



Experiment - 6

Student Name: Mankaran Singh Tandon

UID: 23BCS10204

Branch: BE-CSE

Section/Group: KRG-2B

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Subject Name: Design and Analysis of Algorithms

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Aim: Develop a program and analyse complexity to implement subset-sum problem using Dynamic Programming.

Objective: Objective is to implement subset-sum problem using Dynamic programming.

Input/Apparatus Used: Standard sum value is taken in order to find subsets of that sum using dynamic programming.

Procedure/Algorithm:

Procedure:

So we will create a 2D array of size $(arr.size() + 1) * (target + 1)$ of type boolean. The state $DP[i][j]$ will be true if there exists a subset of elements from $A[0 \dots i]$ with sum value = „j“. The approach for the problem is:

if $(A[i-1] > j)$ $DP[i][j]$

= $DP[i-1][j]$ else

$DP[i][j] = DP[i-1][j] \text{ OR } DP[i-1][j-A[i-1]]$

1. This means that if current element has value greater than „current sum value“ we will copy the answer for previous cases

2. And if the current sum value is greater than the „ith“ element we will see if any of previous states have already experienced the sum=“j” OR any previous states experienced a value „j – A[i]“ which will solve our purpose.

The below simulation will clarify the above approach: set[]={3, 4, 5, 2}

target=6

0 1 2 3 4 5 6

Dynamic Programming – Subset Sum Problem

Given a set of positive integers, and a value *sum S*, find out if there exist a subset in array whose sum is equal to given *sum S*.

Example:

int[] A = { 3, 2, 7, 1 }, S = 6

Output: True, subset is (3, 2, 1)

Recursive Approach:

For every element in the array has two options, either we will include that element in subset or we don't include it.

§ So if we take example as int[] A = { 3, 2, 7, 1 }, S = 6

§ If we consider another int array with the same size as A.

§ If we include the element in subset we will put 1 in that particular index else put 0.

§ So we need to make every possible subsets and check if any of the subset makes the sum as S.

Time Complexity: $O(2^n)$.

Approach: Dynamic Programming

Base Cases:

§ If no elements in the set then we can't make any subset except for 0.

§ If sum needed is 0 then by returning the empty subset we can make the subset with sum 0.

Given – Set = `arrA[]`, Size = `n`, sum = `S`

§ Now for every element in the set we have 2 options, either we include it or exclude it.

§ for any `i`th element-

§ If include it $\Rightarrow S = S - arrA[i]$, $n = n - 1$

§ If exclude it $\Rightarrow S$, $n = n - 1$.

Recursive Equation:

Base Cases:

`SubsetSum (arrA, n, S) = false`, if `sum > 0` and `n == 0` `SubsetSum (arrA, n, S) = true`, if `sum == 0` (return empty set)

Rest Cases

`SubsetSum (arrA, n, S) = SubsetSum (arrA, n-1, S) || SubsetSum (arrA, n-1, S-arrA[n-1])`

How to track the elements.

§ Start from the bottom-right corner and backtrack and check from the True is coming.

§ If value in the cell above is false that means current cell has become true after including the current element. So include the current element and check for the `sum = sum - current elements`.

Time Complexity: $O(\text{sum} * n)$, where `sum` is the „target sum“ and „`n`“ is the size of array.

5. Code and Output:

```
#include <bits/stdc++.h>

using namespace std;
```



```
// Recursive approach bool
subsetSumRecursive(vector<int>& arr, int n, int sum) {
    if (sum == 0) return true; if (n == 0 && sum != 0) return false; if (arr[n - 1] >
    sum) return subsetSumRecursive(arr, n - 1, sum); return subsetSumRecursive(arr,
    n - 1, sum) || subsetSumRecursive(arr, n - 1, sum -
    arr[n - 1]);
}

// Dynamic Programming approach (Bottom-Up) bool
subsetSumDP(vector<int>& arr, int sum, vector<vector<bool>>& dp) {
    int n = arr.size(); for (int
    i = 0; i <= n; i++)
        dp[i][0] = true;
    for (int j = 1; j <= sum; j++)
        dp[0][j] = false;
    for (int i = 1; i <= n; i++) {
        for (int j = 1; j <= sum; j++)
            { if (arr[i - 1] > j)
                dp[i][j] = dp[i - 1][j];
            else
                dp[i][j] = dp[i - 1][j] || dp[i - 1][j - arr[i - 1]];
            } } return dp[n][sum];
}

// Track the subset elements from DP table vector<int> getSubsetElements(vector<int>&
arr, vector<vector<bool>>& dp, int sum) { vector<int> subset; int n = arr.size();
    int i = n, j = sum; while
    (i > 0 && j > 0) {
```

```
        if (dp[i][j] && !dp[i - 1][j])
            { subset.push_back(arr[i -
              1]); j -= arr[i - 1];
            }
        i--;
    } reverse(subset.begin(),
subset.end()); return subset;
}

int main() {
    vector<int> arr = {3, 2, 7, 1};
    int sum = 6;
    int n = arr.size();
    cout << "Subset Sum Problem using Dynamic Programming\n";
    cout << "Set = { "; for (int x : arr) cout << x << " "; cout << "},
    Target Sum = " << sum << "\n\n";
    // Recursive method bool recResult =
    subsetSumRecursive(arr, n, sum);
    cout << "Recursive Approach: " << (recResult ? "Subset Exists" : "Subset Does Not
Exist") << "\n"; cout << "Time
    Complexity: O(2^n)\n\n";
    // Dynamic Programming method vector<vector<bool>> dp(n
    + 1, vector<bool>(sum + 1, false)); bool dpResult =
    subsetSumDP(arr, sum, dp);
    cout << "Dynamic Programming Approach: " << (dpResult ? "Subset Exists" : "Subset
Does Not Exist") << "\n"; cout << "Time
    Complexity: O(n * sum)\n"; cout << "Space
    Complexity: O(n * sum)\n\n";
```

```
if (dpResult) {  
    vector<int> subset = getSubsetElements(arr, dp, sum);  
    cout << "Subset with sum " << sum << " is: { "; for  
    (int x : subset) cout << x << " ";  
    cout << "}\n";  
} else { cout << "No subset found with the given sum.  
    \n";  
}  
return 0;  
}
```

Your Output

```
Subset Sum Problem using Dynamic Programming  
Set = { 3 2 7 1 }, Target Sum = 6
```

```
Recursive Approach: Subset Exists  
Time Complexity:  $O(2^n)$ 
```

```
Dynamic Programming Approach: Subset Exists  
Time Complexity:  $O(n * \text{sum})$   
Space Complexity:  $O(n * \text{sum})$ 
```

```
Subset with sum 6 is: { 3 2 1 }
```