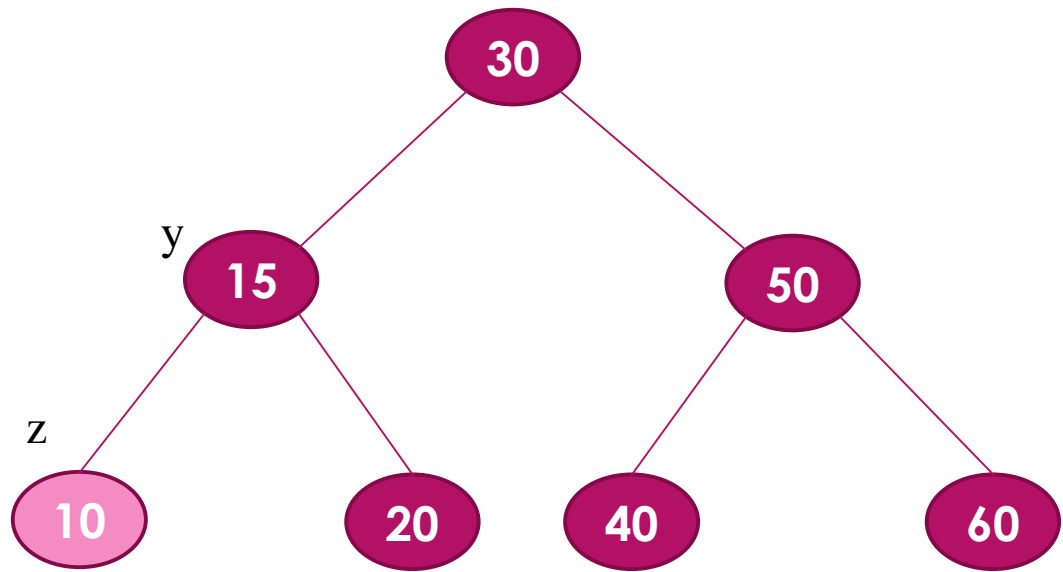


Overview

► Binary Search Tree Deletion

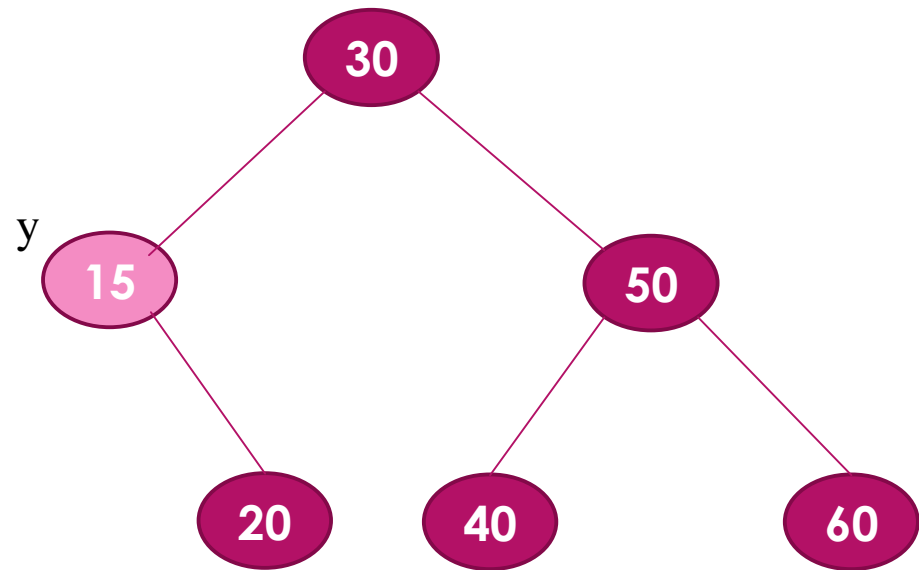
- Examples
- Different cases
- Algorithm

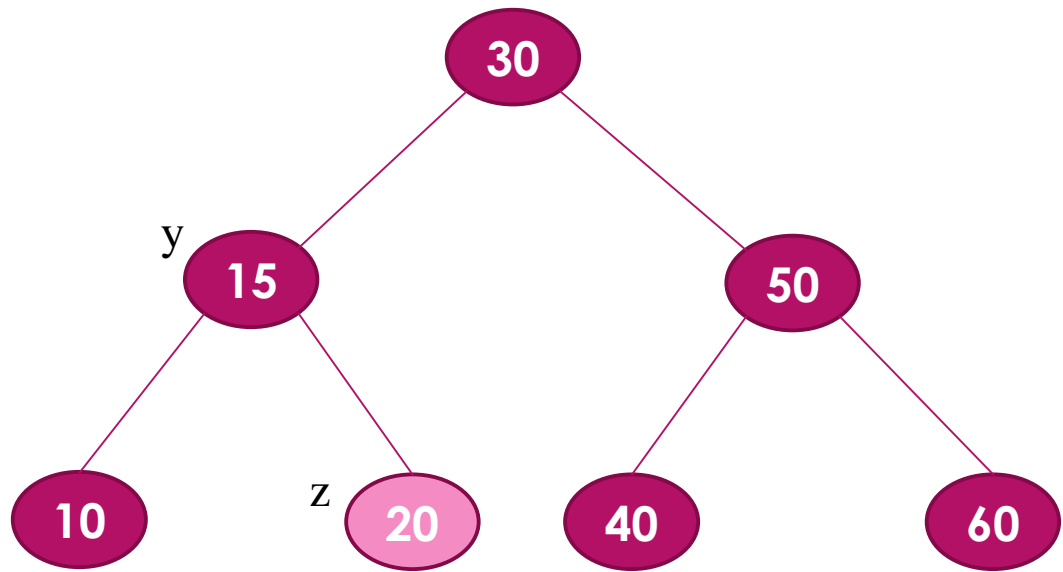


Delete z: a leaf node

Link updates ?

