

Fb | Google | so on ...

### Q1. 2 Sum

Given an array & a no. K. Return true if there exists a pair  $(i, j)$  in array s.t  $A[i] + A[j] = K$   $i \neq j$

$$\{ \begin{smallmatrix} 0 & 1 & 2 & 3 & 4 \\ 2, & 7, & 11, & 15, & 7 \end{smallmatrix} \}$$

$$K=18 \quad (1, 2) \\ \text{True}$$

$$K=14 \quad (1, 4) \\ \text{True}$$

$$K=20 \quad \text{False}$$

$$\{ \begin{smallmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \\ 8, & 9, & 1, & -2, & 4, & 5, & 11, & -6, & 7, & 5 \end{smallmatrix} \}$$

$$K=6 \rightarrow \text{True.}$$

$$\{ \begin{smallmatrix} 0 & 1 & 2 & 3 & 4 & 5 \\ 7, & 1, & 10, & -3, & 9, & 7 \end{smallmatrix} \}$$

$$K = -2 \quad \checkmark$$

$$K = 10 \quad \checkmark$$

$$K = 18 \quad \times$$

$$K = 14 \quad \checkmark$$

Brute Force

Create all the pairs.

```
for (int i=1; i<N; i++) {
```

0 1 2 3

```
    for (int j=0; j<i; j++) {
```

```
        if (i != j && A[i] + A[j] == k) {
```

return true

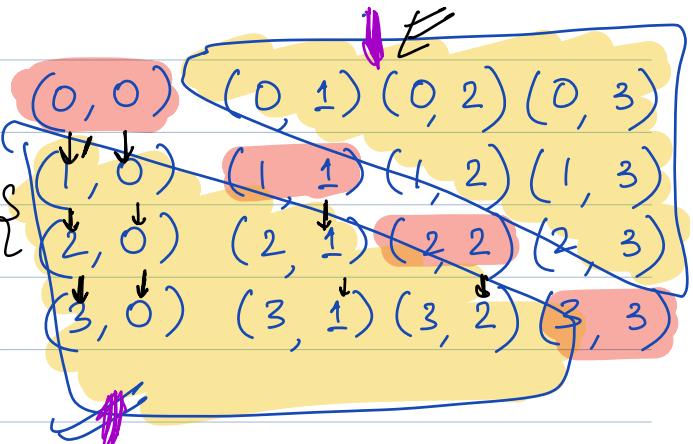
}

}

return false

TC  $\rightarrow O(N^2)$

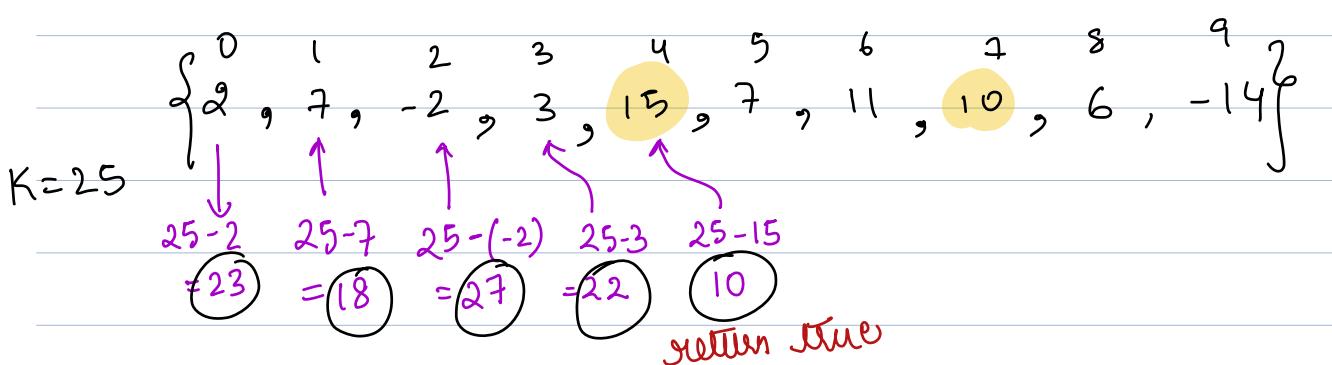
SC  $\rightarrow O(1)$



Observation

$$A[i] + A[j] = K$$

$$A[j] = K - A[i]$$



Check whether a value exists or not

↳ ~~HashMap / HashSet~~

(unique elements)

