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Privacy-Preserving Remote Diagnostics

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Agenda



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- ❑ Introduction
- ❑ Remote Diagnostics
- ❑ Basic knowledge and terminologies
 - Branching Program (BP)
 - Oblivious Transfer (OT)
 - Garbled Circuits (GC)
 - Homomorphic Encryption (HE)

❑ Protocol Digest

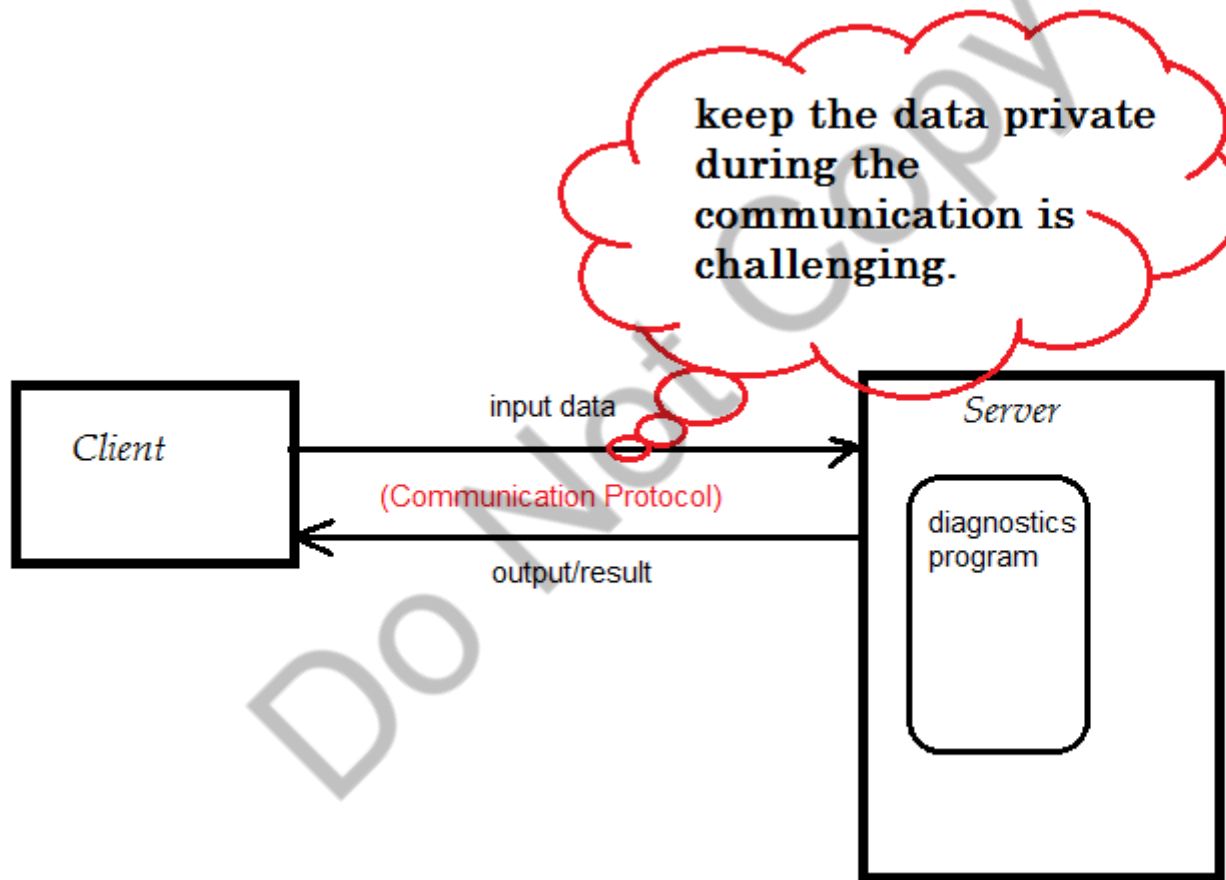
- Secure branching program protocol
- A protocol for secure evaluation of private LBPs -
SecureEvalPrivateLBP