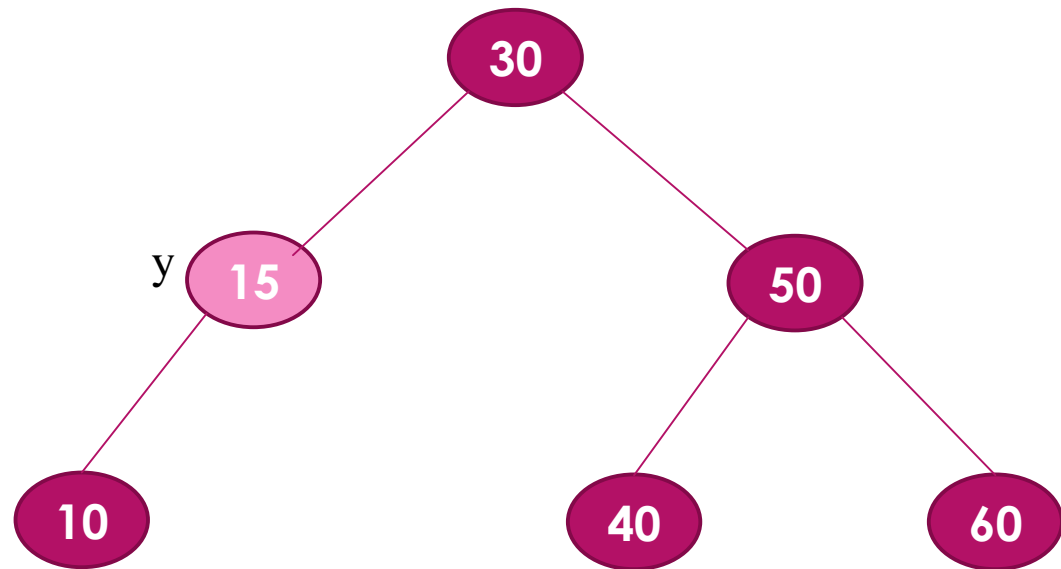
y.lchild to be set to NIL





delete **z**, a leaf
node

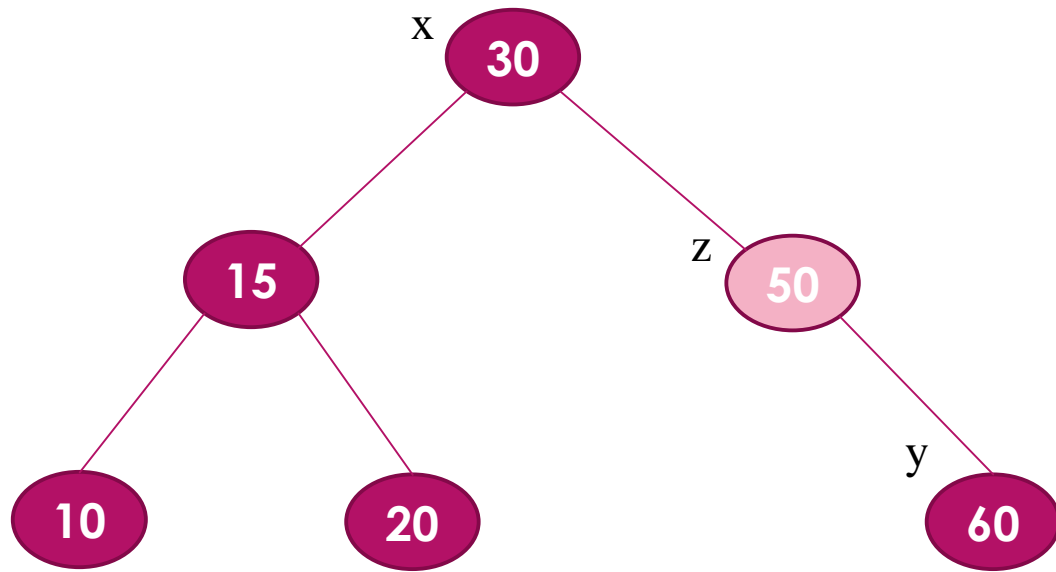
y.rchild to be set to **NIL**



BST Deletion - examples

► Deletion of

- a leaf node
-

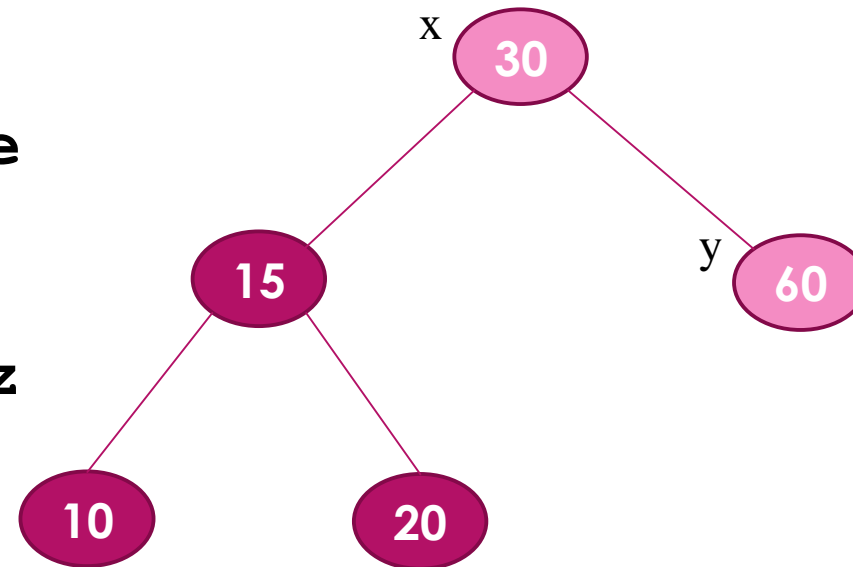


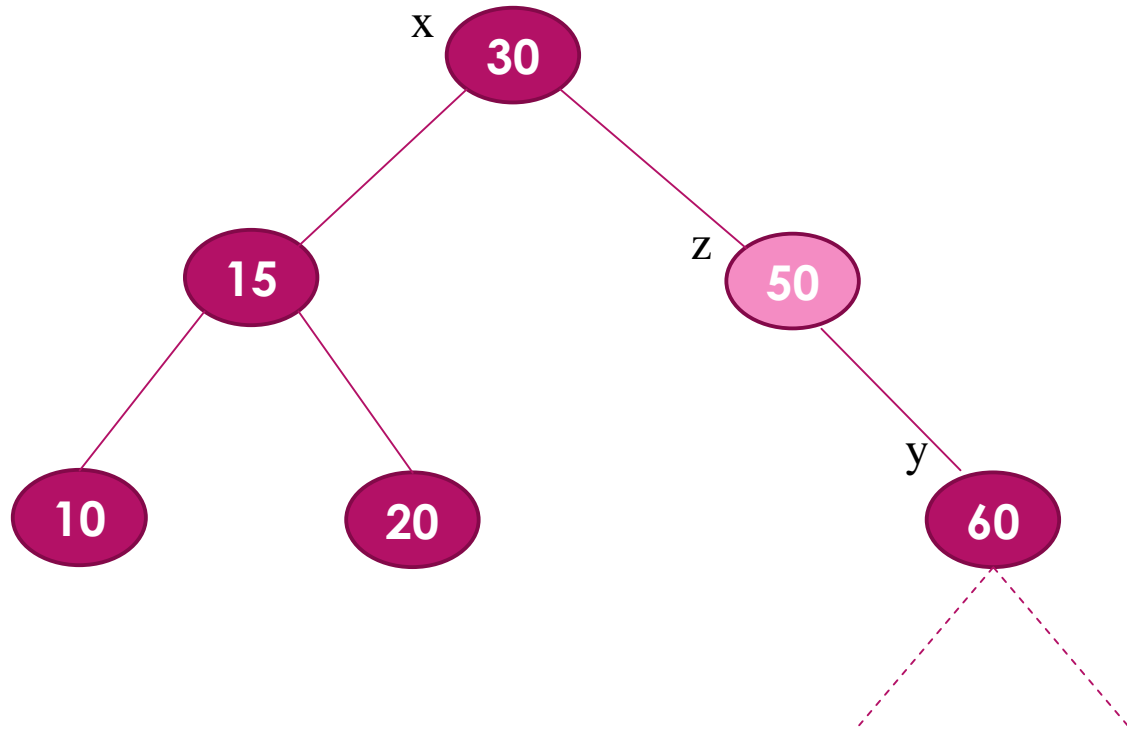
z: internal node with a single child

The lone child y can replace z

y.parent to be set as **x**

x.rchild to be set as **y**

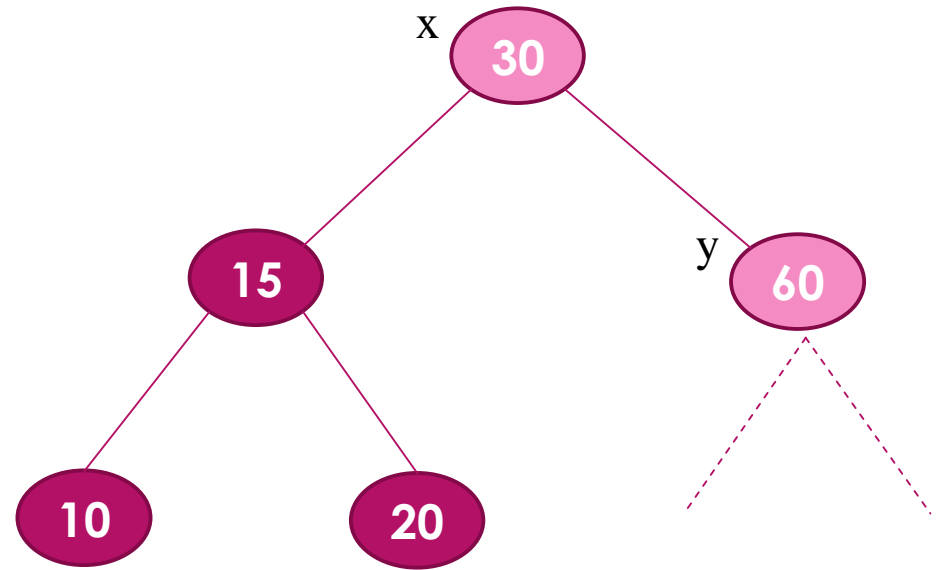