HashSet s

2	7	-2
3	15	6
11	10	
-4		

HashSet<int> s;

for( $i=0$ ;  $i < N$ ;  $i++$ ) {

|     s.insert(A[i])

} for( $i=0$ ;  $i < N$ ;  $i++$ ) {

|      $x = K - A[i]$

|     if ( $x$  exists in  $s$ ) {

|         return true

}

return false

Fails for this case

{ 4, 7, 5, 1 }

K=14

$$14 - 7 \\ = 7$$

4	7
5	1

HashMap → store frequency.

$\{$  0, 1, 2, 3, 4, 5, 6, 7, 8, 9  $\}$   
 ↓  
 2, 7, -2, 3, 15, 7, 11, 10, 6, -14  $\}$

4-2  
= 2

K = 4

HashMap<int, int> mp.

2	→	1
7	→	2
-2	→	1
3	→	1
15	→	1
11	→	1
10	→	1
6	→	1
-14	→	1

HashMap<int, int> mp ;

for(i=0; i<N; i++) {  
 if (A[i] exists in mp) {  
 update (A[i], mp[A[i]]++)

}

else {

    insert (A[i], 1)

}

}

frequency

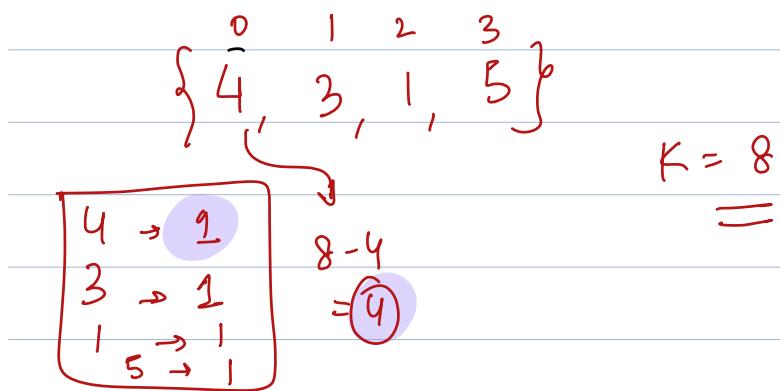
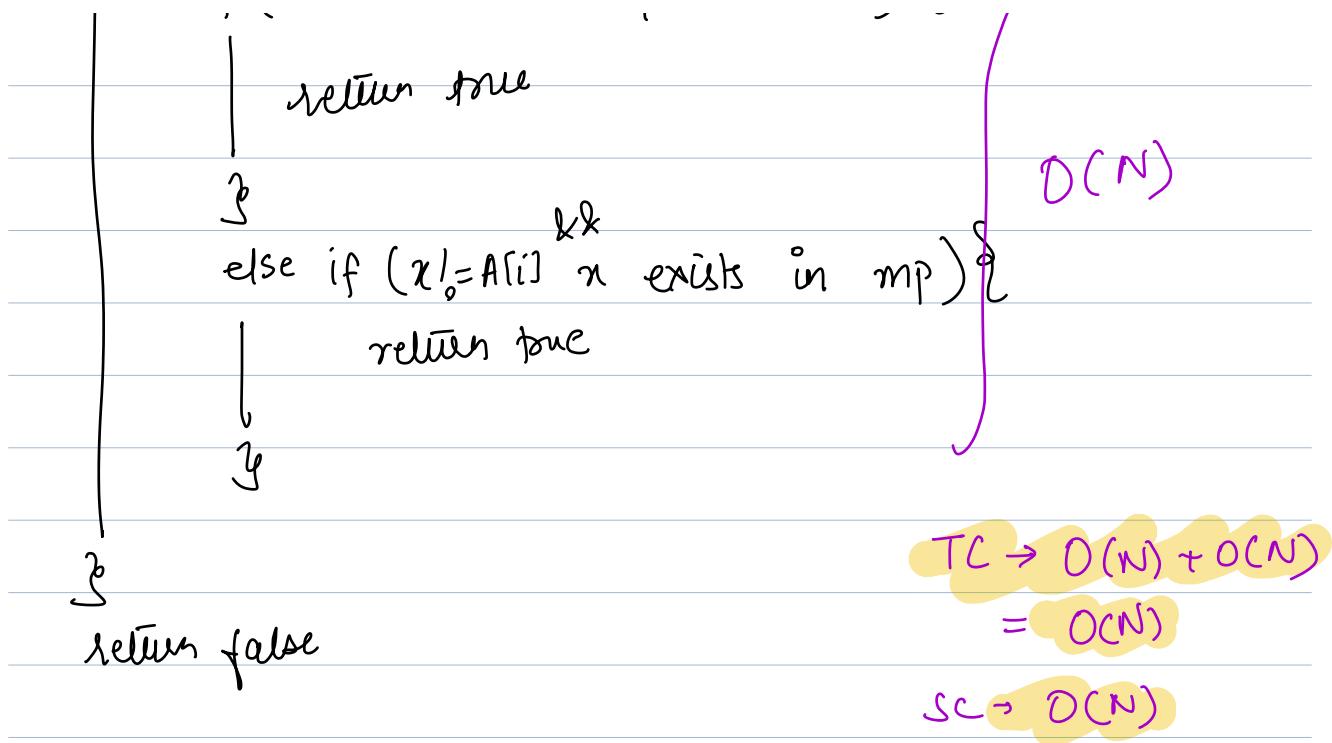
O(N)

