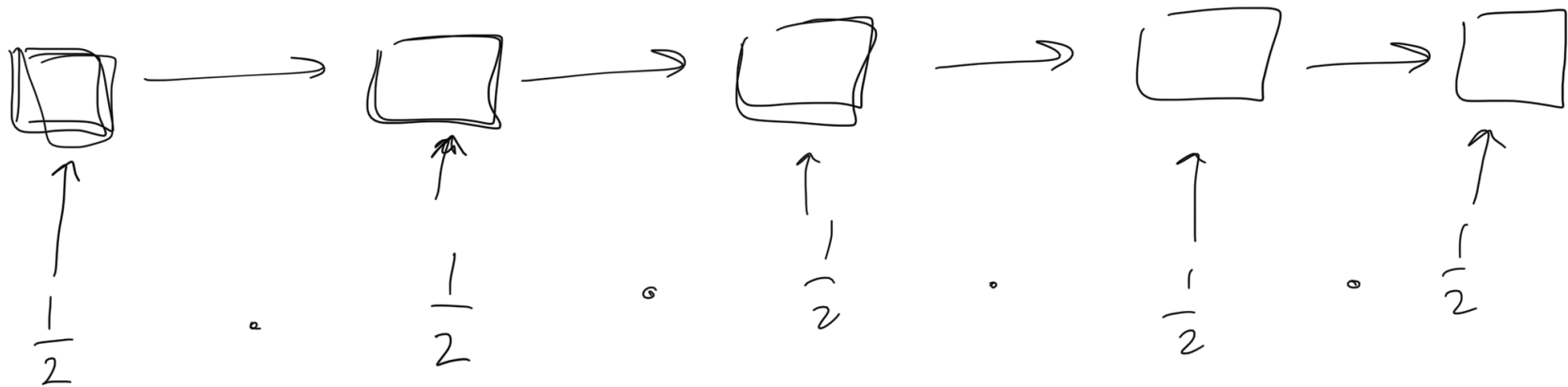
# Inserting & removing nodes...

- Instead, a new node is assigned a level at random where there is
  - a 50% chance that it will have 1 pointer (the lowest level),
  - a 25% chance that it will have 2 pointers (the second lowest level)...
  - and so on...
- This is what makes skiplists *probabilistic*.
- <https://opensa-server.cs.vt.edu/ODSA/Books/CS3/html/SkipList.html#id1>



Worst case :

level 0 : 50%  
level 1 : 25%



$$\frac{1}{32}$$

P(inserting  ~~$N$~~  in a row at level 0) =  $\frac{1}{1024} \left( \frac{1}{2^N} \right)$

# Analysis

- Consider the extremes:
  - All the nodes are at a high level, so lots of pointers. In that case, the insert cost could be quite high (even  $O(N)$ ). (Basically, this would be a linked list with extra pointers).
  - All the nodes are at a low level (e.g. level 0), so they only have 1 pointer each, which is basically just a linked list. In that case, insert and search could be  $O(N)$ .
- The good news: These extremes are unlikely to happen. For example, there is only  $1/1024$  chance that 10 nodes in a row will be in level 0.
- Similar to Quicksort, we accept that the randomization will make sure that the worst case is unlikely to happen.
- Expected memory use:  $O(N)$