❑ Performance Comparison

Introduction



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Remote diagnostics - risk for both the parties



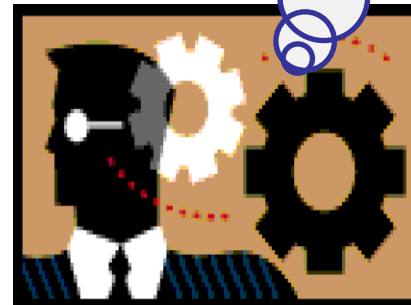
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I want my
personal
information and
secrets protected



client

I want to protect
the diagnostics
program and the
intellectual
property



Server

□ Applications:

- Healthcare (medical treatment sector)
- Remote software fault analysis (fault diagnostics)
- ElectroCardioGram signals

Example: In medical applications:

- **Patient** wants his personal data to be protected from medical applications while receiving an analysis
- The **service provider** wants his algorithms to be protected as it is an intellectual property

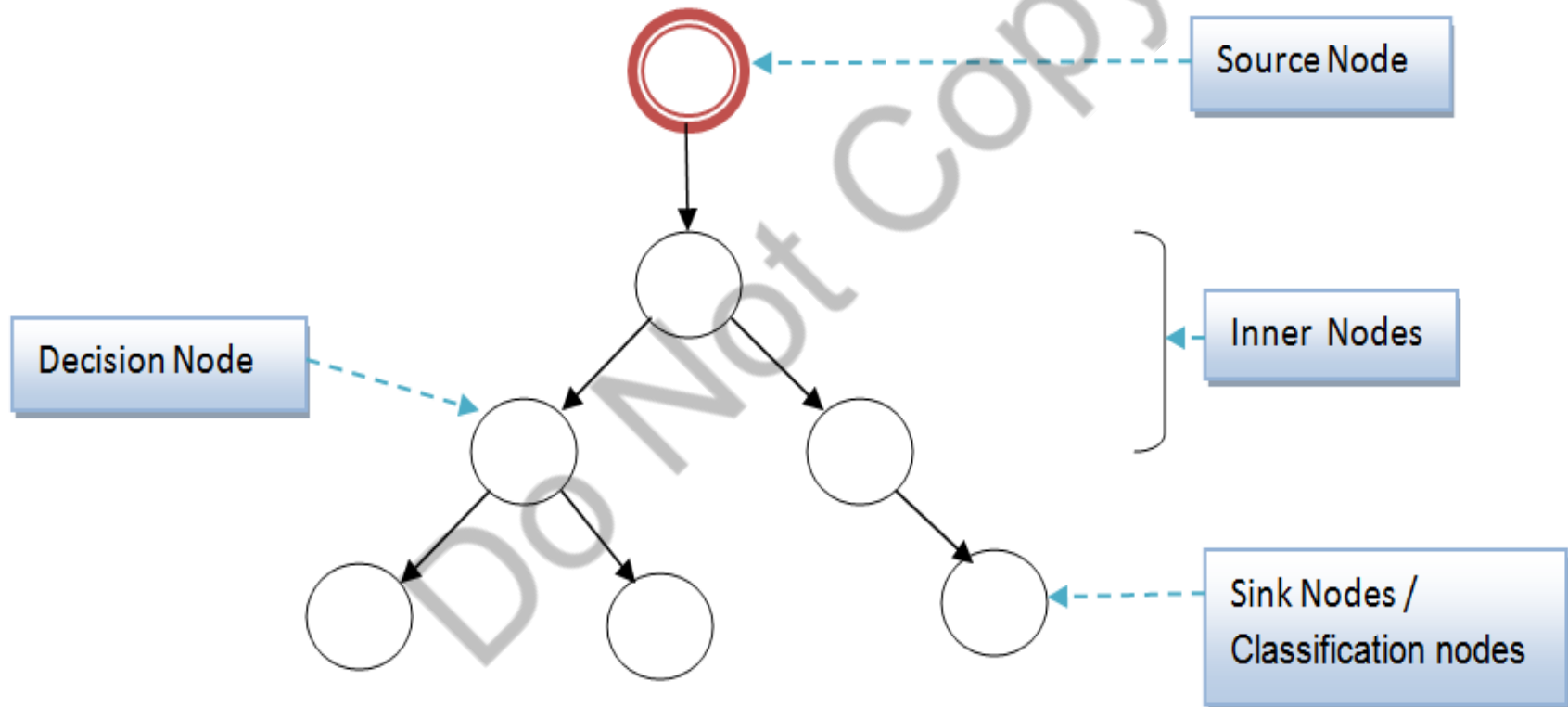
Branching Programs (BP)