for( i=0; i<N ; i++) {

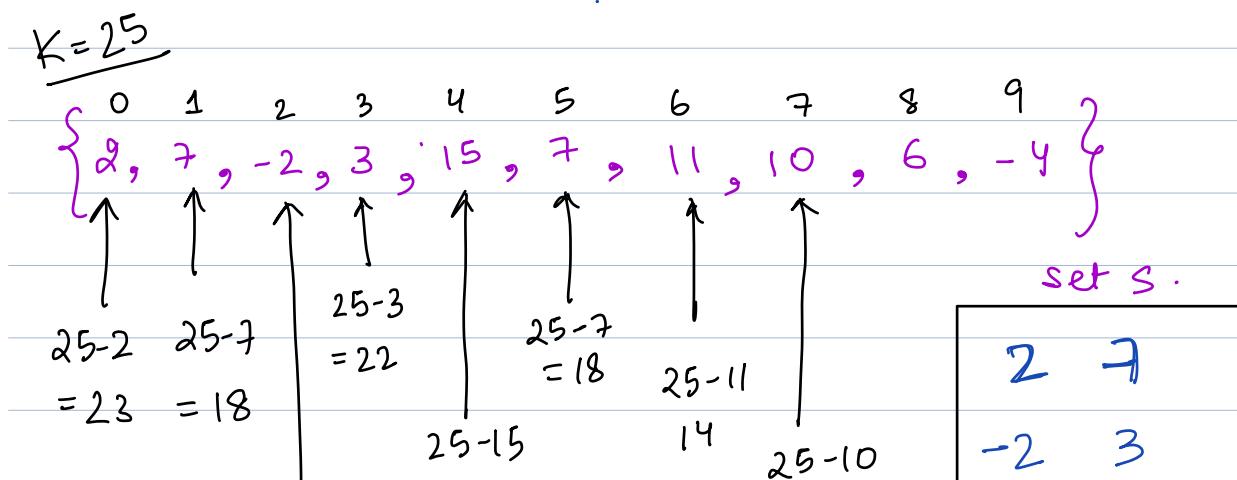
    x = K - A[i]