z: internal node with a single nonempty subtree

y can replace z

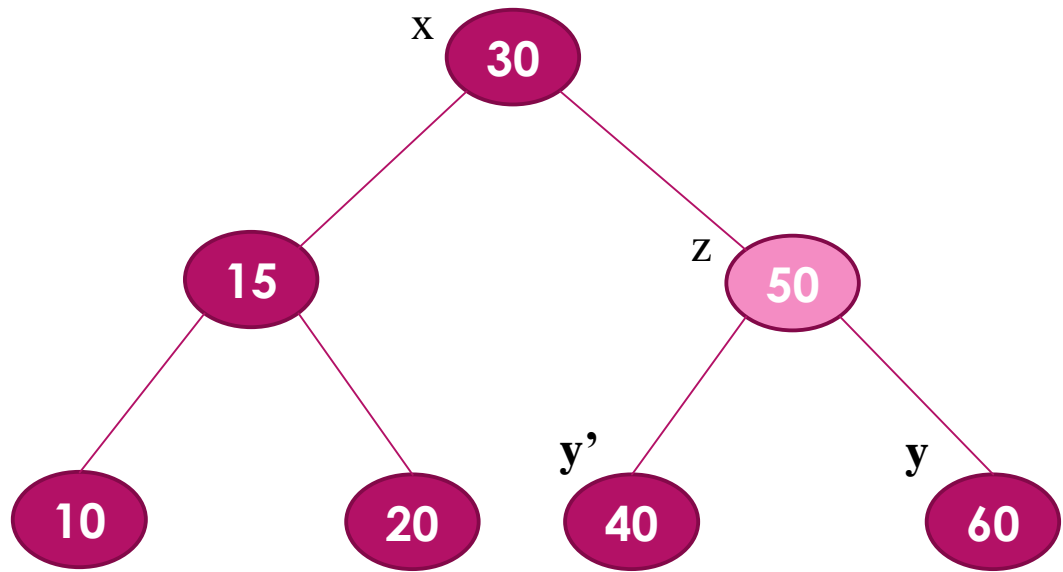
y.parent to be set as **x**
x.rchild to be set as **y**



BST Deletion - examples

► Deletion of

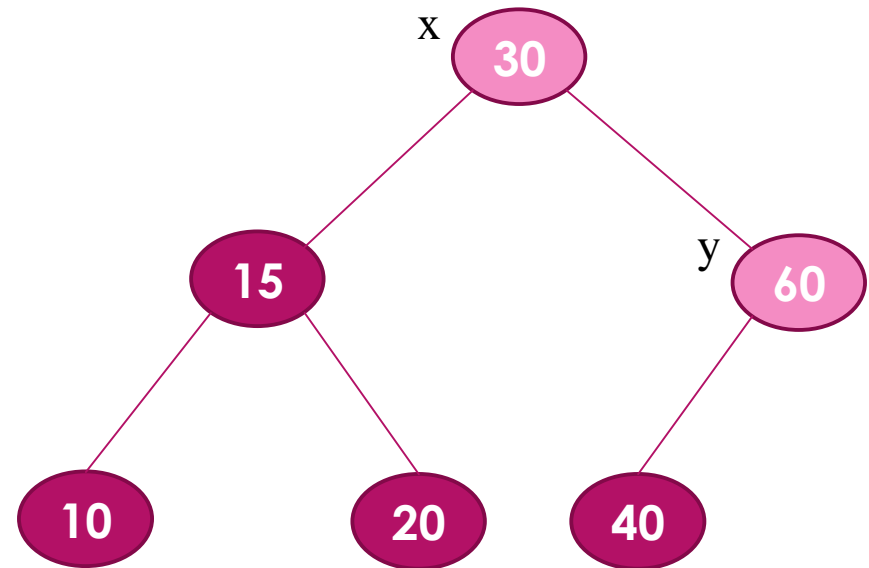
- a leaf node
- a node with one child
-



z: internal node with both left child
and right child nonempty

y.parent to be set as x
x.rchild to be set as y

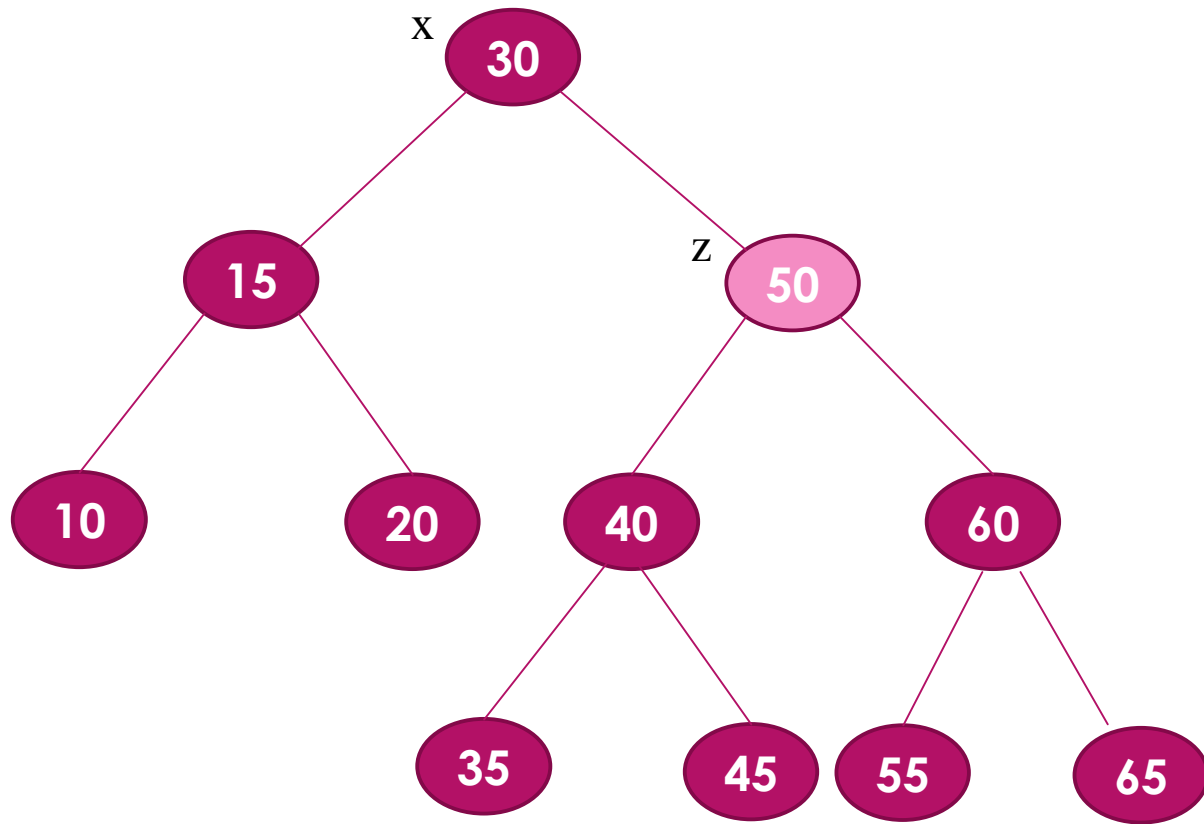
Alternatively can y' replace z ?