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- ❑ Set of decision and classification nodes
- ❑ Decision nodes are inner nodes on which the program branches until a classification node is reached
- ❑ Can represent boolean functions which are compatible with Garbled Circuits

Branching Programs (BP)



Linear Branching Programs



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- ❑ Linear Branching Programs are a generalization BPs
- ❑ While BP use only one attribute for comparison in decision nodes, **LBP use a linear combination of the users attribute vector**

Oblivious Transfer



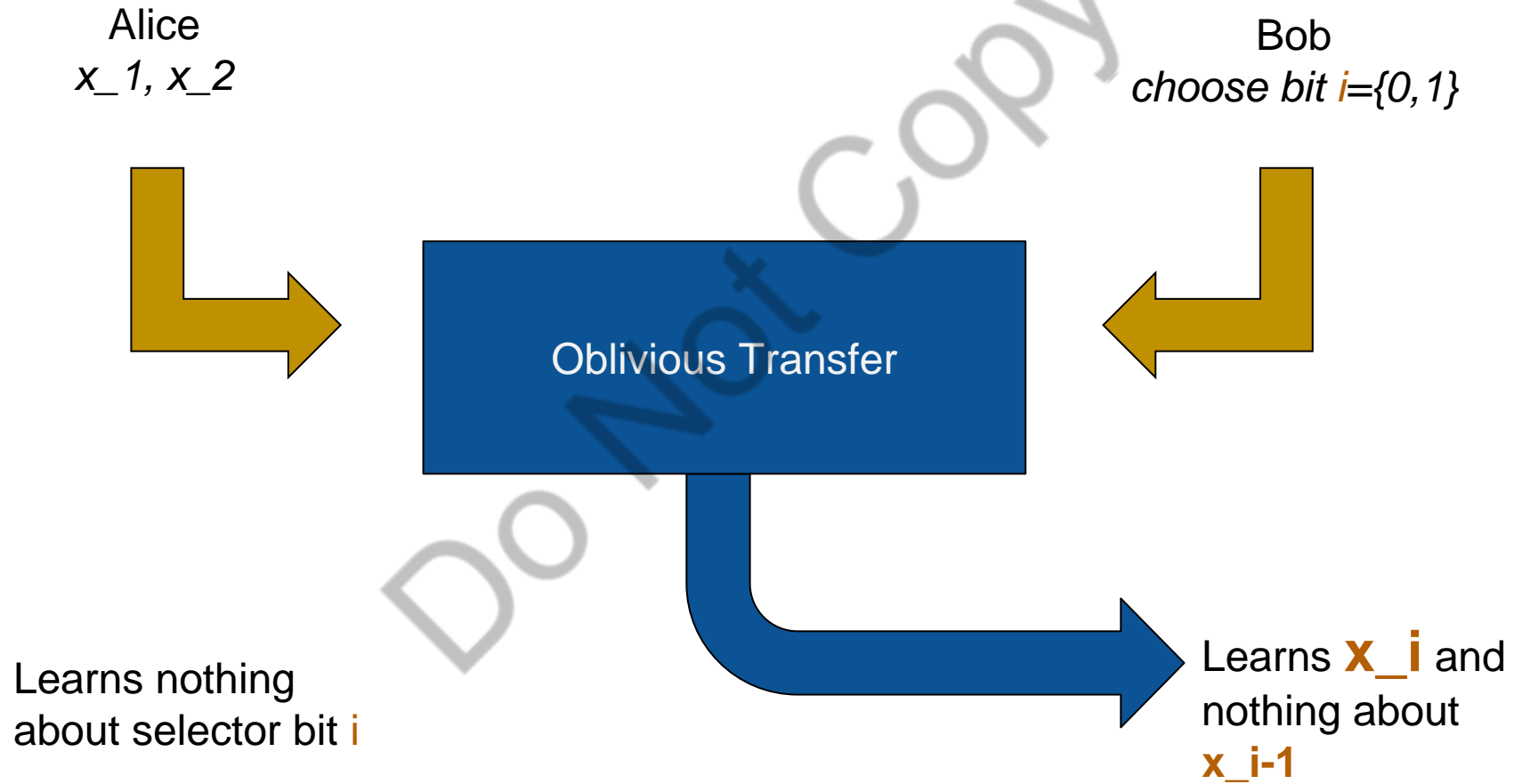
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- ❑ Method for information transfer between a sender and receiver
- ❑ The sender does not know which information was requested
- ❑ While the receiver only learns one of the two inputs
- ❑ Used by Bob to learn his input values for the GC from Alice