    if (x == A[i] && mp[x] >= 2 ) {

}

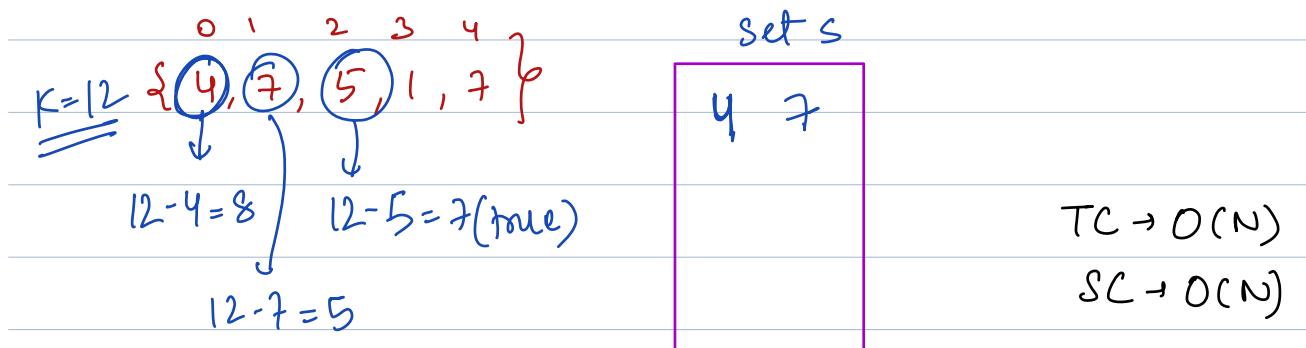


Single Iteration using set



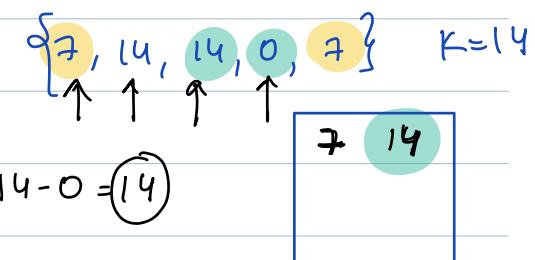
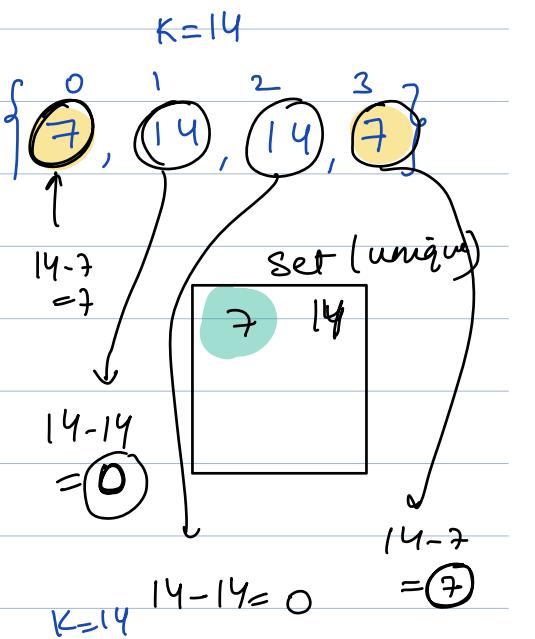
$$25 - \frac{1}{(-2)} = 10$$

= 15  
return true | 15 11



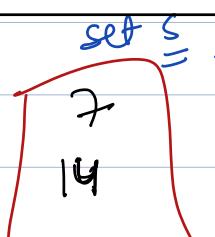
```

HashSet<int> s;
for (i = 0; i < N; i++) {
  n = K - A[i]
  if (n is present in s) {
    return true
  }
  else {
    insert (A[i]) into the s
  }
}
  
```



$A = [7, 14]$

$K=14$



$$\begin{array}{r} \overline{\overline{J}} \\ \downarrow \\ \text{false} \end{array}$$

$$\begin{array}{r} 14-7 \\ = 7 \end{array}$$

$$\begin{array}{r} 14-14 \\ = 0 \end{array}$$

✓

$$A = \{ \begin{array}{c} 0 \\ 7, 14, 5, 6, 9 \\ 1 \\ 2 \\ 3 \\ 4 \end{array} \}$$

$$\begin{array}{l} K=12 \\ \curvearrowright \text{true} \end{array}$$

$$\begin{array}{l} K=18 \\ \curvearrowright \end{array}$$

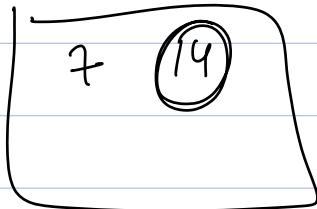
$$\begin{array}{r} 14-0 \\ = 14 \end{array}$$

$$\{ \begin{array}{c} 0 \\ 7, 14, 0 \\ 1 \\ 2 \end{array} \}$$

$$K=14$$

→ True

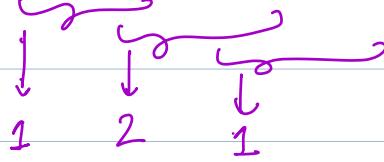
$$\begin{array}{r} 14-7 \\ 14-14 \\ = 7 \\ = 0 \end{array}$$



$$\begin{array}{r} 14+0 \\ = = \\ = 14 \end{array}$$

Ques. Given N array ele, calc the no. of distinct ele  
in every window of size K.