BST Deletion - examples

► Deletion of

- a leaf node
- a node with one nonempty subtree
- a node with both the children

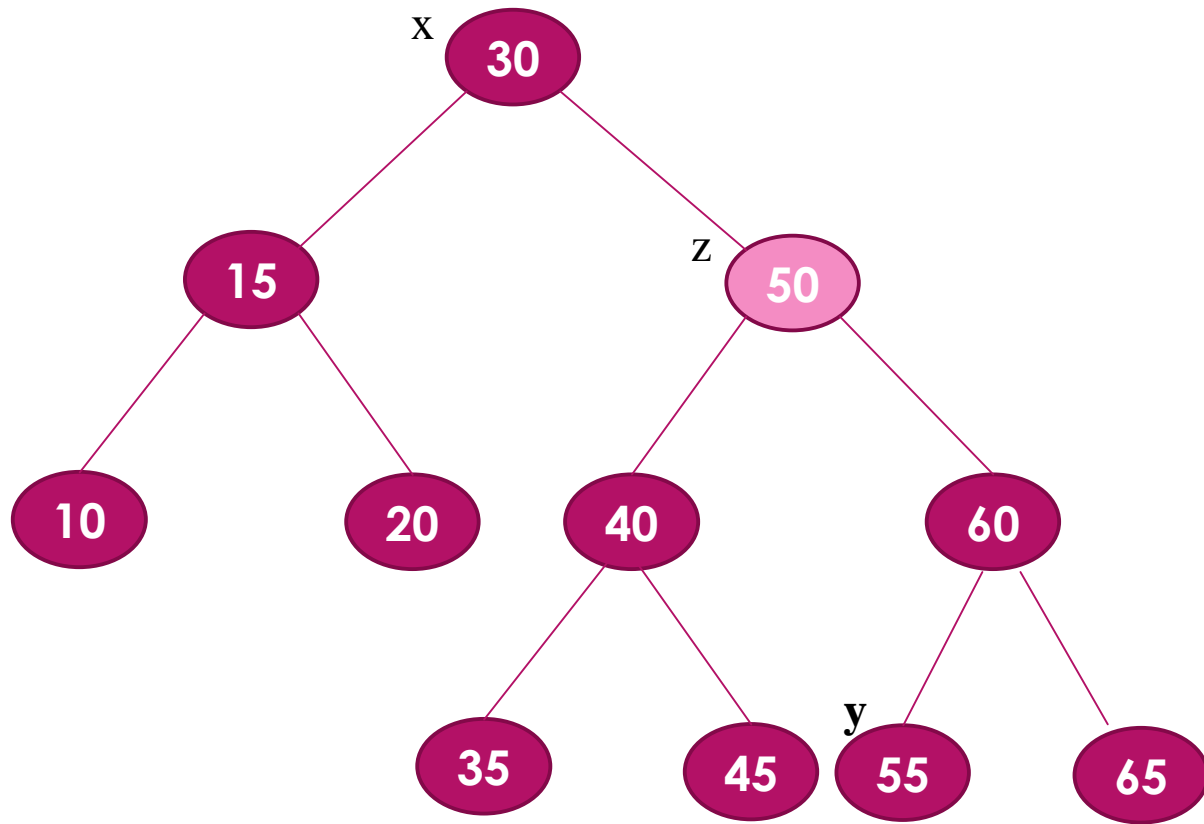


z: internal node, both the subtrees nonempty

Can not simply remove z

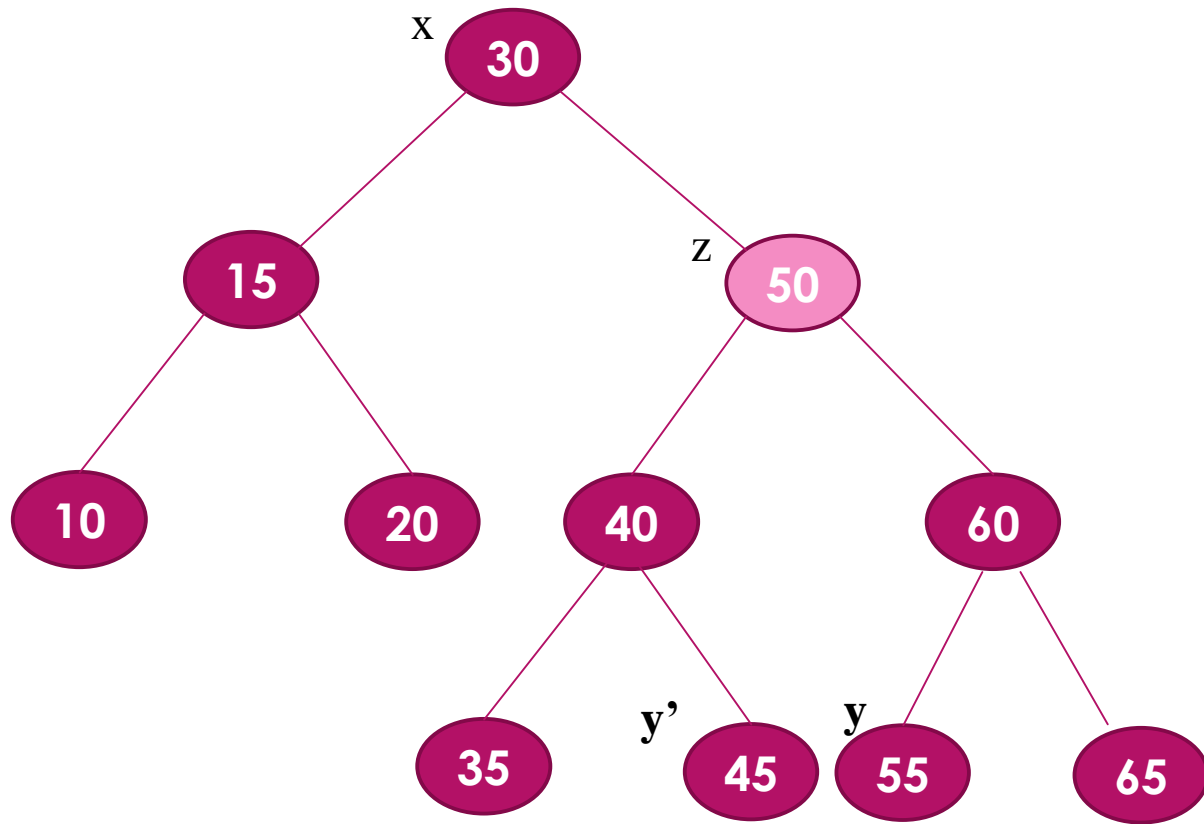
can z be replaced by some other node y ?

Can we select a y from the subtree of z?



z: internal node, both the subtrees nonempty

can z be replaced by y ?

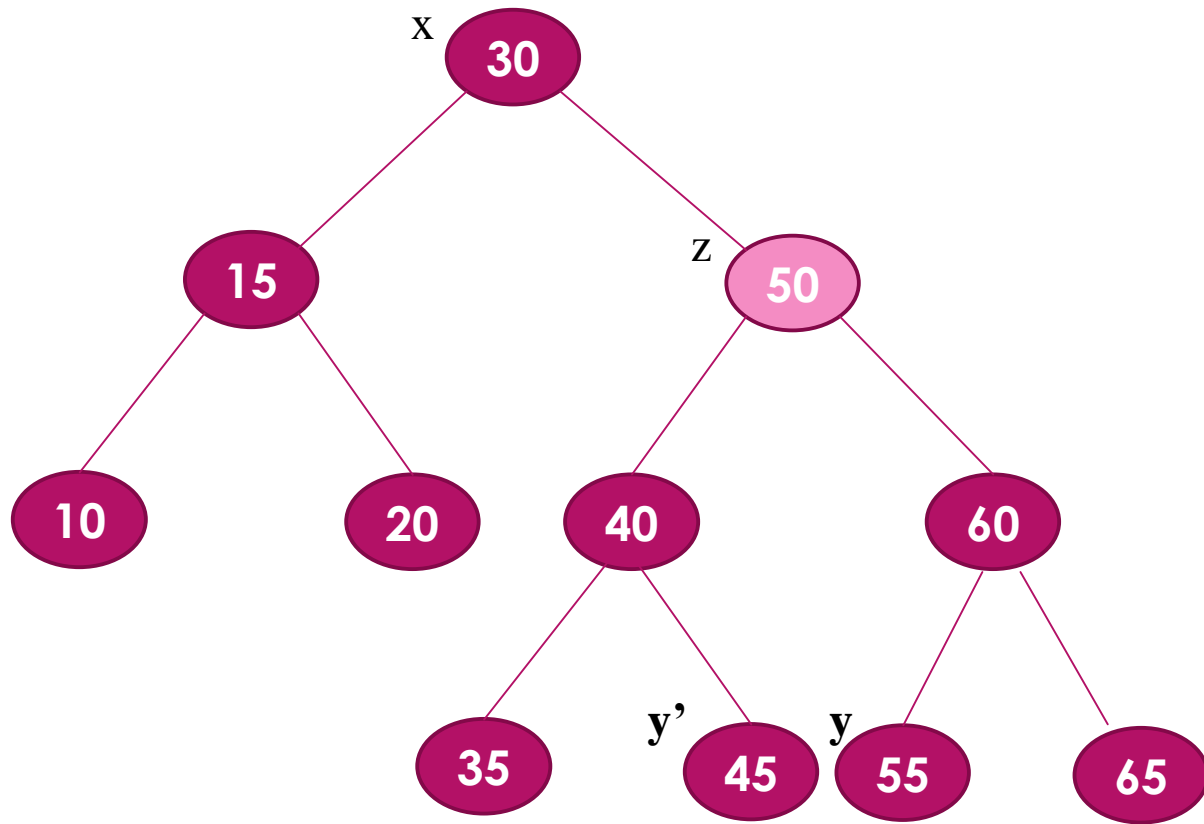


z: internal node, both the subtrees nonempty

can z be replaced by y'?