Oblivious Transfer



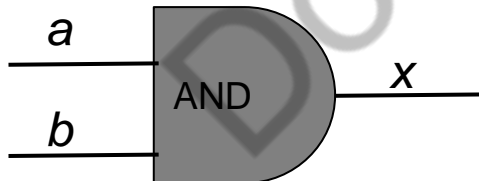
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- ❑ The receiver chooses a bit $b = \{0, 1\}$ which correlates to the input he wants to select
- ❑ The Sender inputs two values between the receiver will choose
- ❑ The receiver receives the chosen value while not learning anything about the other value

Garbled Circuit

- ❑ Encrypted circuit with a pair of wire keys for every input wire.
- ❑ The wire keys enable the computation of the output based on the encrypted input
- ❑ Nothing is learned from the gate computation



a	b	x
a_1	b_1	$Enc_{a_1 \parallel b_1}(x_1)$
a_1	b_0	$Enc_{a_1 \parallel b_0}(x_0)$
a_0	b_1	$Enc_{a_0 \parallel b_1}(x_0)$
a_0	b_0	$Enc_{a_0 \parallel b_0}(x_0)$

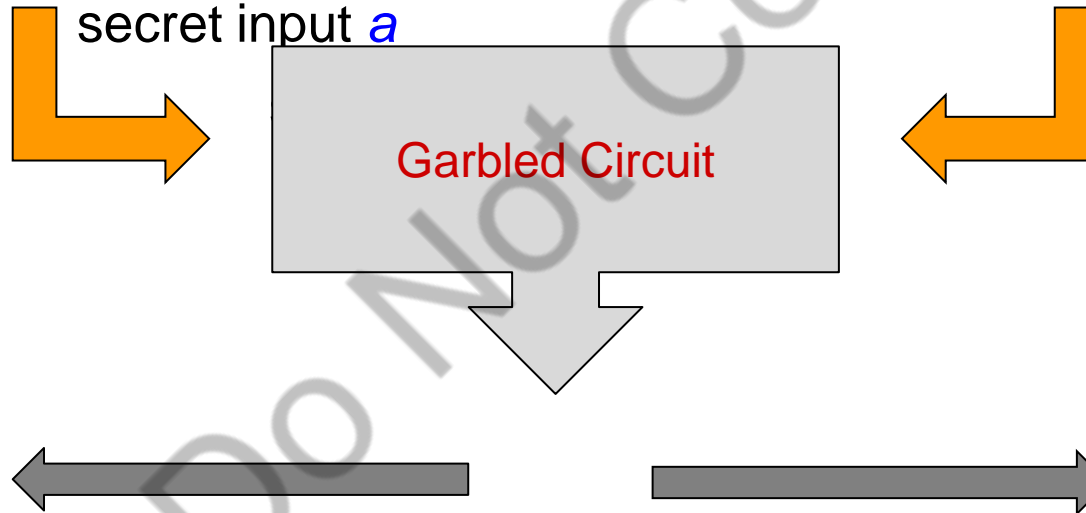
Garbled Circuits (GC) - Yao's garbled circuits protocol

