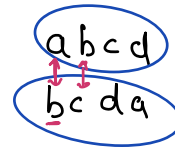
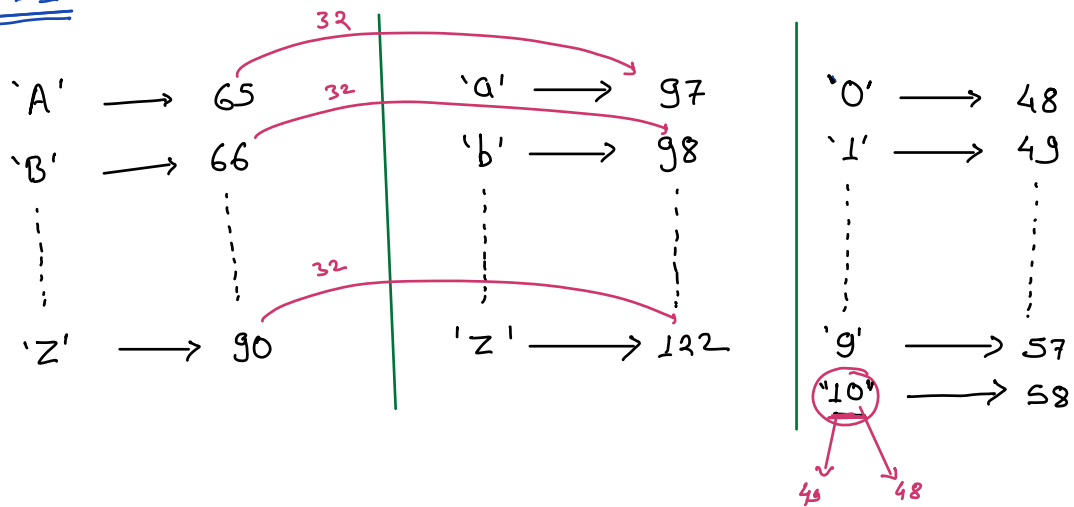


Strings : ~~Group of char~~
 Sequence of char ✓
 Array/List of char ✓



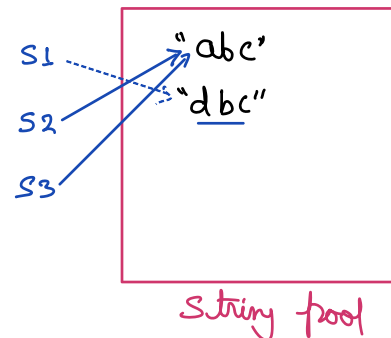
ASCII



Immutable Strings (Java / Python & other languages)

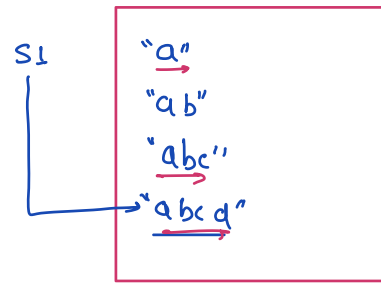
```
String s1 = "abc";
// s1[0] = 'd'
// s1 → "dbc"
```

```
String s2 = "abc";
String s3 = "abc";
s3 = "dbc";
```



Append Char to immutable string

String s1 = "a";
↑
s1 = s1 + "b"; → ~~O(1)~~ O(N)
s1 = s1 + "c"; → ~~O(1)~~ O(N)
s1 = s1 + "d"; → ~~O(1)~~ O(N)



If we append N times

TC : $O(N^2)$

SC : $O(N^2)$

Garbage Collector cleans out useless strings from memory.

String Builder

Java & C#

String Builder sb = new String Builder ();
sb.append('c');
sb.append("abc");
sb.toString() → Java String.

Q Given a string s . Toggle the case of every character.

Upper Case \longrightarrow Lower case

Lower case \longrightarrow Upper case

$s: aBcAE d \longrightarrow AbCa eD$

[Without using any inbuilt method]

toggle(s) {

for($i=0$; $i < s.size()$; $i++$) {

// if $s[i]$ is lower case

if ($s[i] \geq 'a' \ \&\& \ s[i] \leq 'z'$) {

$s[i] -= 32;$

abs('a' - 'A');

}

else if ($s[i] \geq 'A' \ \&\& \ s[i] \leq 'Z'$) {

$s[i] += 32;$

}

}

$s[i] = s[i] \wedge \text{diff}$
(32)

'a' \rightarrow 97

'b' \rightarrow 98

...