Replacing with z's inorder successor/inorder predecessor

BST property is to be maintained



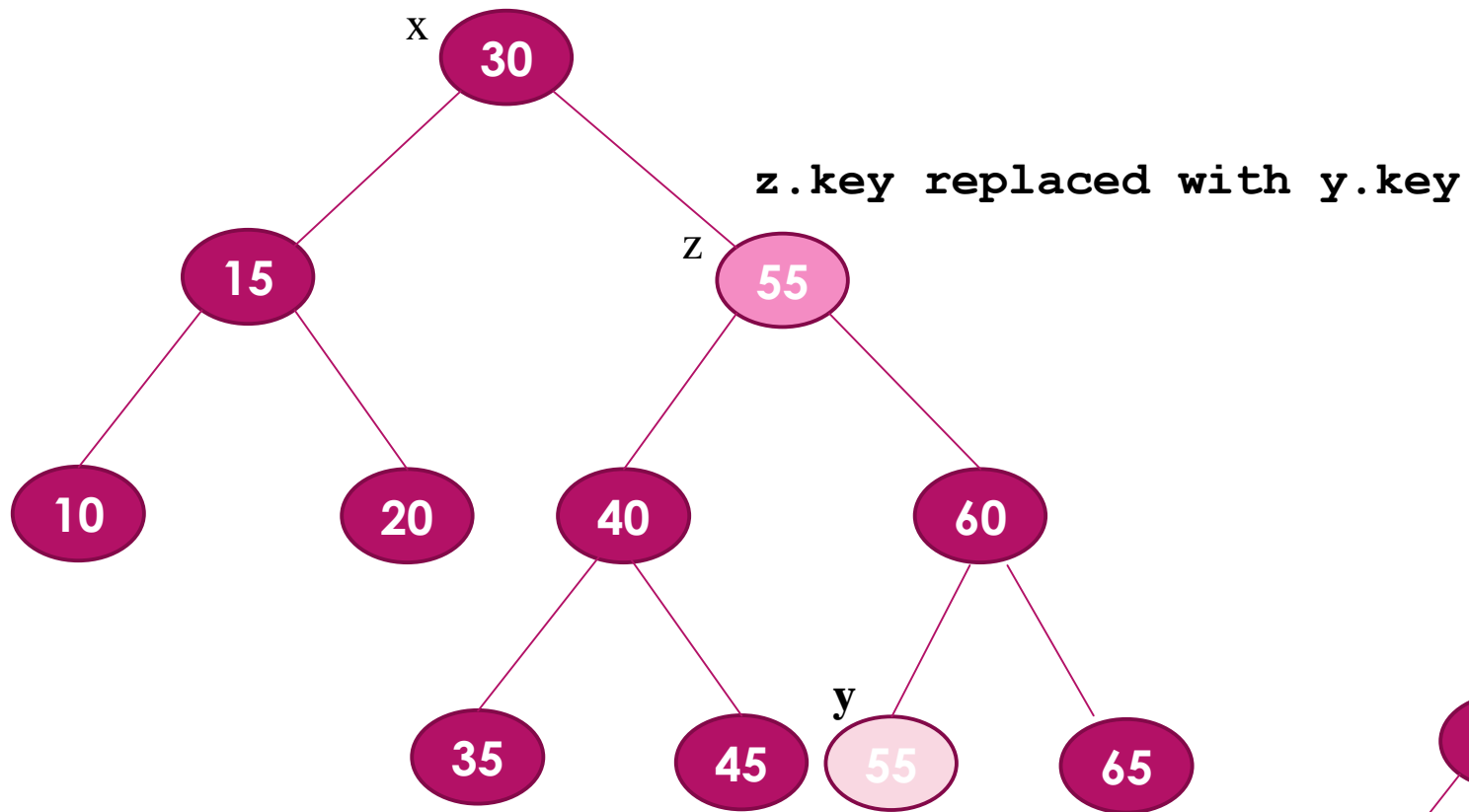
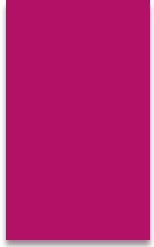
z: internal node, both the subtrees nonempty

Solution1:

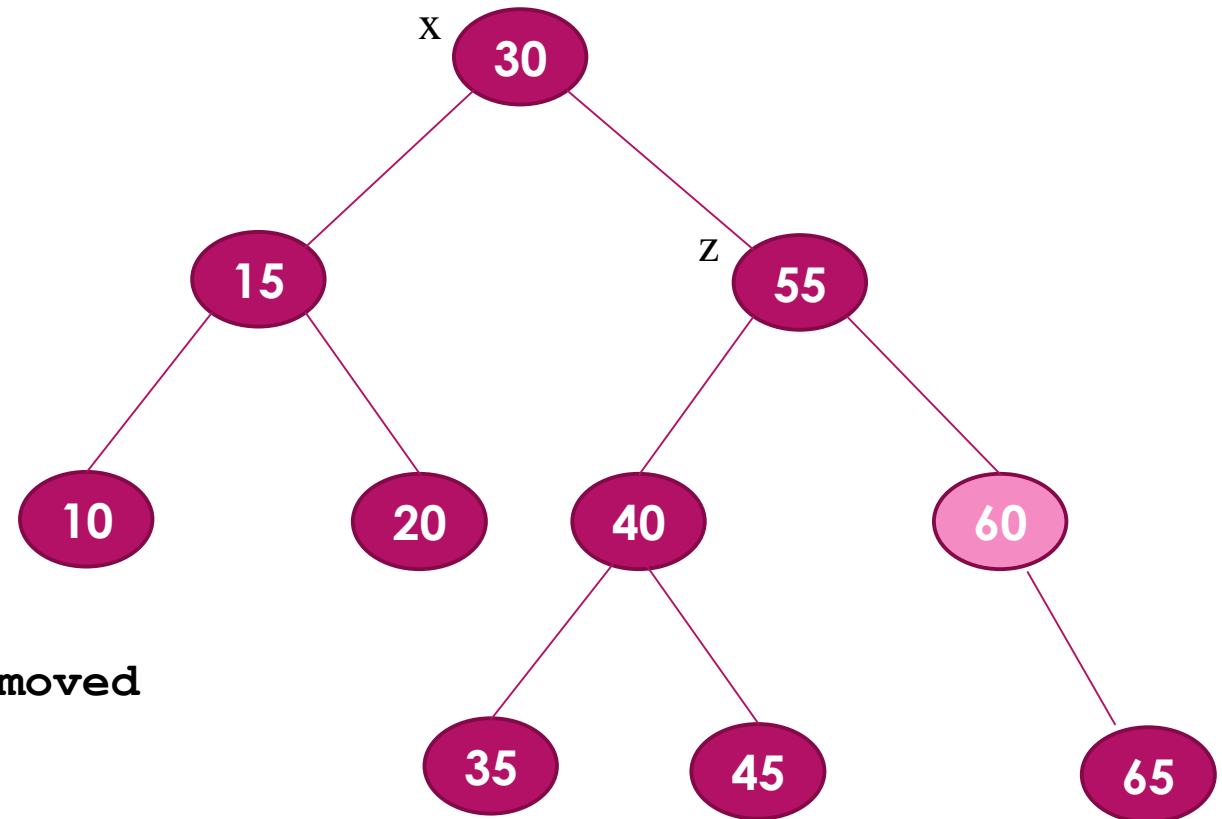
- Let **y** be **z**'s inorder successor
- Replace **z.key** with **y.key**
- Delete **y**