Alice(Server)

agree on function f

Bob(Client)

secret input a



learns nothing about b
 $m = f(a, b)$

learns nothing about a
 $f(a, b)$

$m = f(a, b)$

Yao's garbled circuits protocol



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- ❑ Protocol for secure two-party computation
 - Constant number of rounds
 - Secure against semi honest adversary
- Can compute any function based on a boolean circuit securely
- Builds on garbled circuits and oblivious transfer

Yao's Protocol Steps



- ❑ **Step 1:** Alice generates the garbled circuit C and two encrypted labels for each input wire
- ❑ **Step 2:** Alice sends her encrypted input (wire labels) to Bob
- ❑ **Step 3:** Bob receives his wire labels from Alice using 1-out-of-2 oblivious transfer without revealing his inputs to Alice
- ❑ **Step 4:** Bob uses Alice's wire labels and his own wire labels to compute the output and send it back to Alice

Homomorphic Encryption



- ❑ Manipulation of plaintext without knowing the corresponding ciphertext

- ❑ Additiv homomorphic encryption is :

$$Enc_k(x_1 + x_2) = Enc_k(x_1) + Enc_k(x_2)$$

- ❑ Addition can be done on ciphertext and decrypted plaintext result is correct

❑ Secure branching program protocol

- used for secure evaluation of binary branching programs.