'z' \rightarrow 122

7	6	5	4	3	2	1	0
	↓	↓					
	64	32					

$64 + 32 \rightarrow 96$

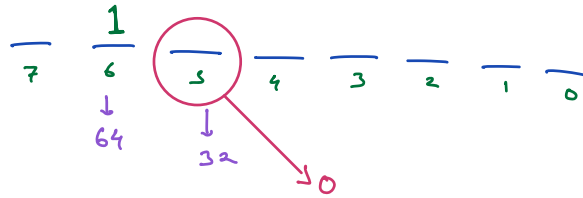
5th bit must be set for $[a, z]$

'A' \rightarrow 65

'B' \rightarrow 66

\vdots

'Z' \rightarrow 90



$$64 + 32 \rightarrow 96$$

5th bit must be unset for [A, Z]

diff ('a' 'A') \Rightarrow 32

'z' \rightarrow <u>122</u>	<u>0</u>	<u>1</u>	<u>1</u>	<u>1</u>	<u>1</u>	<u>0</u>	<u>1</u>	<u>0</u>
'Z' \rightarrow <u>90</u>	<u>0</u>	<u>1</u>	<u>0</u>	<u>1</u>	<u>1</u>	<u>0</u>	<u>1</u>	<u>0</u>

$$'z' - 'Z' \Rightarrow 32$$

In every pair of small & capital char
only 5th bit will be diff.

$$S[i] = S[i] \wedge 32;$$

$$\frac{S[i] \wedge (1 \ll 5)}{2^5}$$

Q Given a string of lower case char. Sort it in dict order.

s: d a b a e c b

a a b b d d e

Use library sorting method $\rightarrow O(N \log N)$

Lower case alphabets $\rightarrow [a-z] \Rightarrow 26$ unique char.

d a b a e c b

d: 2

a: 2

b: 2

e: 1

HashMap/Dict

for a \rightarrow z

a a b b d d e

int count[26] = {0}

Char

index

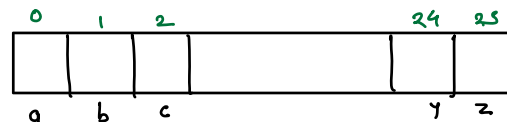
a \rightarrow 97 $\xrightarrow{-97}$ 0

b \rightarrow 98 $\xrightarrow{-97}$ 1

c \rightarrow 99 $\xrightarrow{-97}$ 2

\vdots

z \rightarrow 122 $\xrightarrow{-97}$ 25



S[i] - 'a'

// Count freq.

• `int Count[26] = {0};`

`for (i=0; i<N; i++){`

`// index for S[i] → S[i] - 'a'`

`Count[S[i] - 'a'] ++;`

`}`



// Build the string as per count, in dict order.

`K = 0;`

`for (i=0; i<26; i++){`

`// Count[i] stores freq of i + 'a'`

`for (j=0; j<Count[i]; j++){`

`S[K] = char(i + 'a');`

`K++;`

`}`

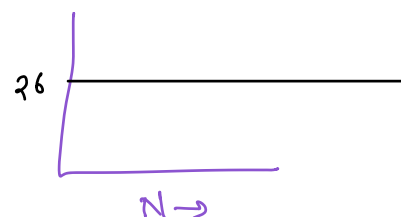
`}`

i	j	# iterations
0	[0, 1a]	a
1	[0, 1b]	b
2	[0, 1c]	c
⋮	⋮	⋮
25	[0, 1z]	z
		N

TC : $O(N)$ + $O(N)$ ⇒ $O(N)$
 Count freq Build the string

SC : $O(1)$ → $O(N)$ (String Build)
 (Extra)
 LC chr 26
 (LC + uc) 52
 All char 256
 if string is mutable

Counting Sort



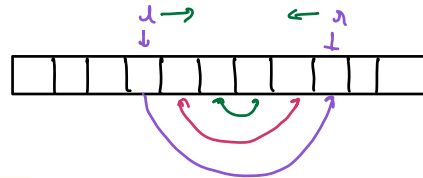
Q Given a string s & two indices l & r .

Reverse the sub-string from l to r

a b d e a g f

a b g a e d f

$l: 2, r: 5$



```
reverse(s, l, r) {
    while (l < r) {
        swap(s[l], s[r]);
        l++;
        r--;
    }
}
```

TC: $O(N)$

SC: $O(1) \rightarrow$ Mutable Strg

$O(N) \rightarrow$ Immutable Strg.

Amazon

Q Given a character array storing a sentence.