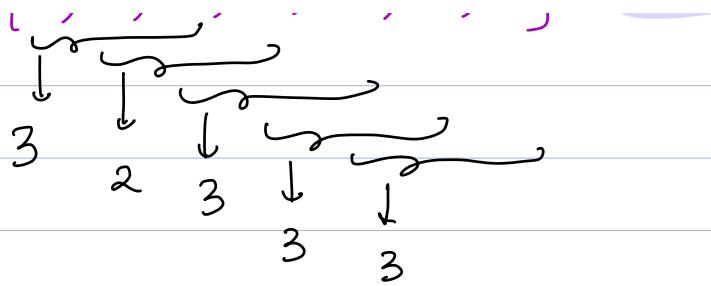
$$A : \{ \begin{array}{c} 1 \\ 1 \\ 2 \\ 2 \end{array} \} \quad K=2$$



result:  $\{1, 2, 1\}$

QUIZ

$$A : \{6, 3, 7, 3, 8, 6, 9\} \quad K=3$$



Result  $\{3, 2, 3, 3, 3\}$

$N$

$$K=1 \quad N$$

$$K=2 \quad N-1$$

$$K=3 \quad N-2$$

$\vdots$

$$K \quad N-K+1$$

$K=3$       start of first window      start of last window

$$1 \quad 0 \quad N-1$$

$$2 \quad 0 \quad N-2$$

$$3 \quad 0 \quad N-3$$

$$4 \quad 0 \quad N-4$$

$\vdots$

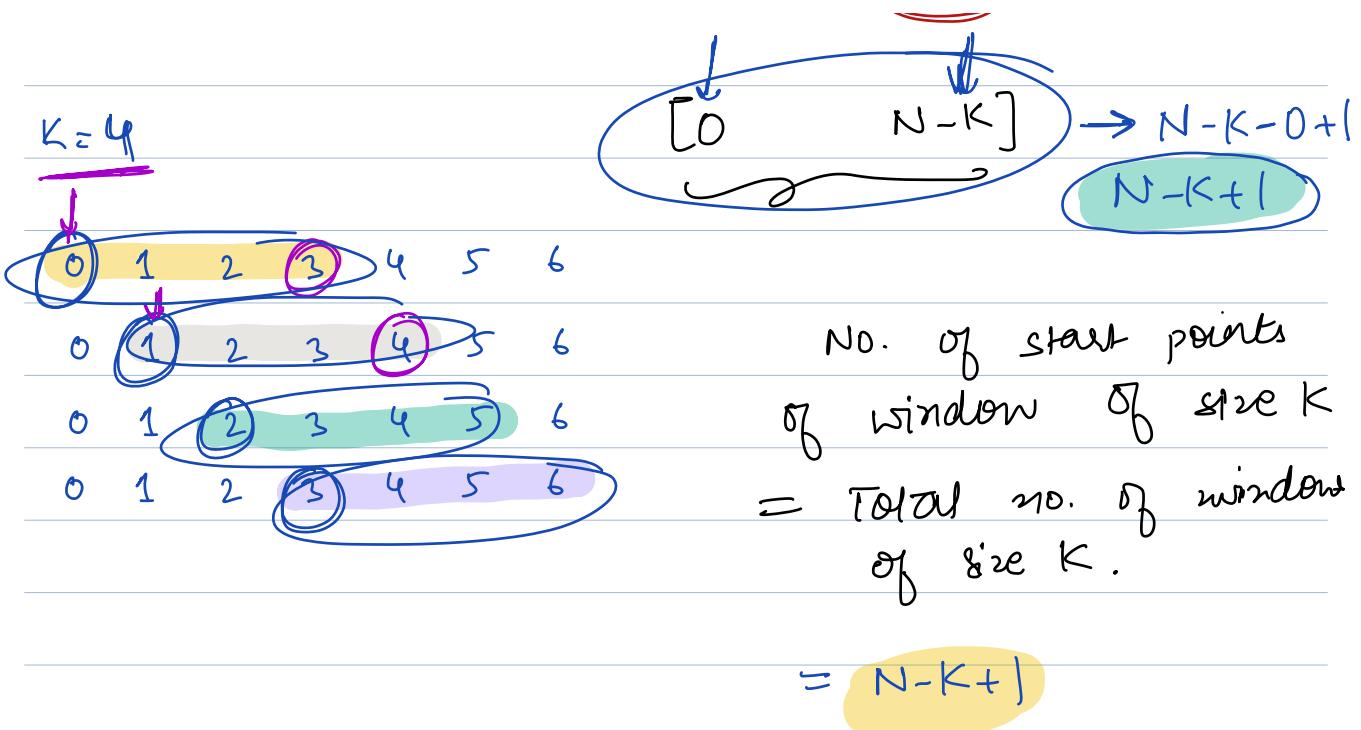
$\vdots$

$\vdots$

$K$

$0$

$N-K$