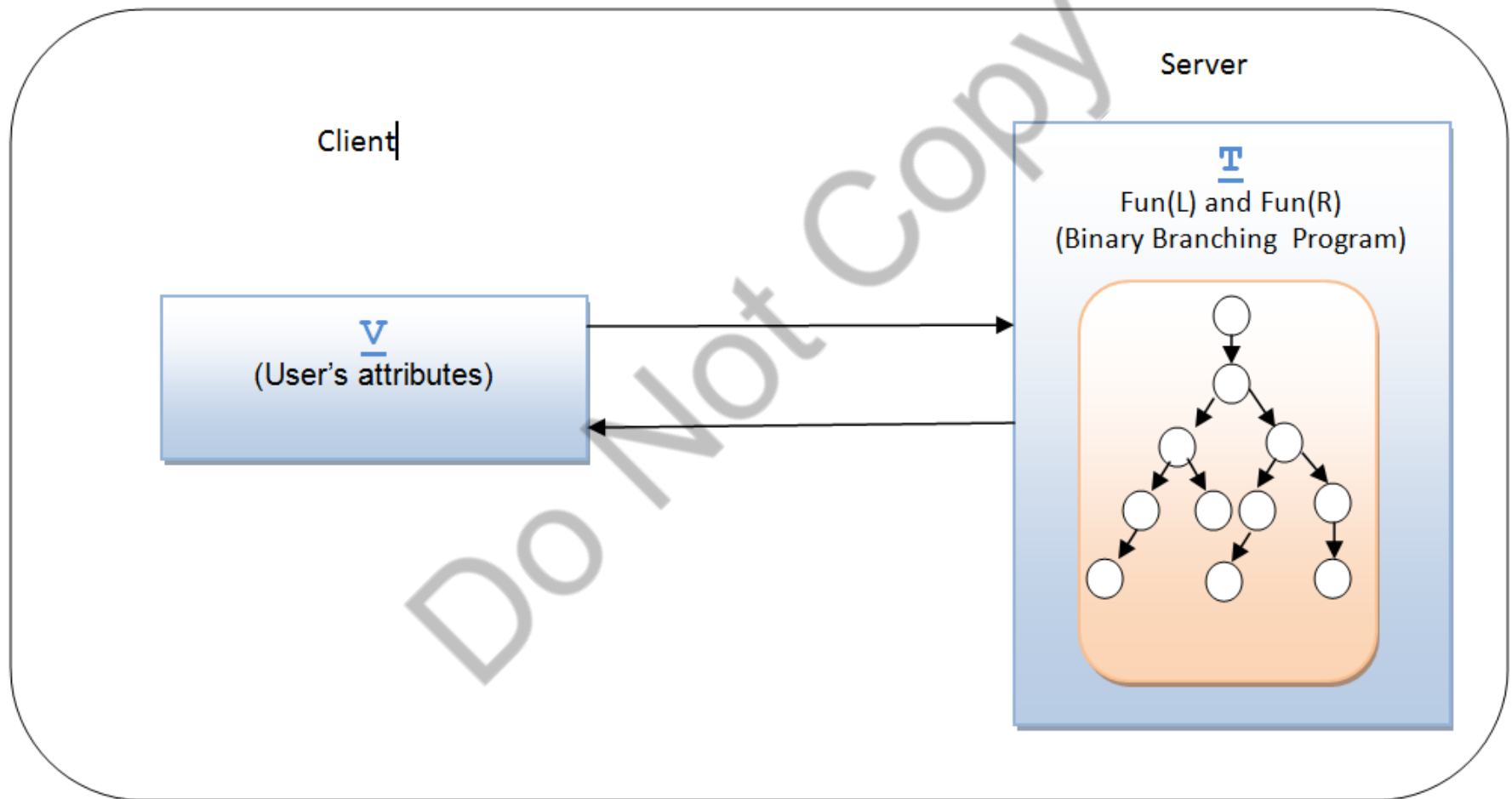
❑ A protocol for secure evaluation of private LBPs - SecureEvalPrivateLBP

- secure evaluation of private linear branching programs.

Secure branching program protocol



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The main goal of the protocol is to securely analyze the T on V

Phase I (Offline): Creation of the secure branching program

- Transforms the nodes in branching program T into secure nodes in branching program T'
- **Classification node** is replaced by encryption of its **classification label**.
- **Decision node** is replaced with a **small garbled circuit** performing offset integer comparison
- User can figure out one of the decryption keys
- the revealed key can help to decrypt the next node on the evaluation path

Secure branching program protocol



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Phase II: Oblivious attribute selection

Phase III: Evaluation of the secure branching program

A protocol for secure evaluation of private LBPs - **SecureEvalPrivateLBP**

- ❑ Secure evaluation of private linear branching programs
- ❑ **Linear Branching Programs (LBP)** generalize binary classification or decision trees and Ordered Binary Decision Diagrams (OBDDs)
- ❑ The protocol **SecureEvalPrivateLBP** is divided into **three phases**:
 - Phase I: CreateGarbledLBP
 - Phase II: ObliviousLinearSelect
 - Phase III: EvalGarbledLBP