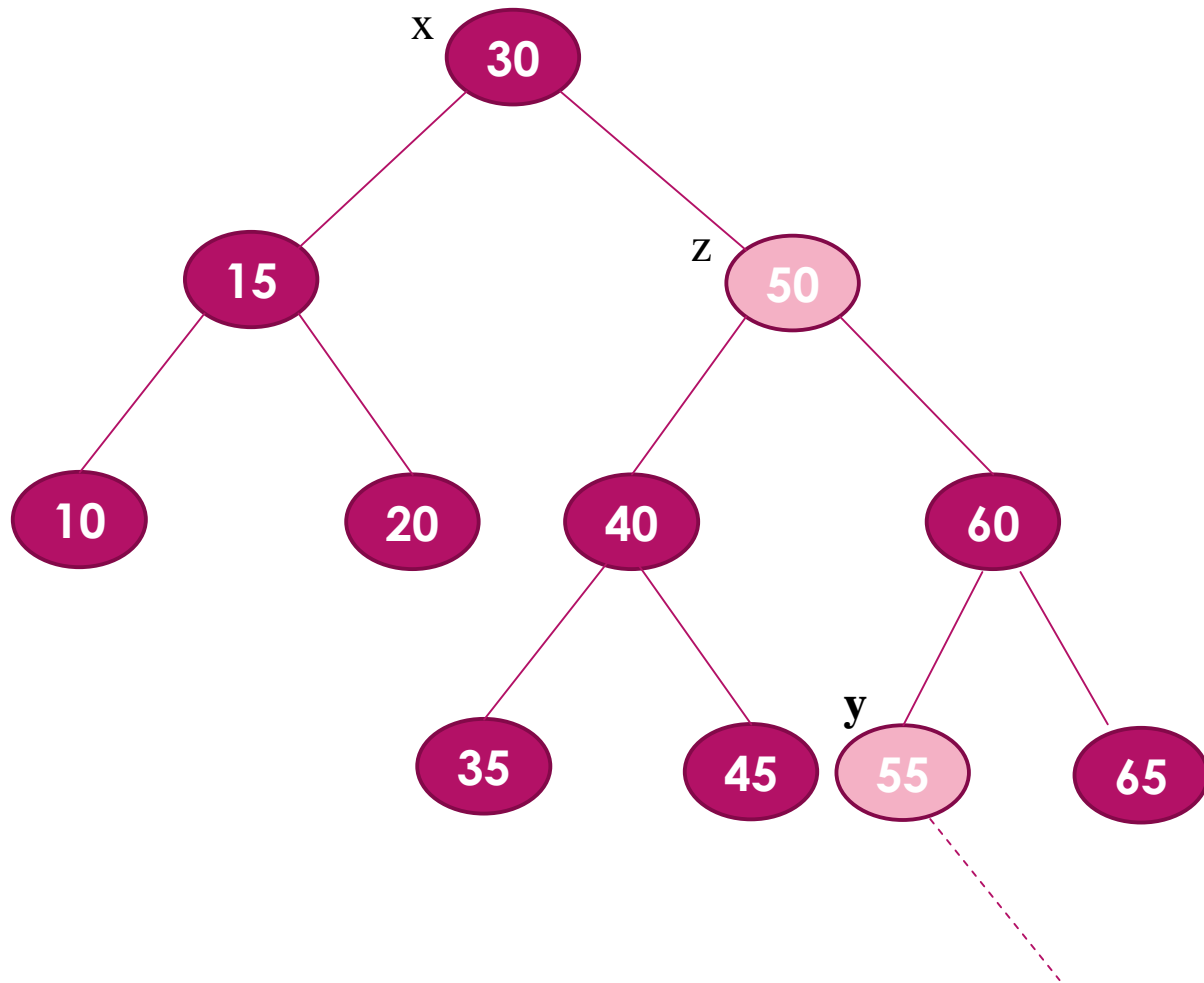
z is not physically removed

Is it correct to replace **z.key** with **y'.key**, where **y'** is the inorder predecessor of **z** and then remove **y'**?



y physically removed

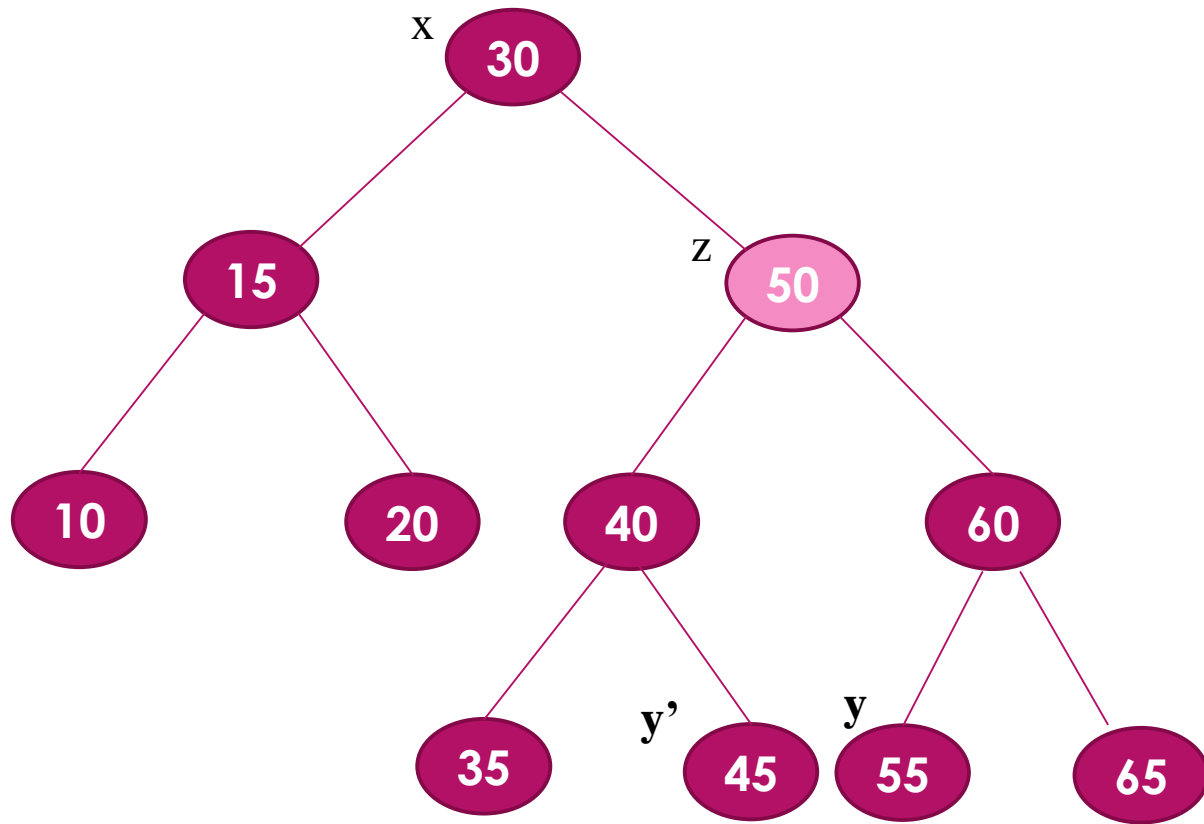




z: internal node with both the subtrees nonempty

z's successor has no left child, can have a nonempty right subtree

Physically deleted node has at most one nonempty subtree



z: internal node, both the subtrees nonempty. y: z's inorder successor)

Solution #1 (CLRS 2nd edn.):

- Copy data in y to z
- Delete y
- z is not physically removed

Solution #2 (CLRS 3rd edn.):

- Node z replaced by node y
- Node y is not deleted

BST Deletion - examples

► Deletion of

- a leaf node
- a node with one nonempty subtree
- a node with both the subtrees non empty