### Brute Force

Iterate over all windows of size K

```
list<int> ans;
for (i=0 ; i<=N-K ; i++) {
```

$j = i+K-1$

```
    HashSet<int> s;
```

```
    for (z = i ; z<=j ; z++) {
```

s.insert (A[z])

}

ans.push (s.size)

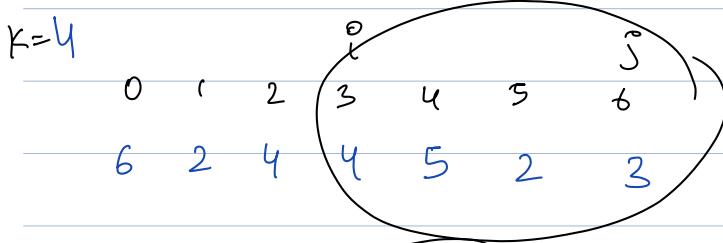
}

Iterations  
 $(N-K+1) (K)$

$K=1$

$$(N-x+y)(1) = N$$

$K=N$



ans  
 $\{3, 3, 3, 4\}$

$$(N-K+1) \cdot N = N$$

$$K \approx N/2$$

$$\left(N - \frac{N}{2} + 1\right) \left(\frac{N}{2}\right)$$

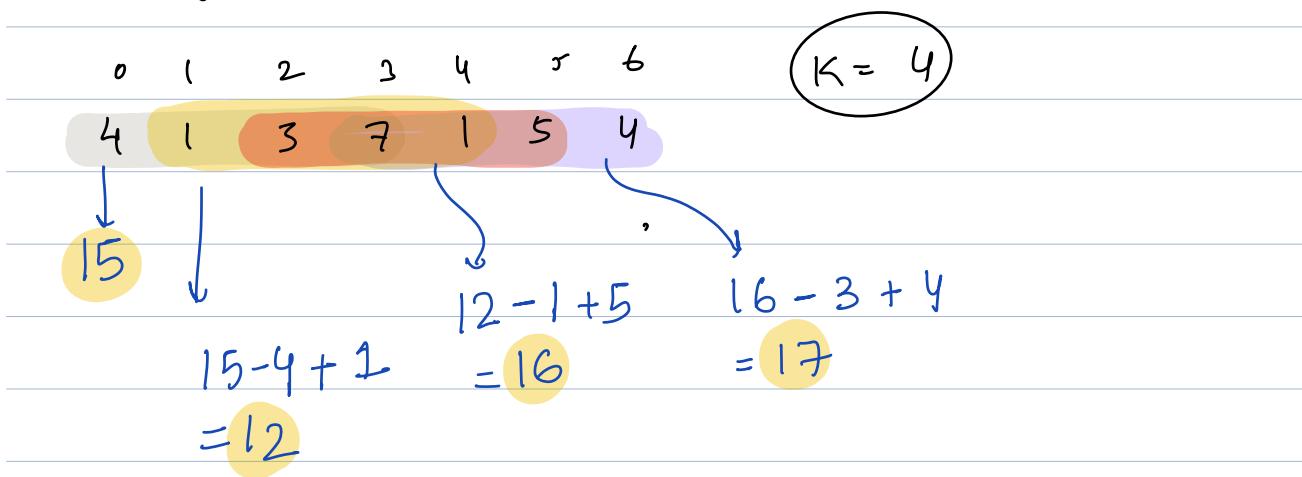
$$\left(\frac{N}{2} + 1\right) \left(\frac{N}{2}\right)$$