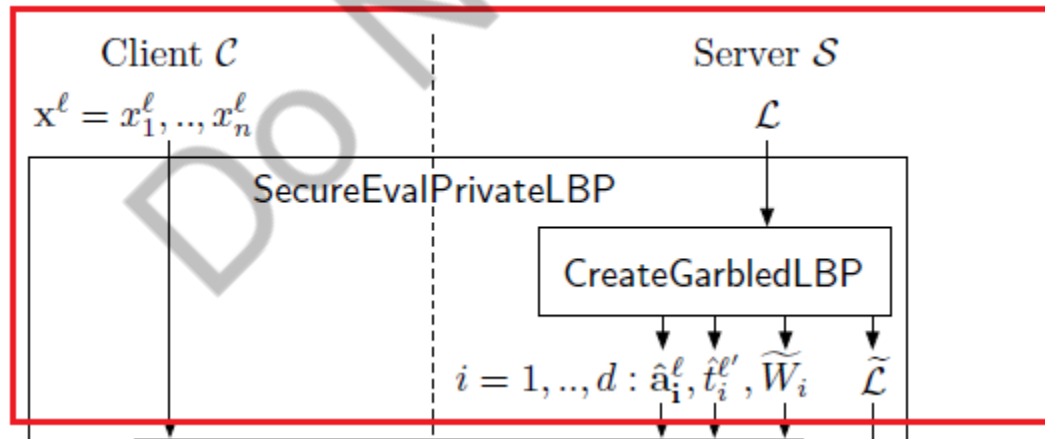
A protocol for secure evaluation of private LBPs - SecureEvalPrivateLBP

C : attribute vector x^ℓ

S : server

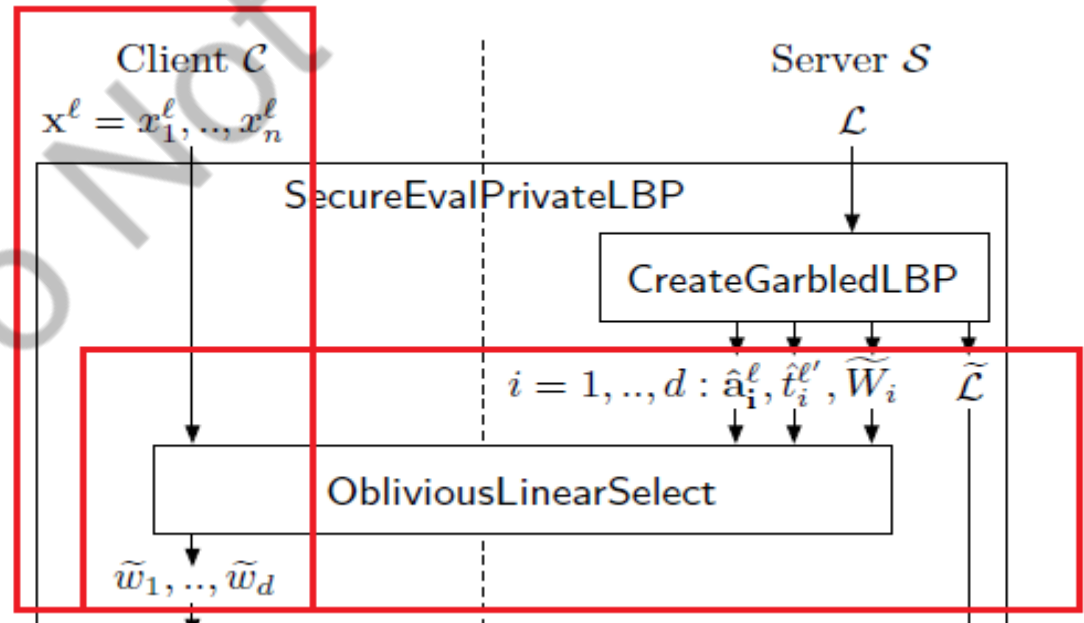
\mathcal{L} : linear branching program

- Phase I: CreateGarbledLBP
 - the server S generates the garbled version of the LBP \mathcal{L}
 - randomize LBP permutation