BST Deletion - Cases

► Let z be the node to be deleted

1. z has no children
2. z has just one child
3. z has both left and right child

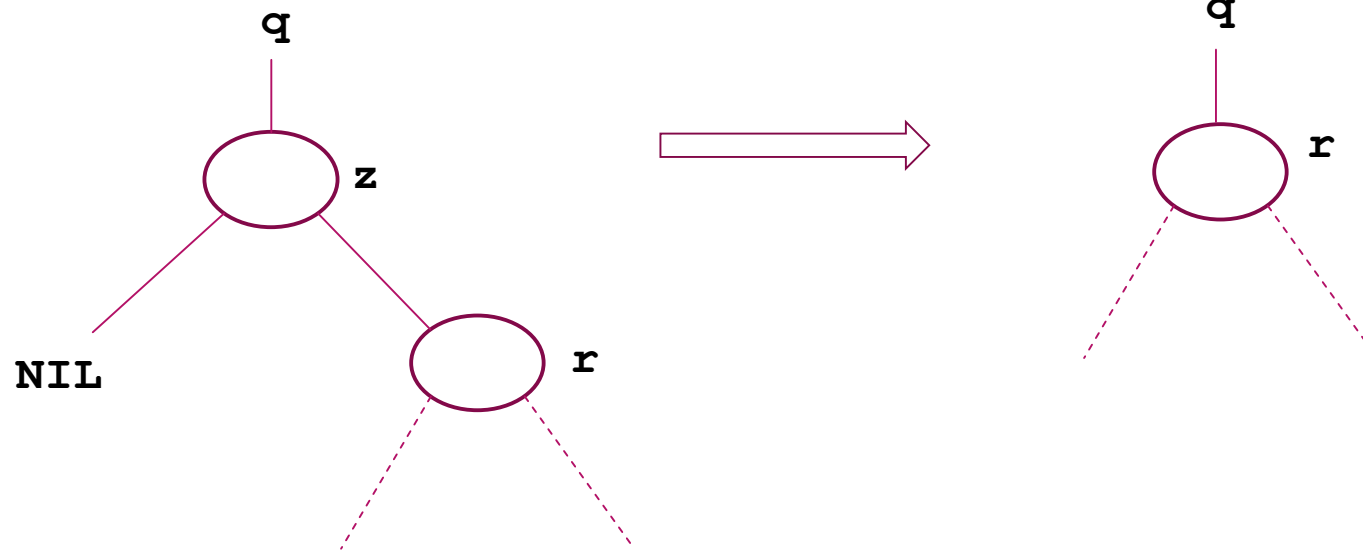
BST Deletion - Algorithm : Cases

► Let z be the node to be deleted

1. z has no left child
2. z has just one child, which is its left child
3. z has both a left and a right child

z being a leaf node - taken care of in case 1 (right child can be empty or non empty)

(a)

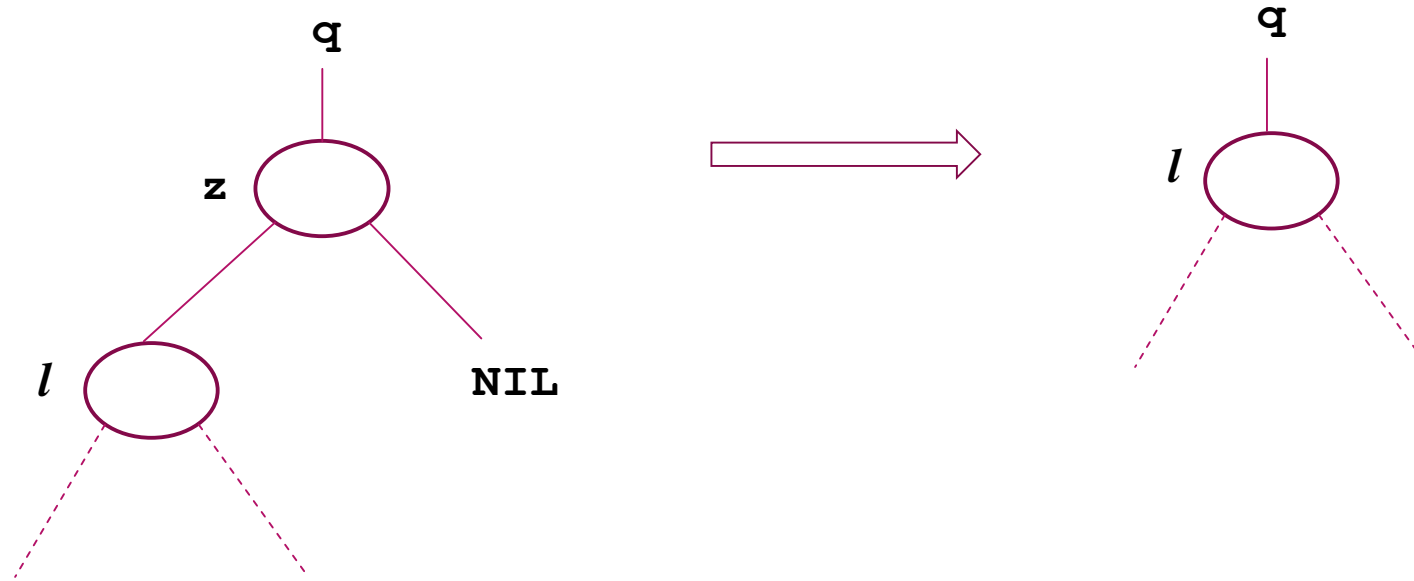


z has no left child

Replace **z** by its right child **r**

r may or may not be empty

(b)



z has left child l no right child
Replace z by its left child

BST Deletion : Cases a and b

```
TREE-DELETE (T, z)
```

```
    if z.left == NIL
```

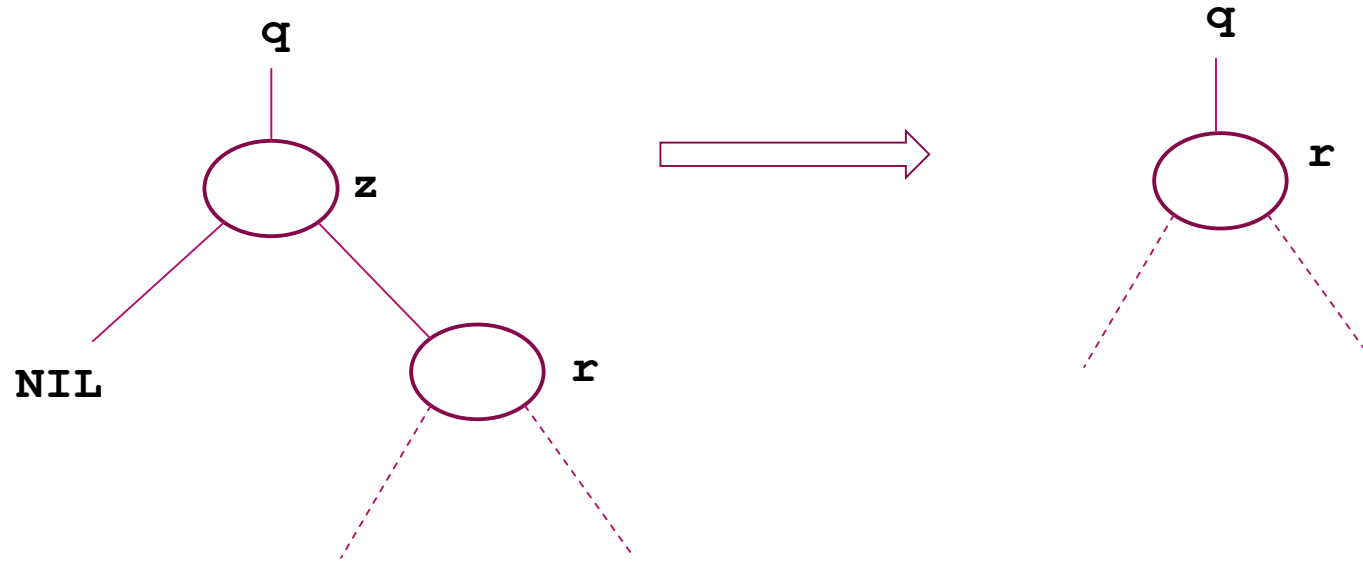
```
        TRANSPLANT (T, z, z.right)
```

```
    elseif z.right == NIL
```

```
        TRANSPLANT (T, z, z.left)
```

```
    else ..... //z has both children, split into two cases c and d
```