$$N^2$$

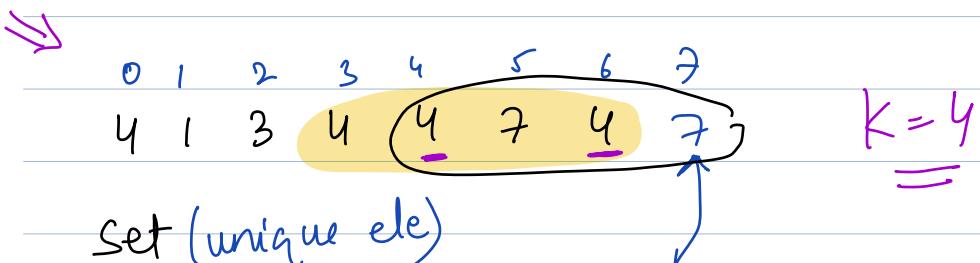
$$TC \rightarrow O(N^2)$$

$$SC \rightarrow O(N)$$

sliding window



Can we apply some logic in this question?



?

L

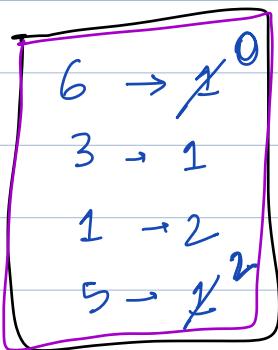
{ 3, 3, 3, 2, 1 }

(set)!

HashMap with frequency.

0 1 2 3 4 5 6 7 8 9 10 11  
6, 3, 1, 1, 5, 5, 5, 6, 9, 14, 9, 9

K = 5



{ 4,

sliding window

① Process window of size K  
⇒ Create the freq map

② Process the rest of the window  
⇒ We will add last ele of next window to map

& remove first ele of prev window from map.

HashMap<int> mp;

for( $i=0$ ;  $i < K$ ;  $i++$ ) {

    if ( $A[i]$  is present) {

$mp[A[i]]++$

} frequency

    else {

        insert ( $A[i]$ , 1)

}

list.push (mp.size())

for( $i=0$ ;  $i \leq n-K-1$ ;  $i++$ ) {

$x = i+K$

$mp[A[i]]--$

    if ( $mp[A[i]] == 0$ ) {

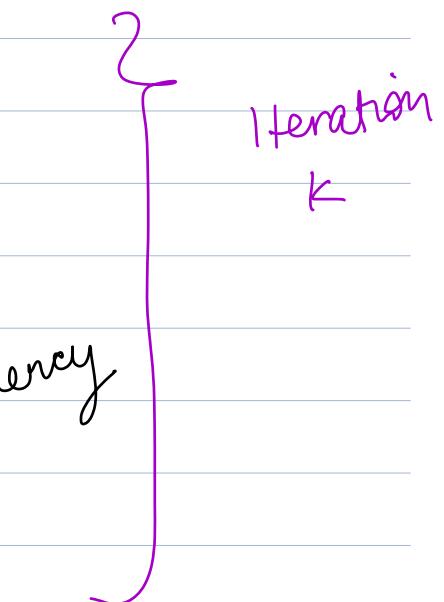
        remove ( $A[i]$  from mp)

}

    else if ( $A[x]$  is present in mp) {

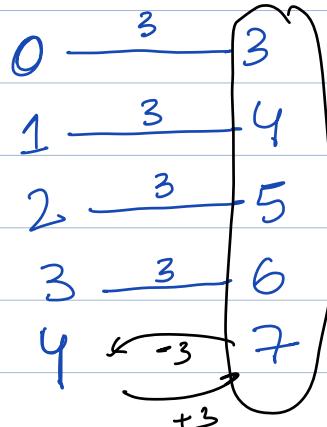
$mp[A[x]]++$

}



$K=3$

Remove



Add

Iteration:  
 $n-K$

else insert ( A[u] , 1 )

list.push ( mp.size() )

$\xrightarrow{k}$   
return list

Total:  $\frac{K+N-K}{N}$

$O(N)$

$K=4$