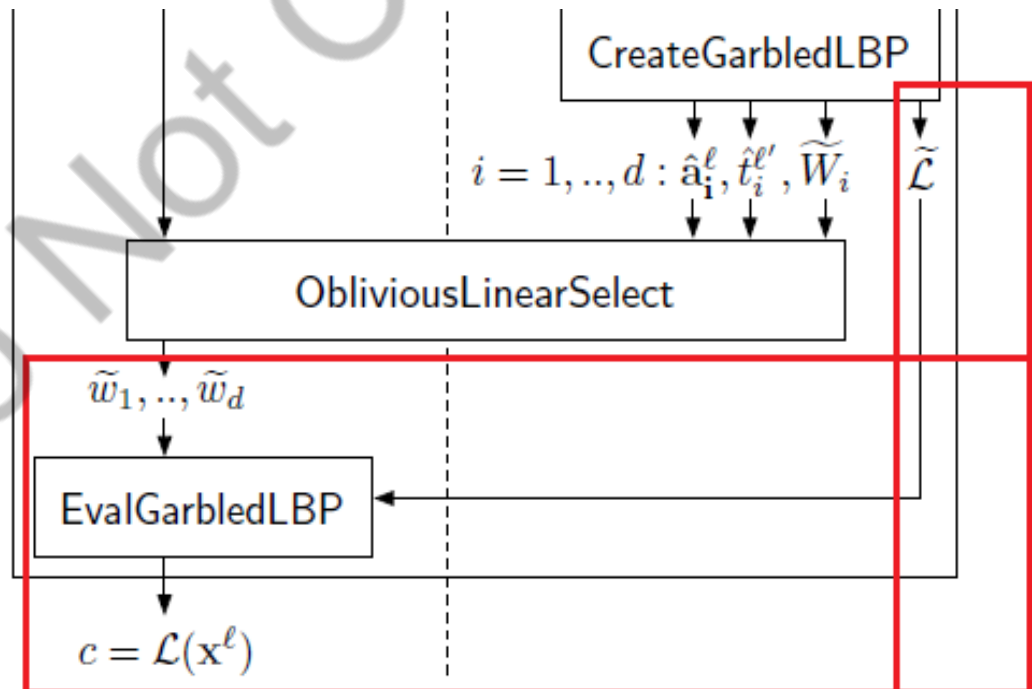
A protocol for secure evaluation of private LBPs - SecureEvalPrivateLBP

- Phase II: ObliviousLinearSelect
 - S blinds the encrypted value in order to hide the encrypted plaintext from C
 - protocol makes sure that S and C should not learn anything about the plaintexts.



A protocol for secure evaluation of private LBPs - SecureEvalPrivateLBP

- Phase III: EvalGarbledLBP
 - This stage takes the garbled values as input and produces the classification label as output



- ❑ Performance is a critical criteria for privacy preserving protocols
- ❑ Efficiency depends on the offline and online computations
- ❑ The garbled circuits can be pre computed on the server
- ❑ The exchange of garbled values happens in an online phase which uses oblivious transfer
- ❑ The evaluation of the Garbled Circuit happens on the client

Performance Improvements of Barni et al.

- ❑ Point and permute
 - Circuit evaluator only needs to decrypt a single ciphertext per garbled gate
- ❑ Incorporate classification nodes into decision nodes
 - Reduces size of LBP and number of Oblivious Transfers by the number of classification nodes
- ❑ Packing
 - Packing multiple ciphertexts into one and thus reducing number of decryptions and communication complexity

Performance Improvements of Barni et al.

□ TinyLBPs

- Constructing the LBP as single Yao gate with d inputs
- only feasible for small d because the size of the LBP grow exponentially in d

□ Key-offsets

Performance Comparison

- ❑ The improved hybrid version of Barni et al. reduces the number of Garbled Circuit and Oblivious Transfer and Homomorphic Encryption computations
- ❑ The reduction results are shown in the following table:

Oblivious Selection Protocol	Private Function	Moves	Asymptotic Communication Complexity		
			GC	OT	HE
[BPSW07]	BP	OT+2	$12zl(t + \kappa)$	OT_t^{zl}	$(n + z)2T$
[BFK ⁺ 09] Hybrid	LBP	OT+2	$12dl't$	$OT_t^{dl'}$	$(n + \frac{l'}{T-\kappa}d)2T$

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References

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Any Questions?