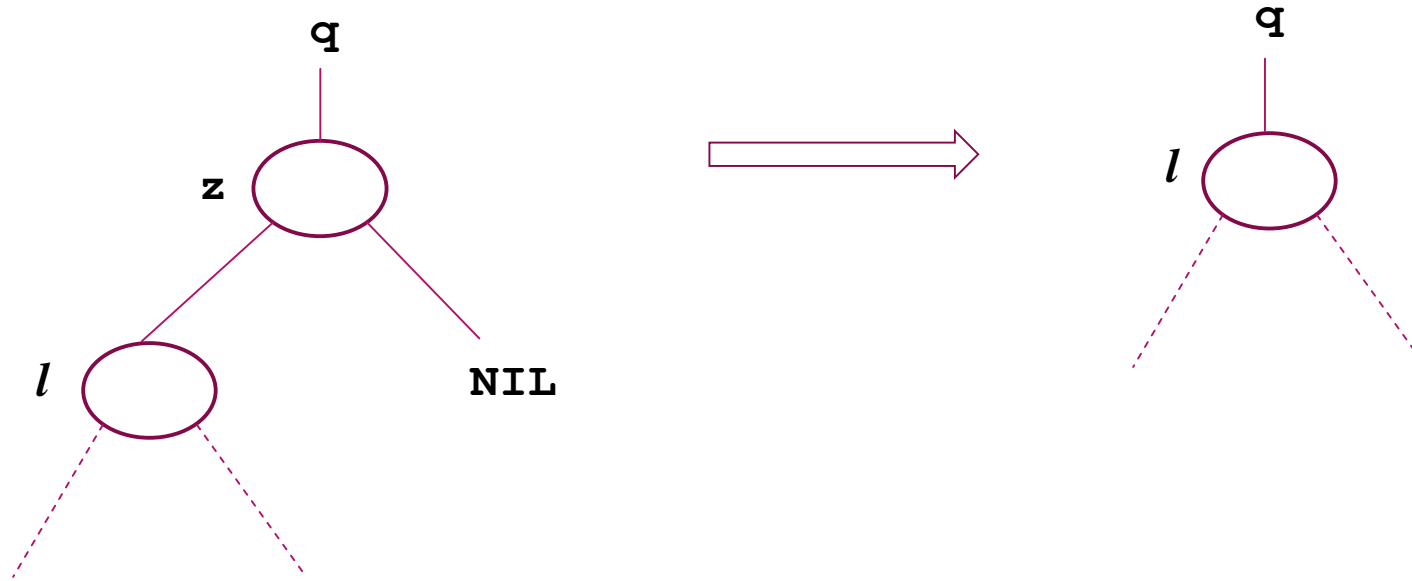
(a)



```
if z.left == NIL  
    TRANSPLANT (T, z, z.right)
```

subtree rooted at **z** replaced with the subtree rooted at **r**

(b)



```
.....elseif z.right == NIL
```

```
    TRANSPLANT (T, z, z.left) //z has left child l
```

subtree rooted at *z* replaced with the subtree rooted at *l*

BST Deletion : Cases a and b

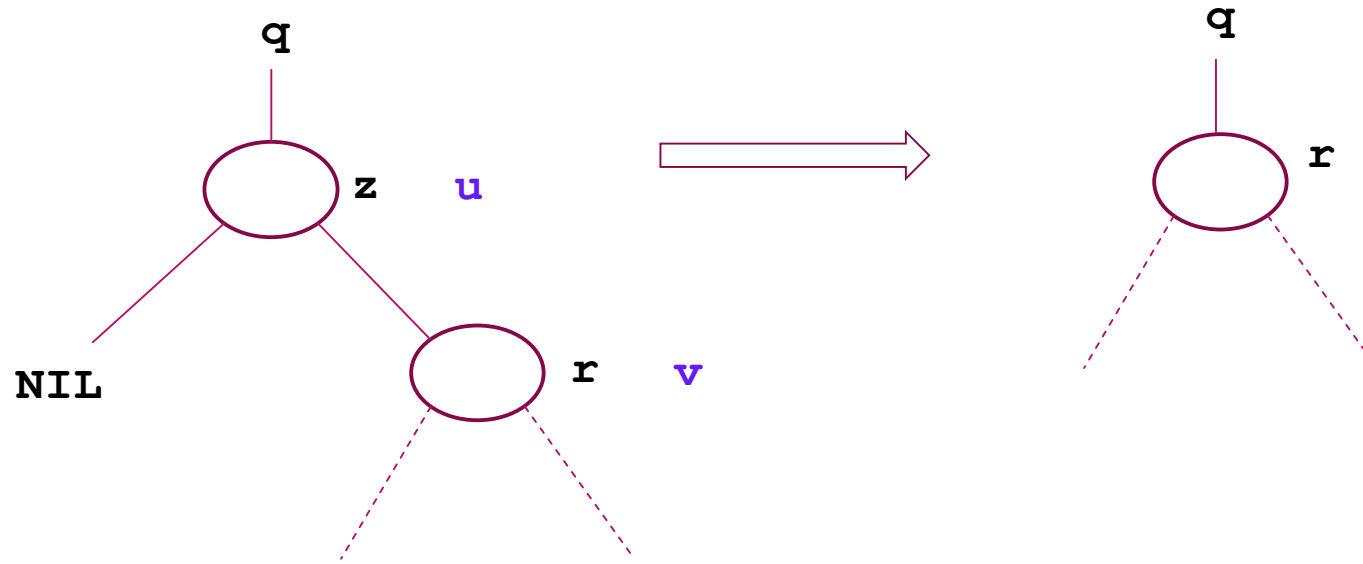
```
TRANSPLANT(T, u, v) // replaces the subtree rooted at u with
                     // the subtree rooted at v

    if u.p == NIL    // u is root
        T.root = v

    elseif u == u.p.left //u is lchild of its parent
        u.p.left = v
    else u.p.right = v

    if v ≠ NIL
        v.p = u.p    // v.left, v.right updations, if required, to
                     // to be done by the caller
```

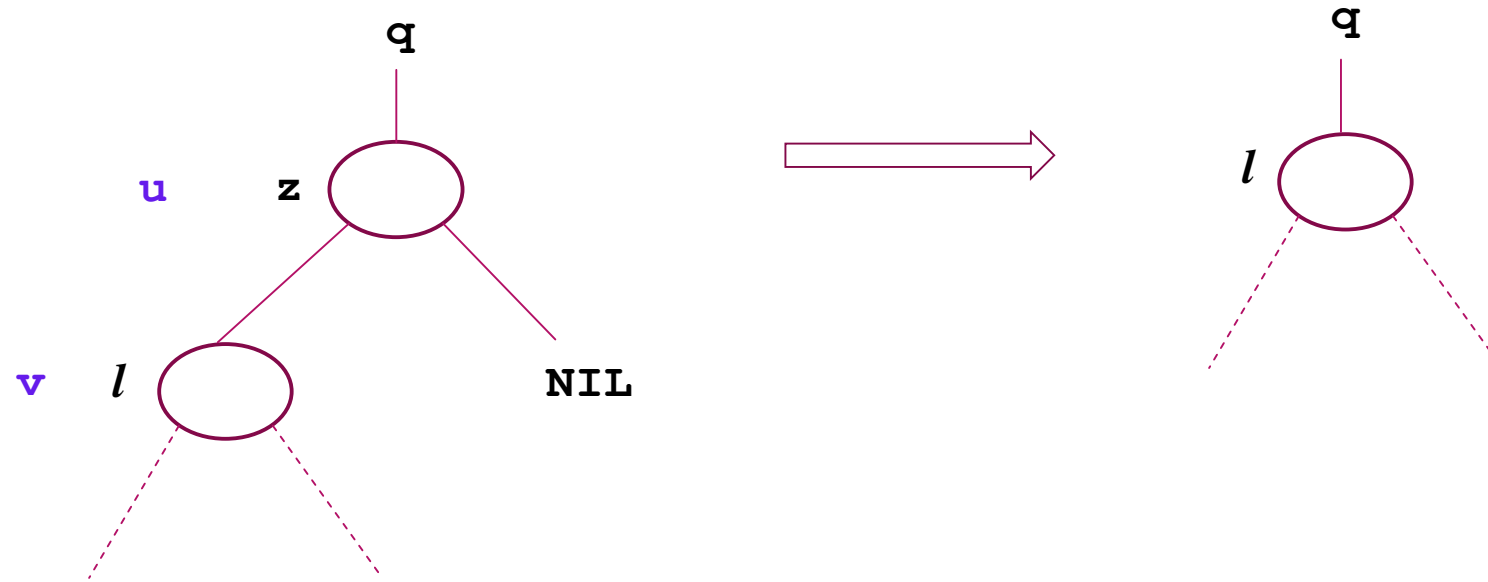
(a)



```
if z.left == NIL
    TRANSPLANT (T, z, z.right)
```

`TRANSPLANT(T, u, v)` replaces the subtree rooted at `u` with the subtree rooted at `v`

(b)



```
.....elseif z.right == NIL  
        TRANSPLANT (T, z, z.left) //z has left child l
```

Reference

1. T H Cormen, C E Leiserson, R L Rivest, C Stein *Introduction to Algorithms*, 3rd ed., PHI, 2010