Reverse it word-by-word

* No extra space is allowed

* Every word is separated by a single white space (" ")

* No inbuilt method (reverse(), split())

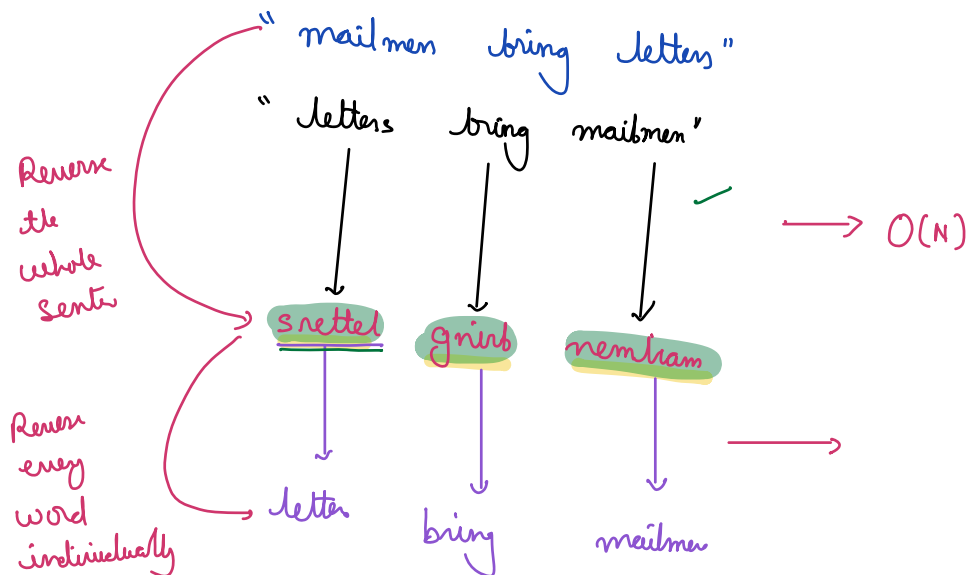
here - is - a - boy



boy a is here

"Are you as clever as I am"

"am I as clever as you Are"



TC : $O(N) + O(N)$

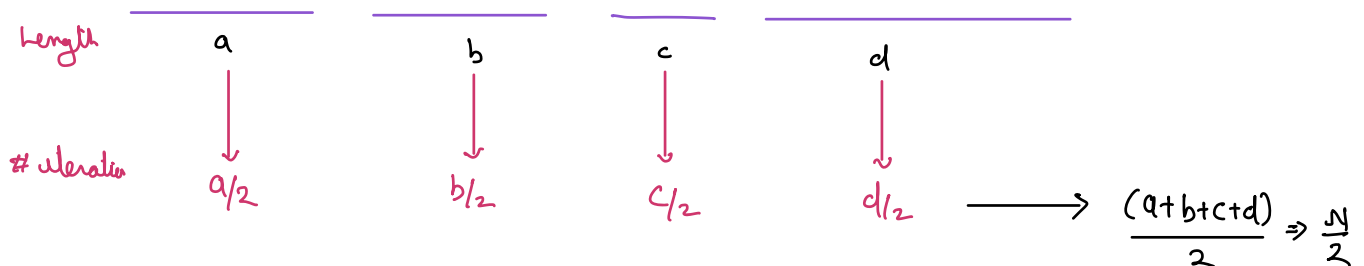
Reverse whole sentence

Revers all words individually

SC : $O(1)$

Tricky

while ($S[l] != ' ' \ \&\& \ l < N$) ?



Amazon
Direct-2
Ola

Q Given a string of size N . (lowercase alphabets)
Return the length of longest palindromic substring.

Palindrome

↓ ↓
→ MOM ←
→ MADAM ←
→ MALAYALAM ←
↑ ↑ ↑ ↑ ↑

s: a b a c a b

⇒ 5

aba → 3
acc → 3
bacab → 5

s: a b c a l m

⇒ 1

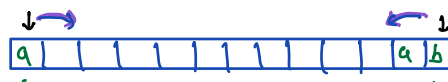
Approach 1

- Iterate over all substring : $O(N^2)$ •
- Check if substring is a palindrome : $O(N)$
- Maintain max length.

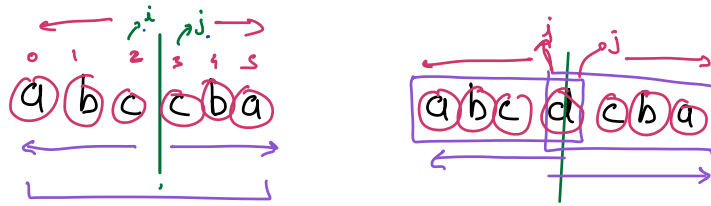
TC: $O(N^2 \times N) \rightarrow O(N^3)$

No of substrings
 $N + (N-1) + (N-2) + (N-3) + \dots + 1$
 $= \frac{N \times (N+1)}{2}$

Approach II



⇒ 2^n



//Return length of longest palindromic centered b/w i & j
 int getPalLength (s, i, j) {

}